

EUROMPI A 25

Oct. 2, 2025, Charlotte, NC, USA

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MPI SESSIONS TUTORIAL

The World Model:

MPI_Init();

[MPI_COMM_WORLD]

MPI_Finalize();

The Sessions Model:

MPI_Session_init();

MPI_Group_from_session_pset();
MPI Comm create from group();

[On-demand communicators]

MPI_Session_finalize();



MPI SESSION TIMELINE



2019

Sep. 25

Hjelm N. et al.: MPI Sessions: Evaluation of an Implementation in Open MPI

2025

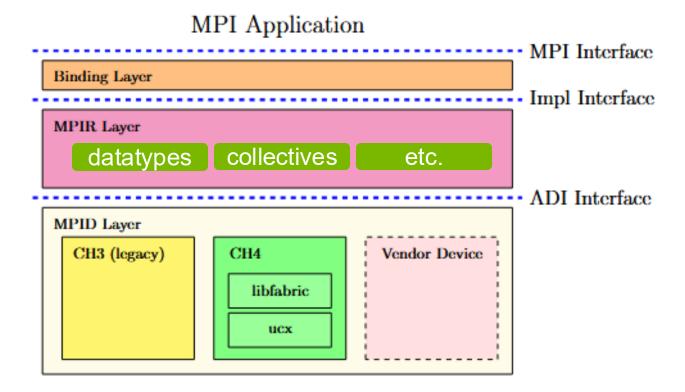
Oct. 2

Zhou H. et al.: Implementing True MPI Sessions and Evaluating MPI Initialization Scalability





MPICH ARCHITECTURE





PROCESS MANAGEMENT INTERFACE (PMI)

PMI Proxy

1. Launching MPI processes

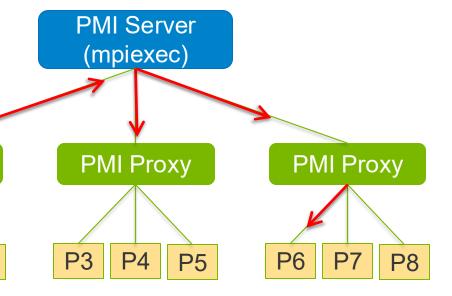
2. Initial process identification

3. Bootstrap communication

1. PMI Put

2. PMI_Barrier

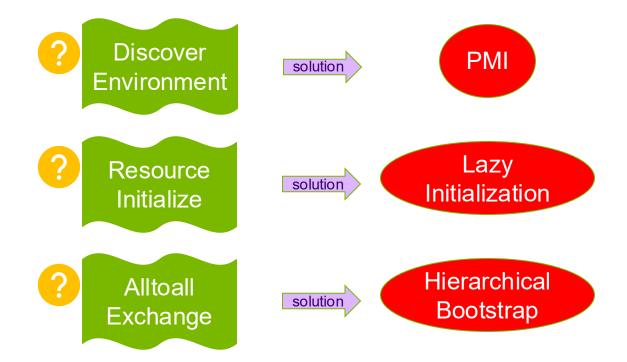
3. PMI_Get





EFFICIENT MPI STARTUP

Bottlenecks In A Large-scale MPI Startup







EFFICIENT ALL-TO-ALL ADDRESS EXCHANGE

Hierarchical bootstrapping

Step 1:

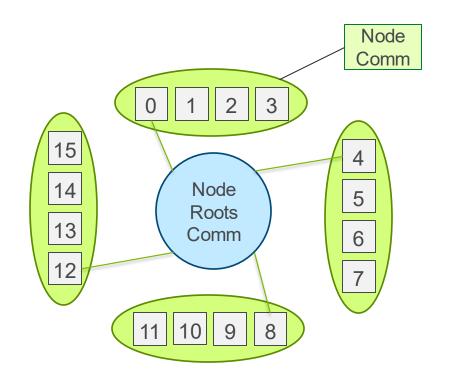
Use PMI to establish node-roots communicator

Step 2:

Use shared memory to establish intra-node communicator

Step 3:

Perform hierarchical allgather using fast MPI communication



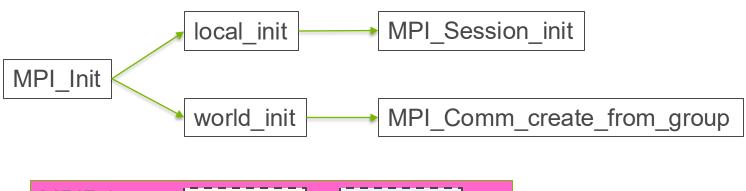


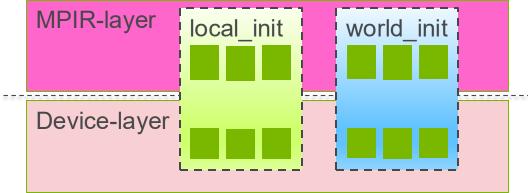






Separating Local and Collective Initialization

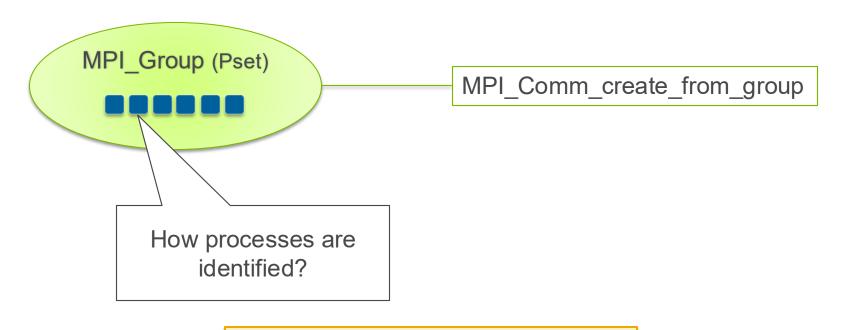








Communicator-Independent Process IDs



world_index:world_rank





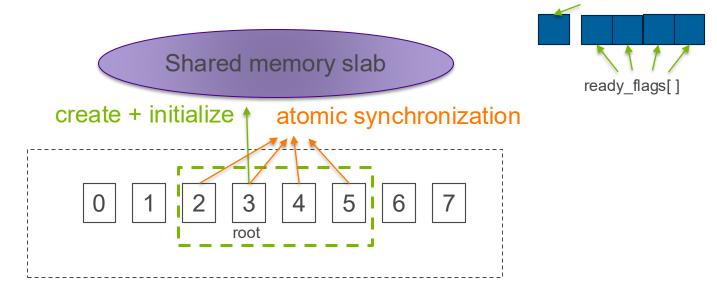
Group-level address exchange via PMI

PMI-1: PMI_Barrier extension (PMI v1.2)

- PMI-2: deprecate
- PMI-x: PMIx Fence



Group-level shared memory initialization



root ready

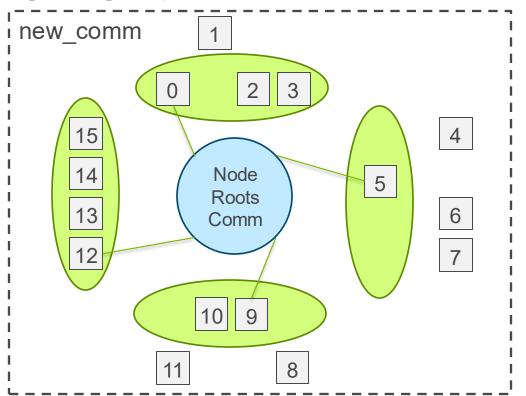
We redesigned shared memory init to an async atomic semantics





Update Hierarchical Bootstraping in a group context

- 1. Node-roots-comm via PMI
- Node-comm via atomic SHM
- 3. Bootstrap new_comm via hierarchical collectives











EXPERIMENTAL EVALUATION

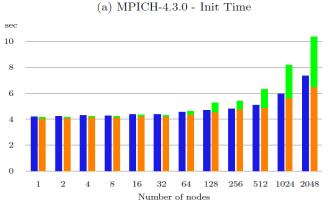
MPICH

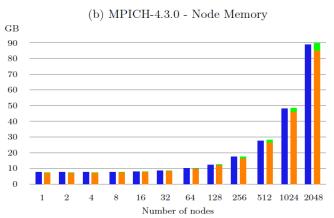
Expectation:

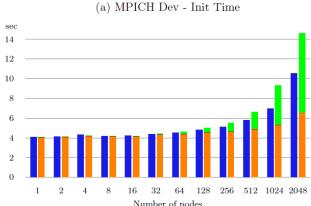
- Equivalency between the world and sessions model.
- Flat local initialization.
- No performance degradation before and after supporting true sessions.

Results:

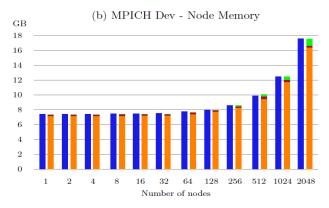
- Internal world comm slightly more efficient.
- Local initialization get slower and takes more memory as num of nodes increase.
- Slight performance degradation







■ MPI Init ■ Session Init ■ Self Comm ■ World Comm

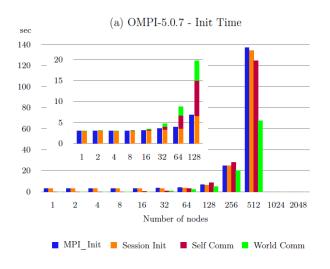


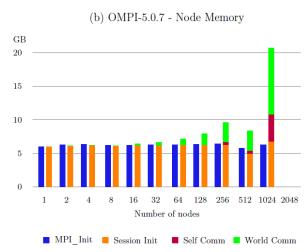






EXPERIMENTAL EVALUATION Open MPI





Aurora, PPN = 96

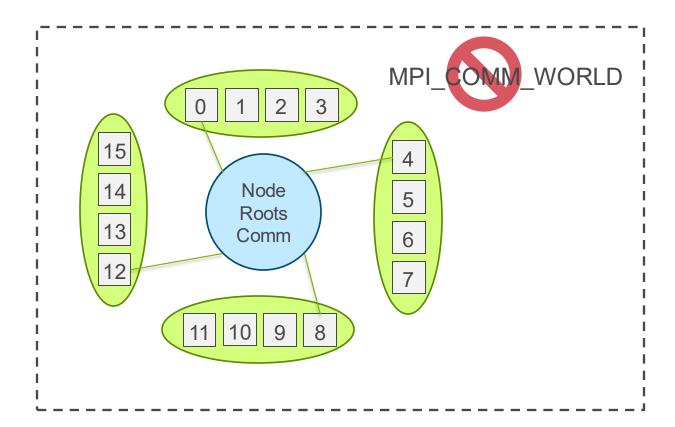




EXPERIMENTAL EVALUATION

Sparse World

- Dense World
 - 192 internode connections
- Sparse World
 - 12 internode connections
 - Reduction by $\frac{1}{(PPN)^2}$

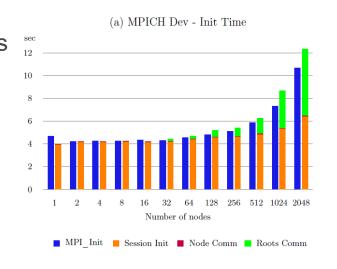


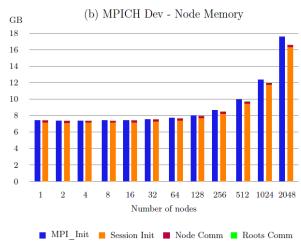


EXPERIMENTAL EVALUATION

Sparse World

- Not much savings over hierarchical bootstrapping
- 2. More significant memory savings
- 3. Require userlayer hierarchical code









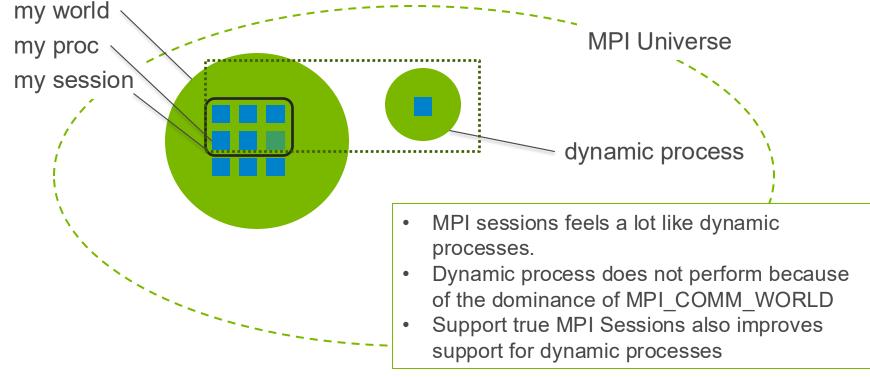






DISCUSSIONS

Lesson 1: MPI Sessions vs. MPI Dynamic Processes



DISCUSSIONS

Lesson 2: The Importance of Implementation

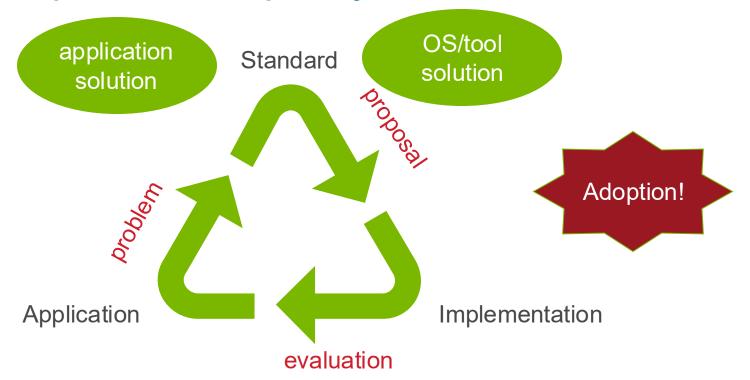






DISCUSSIONS

Lesson 3: The Importance of Complete Cycle







SUMMARY

- MPICH supports MPI Sessions since v4.0 in 2021. However, it relies on an internal comm world to bootstrap communicators in the sessions model.
- We reimplemented in MPICH to support true sessions that does not depend on comm_world.
- Our evaluations show no significant scaling advantage between world model and sessions model if the world communicator is still constructed.
- We show improved initialization time and memory consumption when sparse communicators are used instead.
- Supporting true MPI sessions greatly improves MPICH's dynamic process support.







