Don't shoot down TLB shootdowns!

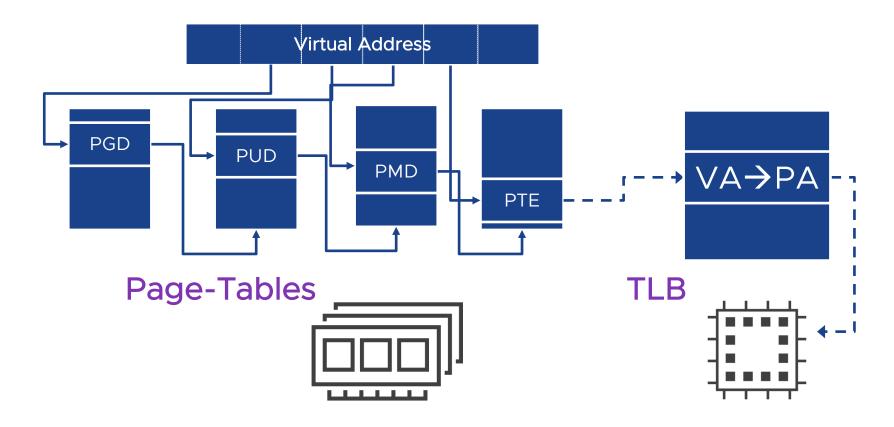
Nadav Amit, Amy Tai, Michael Wei

April 2020





Translation Lookaside Buffer (TLB)

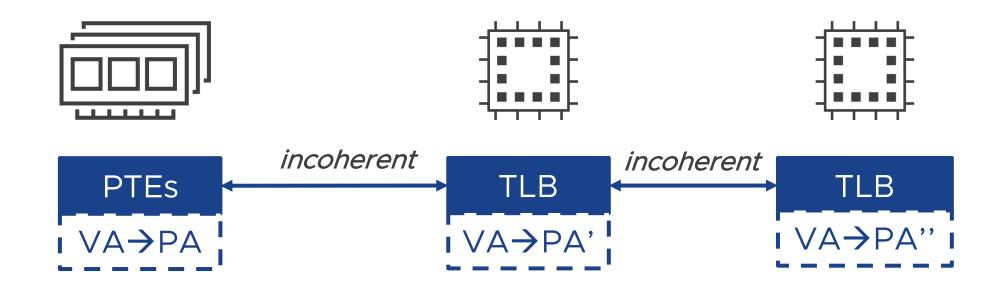


TLB = cache for virtual to physical address translations





TLB Coherency



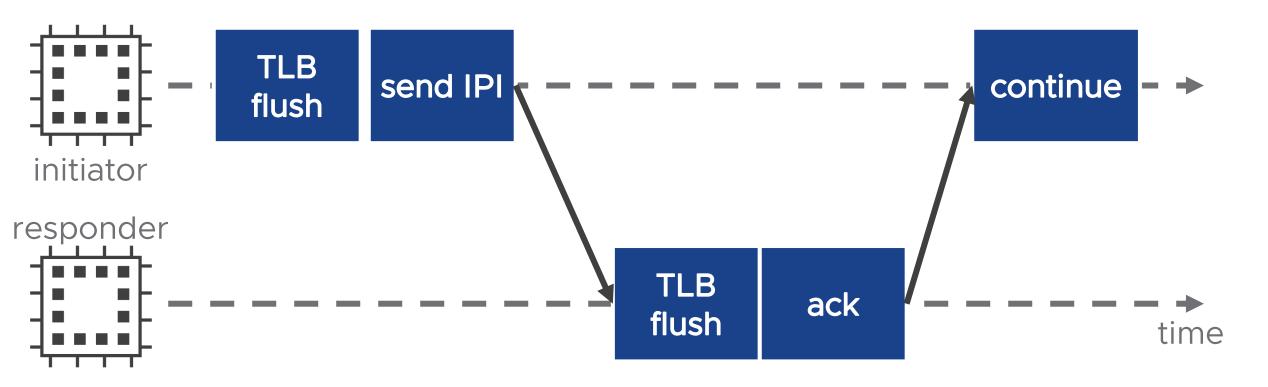
Hardware does not maintain TLBs coherent

The problem is left for software (OS)





TLB Shootdown (in Linux)







Challenge

TLB shootdowns are expensive.

How can we further optimize them?

This work focus on:

- Linux/x86 common lessons
- Userspace mappings common case

Lessons are relevant to other environments



Existing Solutions

Hardware based TLB invalidations

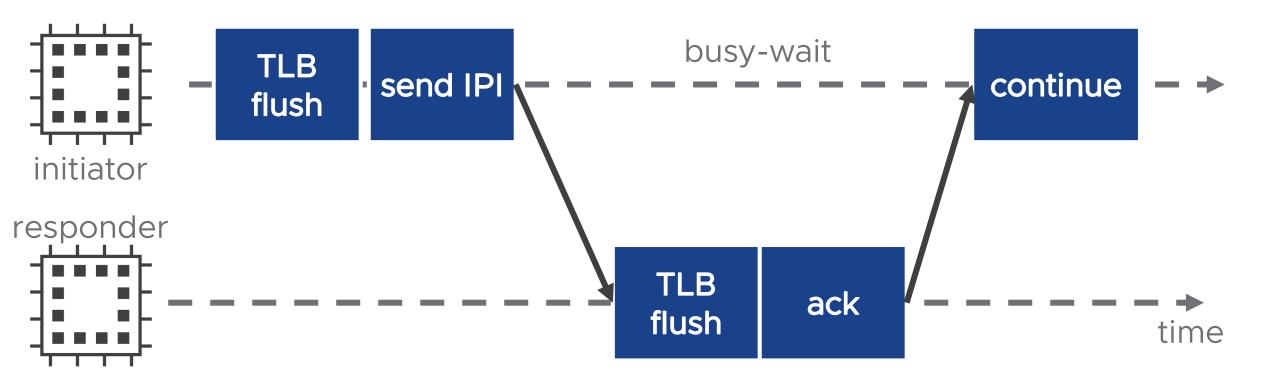
- Not available on all architectures
- Does not coexist (yet) with software techniques:
 - No selective target cores for TLB invalidation

Software solutions

- Replicating page-tables [RadixVM, Clements'13]
 - Can increase overhead with low-latency IPIs
- Aggressive batching [LATR, Kumar'18]
 - Breaks POSIX semantics



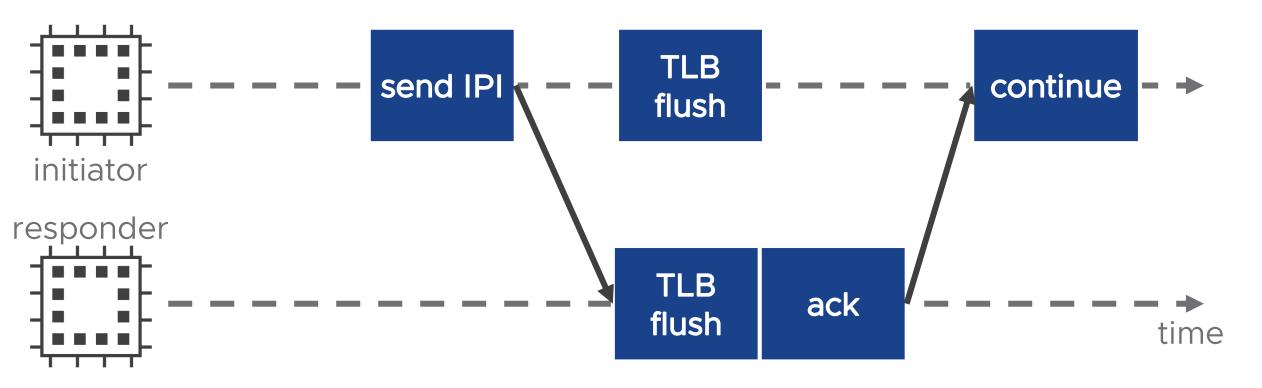
TLB Flushes in Linux and FreeBSD







Optimization 1: Concurrent Flushes (forgotten lesson)



RP3 TLB consistency algorithm [Rosenburg'89]



TLB Shootdown Responder

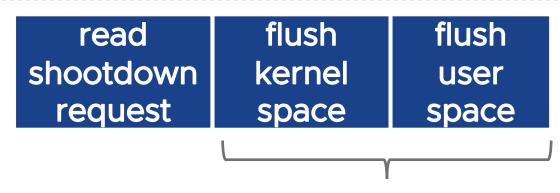
Entry

enter IRQ handler return to user

SMP

read SMP request ack SMP request

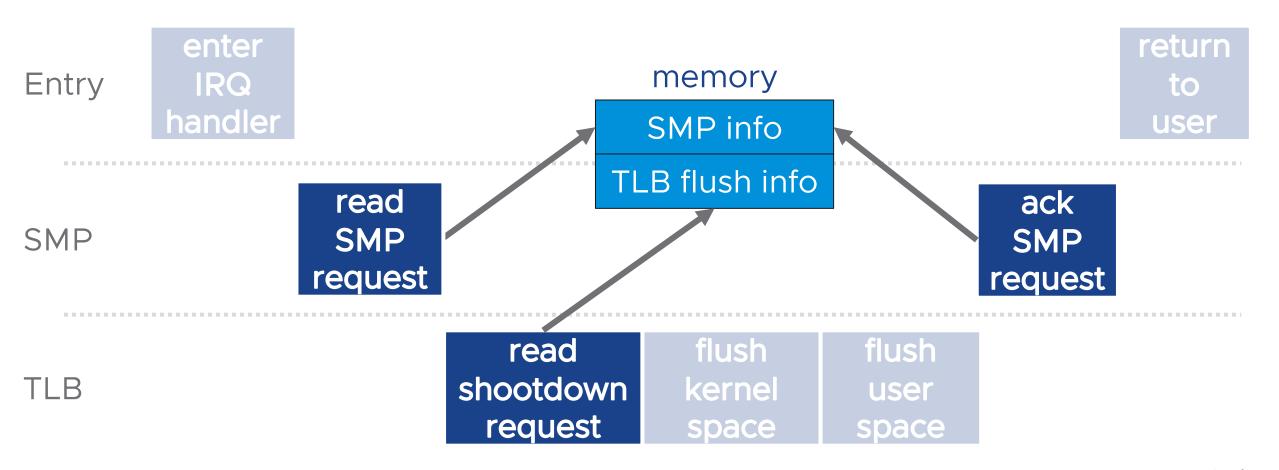
TLB





Page Table Isolation

Optimization 2: Cacheline Consolidation







Optimization 3: Early Acknowledgment

enter
IRQ
handler

return to user

SMP

read SMP request ack SMP request

TLB

read flush flush shootdown kernel user request table





Optimization 3: Early Acknowledgment

Entry

enter IRQ handler Safe: flush will happen

Better: Initiator is faster

return to user

SMP

read SMP request ack SMP request

TLB

read shootdown request flush flush kernel user table





Optimization 4: In-Context Flushes

Entry

enter IRQ handler return to user

SMP

read SMP request

ack SMP request

TLB

read shootdown request

flush kernel table flush user table





Optimization 4: In-Context Flushes

Entry

enter IRQ handler

- 1. Efficient
- 2. Better batching

return to user

SMP

read SMP request

ack SMP request

TLB

read shootdown request flush kernel table flush user table





In the Paper

Userspace-safe batching

Deferring TLB shootdowns while the kernel runs

Avoiding TLB flushes on Copy-on-Write

Special case we can optimize

TLB flushes in virtualization

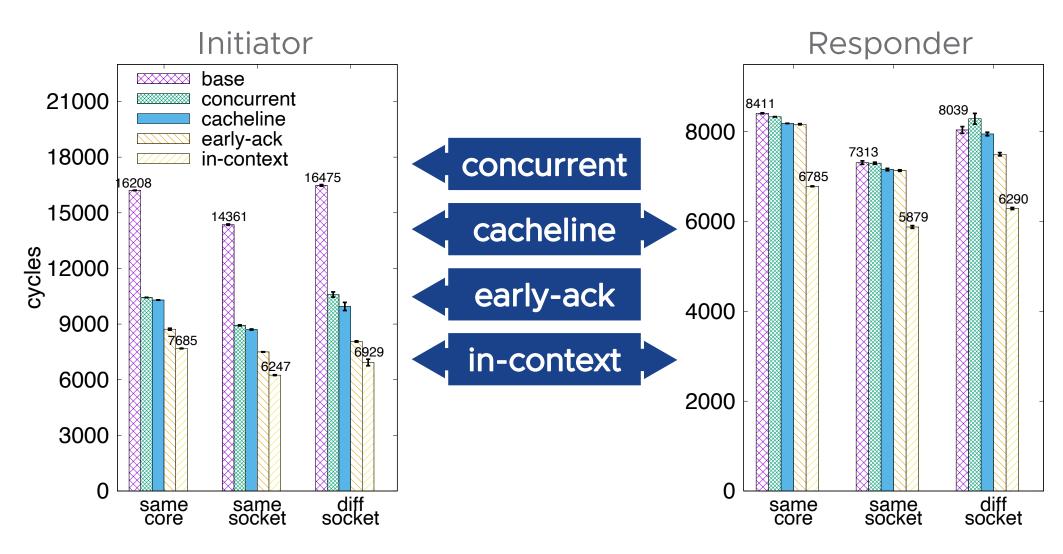
The effect of page size mismatch

Many important and subtle details



Evaluation: Unmapping and Flushing 10 PTEs

madvise(MADV_DONTNEED)







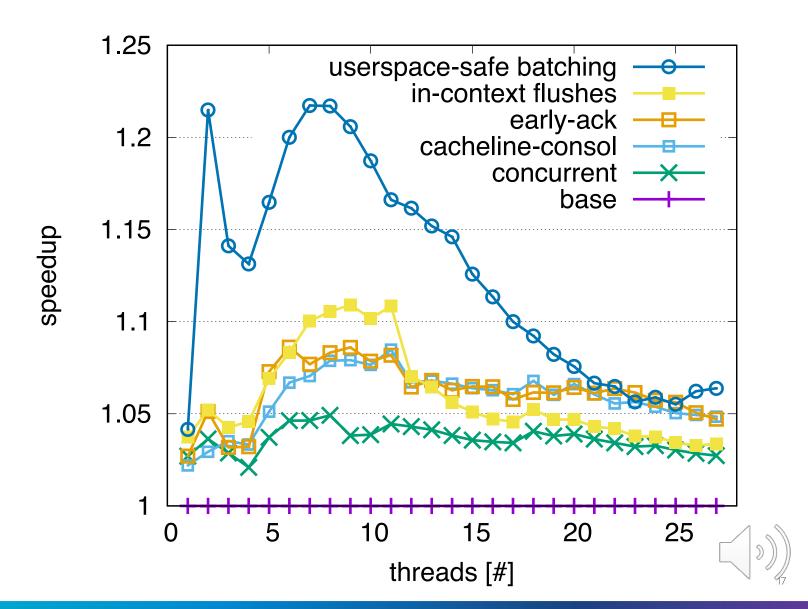
Evaluation: SysBench – Random Writes

Random writes

Periodic flushes

Memory-mapped file

Emulated persistent memory, no write-cache



Conclusions

TLB shootdown can be improved

Doing it well in software → better hardware interfaces

We are working to push these enhancements to Linux

