

Cheat Sheet

Proposition 0.1. *A functor is an equivalence if and only if it is full, faithful, and essentially surjective on objects.*

The naturality axiom of adjoints has two parts:

$$\overline{(F(A) \xrightarrow{g} B \xrightarrow{q} B')} = (A \xrightarrow{\bar{g}} G(B) \xrightarrow{G(q)} G(B')) \quad (1)$$

(that is, $\overline{q \circ g} = G(q) \circ \bar{g}$) for all g and q , and

$$\overline{(A' \xrightarrow{p} A \xrightarrow{f} G(B))} = (F(A') \xrightarrow{F(p)} F(A) \xrightarrow{\bar{f}} B) \quad (2)$$

for all p and f .

Forgetful functors between categories of algebraic structures usually have left adjoints.

Example 0.1. Initial and terminal objects can be described as adjoints. Let \mathcal{A} be a category. There is precisely one functor $\mathcal{A} \rightarrow \mathbf{1}$. Also, a functor $\mathbf{1} \rightarrow \mathcal{A}$ is essentially just an object of \mathcal{A} . Viewing functors $\mathbf{1} \rightarrow \mathcal{A}$ as objects of \mathcal{A} , a left adjoint to $\mathcal{A} \rightarrow \mathbf{1}$ is exactly an initial object of \mathcal{A} .

Similarly, a right adjoint to the unique functor $\mathcal{A} \rightarrow \mathbf{1}$ is exactly a terminal object of \mathcal{A} .

For each $A \in \mathcal{A}$, we have a map

$$(A \xrightarrow{\eta_A} GF(A)) = \overline{(F(A) \xrightarrow{1} F(A))}. \quad (3)$$

Dually, for each $B \in \mathcal{B}$, we have a map

$$(FG(B) \xrightarrow{\epsilon_B} B) = \overline{(G(B) \xrightarrow{1} G(B))}. \quad (4)$$

These define natural transformations

$$\eta : \mathbf{1}_{\mathcal{A}} \rightarrow G \circ F, \quad \epsilon : F \circ G \rightarrow \mathbf{1}_{\mathcal{B}},$$

called the **unit** and **counit** of the adjunction, respectively.

Lemma 0.1. *Given an adjunction $F \dashv G$ with unit η and counit ϵ , the triangles*

$$\begin{array}{ccc} F & \xrightarrow{F\eta} & FGF \\ & \searrow 1_F & \downarrow \epsilon_F \\ & & F \end{array} \quad \begin{array}{ccc} G & \xrightarrow{\eta G} & GFG \\ & \searrow 1_G & \downarrow G\epsilon \\ & & G \end{array}$$

commute.

Remark. These are called the **triangle identities**. An equivalent statement is that the triangles

$$\begin{array}{ccc} F(A) & \xrightarrow{F(\eta_A)} & FGF(A) \\ & \searrow 1_{F(A)} & \downarrow \epsilon_{F(A)} \\ & & F(A) \end{array} \quad \begin{array}{ccc} G(B) & \xrightarrow{\eta_{G(B)}} & GFG(B) \\ & \searrow 1_{G(B)} & \downarrow G(\epsilon_B) \\ & & G(B) \end{array} \quad (5)$$

commute for all $A \in \mathcal{A}$ and $B \in \mathcal{B}$.

Definition 0.1. Given categories and functors

$$\begin{array}{ccc} & \mathcal{B} & \\ & \downarrow Q & \\ \mathcal{A} & \xrightarrow{P} & \mathcal{C}, \end{array}$$

the **comma category** $(P \Rightarrow Q)$ (often written as $(P \downarrow Q)$) is the category defined as follows:

- objects are triples (A, h, B) with $A \in \mathcal{A}$, $B \in \mathcal{B}$, and $h : P(A) \rightarrow Q(B)$ in \mathcal{C} ;
- maps $(A, h, B) \rightarrow (A', h', B')$ are pairs $(f : A \rightarrow A', g : B \rightarrow B')$ of maps such that the square

$$\begin{array}{ccc} P(A) & \xrightarrow{P(f)} & P(A') \\ h \downarrow & & \downarrow h' \\ Q(B) & \xrightarrow{Q(g)} & Q(B') \end{array}$$

commutes.

Example 0.2. Let \mathcal{A} be a category and $A \in \mathcal{A}$. The **slice category** of \mathcal{A} over A , denoted by \mathcal{A}/A , is the category whose objects are maps into A and whose maps are commutative triangles. More precisely, an object is a pair (X, h) with $X \in \mathcal{A}$ and $h : X \rightarrow A$ in \mathcal{A} , and a map $(X, h) \rightarrow (X', h')$ in \mathcal{A}/A is a map $f : X \rightarrow X'$ in \mathcal{A} making the triangle

$$\begin{array}{ccc} X & \xrightarrow{f} & X' \\ & \searrow h & \swarrow h' \\ & A & \end{array}$$

commute.

Slice categories are a special case of comma categories. Functors $\mathbf{1} \rightarrow \mathcal{A}$ are just objects of \mathcal{A} . Now, given an object A of \mathcal{A} , consider the comma category $(\mathbf{1}_{\mathcal{A}} \Rightarrow A)$, as in the diagram

$$\begin{array}{ccc} & \mathbf{1} & \\ & \downarrow A & \\ \mathcal{A} & \xrightarrow{\mathbf{1}_{\mathcal{A}}} & \mathcal{A}. \end{array}$$

An object of $(\mathbf{1}_{\mathcal{A}} \Rightarrow A)$ is in principle a triple (X, h, B) , with $X \in \mathcal{A}$, $B \in \mathbf{1}$, and $h : X \rightarrow A$ in \mathcal{A} ; but $\mathbf{1}$ has only one object, so it is essentially just a pair (X, h) . Hence the comma category $(\mathbf{1}_{\mathcal{A}} \Rightarrow A)$ has the same objects as the slice category \mathcal{A}/A and one can check that it has the same maps too, so $\mathcal{A}/A \cong (\mathbf{1}_{\mathcal{A}} \Rightarrow A)$.

Lemma 0.2. Take an adjunction $\mathcal{A} \xrightleftharpoons[G]{F} \mathcal{B}$ and an object $A \in \mathcal{A}$. Then the unit map $\eta_A : A \rightarrow GF(A)$ is an initial object of $(A \Rightarrow G)$.

Corollary 0.1. Let $G : \mathcal{B} \rightarrow \mathcal{A}$ be a functor. Then G has a left adjoint if and only if for each $A \in \mathcal{A}$, the category $(A \Rightarrow G)$ has an initial object.

Definition 0.2. Let \mathcal{A} be a locally small category and $A \in \mathcal{A}$. We define a functor

$$H^A = \mathcal{A}(A, -) : \mathcal{A} \rightarrow \mathbf{Set}$$

as follows:

- for objects $B \in \mathcal{A}$, put $H^A(B) = \mathcal{A}(A, B)$;
- for maps $B \xrightarrow{g} B'$ in \mathcal{A} , define

$$H^A(g) = \mathcal{A}(A, g) : \mathcal{A}(A, B) \rightarrow \mathcal{A}(A, B')$$

by

$$p \mapsto g \circ p$$

for all $p : A \rightarrow B$.

Definition 0.3. Let \mathcal{A} be a locally small category. A functor $X : \mathcal{A} \rightarrow \mathbf{Set}$ is **representable** if $X \cong H^A$ for some $A \in \mathcal{A}$. A **representation** of X is a choice of an object $A \in \mathcal{A}$ and an isomorphism between H^A and X .

Proposition 0.2. Any set-valued functor with left adjoint is representable.

Definition 0.4. Let \mathcal{A} be a locally small category and $A \in \mathcal{A}$. We define a functor

$$H_A = \mathcal{A}(-, A) : \mathcal{A}^{\text{op}} \rightarrow \mathbf{Set}$$

as follows:

- for objects $B \in \mathcal{A}$, put $H_A(B) = \mathcal{A}(B, A)$;
- for maps $B' \xrightarrow{g} B$ in \mathcal{A} , define

$$H_A(g) = \mathcal{A}(g, A) : \mathcal{A}(B, A) \rightarrow \mathcal{A}(B', A)$$

by

$$p \mapsto p \circ g$$

for all $p : B \rightarrow A$.

Definition 0.5. Let \mathcal{A} be a locally small category. A functor $X : \mathcal{A}^{\text{op}} \rightarrow \mathbf{Set}$ is **representable** if $X \cong H_A$ for some $A \in \mathcal{A}$. A **representation** of X is a choice of an object $A \in \mathcal{A}$ and an isomorphism between H_A and X .

Any map $A \xrightarrow{f} A'$ in \mathcal{A} induces a natural transformation

$$\begin{array}{ccc} \mathcal{A}^{\text{op}} & \begin{array}{c} \xrightarrow{H_A} \\ \Downarrow H_f \\ \xrightarrow{H_{A'}} \end{array} & \mathbf{Set} \end{array}$$

whose component at an object $B \in \mathcal{A}$ is

$$\begin{array}{ccc} H_A(B) = \mathcal{A}(B, A) & \rightarrow & H_{A'}(B) = \mathcal{A}(B, A') \\ p & \mapsto & f \circ p. \end{array}$$

Definition 0.6. Let \mathcal{A} be a locally small category. The **Yoneda embedding** of \mathcal{A} is the functor

$$H_{\bullet} : \mathcal{A} \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$$

defined on objects A by $H_{\bullet}(A) = H_A$ and on maps f by $H_{\bullet}(f) = H_f$.

Here is a summary of the definitions so far.

For each $A \in \mathcal{A}$, we have a functor $\mathcal{A} \xrightarrow{H^A} \mathbf{Set}$.

Putting them all together gives a functor $\mathcal{A}^{\text{op}} \xrightarrow{H^{\bullet}} [\mathcal{A}, \mathbf{Set}]$.

For each $A \in \mathcal{A}$, we have a functor $\mathcal{A}^{\text{op}} \xrightarrow{H_A} \mathbf{Set}$.

Putting them all together gives a functor $\mathcal{A} \xrightarrow{H_{\bullet}} [\mathcal{A}^{\text{op}}, \mathbf{Set}]$.

Fix a small category \mathcal{A} . Take an object $A \in \mathcal{A}$ and a functor $X : \mathcal{A}^{\text{op}} \rightarrow \mathbf{Set}$. The object A gives rise to another functor $H_A = \mathcal{A}(-, A) : \mathcal{A}^{\text{op}} \rightarrow \mathbf{Set}$. The question is: what are the maps $H_A \rightarrow X$? Since H_A and X are both objects of the presheaf category $[\mathcal{A}^{\text{op}}, \mathbf{Set}]$, the 'maps' concerned are maps in $[\mathcal{A}^{\text{op}}, \mathbf{Set}]$. So, we are asking what natural transformations

$$\begin{array}{ccc} \mathcal{A}^{\text{op}} & \begin{array}{c} \xrightarrow{H_A} \\ \Downarrow \\ \xrightarrow{X} \end{array} & \mathbf{Set} \end{array}$$

there are. The set of such natural transformations is called

$$[\mathcal{A}^{\text{op}}, \mathbf{Set}](H_A, X).$$

Are there any other ways to construct a set from the same input data (A, X) ? Yes: simply take the set $X(A)$.

Theorem 0.1 (Yoneda). *Let \mathcal{A} be a locally small category. Then*

$$[\mathcal{A}^{\text{op}}, \mathbf{Set}](H_A, X) \cong X(A) \quad (6)$$

naturally in $A \in \mathcal{A}$ and $X \in [\mathcal{A}^{\text{op}}, \mathbf{Set}]$.

Proposition 0.3. *Let \mathcal{A} be a category.*

- (a) *If \mathcal{A} has all products and equalizers then \mathcal{A} has all limits.*
- (b) *If \mathcal{A} has binary products, a terminal object and equalizers then \mathcal{A} has finite limits.*

Lemma 0.3. *A map $X \xrightarrow{f} Y$ is monic if and only if the square*

$$\begin{array}{ccc} X & \xrightarrow{1} & X \\ 1 \downarrow & & \downarrow f \\ X & \xrightarrow{f} & Y \end{array}$$

is a pullback.

Example 0.3. Take sets and functions $X \xrightarrow[s]{t} Y$. To find the coequalizer of s and t , we must construct in some canonical way a set C and a function $p : Y \rightarrow C$ such that $p(s(x)) = p(t(x))$ for all $x \in X$. Let \sim be the equivalence relation on Y generated by $s(x) \sim t(x)$ for all $x \in X$. Take the quotient map $p : Y \rightarrow Y/\sim$. The maps $Y/\sim \rightarrow B$ correspond one-to-one with maps $f : Y \rightarrow B$, so p is indeed the coequalizer of s and t .

Definition 0.7. (a) Let \mathbf{I} be a small category. A functor $F : \mathcal{A} \rightarrow \mathcal{B}$ **preserves limits of shape \mathbf{I}** if for all diagrams $D : \mathbf{I} \rightarrow \mathcal{A}$ and all cones $\left(A \xrightarrow{p_I} D(I)\right)_{I \in \mathbf{I}}$ on D ,

$$\begin{aligned} & \left(A \xrightarrow{p_I} D(I)\right)_{I \in \mathbf{I}} \text{ is a limit cone on } D \text{ in } \mathcal{A} \\ \Rightarrow & \left(F(A) \xrightarrow{Fp_I} FD(I)\right)_{I \in \mathbf{I}} \text{ is a limit cone on } F \circ D \text{ in } \mathcal{B}. \end{aligned}$$

(b) A functor $F : \mathcal{A} \rightarrow \mathcal{B}$ **preserves limits** if it preserves limits of shape \mathbf{I} for all small categories \mathbf{I} .

(c) **Reflection** of limits is defined as in (a), but with \Leftarrow instead of \Rightarrow .

Definition 0.8. A functor $F : \mathcal{A} \rightarrow \mathcal{B}$ **creates limits (of shape \mathbf{I})** if whenever $D : \mathbf{I} \rightarrow \mathcal{A}$ is a diagram in \mathcal{A} ,

- for any limit cone $\left(B \xrightarrow{q_I} FD(I)\right)_{I \in \mathbf{I}}$ on the diagram $F \circ D$, there is a unique cone $\left(A \xrightarrow{p_I} D(I)\right)_{I \in \mathbf{I}}$ on D such that $F(A) = B$ and $F(p_I) = q_I$ for all $I \in \mathbf{I}$;
- this cone $\left(A \xrightarrow{p_I} D(I)\right)_{I \in \mathbf{I}}$ is a limit cone on D .

The forgetful functors from **Grp**, **Ring**, \dots to **Set** all create limits.

Lemma 0.4. *Let $F : \mathcal{A} \rightarrow \mathcal{B}$ be a functor and \mathbf{I} a small category. Suppose that \mathcal{B} has, and F creates, limits of shape \mathbf{I} . Then \mathcal{A} has, and F preserves limits of shape \mathbf{I} .*

0.1 Bonus

$$\text{Hom}(X, A) \times \text{Hom}(X, B) \cong \text{Hom}(X, A \times B)$$

$$\text{Hom}(A, X) \times \text{Hom}(B, X) \cong \text{Hom}(A + B, X)$$

$$\text{Hom}(X \times A, B) \cong \text{Hom}(X, B^A)$$

$$HA \cong HB \Rightarrow A \cong B$$