

FOOL Programming Language

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Abstract

In this report we show the steps needed to build a simple compiler for an educational programming language called Functional Object Oriented Language (FOOL). Nowadays very few people build compilers or interpreters for mainstream programming languages. On the other hand knowing the theory and the practice behind this discipline should be an important achievement on the life of a computer scientist. Not only because we can understand better the functions and the peculiarities of the languages that we use every day but we can be better developers by knowing all the processes that our code passes through until it become machine code and get executed by the hardware.

Keywords: Fool Programming Languages, ANTLR4, Compiler, Interpreter, Grammar, Assembly Code, C#.

1 Introduction

In the present work we describe the main steps needed to create a compiler for a toy programming language called FOOL. This project is developed as part of the exam *Compilers and Interpreters*. The grammar and some code snippets for the project are provided by the professor. You can see the grammar in the Appendix of this report.

The objective is to construct a compiler that given in input a file written using Fool syntax (.fool) to generate some *assembly code* in a given specific assembly language for a custom virtual machine called *Simple Virtual Machine, SVM*.

The first step of the development is to prepare your development tools and input files.

- **Programming Language:** C# version 8.0 ¹.
- **Development IDE:** We have used Microsoft Visual Studio Community 2017 version:15.9.35² with Microsoft .NET Framework version 4.8.04084 to write, test and debug the project.
- **Lexer/Parser generator:** We used Antlr4.Runtime.Standard version: 4.8.0³ to generate the following files: *FOOL.interp*, *FOOL.tokens*, *FOOLLexer.cs*, *FOOLLexer.interp*, *FOOLLexer.tokens*. This files are automatically generated by ANTLR4 using an input grammar file *FOOL.g4*.
- **FOOL.g4:** Is the input grammar used by ANTLR4 to generate the previous files.
- **Helper Package:** System.Configuration.ConfigurationManager version:v6.0.0⁴ is used to configure the custom project directory paths.

Once prepared the environment and the tools we have to create the project structure and the input files generated by ANTLR4.

In the second step we produce the project directory structure as shown in fig. 1. and populate the folder *Lexer* with the ANTLR4 generated files ⁵. In order to generate those files we have to use the following command from the console:

```
>java -jar antlr-4.8.0.jar -Dlanguage=Csharp FOOL.g4 -visitor
```

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Programming Language.

1. <https://docs.microsoft.com/it-it/dotnet/csharp/whats-new/csharp-8>
2. <https://visualstudio.microsoft.com/it-vs/older-downloads/#visual-studio-2017-and-other-products>
3. <https://www.nuget.org/packages/antlr4.runtime.standard/4.8.0>
4. <https://www.nuget.org/packages/system.configuration.configurationmanager/>
5. <https://hayeol.tistory.com/45>

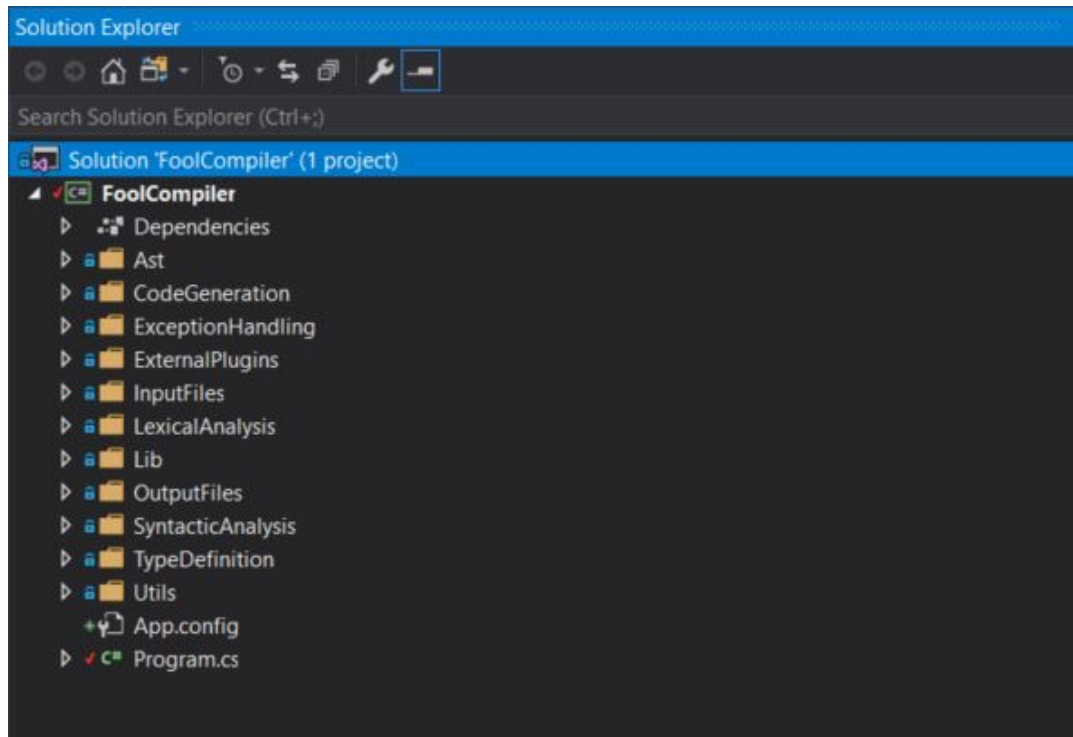


Figure 1. The folder structure of the Project

We have to copy the generated files in the folder *Lexer* of our project.

Now let's have a view in the folder structure of the project: Now according to .NET terminology the project consists of one *solution* called **FoolCompiler** that contains only one *project* called FoolCompiler as well. FoolCompiler project is populated by the following folders:

- **Abstract Syntax Tree (Ast)** contains 22 **.cs**. The first 20 are the implementation of all possible nodes used in the Ast. We have used the following naming convention: *Fool**NodeName**Node*, for example: *FoolMethodNode.cs* or *FoolProgramLetInExpressionNode.cs*. The remaining 2 files are the visitor implementation (*FoolVisitor.cs*) and the interface from which all the nodes inherit (*IFoolNode.cs*).
- **CodeGeneration** folder contains 16 files, some of them are classes that we have write others are files generated in the process of code generation. In this folder are placed all the code that guide the code generation process, the function of the heap and the dispatch table.
- **ExceptionHandling** contains the custom error handling mechanisms. This folder contains 8 files.
- **ExternalPlugins** contains the ANTLR4 jar used to generate the lexer and the parser starting from the grammar.
- **InputFiles** contains 42 **.fool** testing files that we have use to test the compiler.
- **LexicalAnalysis** contain 6 files, the grammar FOOL.g4 and 5 other files generated by ANTLR4.
- **Lib** is a library folder that contain only one helper class.
- **SyntacticAnalysis** contain 3 ANTLR4 generated files.
- **TypeDefinition** contain 9 files, each of them is the definition of one Fool type, such as *FoolIntType.cs*, *FoolBoolType.cs* etc. All this types inherit from the interface *IFoolType.cs*.
- **Utils** contain the utility classes and functions used in the whole project. This file contain 6 files.
Also in the root directory of the project we have 2 very important files:
- **Program.cs** is the starting point of the execution of the project. Here is where the *main* method is defined.

```

1 <?xml version="1.0" encoding="utf-8"?>
2 <configuration>
3   <appSettings>
4     <add key="inputFilePath" value="C:\Users\uzeir\Desktop\FoolCompiler\FoolCompiler\InputFiles\" />
5     <add key="outputFilePath" value="C:\Users\uzeir\Desktop\FoolCompiler\FoolCompiler\OutputFiles\code.svm" />
6   </appSettings>
7 </configuration>

```

Figure 2. Web.config

- **App.config** is the file that we use to set the absolute path to the input and output directories on your local machine. This is needed in order to avoid hard-coding of all the paths in the code. Fig. 2. show the format of this file.

2 How to run the project

The project can be download from the Github⁶ repository and executed locally. Here is a step by step tutorial how to run the project in Windows using Visual Studio 2017. In case you want to run the project using Linux you can use Mono⁷ or install manually .NET⁸ in your computer. Image 3 shows a possible output when the execution terminates.

- **Step 1:** Download and unzip the project folder form GitHub.
- **Step 2:** Open the folder and modify *Web.config* with your own paths.
- **Step 3:** Open the project in the IDE of your choice .
- **Step 4:** Run the project.
- **Step 5:** Choose the input file from the default folder or create your own input file.
- **Step 6:** Click *Enter* to execute the project.

3 Lexical Analysis - Lexer

The whole process of code generation, starting from the input file pass through a number of well defined phases. *Lexical Analysis* or *Lexer* is the first of those steps. In our case the Lexer is generated automatically by the *ANTLR4 tool*. The following listing shows the piece of code in the *Program.cs* that implement the creation and the use of the lexer.

The job of the Lexer is to consume the input file written in the FOOL language and to produce a set of **Tokens**, one for each lexeme in the original input. The Lexer is also called a *tokenizer*.

```

1000 // LEXER
1001 Console.WriteLine("[<<<< LEXICAL ANALYSIS >>>>]");
1002 FOOLLexer lexer = new FOOLLexer(input);
1003 if (lexer.lexicalErrors.Count > 0)
1004 {
1005     foreach (var e in lexer.lexicalErrors)
1006     {
1007         throw new FoolLexerException(e);
1008     }
1009 }
1010 CommonTokenStream tokens = new CommonTokenStream(lexer);
1011 Console.WriteLine("\nOK, DONE!");
1012 Console.WriteLine("*****");

```

Fig. 4. show an example of tokenization for the input code $437 + 743$.

In order to generate the tokens, Lexer make use of the input grammar (FOOL.g4) where we have

6. <https://github.com/euzeir/FoolCompiler>

7. <https://www.mono-project.com/>

8. <https://opensource.com/article/17/11/net-linux>

```

C:\Program Files\dotnet\dotnet.exe
*****
***FOOL COMPILER***
*****
Usage:
  1.[Select one of the provided input files: "input01" up to "input40"]
  2.[Insert your input from Console by adding <file name> + <file text>]
*****
input26

[<<<<< LEXICAL ANALYSIS >>>>>]

OK, DONE!
*****
[<<<<< SYNTAX ANALYSIS >>>>>]

OK, DONE!
*****
[<<<<< SEMANTIC ANALYSIS >>>>>]

OK, DONE!
*****
[<<<<< TYPE CHECKING >>>>>]
Type checking: INT

OK, DONE!
*****
[<<<<< CODE GEN >>>>>]

RESULT: 30
>> Execution Terminated!
>> Please press ESC or X to close the window.

```

Figure 3. Example of Execution

initialize all the needed rules (parser and lexer rules).

We want to point out the fact that the rules are divided in 2 classes, the parser rules that comes first in the grammar and the lexer rules that comes later. Even inside each class the order of the rules is important and helps in case of ambiguity. Another distinction is that lexer rule names are in uppercase letters, and usually expressed in the form of regular expressions or one-to-one matching. Our lexer has also specific rules that ignore comments and white spaces. This code inform the lexer how to generate the tokens and in the case of errors to report. As you can see in the previous code listing the errors are added in a list called *lexicalErrors* and if this list is empty we proceed with the following steps.

The part of lexical rules in the grammar is shown in the Appendix.

```

1000 exp  : left=term ((PLUS | MINUS) right=exp)?
      ;
1002 //some lines further
1004 PLUS  : '+' ;
1006 //Numbers
      fragment DIGIT : '0'..'9';
1008 INTEGER      : DIGIT+;

```

Once the input files is tokenized and the error list is empty we pass to the second phase, the *Parsing*.

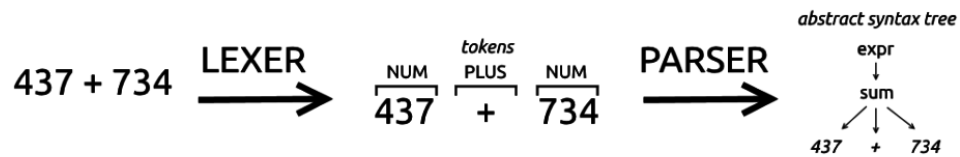


Figure 4. Example of Tokenization and Parsing

4 Syntax Analysis - Parser

Parsing is the second step in the process of input analysis. It is automatically performed by the ANTLR4 plugin and has as final result a tree representation of the input called *Abstract Syntax Tree* or *AST*.

The snippet of code shows the creation on the parser and it's use in out *Program.cs* file

```

1000 // PARSE
1001 Console.WriteLine("[<<<< SYNTAX ANALYSIS >>>>]");
1002 FOOLParser parser = new FOOLParser(tokens);
1003 FOOLParser.ProgContext progContext = parser.prog();
1004 if (parser.NumberOfSyntaxErrors > 0)
1005     throw new FoolParserException("Errors: " + parser.NumberOfSyntaxErrors + "\n");
1006 Console.WriteLine("\nOK, DONE!");
1007 Console.WriteLine("*****");

```

The main purpose of the parsing process is to determine if a given input is part of the language syntax or not based on the rules specified on the grammar. When running ANTLR4 plugin with the *-visitor* option we get back a set of output files. Three of them are particularly interesting in this phase. *FOOLBaseVisitor.cs* and *FOOLVisitor.cs* that represent a blueprint of the visitor pattern applied to our grammar rules. The first is an abstract class that define the signature of all virtual methods that then are implemented by us in the second file. This is the point where we insert the logic of how the process has to progress and which controls has to be made.

Although ANTLR4 allow the use of two patterns, *Listener* and *Visitor*, we used the latest because it permit us to control the manner of how the nodes are parsed and also to get information about each node. This information are then used in the next phase for the *Semantic Analysis* and the *Type checking*. The way the *Visitor* works is very simple, it starts from the root node and execute a depth-first search (DFS) on the AST of the chosen rule based on the input and the grammar. Keep in mind that the visitor produced by ANTLR4 has one method for each rule in the grammar file.

The following listing is the virtual default method produced by ANTLR4 in the *FOOLBaseVisitor.cs* file for the *LetInExp* rule in the grammar.

```

1000 public virtual Result VisitLetInExp([NotNull] FOOLParser.LetInExpContext context)
1001 { return VisitChildren(context); }

```

Here is the starting rule in the grammar from which ANTLR4 generated the previous method. As you can see is the second option in the starting rule.

```

1000 prog : exp SEMIC #singleExp
1001      | let IN (exp SEMIC | stms) #letInExp
1002      | (classdec)+ let? IN (exp SEMIC | stms) #classExp
1003      ;

```

Now in our visitor implementation we have to insert the logic of how the visitor will behave in each node it reaches.

The following listing represents the code we implemented for our project for the *LetInExp* rule. As you can see in the code snippet the method get as parameter a *context* object of the type *LetInExpContext* in this case. The logic of how the analysis has to proceed is explicitly express in the code. Inside the method we check different conditions, in this case if we are dealing with an *expression* of a *statement* and base on it we create an instance of *FoolProgramLetInExpressionNode* class with different parameters.

```
1000     public override IFoolNode VisitLetInExp(LetInExpContext context)
1001     {
1002         List<IFoolNode> declarations = new List<IFoolNode>();
1003
1004         foreach (DecContext dec in context.let().dec())
1005         {
1006             declarations.Add(Visit(dec));
1007         }
1008
1009         if (context.exp() != null)
1010         {
1011             return new FoolProgramLetInExpressionNode(declarations, Visit(
1012 context.exp()));
1013         }
1014         else
1015         {
1016             List<IFoolNode> statements = new List<IFoolNode>();
1017             foreach (StmContext stm in context.stms().stm())
1018             {
1019                 statements.Add(Visit(stm));
1020             }
1021             return new FoolProgramLetInExpressionNode(declarations, statements)
1022         }
1023     }
```

Keep in mind that all the classes defined inherit by the single interface *IFoolNode* that define three methods as shown in the following snippet.

```
1000     public interface IFoolNode
1001     {
1002         IFoolType TypeCheck();
1003         List<string> CheckSemantics(FoolEnvironment environment);
1004         string CodeGeneration();
1005     }
```

If the parsing finish correctly; no errors found the flow goes throw the *Semantic Analysis* that check the logical meaning of our input code.

5 Semantic Analysis

Semantic Analysis is a fundamental part on the whole code analysis process that try to find that kind of errors that the previous two types of analyser couldn't find, semantic/meaning errors. During this process, the semantic analyser makes use of an important structure called *Symbol Table*. The symbol table in our case is implemented as a list of *dictionaries* in the C# terminology.

More formally. we are going to list the four type of checks that semantic analyser perform:

- check for multiple declarations of the same *name* in a given *scope*. In case of error an exception is raised.
- check for use of undeclared variables. As general rule in Fool a variable has to be declared before being used.
- check for type mismatch. Errors of type: *int**x*; *bool**y*; *x* = *y*;
- check if the methods are called with the right type and number of parameters.

This checks are done in two sequential steps by traversing the AST. The first step is called *Scope Checking* and the second is called *Type Checking*. We know that the scoping rules are fundamental for the correctness of the code. Fool language uses *static scoping*. The following listing show the entry point in our *Program.cs* file where we perform semantic checking.

```

1000 // SEMANTIC
      Console.WriteLine("[<<<< SEMANTIC ANALYSIS >>>>]");
1002 FoolVisitor visitor = new FoolVisitor();
      IFoolNode ast = visitor.Visit(progContext);
1004 FoolEnvironment environment = new FoolEnvironment();
      List<string> error = ast.CheckSemantics(environment);
1006 if (error.Count > 0) throw new FoolSemanticException(error);
      Console.WriteLine("\nOK, DONE!");
1008 Console.WriteLine("*****");
      // TYPE CHECKING
1010 Console.WriteLine("[<<<< TYPE CHECKING >>>>]");
      IFoolType type = ast.TypeCheck(); //type-checking bottom-up
1012 Console.WriteLine("Type checking: " + type.GetFoolType().ToUpper());
      Console.WriteLine("\nOK, DONE!");
1014 Console.WriteLine("*****");

```

As you can see in the code snippet, type checking is done by calling *ast.TypeCheck()* on the tree. Is important to mention that we have defined 8 different types in the project. All the types derive and implement *IFoolType.cs* interface that defines the signature of 3 methods.

```

1000     public interface IFoolType
      {
1002         string GetId();
            bool IsSubType(IFoolType type);
1004         string GetFoolType();
      }

```

The logic of *GetId()* and *GetFoolType()* is very simple. They just return the type name and the type.

IsSubType() is more elaborate. It takes as input a given type and determines if it is a sub-type of current type.

At the end if there are not errors found we go forward with the *Code Generation*. Keep in mind that at this point the helper structures that we used previously such as Symbol Tables are no needed anymore.

6 Code Generation

The purpose of this step is to generate an assembly-like code that later will be executed by our stack machine called SVM (Simple Virtual Machine). Stack machines are very simple models where the operand are always on the top of the stack.

The grammar for the assembly language, *SVM.g4* is provided by the professor and further enriched by our work adding some more instructions to it, such as *new*, *loadm* etc. Keep in mind that in this

grammar we have removed *actions* and *attributes*. We perform for this grammar all the previous analysis except for the semantic analysis and return the assembly code that has to be executed later on. The code snippet below show the initialization of the *virtualMachine* object. We call on it the *cpu()* method that has the task to execute the assembly code and produce a value as input or an error if any.

As we mention previously the virtual machine used in the project is a stack machine, where the operators are always in the top of the stack and the common *push* and *pop* stack operations are defied. In the case of the objects, this memory arrangement is not enough. We need another data structure called *heap* to deal with object. As in all other OOP languages even in Fool the object are instantiate using *new* keyword. The logic of how heap is implemented is defined in the file *FoolHeap.cs*.

Another important part of the OOP is the inheritance. Fool programming language permits inheritance and method overriding. This feature is implemented using the *dispatch table*. Itself the dispatch table is organized as a dictionary where the Keys are the name of the classes and the Values are a list of dispatch table entries called *FoolDispatchTableEntry*. The following code snippet show the structure of each DispatchTableEntry.

```
1000     public class FoolDispatchTableEntry
1001     {
1002         private string _methodId;
1003         private string _methodLabel;
1004
1005         public FoolDispatchTableEntry(string methodId, string methodLabel)
1006         {
1007             _methodId = methodId;
1008             _methodLabel = methodLabel;
1009         }
1010         public string GetMethodId()
1011         {
1012             return _methodId;
1013         }
1014         public string GetMethodLabel()
1015         {
1016             return _methodLabel;
1017         }
1018     }
```

The two methods that are defined in this class are used to get the the name of the method (GetMethodId()) and to get the label (GetMethodLabel). The labels are used to show the part of the assembly file where the methods code start.

At the end of the executing we will be able to see in the terminal a value or an error.

The resource that i most used to prepare this project is mentioned in the bibliography and is the book "The definitive ANTLR 4 Reference" and the slides of the course.

7 Apendix

The following listing shows the Parser rules used in the Project.

```
1000 /*-----
1001  *  PARSER RULES
1002  *-----*/
1003 prog   : exp SEMIC
1004         | let IN (exp SEMIC | stms)
1005         | (classdec)+ let? IN (exp SEMIC | stms) ;
1006 let    : LET (dec SEMIC)+ ;
1007 classdec : CLASS ID (EXTENDS ID)?
1008           (LPAR vardec (COMMA vardec)* RPAR)?
1009           (CLPAR met* CRPAR)? SEMIC ;
1010 letVar : LET (varasm)+ ;
1011 vardec : type ID ;
1012 varasm : vardec ASM (exp | NULL) ;
1013 fun    : (VOID | type) ID LPAR ( vardec ( COMMA vardec)* )?
1014         RPAR CLPAR (letVar IN)? (exp SEMIC | stms) CRPAR ;
1015 dec    : varasm
1016         | fun ;
1017 met    : fun ;
1018 type   : INT
1019         | BOOL
1020         | ID ;
1021 exp    : left=term ((PLUS | MINUS) right=exp)? ;
1022 term   : left=factor ((TIMES | DIV) right=term)? ;
1023 factor : left=value ( logicoperator=
1024                 (EQ|
1025                 GREATERTHAN|
1026                 LESSERTHAN|
1027                 GREATEREQUAL|
1028                 LESSEREQUAL|AND|OR) right=factor)? ;
1029 value  : (MINUS)? INTEGER
1030         | (NOT)? ( TRUE | FALSE )
1031         | LPAR exp RPAR
1032         | IF cond=exp THEN CLPAR thenBranch=exp
1033           SEMIC CRPAR (ELSE CLPAR elseBranch=exp SEMIC CRPAR)?
1034         | (MINUS| NOT)? ID
1035         | functioncall
1036         | ID DOT functioncall
1037         | NEW ID LPAR (exp (COMMA exp)* )? RPAR ;
1038 functioncall : ID LPAR (exp (COMMA exp)* )? RPAR ;
1039 stm    : ID ASM (NULL | exp)
1040         | IF LPAR cond=exp RPAR THEN CLPAR thenBranch=stms CRPAR
1041           (ELSE CLPAR elseBranch=stms CRPAR)? ;
1042 stms   : (stm SEMIC)+ ;
```

The following listing shows the Lexer rules used in the Project.

```
1000 /*-----
1001  * LEXER RULES
1002  *-----*/
1003 SEMIC : ';' ;
1004 COLON : ':' ;
1005 COMMA : ',' ;
1006 EQ : '==';
1007 ASM : '=' ;
1008 PLUS : '+' ;
1009 MINUS : '-' ;
1010 TIMES : '*' ;
1011 DIV : '/' ;
1012 TRUE : 'true' ;
1013 FALSE : 'false' ;
1014 LPAR : '(' ;
1015 RPAR : ')' ;
1016 CLPAR : '{' ;
1017 CRPAR : '}' ;
1018 IF : 'if' ;
1019 THEN : 'then' ;
1020 ELSE : 'else' ;
1021 LET : 'let' ;
1022 IN : 'in' ;
1023 VAR : 'var' ;
1024 FUN : 'fun' ;
1025 INT : 'int' ;
1026 BOOL : 'bool' ;
1027
1028 //ADDED FOR PROJECT.
1029
1030 OR : '||' ;
1031 AND : '&&' ;
1032 NOT : 'not' ;
1033 GREATERTHAN : '>' ;
1034 LESSERTHAN : '<' ;
1035 GREATEREQUAL : '>=' ;
1036 LESSEREQUAL : '<=' ;
1037
1038 VOID : 'void' ;
1039 CLASS : 'class' ;
1040 THIS : 'this' ;
1041 NEW : 'new' ;
1042 DOT : '.' ;
1043 EXTENDS : 'extends' ;
1044 NULL : 'null' ;
1045
1046 //Numbers
1047 fragment DIGIT : '0'..'9';
1048 INTEGER : DIGIT+;
1049
1050 //IDs
1051 fragment CHAR : 'a'..'z' | 'A'..'Z' ;
1052 ID : CHAR (CHAR | DIGIT)* ;
1053
1054 //ESCAPED SEQUENCES
1055 WS : (' ' | '\t' | '\n' | '\r')-> skip;
1056 LINECOMMENTS : '//' (~('\n' | '\r'))* -> skip;
```