

Computing eigenvalues and eigenvectors

Power iteration

Power iteration

Let $\mathbf{A} \in \mathbb{R}^{m \times m}$ has a full set of eigenvalues

$$\lambda_k, \quad k = 1, \dots, m$$

Suppose that

$$|\lambda_1| > |\lambda_2| \geq |\lambda_3| \geq \dots \geq |\lambda_m|$$

NB: $\lambda_1 \in \mathbb{R}$, otherwise λ_1^* is also an eigenvalue.

Power iteration

Only compute λ_1 .

Iteration:

1. Take arbitrary $\mathbf{x}_0 \in \mathbb{R}^m$.
2. Iterate:

$$\begin{cases} \mathbf{x}_{k+1} = \mathbf{A} \mathbf{x}_k \\ \lambda^{(k)} = \frac{\langle \mathbf{x}_k | \mathbf{x}_{k+1} \rangle}{\langle \mathbf{x}_k | \mathbf{x}_k \rangle} \quad \left(\equiv \frac{\langle \mathbf{x}_k | \mathbf{A} | \mathbf{x}_k \rangle}{\langle \mathbf{x}_k | \mathbf{x}_k \rangle} \right), \end{cases} \quad k = 0, 1, \dots$$

Then, $\lambda^{(k)} \rightarrow \lambda_1, \quad k \rightarrow \infty$.

Power iteration: why does it work

Consider

$$\mathbf{x}_0 = c_1 \mathbf{e}_1 + c_2 \mathbf{e}_2 + \cdots + c_m \mathbf{e}_m$$

where $\{\mathbf{e}_k\}$ are normalized eigenvectors of \mathbf{A} :

$$\mathbf{A} \mathbf{e}_k = \lambda_k \mathbf{e}_k \quad k = 1, \dots, m$$

Then,

$$\begin{aligned} \mathbf{x}_k &= \mathbf{A}^k \mathbf{x}_0 = \lambda_1^k c_1 \mathbf{e}_1 + \lambda_2^k c_2 \mathbf{e}_2 + \cdots + \lambda_m^k c_m \mathbf{e}_m \\ &= \lambda_1^k \left[c_1 \mathbf{e}_1 + \underbrace{\left(\frac{\lambda_2}{\lambda_1} \right)^k c_2 \mathbf{e}_2 + \cdots + \left(\frac{\lambda_m}{\lambda_1} \right)^k c_m \mathbf{e}_m}_{\rightarrow 0 \text{ as } k \rightarrow \infty} \right] \end{aligned}$$

Power iteration

Questions:

1. What if $c_1 = 0$ (i.e. $\mathbf{x}_0 \perp \mathbf{e}_1$)?

numerical errors will generate $\neq 0$ components in the direction of \mathbf{e}_1

2. Numerical stability?

3. Rate of convergence?

LOW, *a priori* error estimate?

4. *A posteriori* error estimate?

Stability

$$\begin{array}{ll} \text{If } |\lambda_1| > 1, & \lambda_1^k \rightarrow \infty \quad \text{overflow} \\ |\lambda_1| < 1, & \lambda_1^k \rightarrow 0 \quad \text{underflow} \end{array}$$

\Rightarrow Need to keep \mathbf{x}_k normalized: $\|\mathbf{x}_k\| = 1, \quad k = 0, 1, \dots$

Modified iteration:

$$\begin{cases} \mathbf{y}_{k+1} = \mathbf{A} \mathbf{x}_k \\ \mathbf{x}_{k+1} = \frac{\mathbf{y}_{k+1}}{\|\mathbf{y}_{k+1}\|} \leftarrow \text{normalize} \\ \lambda^{(k)} = \langle \mathbf{x}_k | \mathbf{y}_{k+1} \rangle \end{cases}$$

$$k = 0, 1, \dots$$

A priori error estimate

$$\frac{|\lambda^{(k)} - \lambda_1|}{|\lambda_1|} \leq \text{const} \times \left(\frac{\lambda_2}{\lambda_1} \right)^k$$

АКД §8.2

Convergence

$\lambda^{(k)}$ converges to λ_1 .

Moreover, \mathbf{x}_k converges to the eigenvector \mathbf{e}_1 .

Convergence: $\sin \varphi_k \rightarrow 0, \quad k \rightarrow \infty$

$$\cos \varphi_k = \frac{\langle \mathbf{x}_k | \mathbf{e}_1 \rangle}{\|\mathbf{x}_k\| \cdot \|\mathbf{e}_1\|}$$

A posteriori error estimate

There is no general solution.

For symmetric matrices:

Lemma: Let $\lambda_* \in \mathbb{R}$ and $\mathbf{x}_* \in \mathbb{R}^m$. Let $\mathbf{A} \in \mathbb{R}^{m \times m}$ is symmetric.

Then, \exists an eigenvalue λ of \mathbf{A} , such that

$$|\lambda - \lambda_*| \leq \frac{\|\mathbf{A}\mathbf{x}_* - \lambda_*\mathbf{x}_*\|}{\|\mathbf{x}_*\|}.$$

A posteriori error estimate

Therefore, an *a posteriori* error bound for the power iteration for symmetric matrices:

$$|\lambda_1 - \lambda^{(k)}| \leq \frac{\|\mathbf{y}_k - \lambda^{(k)} \mathbf{x}_{k-1}\|}{\|\mathbf{x}_{k-1}\|}$$

i.e. stop the iteration when the r.h.s is smaller then the predefined tolerance.

The rate of convergence; shifts

Convergence is slow if the spectral gap is small, $|\lambda_2 - \lambda_1| \ll \lambda_1$,
since then $\left| \frac{\lambda_2}{\lambda_1} \right| \approx 1$.

Shifts

Consider

$$\mathbf{B} = \mathbf{A} - \sigma \hat{\mathbf{1}}$$

with some $\sigma \in \mathbb{R}$. $\text{eigvals}(\mathbf{B}) = \text{eigvals}(\mathbf{A}) - \sigma$

Use the power iteration for \mathbf{B} . Convergence is controlled by $\lambda_k - \sigma$.
If $\lambda_1 - \sigma$ is the max eigenvalue of \mathbf{B} , then select σ such that

$$\max_{k=2, \dots, m} \left| \frac{\lambda_k - \sigma}{\lambda_1 - \sigma} \right| \Rightarrow \min$$

Krylov subspaces

Krylov subspaces

$$\mathcal{K}_r(\mathbf{A}, \mathbf{x}_0) = \text{span}(\mathbf{x}_0, \mathbf{A}\mathbf{x}_0, \mathbf{A}^2\mathbf{x}_0, \dots, \mathbf{A}^{r-1}\mathbf{x}_0)$$

Power iteration generates a sequence of *Krylov subspaces*

$$\mathcal{K}_1 \subset \mathcal{K}_2 \subset \dots \subset \mathbb{R}^m$$

Note that also

$$\mathbf{A}\mathcal{K}_r \subset \mathcal{K}_{r+1}$$

Krylov subspace iterations

Power iteration generated vectors,

$$\mathbf{A}^s \mathbf{x}_0$$

are progressively closer to being linear dependent with
 $s = 1, 2, \dots$.

Krylov subspace methods: form a basis of \mathcal{K}_r by reorthogonalization.

Inverse iteration

Inverse iteration

Suppose we know $\lambda_* \approx \lambda_s$, a good approximation to λ_s . How to compute \mathbf{e}_s , its eigenvector?

Idea 1

By definition, $\mathbf{A}\mathbf{e}_s = \lambda_s\mathbf{e}_s$, so try looking for a non-trivial solution of

$$(\mathbf{A} - \lambda_*\hat{\mathbf{1}})\mathbf{x} = 0$$

Since $\lambda_* \neq \lambda_s$, $\mathbf{A} - \lambda_*\hat{\mathbf{1}}$ is non-degenerate \Rightarrow there is only the trivial solution $\mathbf{x} = 0$.

Inverse iteration

The algorithm:

$$\begin{aligned}(\mathbf{A} - \lambda_* \hat{\mathbf{1}}) \mathbf{y}_{k+1} &= \mathbf{x}_k \\ \mathbf{x}_{k+1} &= \frac{\mathbf{y}_{k+1}}{\|\mathbf{y}_{k+1}\|}\end{aligned}\tag{1}$$

for $k = 0, 1, \dots$ and some initial condition \mathbf{x}_0 .

Never invert a matrix explicitly!

Inverse iteration

Formally:

$$\mathbf{y}_{k+1} = (\mathbf{A} - \lambda_* \hat{\mathbf{1}})^{-1} \mathbf{x}_k$$

\Rightarrow inverse iteration is nothing but a power iteration with the matrix $(\mathbf{A} - \lambda_* \hat{\mathbf{1}})^{-1}$

How and why does it work?

Inverse iteration

Consider symmetric \mathbf{A} . Let $\{\mathbf{e}_\alpha\}$, $\alpha = 1, \dots, m$ is an orthonormal eigensystem of \mathbf{A} .

$$\mathbf{x}_0 = \sum_{\alpha=1}^m c_\alpha \mathbf{e}_\alpha; \quad \mathbf{y}_1 = \sum_{\alpha=1}^m b_\alpha \mathbf{e}_\alpha$$

The iteration (1) becomes:

$$\sum_{\alpha=1}^m (\lambda_\alpha - \lambda_*) b_\alpha \mathbf{e}_\alpha = \sum_{\alpha=1}^m c_\alpha \mathbf{e}_\alpha$$

Due to orthonormality of $\mathbf{e}_\alpha \Rightarrow$

$$b_\alpha = \frac{c_\alpha}{\lambda_\alpha - \lambda_*}$$

Inverse iteration

Therefore,

$$\mathbf{y}_1 = \sum_{\alpha=1}^m \frac{c_{\alpha}}{\lambda_{\alpha} - \lambda_*} \mathbf{e}_{\alpha} = \frac{1}{\lambda_s - \lambda_*} \left[c_s \mathbf{e}_s + \underbrace{\sum_{\alpha \neq s} \frac{\lambda_s - \lambda_*}{\lambda_{\alpha} - \lambda_*} \mathbf{e}_{\alpha}} \right]$$

"small" if $|\lambda_s - \lambda_*| \ll |\lambda_{\alpha} - \lambda_*|, \quad \alpha \neq s$

Each iteration multiplies $\mathbf{e}_{\alpha \neq s}$ by

$$\frac{\lambda_s - \lambda_*}{\lambda_{\alpha} - \lambda_*}.$$

Modified inverse iteration

Inverse iteration requires a good approximation to the eigenvalue.

Try using the Rayleigh ratio

$$\lambda^{(k)} = \langle \mathbf{x}_k | \mathbf{A} \mathbf{x}_k \rangle ,$$

$$\left(\mathbf{A} - \lambda^{(k)} \hat{\mathbf{1}} \right) \mathbf{y}_{k+1} = \mathbf{x}_k ,$$

$$\mathbf{x}_{k+1} = \frac{\mathbf{y}_{k+1}}{\|\mathbf{y}_{k+1}\|}$$

for $k = 0, 1, \dots$ and $\|\mathbf{x}_0\| = 1$

Modified inverse iteration

Modified inverse iteration converges fast if the angle $\widehat{(\mathbf{x}_0, \mathbf{e}_s)}$ is small.

How to get $\mathbf{x}_0 \approx \mathbf{e}_s$? E.g. a couple of power iterations.