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# DNS, Web

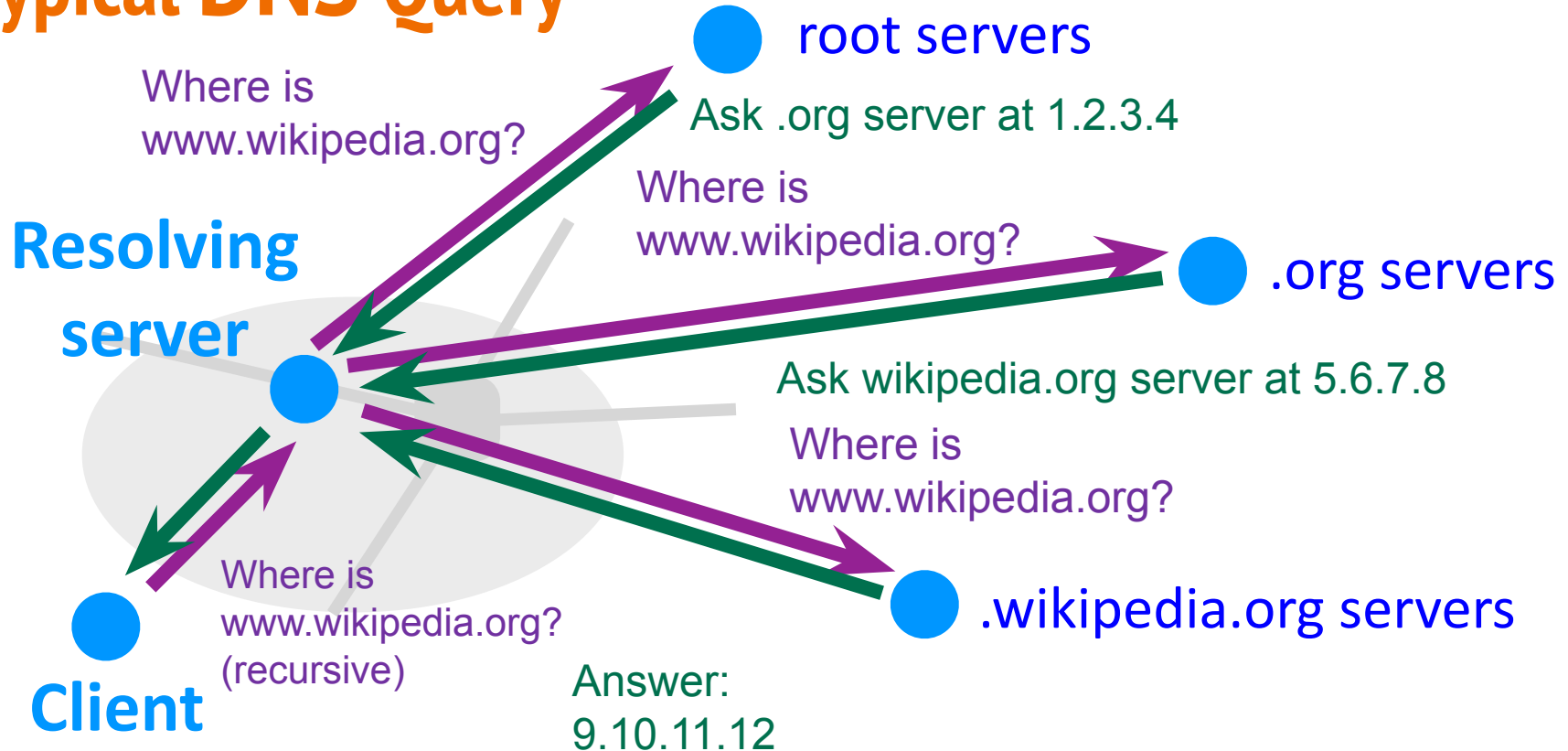
— CS 168 – Fall 2022 – Discussion 11 —

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# Agenda

- DNS Caching
- Anatomy of a URL
- HTTP
- CDN Caching

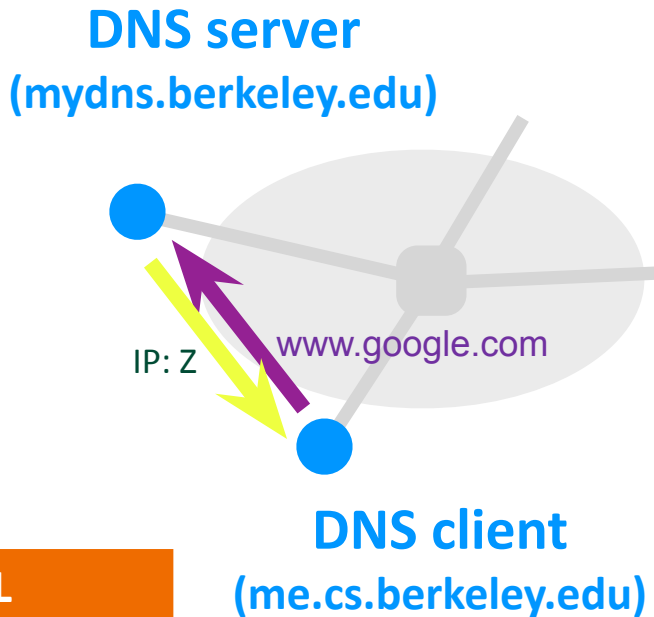
# Typical DNS Query



# Caching DNS Responses

- DNS responses cached in DNS servers
  - Expire after TTL (time-to-live)
- Most popular sites visited often
  - Top sites frequently cached. Fast!

Hostname	IP	TTL
www.google.com	Z	60 min



# Anatomy of a URL

**scheme://host[:port]/path/resource**

<i>Scheme</i>	Usually a protocol like (http, ftp, https, smtp, rtsp, etc.)
<i>host</i>	DNS hostname or an IP address
<i>port</i>	Defaults to protocol's standard port e.g. http: 80 https: 443
<i>path</i>	Traditionally reflects the file system
<i>resource</i>	Identifies the desired resource

Can also extend to program executions:

```
http://us.f413.mail.yahoo.com/ym/ShowLetter?box=
%40B%40Bulk&MsgId=2604_1744106_29699_1123_1261_0
_28917_3552_1289957100&Search=&Nhead=f&YY=31454&
order=down&sort=date&pos=0&view=a&head=b
```

# HTTP

- Client-server architecture
  - Server is “always on” and “well known”
  - Clients initiate contact to server
- Synchronous request/reply protocol
  - “Synchronous” means same HTTP session used for request and reply
  - Runs over TCP, Port 80

# Requests and Responses

## Request

```
GET /somedir/page.html HTTP/1.1
Host: www.someschool.edu
User-agent: Mozilla/4.0
Connection: close
Accept-language: fr
(blank line)
```

## Response

```
HTTP/1.1 200 OK
Connection close
Date: Thu, 06 Aug 2006 12:00:15 GMT
Server: Apache/1.3.0 (Unix)
Last-Modified: Mon, 22 Jun 2006 ...
Content-Length: 6821
Content-Type: text/html
(blank line)
data data data data data ...
```

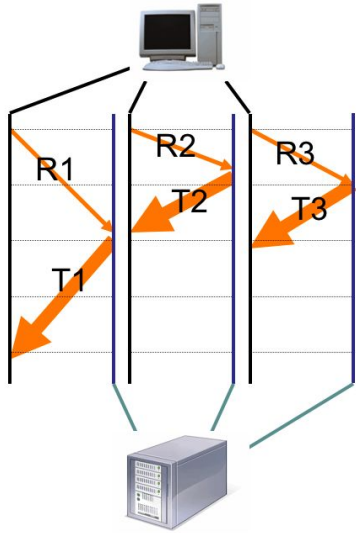
# Performance!

- A wise person once said “Architect for flexibility, engineer for performance”
- We have an architected solution – let’s explore ways to make it go fast

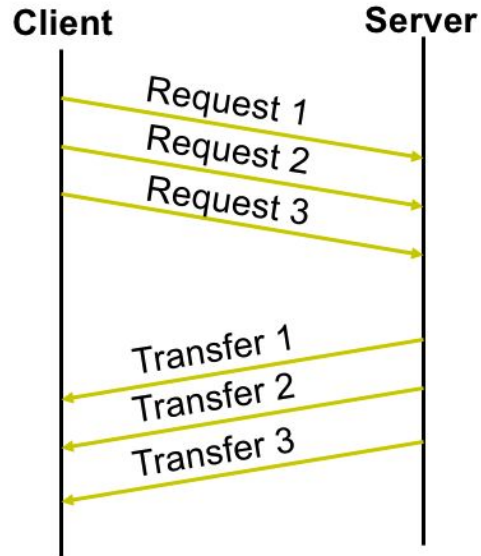


# Request Patterns

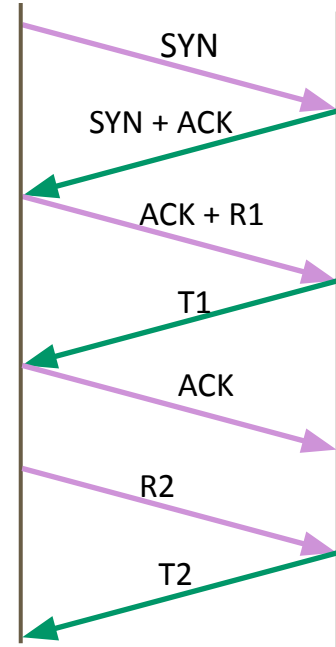
## Concurrent



## Pipelined



## Persistent



# Caching: How

- Idea: Replication
  - Replicate the content across multiple copies to reduce bottlenecks
- Implementation: content distribution networks (CDNs)
  - The client content provider modifies its content so that embedded URLs reference the new domains.
  - “Akamaize” content
  - e.g.: [http://www.netflix.com/tiger\\_king.jpg](http://www.netflix.com/tiger_king.jpg) might become [http://a1386.g.akamai.net/tiger\\_king.jpg](http://a1386.g.akamai.net/tiger_king.jpg)

# Caching: CNAMEs

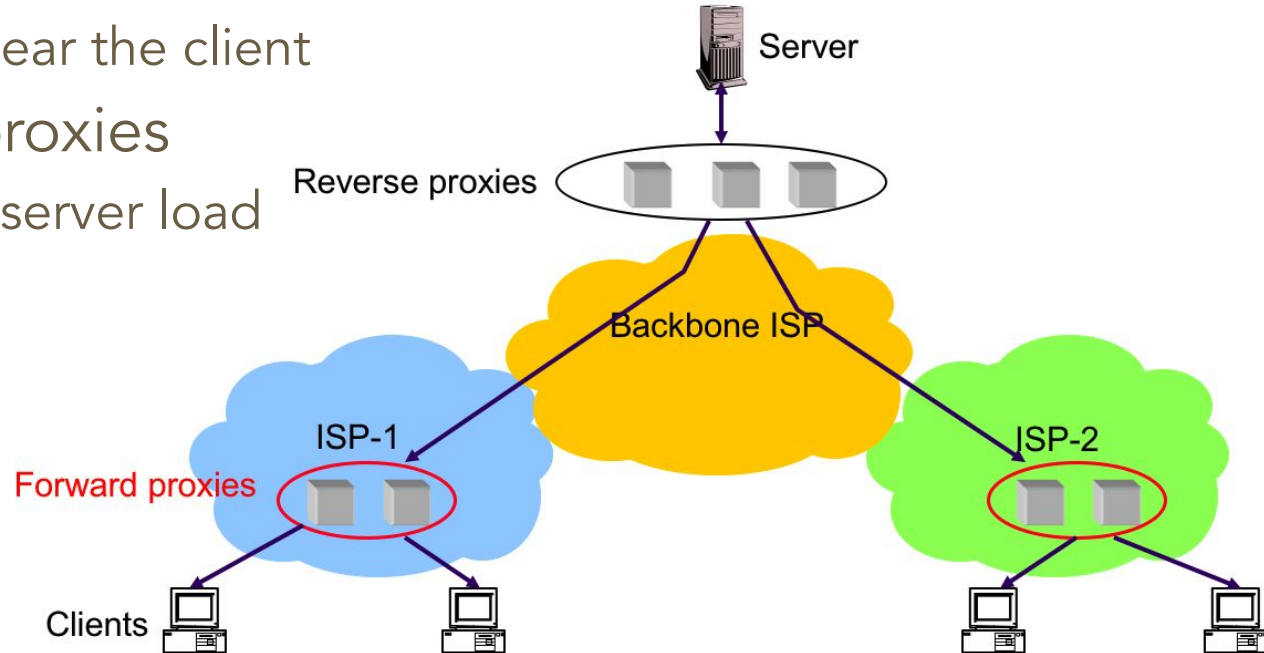
- Instead of using a weird URL, use CNAMEs!
  - List picture as [http://cdn.netflix.com/tiger\\_king.jpg](http://cdn.netflix.com/tiger_king.jpg)
  - The authoritative server for this is controlled by netflix
  - CNAME record aliases [cdn.netflix.com](http://cdn.netflix.com) to [a1386.g.akamai.net](http://a1386.g.akamai.net)
- Akamai handles sending to the closest server because [a1386.g.akamai.net](http://a1386.g.akamai.net) is under an Akamai nameserver
- Only change to the HTML is [http://www.netflix.com/tiger\\_king.jpg](http://www.netflix.com/tiger_king.jpg) became [http://cdn.netflix.com/tiger\\_king.jpg](http://cdn.netflix.com/tiger_king.jpg)

# Caching: Specifics in HTTP

- GET Request header:
  - `If-Modified-Since` – returns “not modified” if resource not modified since specified time
  - `Cache-Control: no-cache` – ignore all caches; always get resource directly from server (think force refresh)
- Response header:
  - `Cache-Control: max-age=<ttl>` – TTL: how long to cache the resource
  - `Cache-Control: no-store` – Don’t cache this
- When making request, if within the TTL, just load cached resource... otherwise, send with *if-modified*.
  - Server will either send a HTTP 304 (“Not Modified”) or HTTP 200 (changed, and here’s the new data)

# Caching: Where

- Forward Proxies
  - Cache near the client
- Reverse proxies
  - Reduce server load



# Worksheet

# Question 1

1. A URL only has information pertaining to the application layer (L7).
2. CDNs lower the latency for users.
3. HTTP 1.0 Requests are not human readable.
4. A server responds with only a header for an "If-Modified-Since" request on an object that has not been recently modified.
5. Pipelined connections are frequently used in practice.



## Question 2 Hints

- How long does it take to establish the connection?
- How long does it take to receive the webpage / media files?



## Question 2

- (1) Sequential requests with non-persistent TCP connections.

## Question 2

(2) Concurrent requests with non-persistent TCP connections.

## Question 2

(3) Sequential requests with a single persistent TCP connection.

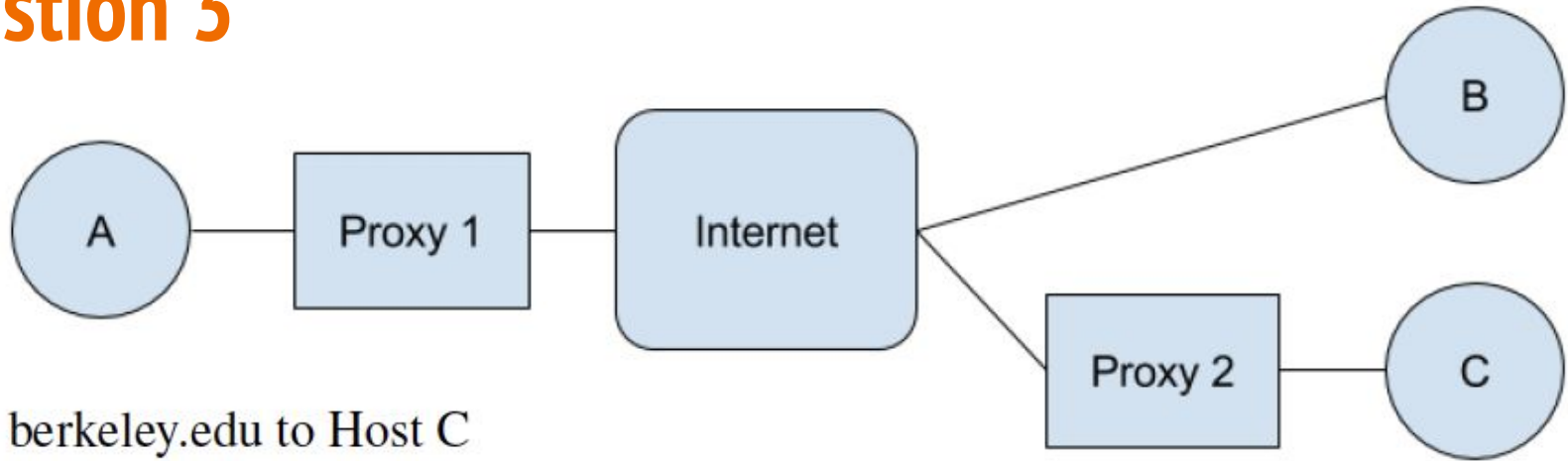
## Question 2

(4) Pipelined requests within a single persistent TCP connection.

## Question 2

- (5) We have been assuming that the throughput for sending media files is  $T$  for a single connection, and  $\frac{T}{n}$  for  $n$  concurrent connections. However, depending on the size of the media files, we can make more inferences about how fast we can send the media files. If the media files are very small, what kind of delay would dominate the time it would take to send them? What if the files are very large?

## Question 3

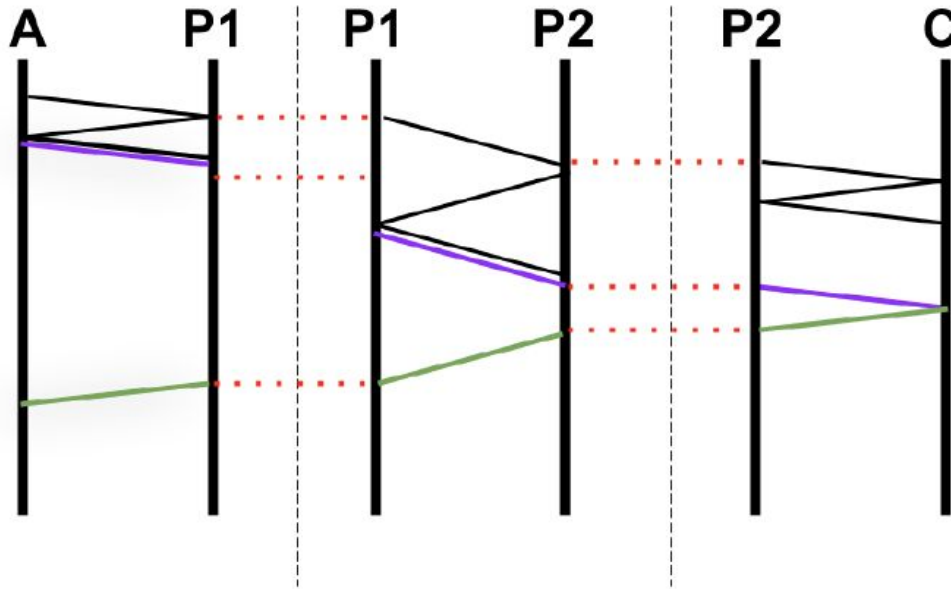


- berkeley.edu to Host C
- eecs.berkeley.edu to Host C
- stanford.edu to Host B
- mit.edu to Host B
- stanford.edu to Host B
- berkeley.edu to Host C

1. Sent Sequentially
2. Sent concurrently

# 3.1

$$\begin{aligned} &2(S + 2L_C + L) + 2(S + 2L_B + L) + 4L + S + 2(L_C - L) + L = \\ &2S + 4L_C + 2L + 2S + 4L_B + 2L + 4L + S + 2L_C - 2L + L = \\ &5S + 6L_C + 4L_B + 7L = \\ &5(5L + 2I) + 6(3L + I) + 4(2L + I) + 7L = \\ &25L + 10I + 18L + 6I + 8L + 4I + 7L = \\ &58L + 20I \end{aligned}$$



## 3.2

**Solution:** Since all the requests occur concurrently, there is no opportunity for caching by the proxies. Therefore the total time is just the max of all the times for the individual requests, which is the first request. It takes:

$$\begin{aligned}S + 2L_C + L &= \\5L + 2I + 2(3L + I) + L &= \\5L + 2I + 6L + 2I + L &= \\12L + 4I\end{aligned}$$



## Question 4: DNS

