

## STA 372-6 Homework 6: Dynamic Programming

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### 1. Code to find shortest path:

```
4 - #####
5 #Problem 1
6 |
7 # create weighted adjacency matrix
8 #      To: 1  2  3  4  5  6  7  8  9 10
9 # From: 1  0  0  3  2  0  0  0  0  0
10 #      2  0  0  4  2  4  0  0  0  0
11 #      3  0  0  0  0  0  3  2  0  0
12 #      4  0  0  0  0  0  4  0  2  5
13 #      5  0  0  0  0  0  0  2  2  0
14 #      6  0  0  0  0  0  0  0  0  3
15 #      7  0  0  0  0  0  0  0  0  4
16 #      8  0  0  0  0  0  0  0  0  2
17 #      9  0  0  0  0  0  0  0  0  3
18 #     10  0  0  0  0  0  0  0  0  0
19
20 # for simplicity, apply igraph package to aid in path calculation
21 require ("igraph")
22
23 adj_mat1<-matrix(c(0, 0, 3, 2, 0, 0, 0, 0, 0, 0,
24                   0, 0, 4, 2, 4, 0, 0, 0, 0, 0,
25                   0, 0, 0, 0, 0, 3, 2, 0, 0, 0,
26                   0, 0, 0, 0, 0, 4, 0, 2, 5, 0,
27                   0, 0, 0, 0, 0, 0, 2, 2, 0, 0,
28                   0, 0, 0, 0, 0, 0, 0, 0, 0, 3,
29                   0, 0, 0, 0, 0, 0, 0, 0, 0, 4,
30                   0, 0, 0, 0, 0, 0, 0, 0, 0, 2,
31                   0, 0, 0, 0, 0, 0, 0, 0, 0, 3,
32                   0, 0, 0, 0, 0, 0, 0, 0, 0, 0), nrow=10, byrow=T)
33
34 # create graph object
35 network<-graph_from_adjacency_matrix(adj_mat1, mode="directed", weighted=T)
36 plot(network, layout=layout_nicely(network, dim = 2))
37
38 # all paths from 1 to 10
39 paths1<-all_simple_paths(network, 1, 10)
40
41 # shortest path from 1 to 10
42 spath1<-shortest_paths(network, 1, 10)
43 spath1
44 # weight of shortest path from 1 to 10
45 distances(network, 1, 10)
46
47 # all paths from 2 to 10
48 paths2<-all_simple_paths(network, 2, 10)
49
50 # shortest path from 2 to 10
51 spath2<-shortest_paths(network, 2, 10)
52 spath2
53 distances(network, 2, 10)
54
```

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Output from above code:

```
> # shortest path from 1 to 10
> spath1<-shortest_paths(network, 1, 10)
> spath1
$vpath
$vp[1]
+ 4/10 vertices:
[1] 1 4 8 10

> # shortest path from 2 to 10
> spath2<-shortest_paths(network, 2, 10)
> spath2
$vp[1]
+ 4/10 vertices:
[1] 2 4 8 10
```

Thus the shortest path from 1 to 10 is  $1 \rightarrow 4 \rightarrow 8 \rightarrow 10$  with a total degree of 6, and the shortest path from 2 to 10 is  $2 \rightarrow 4 \rightarrow 8 \rightarrow 10$  also with a degree of 6.

### 2. Dynamic approach to car valuation

- States: time
- Actions: sell or operate
- Bellman Equation:

$$V(t) = \begin{cases} \min\{-resale_t + operating\ cost_t, 0\} & \text{if } t = 6 \\ \min\{-resale_t + operating\ cost_t, V(t+1)\} & \text{if } 0 < t < 6 \end{cases}$$

- Output from code:

Decisions	Sell	Sell	Sell	Sell	Sell	Operate
Values (Costs)	-\$13,400	-\$11,000	-\$6,400	-\$3,600	-\$800	0

Note: values of costs does not consider the original purchase price.

### 3. Travelling Salesman Variation

- Net gains from travelling matrix  $M$ :

		To:		
		Indianapolis	Bloomington	Chicago
From:	Indianapolis	\$120	\$110	\$150
	Bloomington	\$70	\$160	\$100
	Chicago	\$100	\$90	\$170

- States: (time, place) where time  $t \in \{1, 2, 3, 4\}$  and place  $c \in \{1, 2, 3\}$  corresponding to Indianapolis, Bloomington, and Chicago, respectively
- Actions: choose place  $c \in \{1, 2, 3\}$
- Bellman Equation:

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$$V(t, c) = \begin{cases} \max\{M_{2,1}, M_{2,2}, M_{2,3}\} & \text{if } t = 1 \\ \max\{M_{c(t-1),1}, M_{c(t-1),2}, M_{c(t-1),3}\} & \text{if } 1 < t < 4 \\ \max\{M_{c(t-1),1}\} & \text{if } t = 4 \end{cases}$$