# CSCI 200: Foundational Programming Concepts & Design Lecture 38



Searching Algorithms

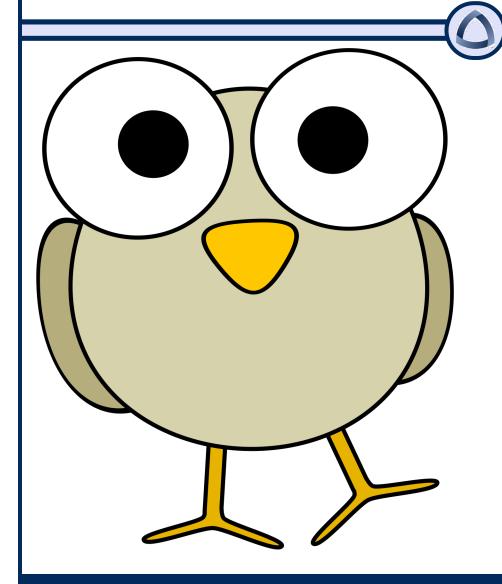
## Previously in CSCI 200

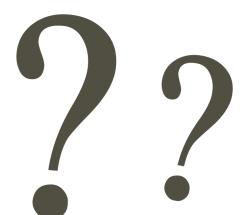
- Merge Sort
  - split()
  - merge()
- Recursion
  - Defined in terms of self
  - Solve smaller version of same problem
    - Divide-and-Conquer
    - Decrease-and-Conquer

## Sorting Complexities

Algorithm	Worst Case	Best Case	Average Case
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n^2)$	O( <i>n</i> )	$O(n^2)$
Bubble Sort	$O(n^2)$	O( <i>n</i> )	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$

#### Questions?





## Learning Outcomes For Today

- Explain how sorting a list affects the performance of searching for a value in a list.
- Implement linear and binary search.

## On Tap For Today

- Searching
  - Linear Search
  - Binary Search
- Practice

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## Searching

Different Types of Searches







Binary (ordered list)

## On Tap For Today

- Searching
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#### Linear Search

- No knowledge of list contents
  - No requirement of list ordering
    - List is unsorted (or sorted)

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Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
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Algorithm	Worst Case	Best Case	Average Case
Linear Search			

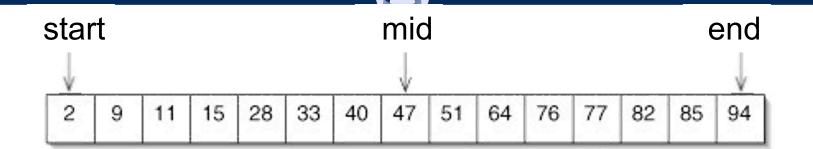
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Algorithm	Worst Case	Best Case	Average Case
Linear Search	O( <i>n</i> )	O(1)	O(n)

## On Tap For Today

- Searching
  - Linear Search
  - Binary Search
- Practice

• List must be sorted



- Values to keep track of
  - The list to search
  - Target value
  - Index to start search
  - Index to end search

#### Binary Search Pseudocode

- Examine middle element
  - If equals target
    - item found, return location
  - If greater than target
    - Ignore top half of list, continue search on bottom half
  - If less than target
    - Ignore bottom half of list, continue search on top half
- Repeat until
  - Target is found
  - start and end cross (not found)

#### Binary Search Pseudocode

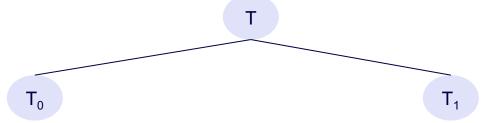
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- T(1) = 1
  - Follows the form of the "Master Theorem"

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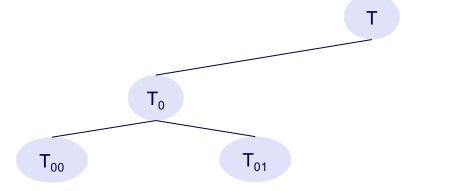


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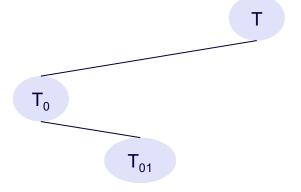
 $\mathsf{T}_0$ 



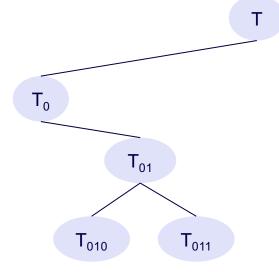
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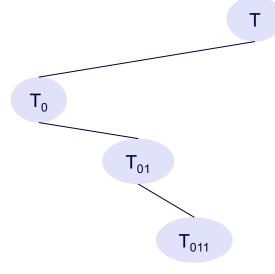
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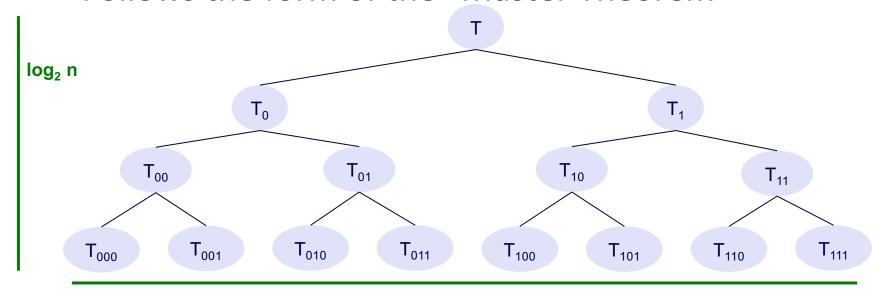
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• O(log n)

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Algorithm	Worst Case	Best Case	Average Case
Linear Search	O(n)	O(1)	O(n)
Binary Search	O(log n)	O(1)	O(log n)

Other search algorithms exist
 (just as other sort algorithms exist too)

- Values to keep track of
  - The array to search, Target value, Index to start search, Index to end search
- Examine middle element
  - If equals target
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  - If end < start</p>
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- Decrease-and-Conquer

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  - If less than target
    - Ignore bottom half of list, continue search on top half
- Decrease-and-Conquer
- So tempting to make recursive function...

### Recursive Binary Search

```
template<typename T>
int binary search(const List<T>& LIST, const T TARGET,
                  const int START POS, const int END POS) {
  const int MIDDLE POS = (END POS - START POS) / 2 + START POS;
  if(END POS < START POS)</pre>
    return -1;
  if(LIST[MIDDLE POS] == TARGET)
    return MIDDLE POS;
  if(LIST[MIDDLE POS] > TARGET)
    return binary search(LIST, TARGET, START POS, MIDDLE POS - 1);
  if(LIST[MIDDLE POS] < TARGET)</pre>
    return binary search(LIST, TARGET, MIDDLE POS + 1, END POS);
int targetPos = binary_search(myList, target, 0, myList.size() - 1);
```

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int targetPos = binary_search(myList, target, 0, myList.size() - 1);
```

Concern/danger of recursion?

#### The Call Stack

```
void print space(const int N) {
  for(int i = 0; i < N; i++) cout << " ";</pre>
void recurse( const int N ) {
  if(N \le 1)
    print space(N); cout << "Done!" << endl;</pre>
  } else {
    print_space(N); cout << "Start " << N << endl;</pre>
    recurse (N-1);
    print_space(N); cout << "End " << N << endl;</pre>
int main() {
  recurse (6);
  return 0;
```

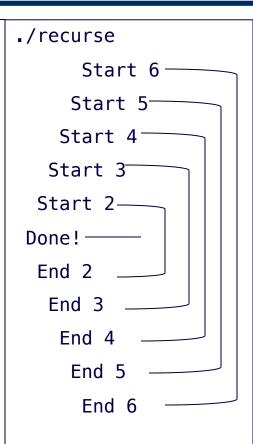
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```

```
./recurse
      Start 6
     Start 5
   Start 4
  Start 3
 Start 2
Done!
 End 2
  End 3
    End 4
     End 5
      End 6
```

### The Call Stack

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  return 0;
```



What is the run time?

```
template<typename T>
int binary search(const List<T>& LIST, const T TARGET,
                  const int START POS, const int END POS) {
    return binary search(LIST, TARGET, START POS, MIDDLE POS - 1);
```

• What is the run time? O(log n)

- What is the run time? O(log n)
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- What is the extra memory usage? O(log n)
- Is recursion necessary? No
  - Is there any backtracking going on? No
  - Any post-recursive work? No

## Iteration v Recursion

If Task() definition is of form
 A()
 Task()
 B()

use recursion

• If Task () definition is of form

A()

Task()

use iteration in place of tail recursion

## Iteration v Recursion

• If **Task()** definition is of form

```
A()
Task()
B()
```

use recursion

• If Task () definition is of form

```
A()
Task()
```

use iteration in place of tail recursion

• (NOTE: general rule of thumb...iteration can always replace recursion...)

## Iterative Binary Search

- Run Time is still O(log n)
- Extra memory usage is now O(1)

```
template<typename T>
int binary_search(const List<T>& LIST, const T TARGET) {
  int startPos = 0, endPos = myList.size() - 1;
  int targetPos = -1;
  while( true ) {
    // perform search...
  }
  return targetPos;
}
```

# Algorithm Complexities

- Scenario A
  - Unsorted list of n elements
  - Need to check if m values exist
  - Total Cost?

- Scenario B
  - Sort list of n elements
  - Need to check if m values exist
  - Total Cost?

# Algorithm Complexities

#### Scenario A

- Unsorted list of n elements O(1)
- Need to check if m values exist O(mn)
- Total Cost?  $O(mn) \rightarrow O(n^2)$

#### Scenario B

- Sort list of n elements  $O(n \log n)$
- Need to check if m values exist  $O(m \log n)$
- Total Cost?  $O(max(m,n) \log n)$  →  $O(n \log n)$

# On Tap For Today

- Searching
  - Linear Search
  - Binary Search
- Practice

## To Do For Next Time

- Rest of semester
  - W 11/29: 2D Lists + BFS/DFS
  - F 12/01: Stack & Queue

- M 12/04: Trees & Graphs, Quiz 6
- W 12/06: Exam Review
- R 12/07: Set6, SetXP, Final Project due

M 12/11: Final Exam