# CSCI 200: Foundational Programming Concepts & Design Lecture 41



**Trees & Graphs** 

## Previously in CSCI 200

- BFS and DFS --  $O(n^2)$ 
  - Same pseudocode except for Queue or Stack implementation
    - Explore neighbors recursively via iterative implementation
  - Reach same conclusion, potentially at varying speeds
  - "Graph" algorithms to discover paths between connected nodes

# On Tap For Today

Trees

Graphs

Practice

# On Tap For Today

Trees

Graphs

Practice

## Data Structure Operations

Operation		Array	Doubly-Linked List
Element Access		O(1)	O(n)
Traversal	Forwards	0/n	() (n)
Traversal	Backwards	O(n)	O( <i>n</i> )
	Front		O(1)
Add	Middle	O(n)	O(n)
	Back		O(1)
	Front	O(n)	O(1)
Delete	Middle		O(n)
	Back		O(1)
Search		O( <i>n</i> )	O( <i>n</i> )
Min / Max		O( <i>n</i> )	O(n)
Memory		n*sizeof(T) contiguous	n*(sizeof(T)+16) fragmented

# Algorithm Complexities

Algorithm	Worst Case	Best Case	Average Case
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n^2)$	O( <i>n</i> )	$O(n^2)$
Bubble Sort	$O(n^2)$	O( <i>n</i> )	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$

Algorithm	Worst Case	Best Case	Average Case
Linear Search	O( <i>n</i> )	O(1)	O(n)
Binary Search	O(log n)	O(1)	O(log n)

Data	Struct	ure	Opera	tions
				tioi 13

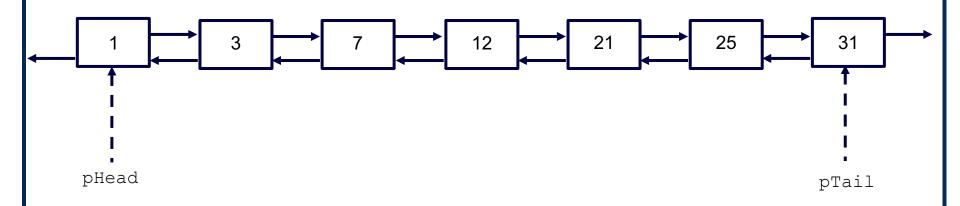
Operation		Array	Doubly-Linked List
Element Access		O(1)	O(n)
<b>-</b>	Forwards	O(n)	O(n)
Traversal	Backwards	O( <i>n</i> )	O(n)
	Front		O(1)
Add	Middle	O( <i>n</i> )	O(n)
	Back		O(1)
	Front		O(1)
Delete	Middle	O( <i>n</i> )	O(n)
	Back		O(1)
Sort		O(n log n)	O(n log n)
Search	Linear	O( <i>n</i> )	O(n)
Search	Binary		
Min / Max	Unsorted	O( <i>n</i> )	O(n)
	Sorted	O(1)	O(1)
Memory		n*sizeof(T) contiguous	n*(sizeof(T)+16) fragmented

CS @ Mines

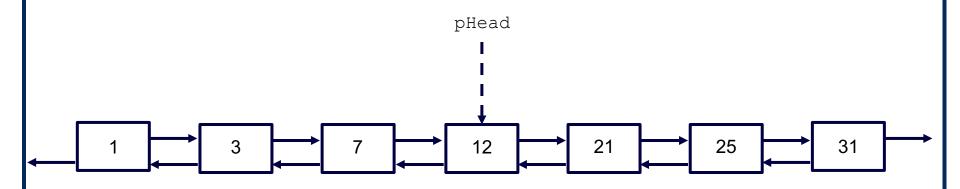
Data	Struct	ure	Opera	tions
				tioi 13

Operation		Array	Doubly-Linked List
Element Access		O(1)	O(n)
<b></b>	Forwards	O(n)	O(n)
Traversal	Backwards	O(n)	O(n)
	Front		O(1)
Add	Middle	O( <i>n</i> )	O(n)
	Back		O(1)
	Front		O(1)
Delete	Middle	O(n)	O(n)
	Back		O(1)
Sort		O(n log n)	O(n log n)
Search	Linear	O( <i>n</i> )	O(n)
Search	Binary	$O(\log n)$	$O(n \log n)$
Min / Max	Unsorted	O( <i>n</i> )	O(n)
	Sorted	O(1)	O(1)
Memory		n*sizeof(T) contiguous	n*(sizeof(T)+16) fragmented

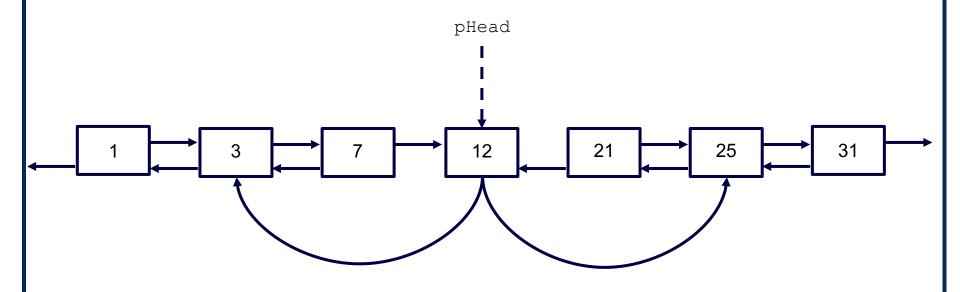
CS @ Mines



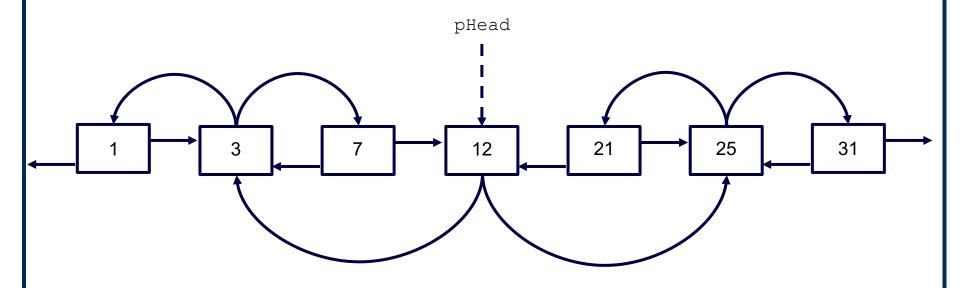
• Point to midpoint



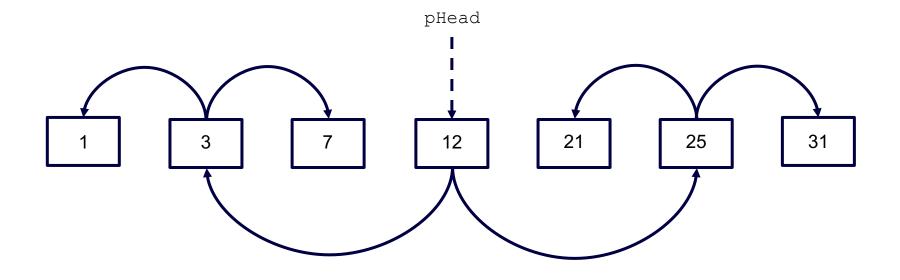
- Point to midpoint
  - Have that point to midpoints

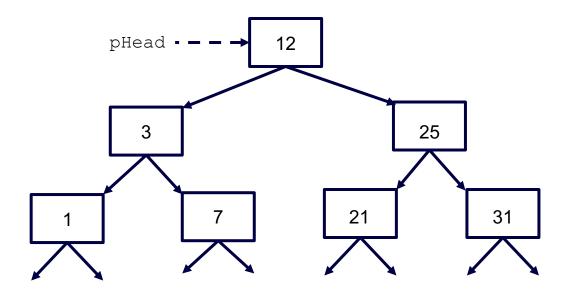


- Point to midpoint
  - Have that point to midpoints
    - And so on

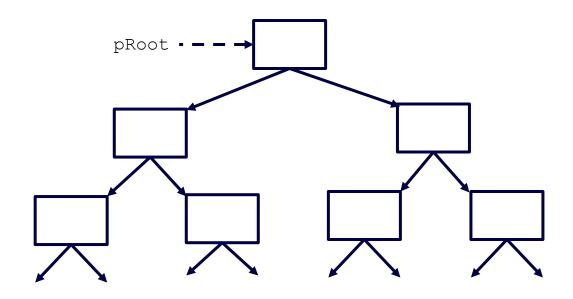


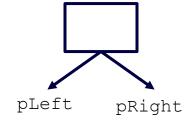
Cost of "Binary Search"?





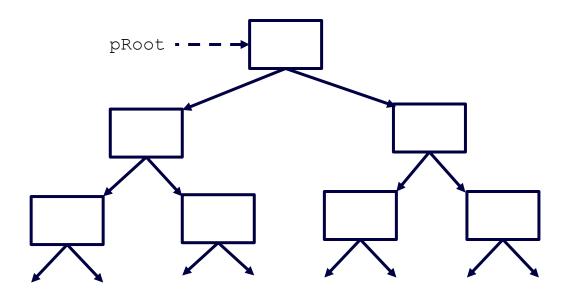
## **Binary Tree**

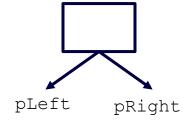




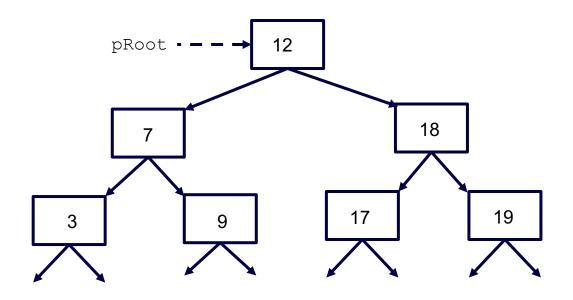
### **Binary Tree**

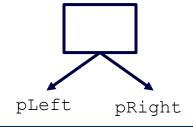
• When adding – apply insertion sort



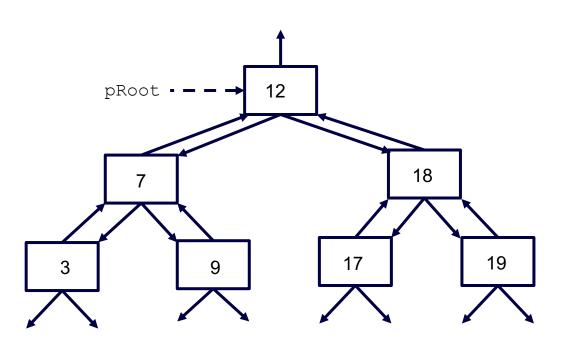


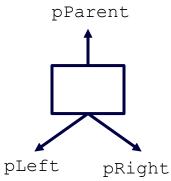
## Binary Search Tree





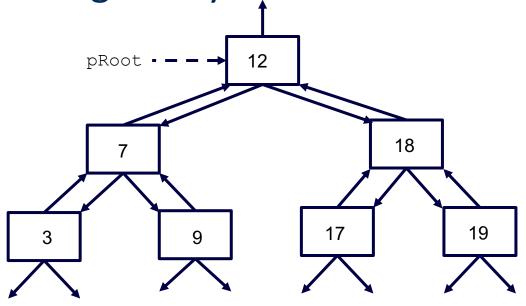
# Binary Search Tree



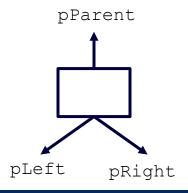


#### Binary Search Tree

When adding – may need to rebalance



 Note: For AXC - not doing a self-balancing tree

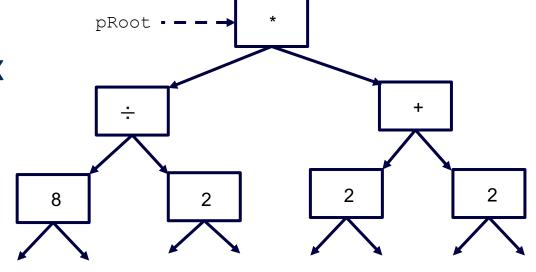


#### Tree Traversal

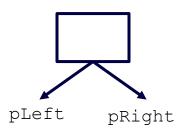
- Solve:  $8 \div 2(2+2)$ 
  - BODMAS / PEMDAS

#### Tree Traversal

- Infix
- Prefix
- Postfix



- Solve:  $8 \div 2(2+2)$ 
  - BODMAS / PEMDAS



Data Structura Oparationa					
Operation		Array	Doubly-Linked List	Balanced BST	
Element Acc	ess	O(1)	O(n)	O(log n)	
Traversal	Forwards	O(n)	O(n)	O(n)	
Iraversai	Backwards	O(n)			
	Front		O(1)		
Add	Middle	O( <i>n</i> )	O(n)	O(log n)	
	Back		O(1)		
	Front	O(n)	O(1)	O(log n)	
Delete	Middle		O(n)		
	Back		O(1)		
Sort		O(n log n)	O(n log n)	N/A	
Search	Linear	O(n)	O(n)	N/A	
Search	Binary	O(log n)	$O(n \log n)$	O(log n)	
Min / Max	Unsorted	O(n)	O(n)	Ollog p)	
	Sorted	O(1)	O(1)	O(log n)	
Memory		n*sizeof(T) contiguous	n*(sizeof(T)+16) fragmented	n*(sizeof(T)+16) fragmented	

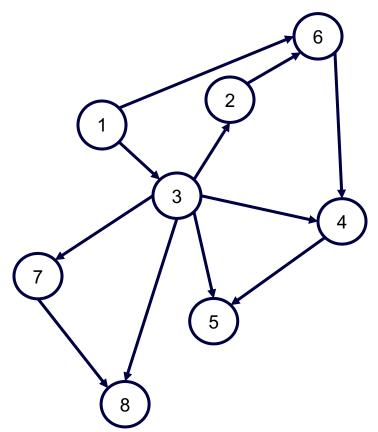
# On Tap For Today

Trees

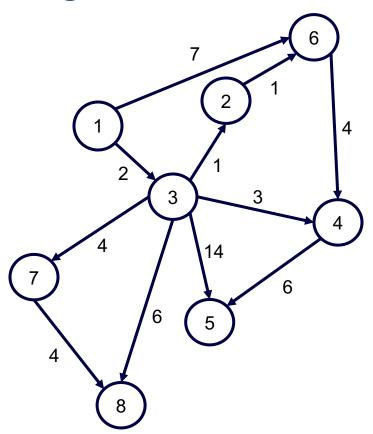
Graphs

Practice

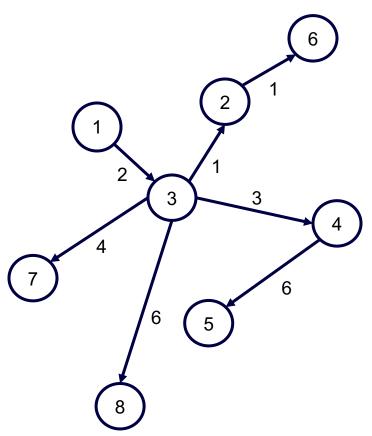
• Directed Acyclic Graph (DAG)



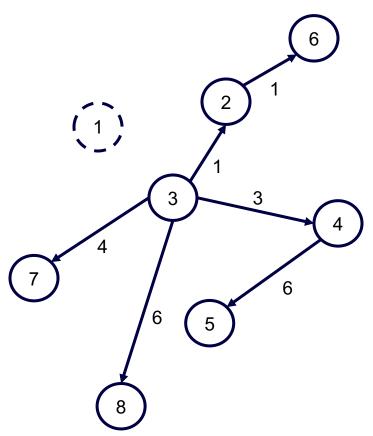
Weighted Directed Acyclic Graph



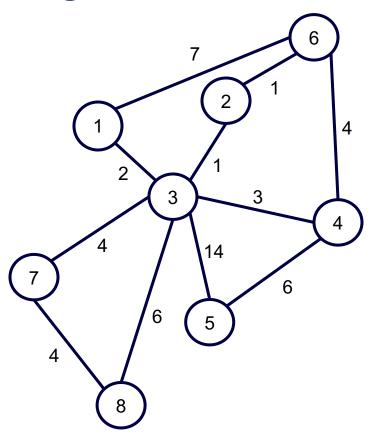
Shortest Path Tree from 1 to ? (Djikstra's)



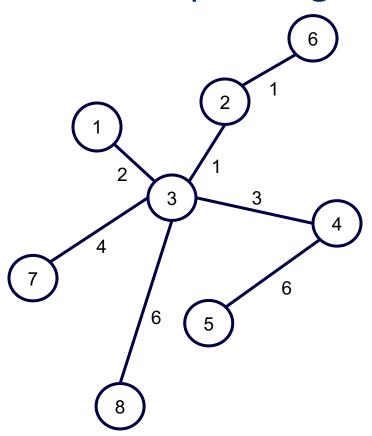
Shortest Path Tree from 3 to ? (Djikstra's)



Weighted Undirected Graph



Minimum Spanning Tree



# On Tap For Today

Trees

Graphs

Practice

#### To Do For Next Time

- Rest of semester
  - W 12/06: Exam Review
  - R 12/07: Set6, SetXP, Final Project due

- M 12/11: Final Exam

#### List Quiz

- Make Canvas Full Screen
- Access Code:
- 12 Minutes

