

# CENG 424

## Logic For Computer Science

Spring 2022-2023

### Assignment 3

Name Surname: Andaç Berkay Seval  
Student ID: 2235521

#### Answer for Q1

1.  $G((a \cup b) \Rightarrow \neg c)$
2.  $G\neg(c \wedge Xc)$
3.  $G(a \Rightarrow (Xb \Rightarrow (c \cup (a \wedge b))))$

#### Answer for Q2

1. Consider a transition system consists of two states  $s1$  and  $s2$  such that just  $a$  holds in  $s1$  and just  $b$  holds in  $s2$ .  $s1$  has a transition to  $s2$  and  $s2$  has a transition to  $s1$  ( $s1 \rightarrow s2$  and  $s2 \rightarrow s1$ ). If we take a path  $s1 \rightarrow s2 \rightarrow s1 \rightarrow s2 \rightarrow \dots$  we can see that  $G(a \vee b)$  holds. However,  $G_a$  does not hold since  $a$  does not hold in every state, just in  $s1$ . Also,  $G_b$  does not hold since  $b$  does not hold in every state, just in  $s2$ . Thus,  $G_a \vee G_b$  does not hold. Hence, equivalence does not hold.

2. We can write the right hand side of the equivalence as  $\neg G_a \wedge \neg G_b \equiv F\neg a \wedge F\neg b$ . Also, we know that  $F$  does not distribute over  $\wedge$ . Consider a transition system consists of two states  $s1$  and  $s2$ . Let atomic propositions of the transition system be  $x$  and  $y$  such that  $x \equiv \neg a$  and  $y \equiv \neg b$ . Just  $x$  holds in  $s1$  and just  $y$  holds in  $s2$ .  $s1$  has a transition to  $s2$  and  $s2$  has a transition to  $s1$  ( $s1 \rightarrow s2$  and  $s2 \rightarrow s1$ ). If we take a path  $s1 \rightarrow s2 \rightarrow s1 \rightarrow s2 \rightarrow \dots$  we can see that the path satisfies  $Fx \wedge Fy \equiv F\neg a \wedge F\neg b$  but not  $F(x \wedge y) \equiv F(\neg a \wedge \neg b)$ . Hence, equivalence does not hold.

3. Let  $\mathcal{M} = (S, \rightarrow, L)$  be a model.  $G_a$  means that for all  $i \geq 1$ ,  $\pi^i \models a$  for all paths  $\pi$  in the model  $\mathcal{M}$ .  $XFa$  means that  $\pi^2 \models Fa$  for all paths  $\pi$  in the model  $\mathcal{M}$ , which means there is some  $i \geq 2$  such that  $\pi^i \models a$  for all paths  $\pi$  in the model  $\mathcal{M}$ . Since  $G_a$  covers  $XFa$ ,  $G_a \wedge XFa \equiv G_a$ .

4. Let  $\mathcal{M} = (S, \rightarrow, L)$  be a model.  $Ga$  means that for all  $i \geq 1$ ,  $\pi^i \models a$  for all paths  $\pi$  in the model  $\mathcal{M}$ .  $Fb$  means that there is some  $j \geq 1$  such that  $\pi^j \models b$  for all paths  $\pi$  in the model  $\mathcal{M}$ .  $a \cup (\neg a \vee b)$  means that there is some  $k \geq 1$  such that  $\pi^k \models (\neg a \vee b)$  and for all  $l = 1, \dots, k-1$   $\pi^l \models a$ . If we look at the left hand side of the equivalence, if  $a$  holds in every state, then  $b$  will hold eventually. In the right hand side, in some state in the future,  $\neg a \vee b$  will hold and  $a$  holds until that state. For the left hand side, since  $b$  will hold eventually,  $\neg a \vee b$  will hold too even  $\neg a$  is false since  $Ga$  holds. Also, until  $\neg a \vee b$  will hold,  $a$  holds until that state. Hence  $Ga \Rightarrow Fb \equiv a \cup (\neg a \vee b)$ .

## Answer for Q3

1. It does not hold. Consider the path  $s1 \rightarrow s4 \rightarrow s2 \rightarrow s4 \rightarrow s2 \rightarrow \dots$ . According to this path,  $c$  will not eventually hold in all states ( $a \rightarrow b \rightarrow c \rightarrow b \rightarrow c \rightarrow \dots$ )
2. It holds. Since for all paths  $c$  holds infinitely often. After the starting state  $s1$  there is just one state where  $c$  does not hold which is  $s4$ . Also,  $s4$  does not have a transition to itself. Thus, it has to go to other states continually which include  $c$ . Hence  $GFc$  holds.
3. It does not hold since  $a$  does not hold in every state. Consider the path  $s1 \rightarrow s3 \rightarrow s4 \rightarrow s3 \rightarrow s4 \rightarrow \dots$ . According to this path,  $a$  holds just in the first state  $s1$ .
4. It holds. In order to  $X\neg c$  be true,  $s4$  is the only state where  $s1$  has a transition to since  $c$  holds in  $s3$  but not in  $s4$ .  $c$  holds in all the next states of  $s4$  ( $s2, s3, s5$ ). Thus,  $XXc$  also true. Hence,  $X\neg c \Rightarrow XXc$  holds. (I assume  $s1$  is the only initial state)
5. It holds. The formula states that there is some  $i \geq 1$  such that  $\pi^i \models G(b \vee c)$  and  $j = 1, \dots, i-1$   $\pi^j \models a$ . After the initial state  $s1$ ,  $b$  or  $c$  hold in every state ( $b$  and  $c$  hold in  $s3$ ,  $b$  holds in  $s4$ ).  $c$  holds in  $s2$ ,  $b$  and  $c$  hold in  $s5$  which are reachable from  $s4$ . We can say that after the initial state  $s1$ ,  $b \vee c$  becomes true in all states of all paths and  $a$  holds in  $s1$ . Hence,  $a \cup G(b \vee c)$  holds. (I assume  $s1$  is the only initial state)

## Answer for Q4

1.  $\forall G \forall F(a \wedge b)$
2.  $\forall G \neg(a \wedge b)$
3.  $\forall G(b \Rightarrow \exists Fa)$

## Answer for Q5

1.  $\text{Sat}(\phi_1) = \{s0, s1, s2, s3, s4\}$ .  $\forall G \forall Fb$  states that for all paths globally, eventually  $b$  will hold for all paths, so  $b$  holds infinitely often in any path along the path. If we look at  $TS'$ , we can see that for every individual state (we choose them to be initial state),  $\phi_1$  holds. Hence,  $TS' \models \phi_1$  holds since  $s0$  (initial state) is in  $\text{Sat}(\phi_1)$ .

2.  $\text{Sat}(\phi_2) = \{\}$ .  $\forall G \exists F a$  states that for all paths globally, eventually  $a$  will hold for some path, so  $a$  holds infinitely often in some path along all paths. If we look at  $\text{TS}'$ , we can see that none of the states holds  $\phi_2$ . After reaching states  $s_3$  or  $s_4$ ,  $\text{TS}'$  continues without changing its state and in these states just  $b$  holds. Thus,  $\text{TS}' \models \phi_2$  does not hold.

3.  $\text{Sat}(\phi_3) = \{s_1, s_2\}$ .  $\forall(a \text{ U } b)$  states that for all paths,  $a$  holds until  $b$  holds. If we look at the  $\text{TS}'$ , we can see that for the states  $s_1$  and  $s_2$ ,  $a$  holds in both.  $s_1$  has just one transition which is to  $s_2$  and  $b$  holds in  $s_2$ .  $s_2$  has just one transition which is to  $s_3$  and  $b$  holds in  $s_3$ . However,  $\text{TS}' \models \phi_3$  does not hold since  $s_0$  (initial state) is not in  $\text{Sat}(\phi_3)$ .

4.  $\text{Sat}(\phi_4) = \{s_0, s_1, s_2, s_3, s_4\}$ .  $\exists X(\forall G b)$  states that there is some path in the next state,  $b$  globally hold for all paths along the path. If we look at the  $\text{TS}'$ , we can see that for all states as initial states,  $\phi_4$  holds. Therefore,  $\text{TS}' \models \phi_4$  holds since  $s_0$  (initial state) is in  $\text{Sat}(\phi_4)$ .

5.  $\text{Sat}(\phi_5) = \{s_1, s_2\}$ .  $\forall G \forall(a \text{ U } b)$  states that for all paths globally,  $a$  holds until  $b$  holds for all paths along the path. We can see that for states  $s_1$  and  $s_2$  as initial states,  $a$  holds in both.  $s_1$  has just one transition which is to  $s_2$  and  $b$  holds in  $s_2$  and  $s_2$  has just one transition which is to  $s_3$  and  $b$  holds in  $s_3$  and  $s_3$  has just one transition to itself so  $b$  holds continuously in  $s_3$ . However,  $\text{TS}' \models \phi_5$  does not hold since  $s_0$  (initial state) is not in  $\text{Sat}(\phi_5)$ .