

University of California, Berkeley – College of Engineering

Department of Electrical Engineering and Computer Sciences

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CS61C MIDTERM 1



After the exam, indicate on the line above where you fall in the emotion spectrum between “sad” & “smiley”...

<i>Last Name</i>	
<i>First Name</i>	
<i>Student ID Number</i>	
<i>CS61C Login</i>	cs61c-
<i>The name of your SECTION TA (please circle)</i>	David Donggyu Fred Jeffrey Martin Nolan Sagar Shreyas William
<i>Name of the person to your Left</i>	
<i>Name of the person to your Right</i>	
<i>All the work is my own. I had no prior knowledge of the exam contents nor will I share the contents with others in CS61C who have not taken it yet. (please sign)</i>	

Instructions (Read Me!)

- This booklet contains 7 numbered pages including the cover page. **The back of each page is blank and can be used for scratch-work, but will not be looked at for grading.** (i.e. the sides of pages without the printed “SID: _____” header will not even be scanned into gradescope).
- Please turn off all cell phones, smartwatches, and other mobile devices. Remove all hats & headphones. Place your backpacks, laptops and jackets under your seat.
- You have 80 minutes to complete this exam. The exam is closed book; no computers, phones, or calculators are allowed. You may use one handwritten 8.5”x11” page (front and back) of notes in addition to the provided green sheet.
- There may be partial credit for incomplete answers; write as much of the solution as you can. We will deduct points if your solution is far more complicated than necessary. When we provide a blank, please fit your answer within the space provided. “IEC format” refers to the mebi, tebi, etc prefixes.

	Q1	Q2	Q3	Q4	Q5	Q6	Q7	Total
Points Possible	10	12	6	12	18	16	16	90
Points Earned								

Q1: Number Representation (10 points)

1) Convert the following 8-bit two's-complement numbers from hexadecimal to decimal:

0x80 = -128

0xFF = -1

0x0F = 15

2) For two n -bit numbers, what is the difference between the largest unsigned number and the largest two's-complement number? In other words, what is **MAX_UNSIGNED_INT** – **MAX_SIGNED_INT**? Write your answer in terms of n .

$$2^n - 1 - (2^{n-1} - 1) = 2^{n-1}$$

3) Fill in the blanks to return the largest positive number a 32-bit two's-complement number can represent.

```
unsigned int max_twos() {
    return ((1 << 31) - 1);
}
```

4) Consider a new type of notation for representing signed numbers, *biased* notation. The formula for obtaining the value from a number written in biased notation is:

$$\text{value} = \text{value_as_unsigned} - b$$

Where b is a constant called the *bias*. Example with 4 bits and a bias of 4:

$$\begin{aligned} 0b0011 &= 3 - 4 = -1 \\ 0b1110 &= 14 - 4 = 10 \end{aligned}$$

If we wanted an n -bit biased system to represent the same range as two's complement numbers, what is the value of b ?

The most negative number needs to be -2^{n-1} . The most negative number in biased notation is $0 - b$ so our bias is 2^{n-1} .

Q2: Pointers and Memory (12 points)

1) Assume you are given an int array `arr`, with a pointer `p` to its beginning:

```
int arr[] = {0x61c, 0x5008, 0xd, 0x4, 0x3, 0x4ffc};
int *p = arr;
```

Suppose `arr` is at location `0x5000` in memory, i.e., the value of `p` if interpreted as an integer is `0x5000`. To visualize this scenario:

0x61c	0x5008	0xd	0x4	0x3	0x4ffc
↑ <i>arr[0]</i>		...			<i>arr[5]</i>
p					

Assume that integers and pointers are both 32 bits. What are the values of the following expressions? If an expression may cause an error, write “Error” instead.

- a) `*(p+3)` = 0x4
- b) `p[4]` = 0x3
- c) `*(p+5) + p[3]` = 0x5000
- d) `*(int*)(p[1])` = 0xd(13)
- e) `*(int*)(*(p+5))` = error(out of bounds)

2) Consider the following code and its output. Fill in the blanks.

```
void foo1( __int*__ a, int n) {
    int i;
    for (i = 0 ; i < n ; i++) {
        (*(a+i)) += 3;
    }
}

void foo2( __int**__ p) { p++; }

int main() {
    int x = __4__ ;
    int a[] = {1, 2, 3, 4, 5};
    int *p = &a[1];
    foo1(a, sizeof(a) / sizeof(int));
    foo2(&p);
    printf("%d, %d, %d\n", a[1], *(++p), a[x]);
}
```

The output of this code is:

__5__ , __6__ , 8

Q3: C Memory Model (6 points)

For each of the following functions, answer the questions below in the corresponding box to the right:

- 1) Does this function return a usable pointer to a string containing "asdf"?
- 2) Which area of memory does the returned pointer point to?
- 3) Does this function leak memory?

You may assume that `malloc` calls will always return a non-NULL pointer.

```
char * get_asdf_string_1() {
    char *a = "asdf";
    return a;
}
```

get_asdf_string_1
1)
2)
3)

```
char * get_asdf_string_2() {
    char a[5];
    a[0]='a';
    a[1]='s';
    a[2]='d';
    a[3]='f';
    a[4]='\0';
    return a;
}
```

get_asdf_string_2
1)
2)
3)

```
char * get_asdf_string_3() {
    char * a = malloc(sizeof(char) * 5);
    a = "asdf";
    return a;
}
```

get_asdf_string_3
1)
2)
3)

```
char * g = "asdf";

char * get_asdf_string_4() {
    return g;
}
```

get_asdf_string_4
1)
2)
3)

Q4: Linked Lists (12 points)**1) Fill out the declaration of a singly linked linked-list node below.**

```
typedef struct node {
    int value;
    __struct node*__ next; // pointer to the next element
} sll_node;
```

2) Let's convert the linked list to an array. Fill in the missing code.

```
int * to_array(sll_node *sll, int size) {
    int i = 0;
    int *arr = __malloc(size * sizeof(int))__;
    while (sll) {
        arr[i] = __sll->value____;
        sll = __sll->next____;
        __i++____;
    }
    return arr;
}
```

3) Finally, complete delete_even() that will delete every second element of the list. For example, given the lists below:Before: Node 1 → Node 2 → Node 3 → Node 4After: Node 1 → Node 3**Calling delete_even() on the list labeled “Before” will change it into the list labeled “After”. All list nodes were created via dynamic memory allocation.**

```
void delete_even(sll_node *sll) {
    sll_node *temp;
    if (!sll || !sll->next) return;
    temp = __sll->next____;
    sll->next = __temp->next (or sll->next->next)____;
    free(__temp____);
    delete_even(__sll->next____);
}
```

Q5: MIPS with FUNctions (18 points)

The function `countChars(char *str, char *target)` returns the number of times characters in `target` appear in `str`. For example:

```
countChars("abc abc abc", "a") = 3
countChars("abc abc abc", "ab") = 6
countChars("abc abc abc", "abcd") = 9
```

The C code for `countChars` is given to you in the box on right. The helper function `isCharInStr(char *target, char c)` returns 1 if `c` is present in `target` and 0 if not.

```
int countChars(char *str, char *target) {
    int count = 0;
    while (*str) {
        count += isCharInStr(target, *str);
        str++;
    }
    return count;
}
```

Finish the implement of `countChars` in TAL MIPS below. You may not need every blank.

`countChars:`

```
addiu $sp, $sp, ___-16___
__sw $ra, 0($sp)_____ # Store onto the stack if needed
__sw $s0, 4($sp)_____
__sw $s1, 8($sp)_____
__sw $s2, 12($sp)_____
```

```
_____
addiu $s0, $zero, 0      # We'll store the count in $s0
addiu $s1, $a0, 0
addiu $s2, $a1, 0
```

`loop:`

```
addiu $a0, $s2, 0
__lb $a1, 0($s1)_____
beq __$a1, $zero, done___
jal isCharInStr
__addu $s0, $s0, $v0_____
__addiu $s1, $s1, 1_____
__j loop_____
```

`done:`

```
__addiu $v0, $s0, 0_____ # Load from the stack if needed
__lw $ra, 0($sp)_____
__lw $s0, 4($sp)_____
__lw $s1, 8($sp)_____
__lw $s2, 12($sp)_____
_____
addiu $sp, $sp, ___16___
jr $ra
```

Q6: MIPS Instruction Formats (16 points)

Convert the following TAL MIPS instructions into their machine code representation (binary format) or vice versa. For rows where you convert instructions to machine code, we've provided boxes to the right that you should fill in with the appropriate fields (in binary):

<u>MIPS</u>	<u>Machine Code</u>						
foo_bar:	0b000000001000000000100000100001						
loop: beq \$a1 \$0 end	<table border="1"><tr><td></td><td></td><td></td><td></td></tr></table>						
	0b00000000000000010000100000100000						
j loop	0b00001000000000000000000000000001						
end: jr \$ra	<table border="1"><tr><td></td><td></td><td></td><td></td><td></td><td></td></tr></table>						

Q7: MIPS Addressing Modes (16 points)

We have a function that, when given a branch instruction, returns the number of bytes that the Program Counter (PC) would change by, i.e. (PC_of_branch_target - PC_of_branch_instruction).

```
branchAmount(branch_inst):
    calculate the instruction offset from branch_inst
    convert the offset to byte addressing
    return PC_of_branch_target - PC_of_branch_instruction
```

Write branchAmount in TAL MIPS (no pseudoinstructions) .You may not need all the blanks. Assume that register \$a0 contains a valid branch instruction.

```
branchAmount:
    andi $t0, $a0, 0x8000      # Mask out a certain bit
    bne $t0, $zero, label1
    __andi $v0, $a0, 0xFFFF__
    _____
    j label2
label1:
    __lui $t1, 0xFFFF__
    _____
    or $v0, $a0, $t1
label2:
    sll __$v0__, __$v0__, __2__ # Convert to byte addressing
    __addiu $v0, $v0, 4__
label3:
    jr $ra
```