



Iterated Revelation: How to incentivize experts to reveal novel actions

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Ph.D student	supervisor	prob. of success

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Ph.D student	supervisor	prob. of success	research ideas
homeowner	architect	???	house design

Project Choice

Armstrong and Vickers (2010), Guo and Shmaya (2023) study project choice:

- ◊ A manager (decision maker) is uncertain about which projects are feasible
- ◊ A subordinate (expert) makes recommendations
- ◊ The manager commits to a selection rule
 - ◊ how to choose a project from the subordinate's recommendations

But what if ex-ante commitment to a selection rule is infeasible?

- ◊ Unawareness: a student unaware of state-of-the-art research ideas
- ◊ Inexpressibility: impractical to express every possible house design
- ◊ Enforceability: A regulator may be unable to make reasonable commitments

Literature

- ◊ Delegation / Project Choice
 - ◊ Holmstrom (1980); Armstrong and Vickers (2010); Guo and Shmaya (2023)
- ◊ Incomplete Contracting / Unawareness in Contracting
 - ◊ Grossman and Hart (1986); Maskin and Tirole (1999); Tirole (2009); Hart (2017); Piermont (2017); Lei and Zhao (2021); Francetich and Schipper (2021)
- ◊ Strategic Information Transmission
 - ◊ Milgrom (1981), Crawford and Sobel (1982); Seidmann and Winter (1997); Aumann and Hart (2003); Chakraborty and Harbaugh (2010)
- ◊ Evidentiary disclosure
 - ◊ Dye, 1985; Green and Laffont, 1986; Grossman and Hart, 1986; Bull and Watson, 2007; Ben-Porath et al., 2019



Model





The environment is described by

\mathcal{A} – a (compact) set of actions

(u_d, u_e) – (continuous) utility functions $\mathcal{A} \rightarrow \mathbb{R}$

\mathcal{R} – a collection of non-empty (compact) subsets of \mathcal{A} such that

- ◊ Exists $r^\dagger \in \mathcal{R}$ such that $r^\dagger \subseteq r$ for all $r \in \mathcal{R}$
- ◊ for $r \in \mathcal{R}$, there are a finite $r' \in \mathcal{R}$ with $r' \subseteq r$.

Revelation Types

- ◊ A **revelation type** $r \in \mathcal{R}$ is a set of actions / projects that an agent can express
- ◊ Say that r is **more expressive** than r' , if $r' \subseteq r$
- ◊ r^\dagger is the **dm**'s type

Key Assumptions



Let the **expert** be of type $r \in \mathcal{R}$:

Voluntary Disclosure: the **expert** can always masquerade as a type $r' \subseteq r$

Information Spillover: Only subsets of actions $r' \in \mathcal{R}$ can be revealed

Hard Evidence: If the **expert** reveals r , then any $a \in r$ is ‘real’

Selection Rules

An **selection rule** is a function from types to actions:

$$\begin{array}{ccc} f: & r & \mapsto a \\ & \cap & \cap \\ \mathcal{R} & \rightarrow & r \end{array}$$

- ◊ Selection rules are *direct mechanisms*
- ◊ These are the object of study in the project choice literature
- ◊ Inexpressibility precludes direct mechanisms / revelation principle

Call f **monotone** if ex's payoff is monotone in her type

$$u_{\text{e}}(f(r')) \leq u_{\text{e}}(f(r)) \quad (1)$$

whenever $r' \subseteq r$, and **strongly monotone** if in addition (2) holds strictly whenever $f(r) \neq f(r')$.

- ◊ If direct mechanisms existed, monotonicity is incentive compatibility
- ◊ Direct mechanisms don't exist: even with monotonicity, there need not be any 'strategic' way of enacting a selection rule.

Call f **efficient** if outcome is Pareto undominated:

for all $r \supseteq r_{\textcolor{blue}{d}}$, is no $a \in r$ such that

$$u_{\textcolor{blue}{d}}(a) \geq u_{\textcolor{blue}{d}}(f(r)) \quad \text{and} \quad u_{\textcolor{red}{e}}(a) \geq u_{\textcolor{red}{e}}(f(r)), \quad (2)$$

with at least one inequality holding strictly.

So, what can we do without ex-ante commitment to a selection rule?

Example

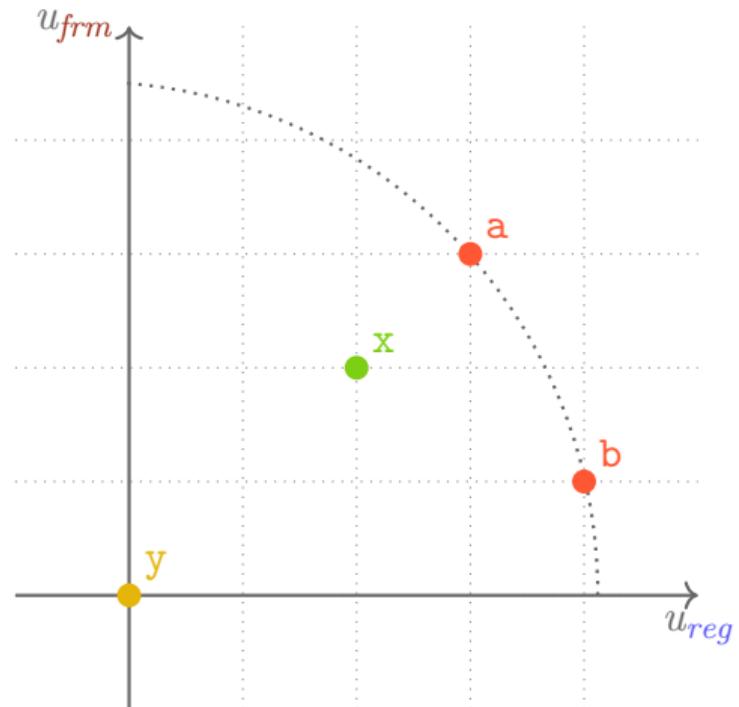
- ◊ A **regulator** (the **decision maker**) is evaluating mergers (i.e., projects) :
 - ◊ can only choose a merger structure it is aware of
- ◊ A **firm** (the **expert**) may be aware of novel ways of structuring a merger
- ◊ choice of merger structure determines payoffs for both players
 - ◊ the **firm** cares about producer welfare
 - ◊ the **regulator** cares about consumer welfare foremost, but
 - ◊ also cares about efficiency

Example



- ◊ Each merger yields (x_{reg}, x_{frm}) :
 - ◊ x_{reg} is regulator's payoff (consumer welfare)
 - ◊ x_{frm} is firms's (producer welfare)
- ◊ There are four ways to structure the merger:

$\mathbf{x} = (2, 2)$	$\mathbf{y} = (0, 0)$
$\mathbf{a} = (3, 3)$	$\mathbf{b} = (4, 1)$
- ◊ \mathbf{a} and \mathbf{b} must be revealed together



Example



- ◊ Initially the **regulator** is aware of

$$\mathbf{x} = (2, 2) \quad \mathbf{y} = (0, 0)$$

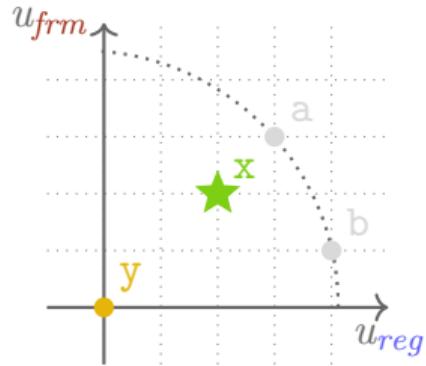
and the **firm** is also aware of

$$\mathbf{a} = (3, 3) \quad \mathbf{b} = (4, 1)$$

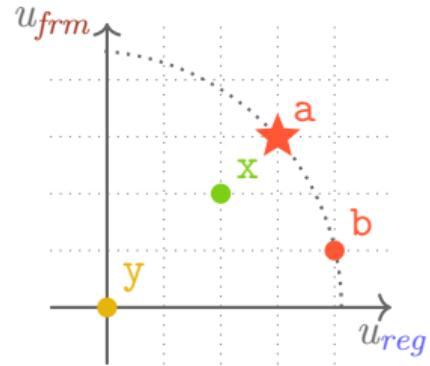
- ◊ Since **a** and **b** must be revealed together

$$\mathcal{R} = \{\{\mathbf{y}, \mathbf{x}\}, \{\mathbf{y}, \mathbf{x}, \mathbf{a}, \mathbf{b}\}\}$$

A monotone and efficient selection rule:



$$f(\{y, x\}) = x$$

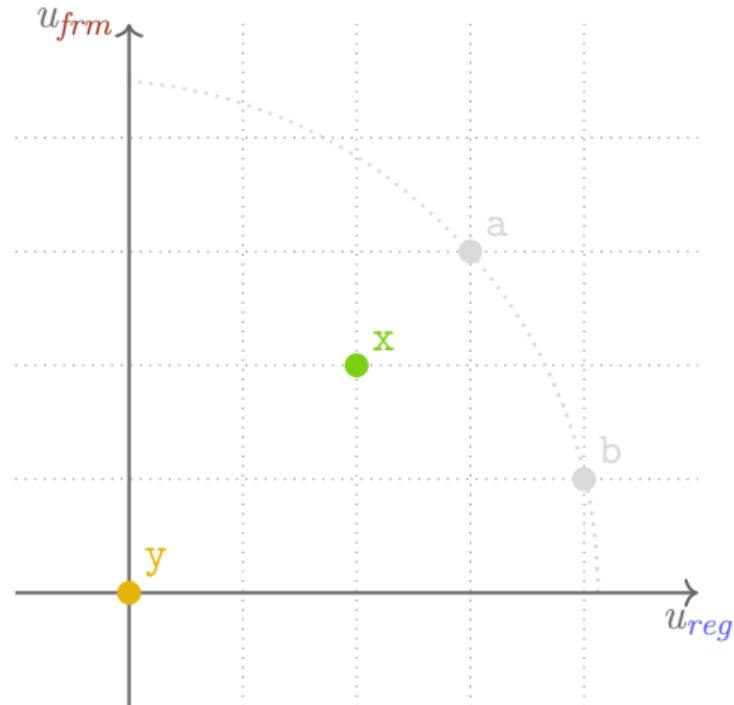


$$f(\{y, x, a, b\}) = a$$

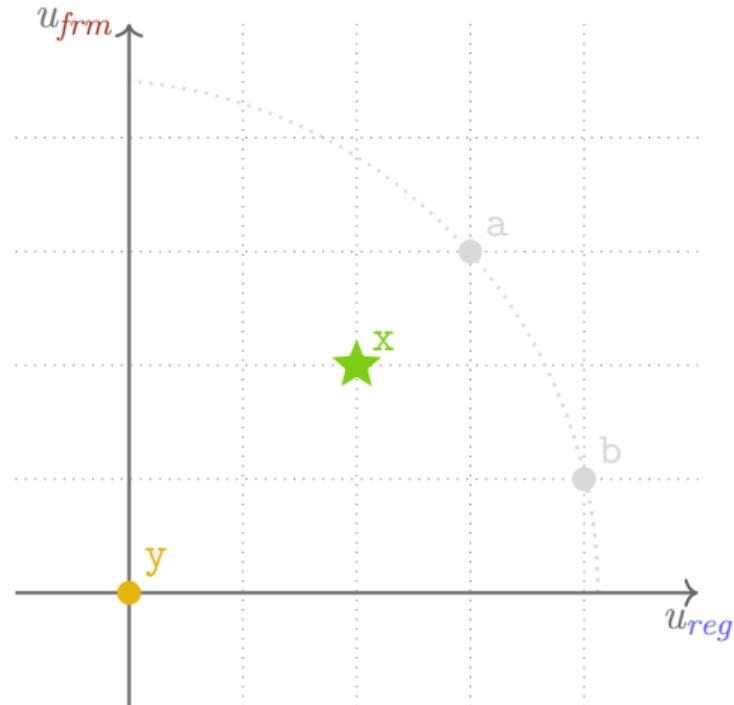
Example A

- ◊ If the **firm** had full control over what to reveal: simply reveal $\mathbf{a} = (3, 3)$
- ◊ However, not all mergers can be independently revealed:
 - ◊ Revealing one merger in a ‘class’ reveals the existence of the whole class, etc
- ◊ **Regulator** cannot commit to \mathbf{a} (or \mathbf{b}) until it is revealed
 - ◊ Cannot express the direct selection rule above

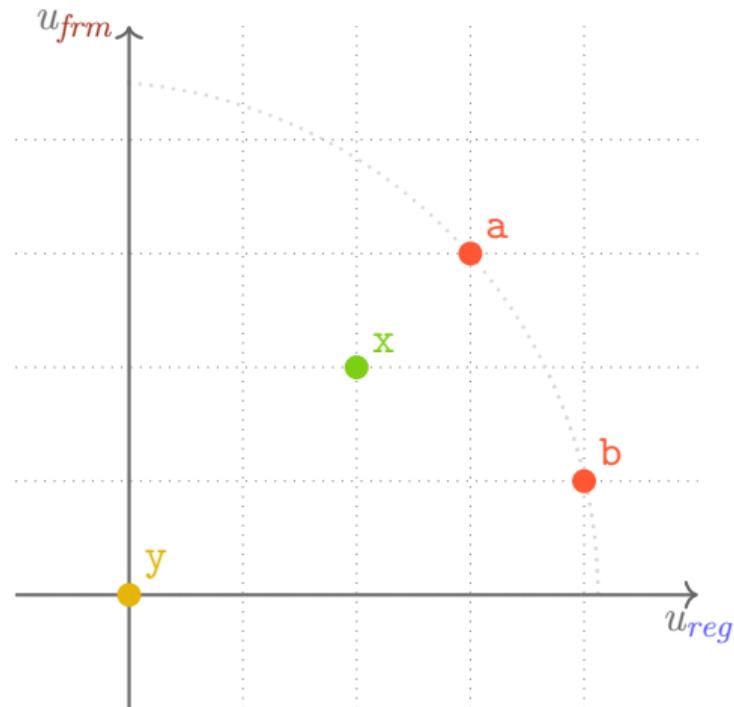
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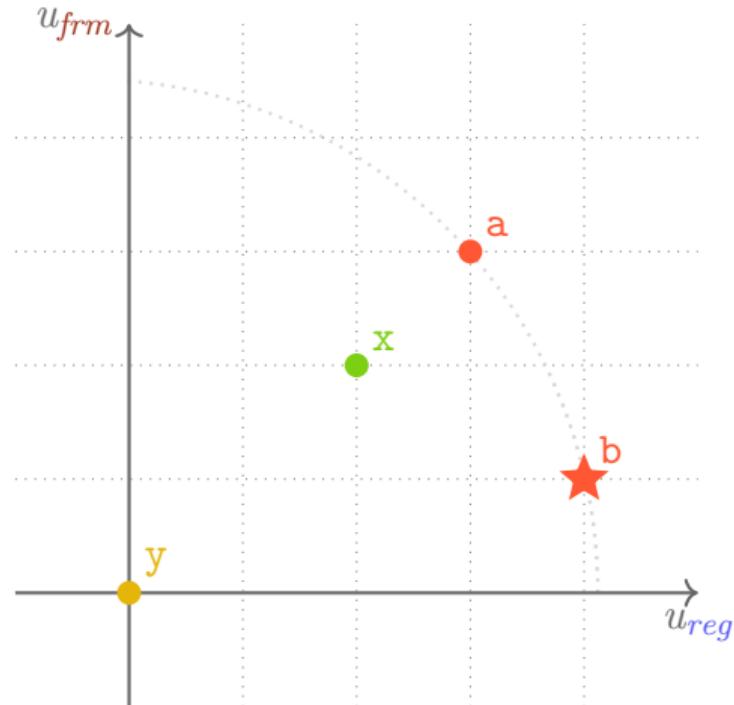
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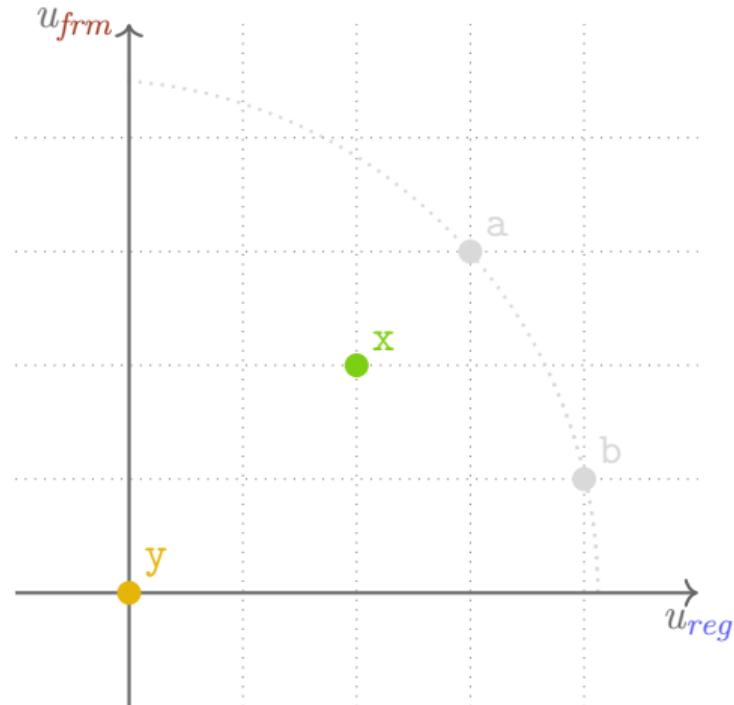
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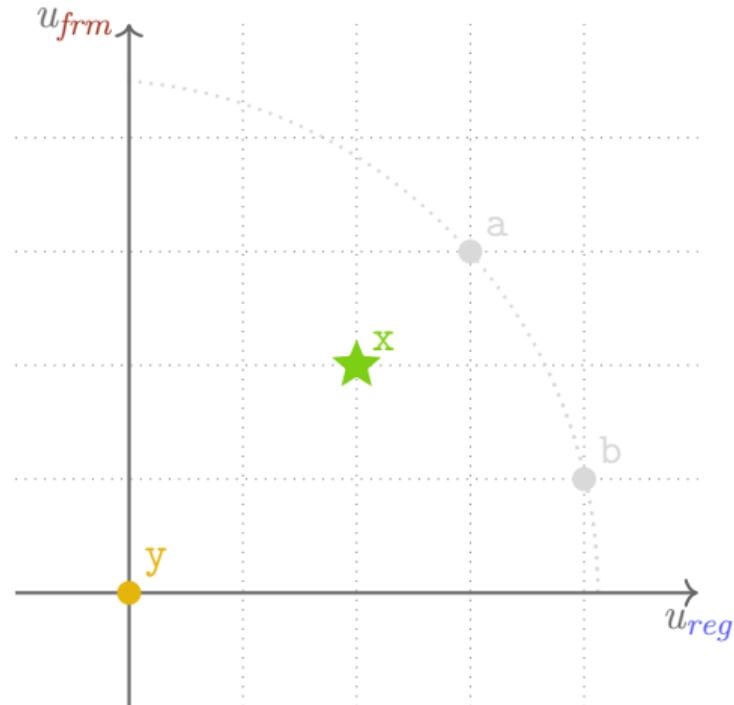
Example A

- ◊ Since the **firm** prefers **x** to **b**, she would choose not to reveal.
- ◊ This is Pareto Inefficient: **a** dominates **x**
- ◊ What if **regulator** and **firm** can create the following contract (before revelation):
 - ◊ shortlist an ‘outside option’ (that the **regulator** is aware of)
 - ◊ this can be re-negotiated after revelation
 - ◊ the **regulator** can propose a new merger, but the **firm** can veto (therefore implement outside option)

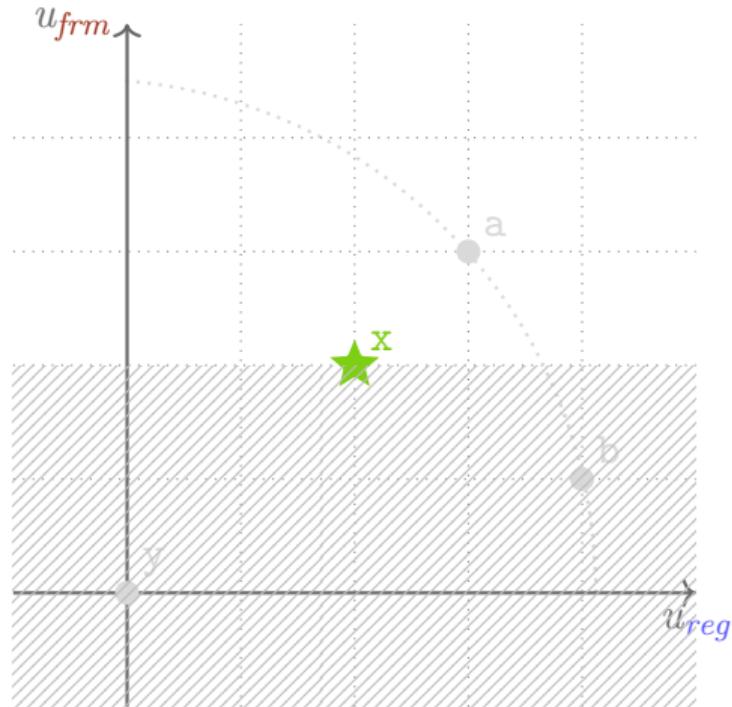
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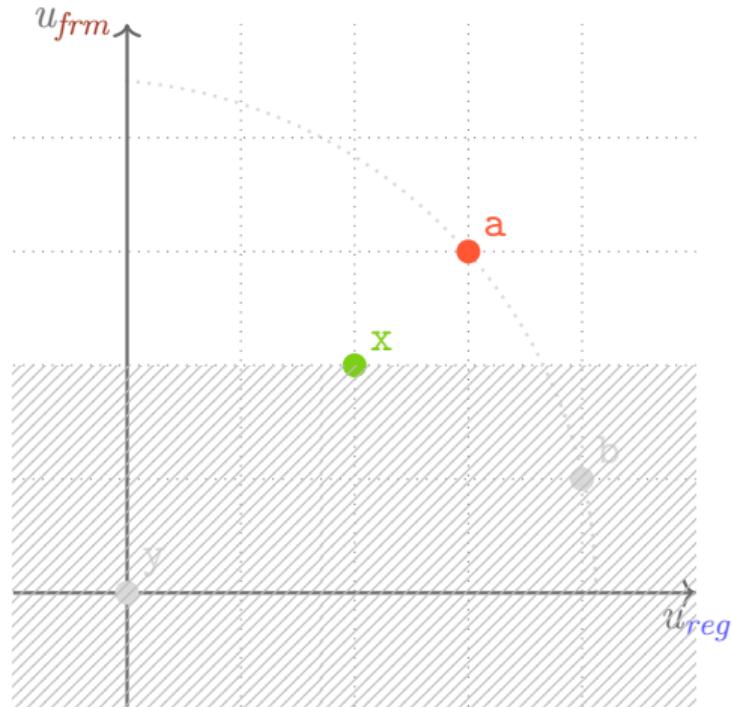
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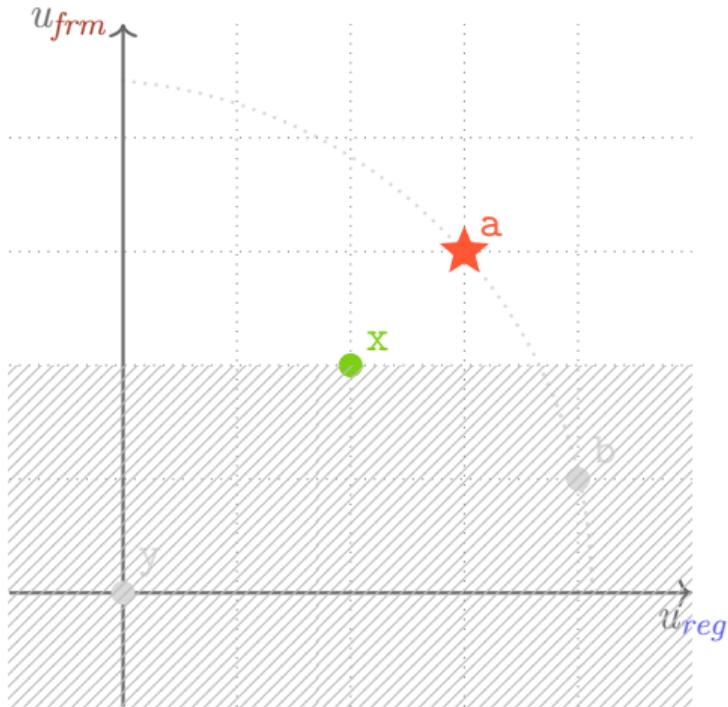
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So a two stage game with commitment to not revoke the prior proposal resulted in

- ◊ full revelation
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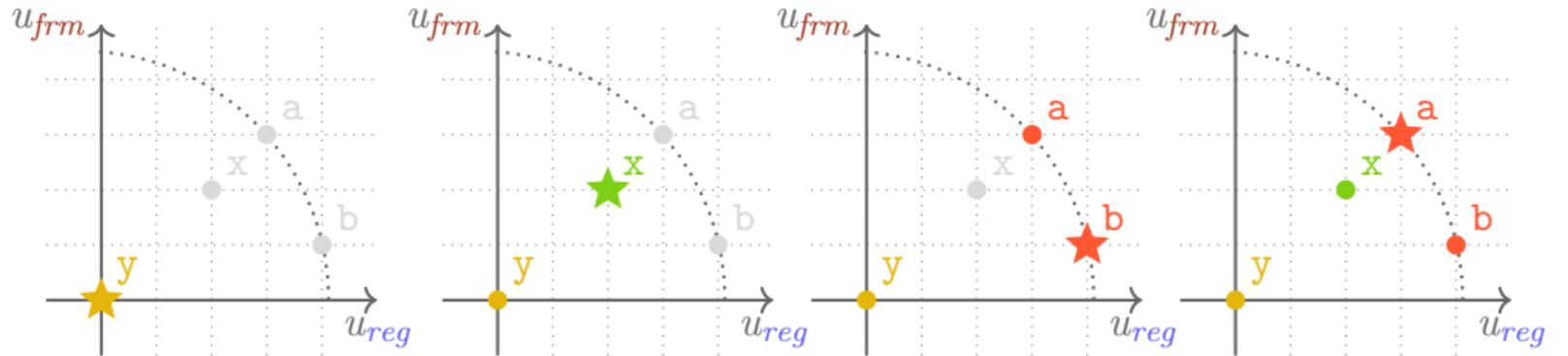
Does this always work? No

Example B

- ◊ What if the regulator was also initially unaware of x
- ◊ x and $\{a, b\}$ can be revealed independently:

$$\mathcal{R} = \{\{y\}, \{y, x\}, \{y, a, b\}, \{y, x, a, b\}\}$$

A monotone and efficient selection rule:



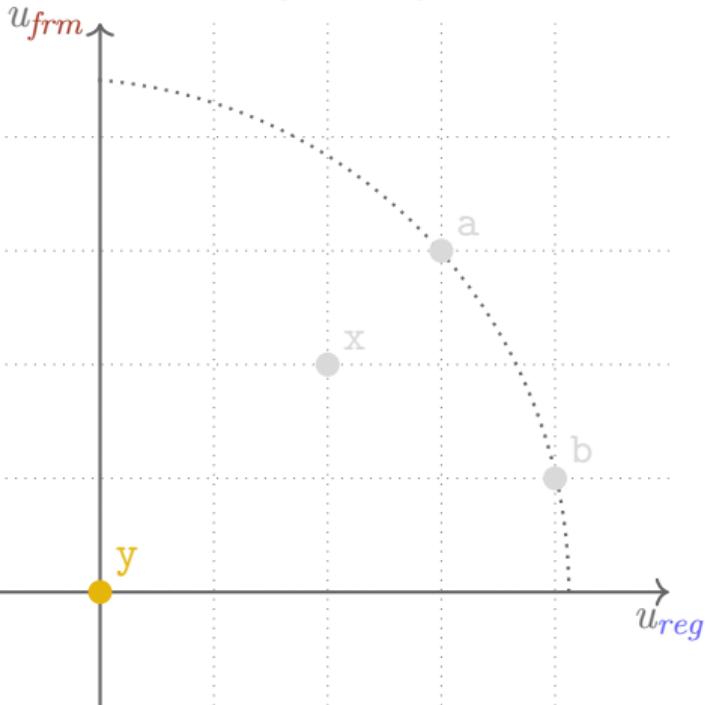
$$f(\{y\}) = y$$

$$f(\{y, x\}) = x$$

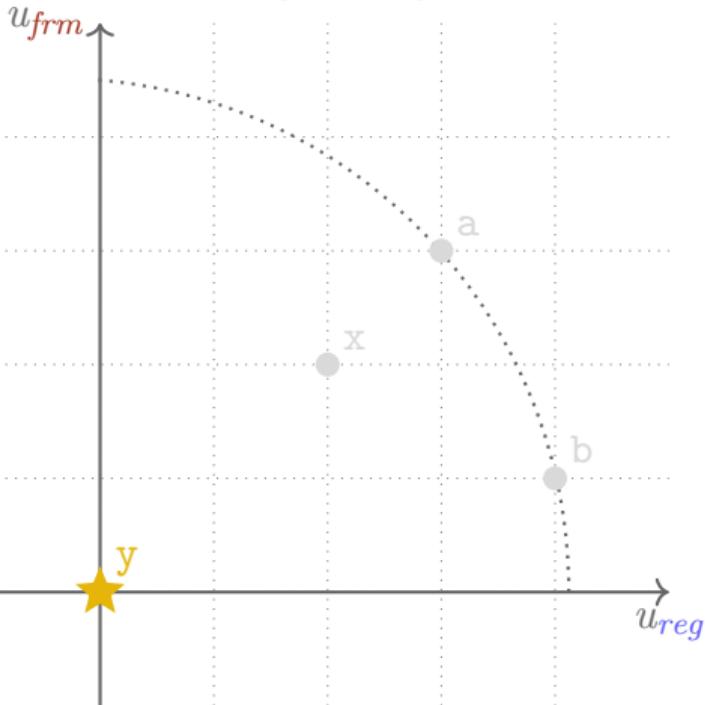
$$f(\{y, a, b\}) = b$$

$$f(\{y, x, a, b\}) = a$$

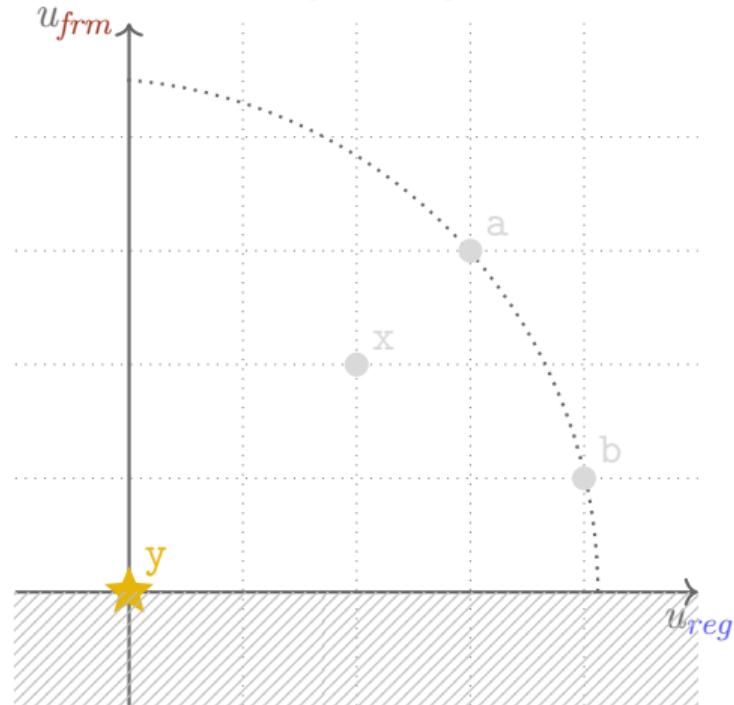
Example B



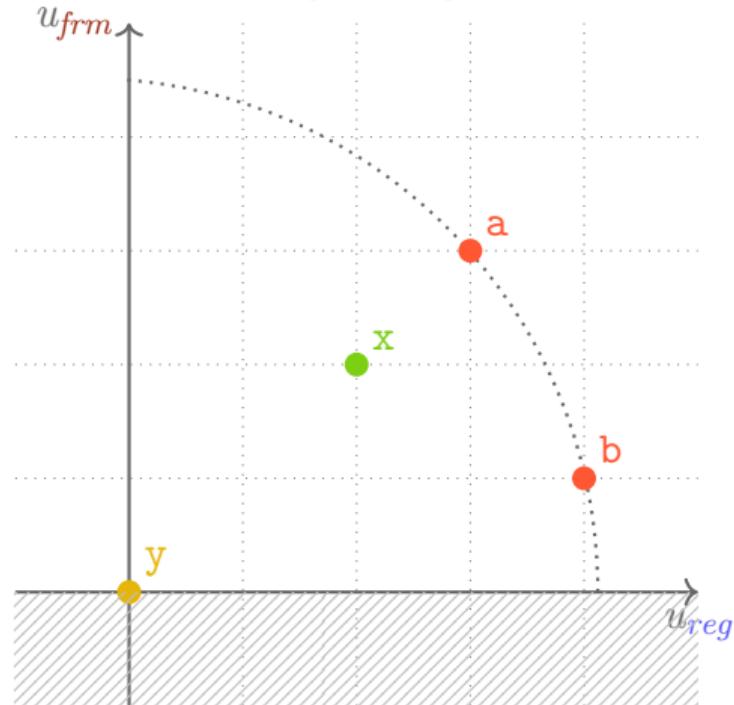
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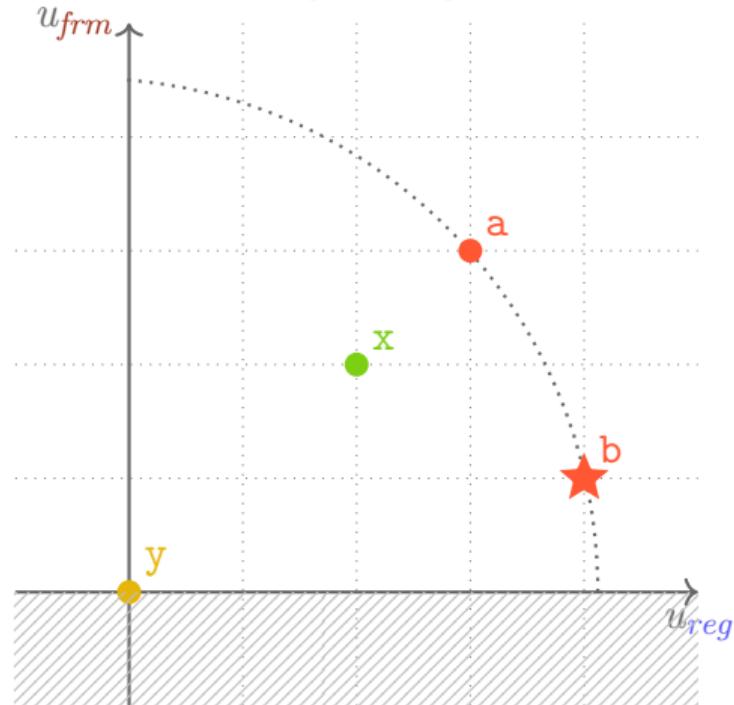
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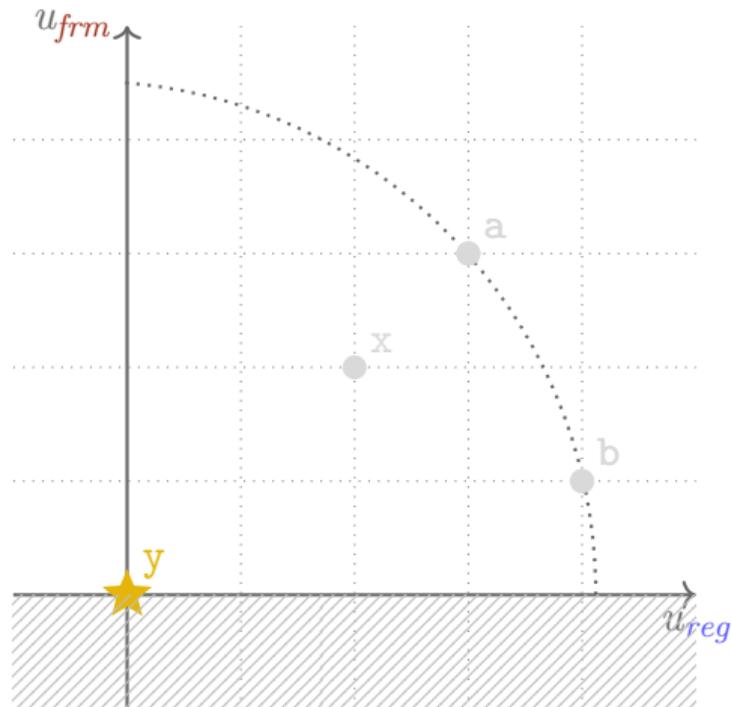


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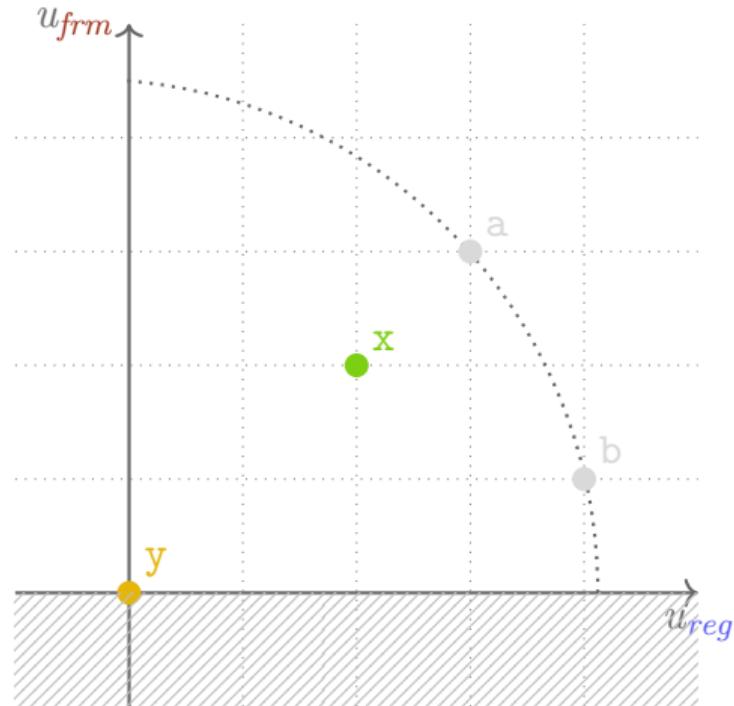


But the **firm** does not have to reveal everything! Instead, reveal only **x**

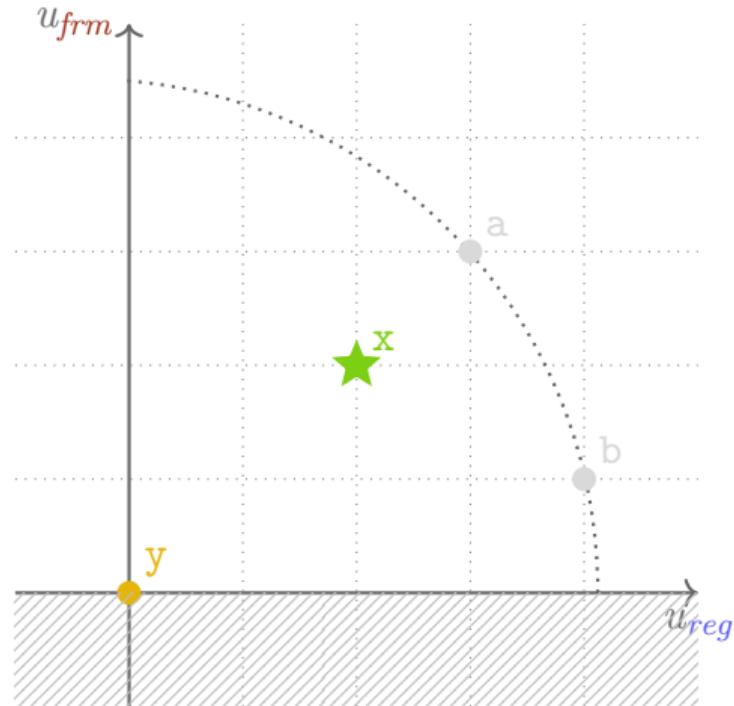
Example B



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Example B



The **firm** prefers **x** to **b**:

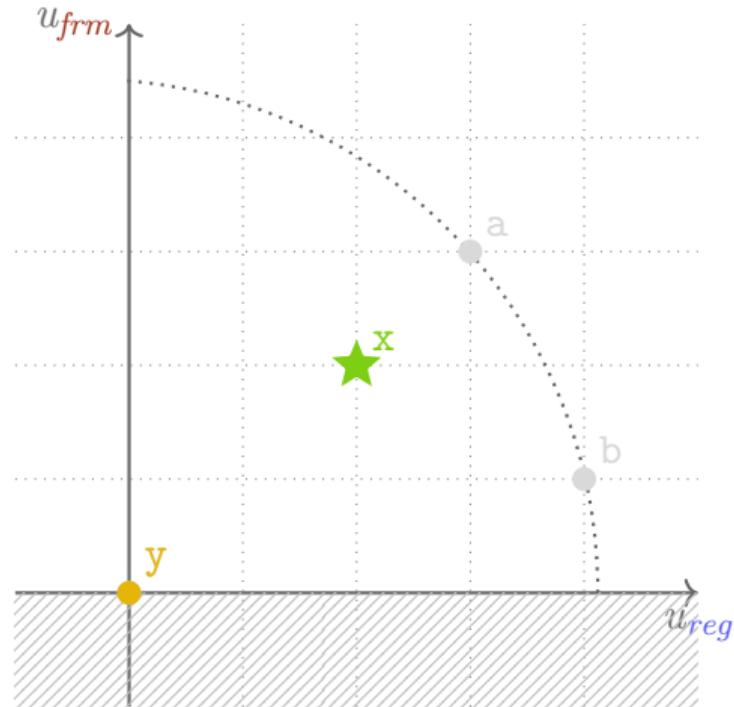
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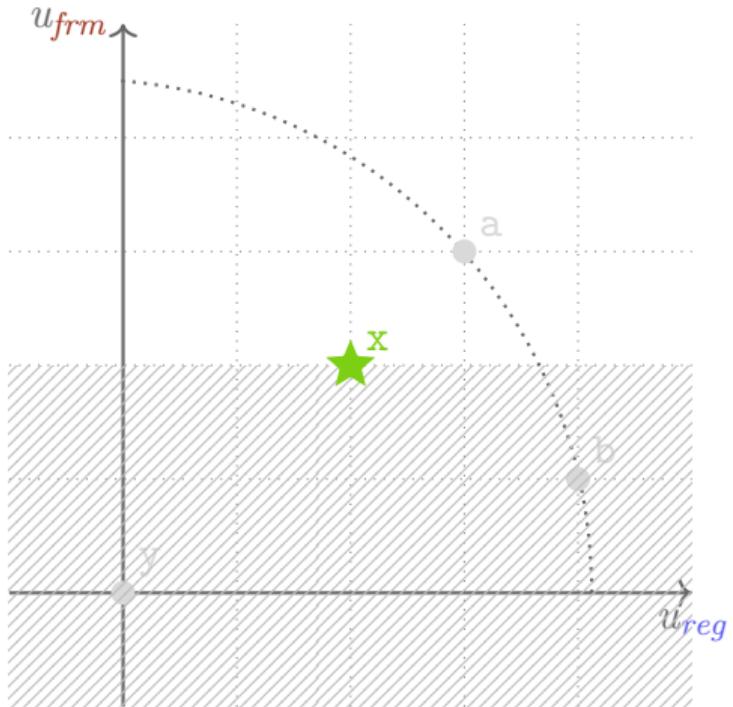
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But, we can repeat!

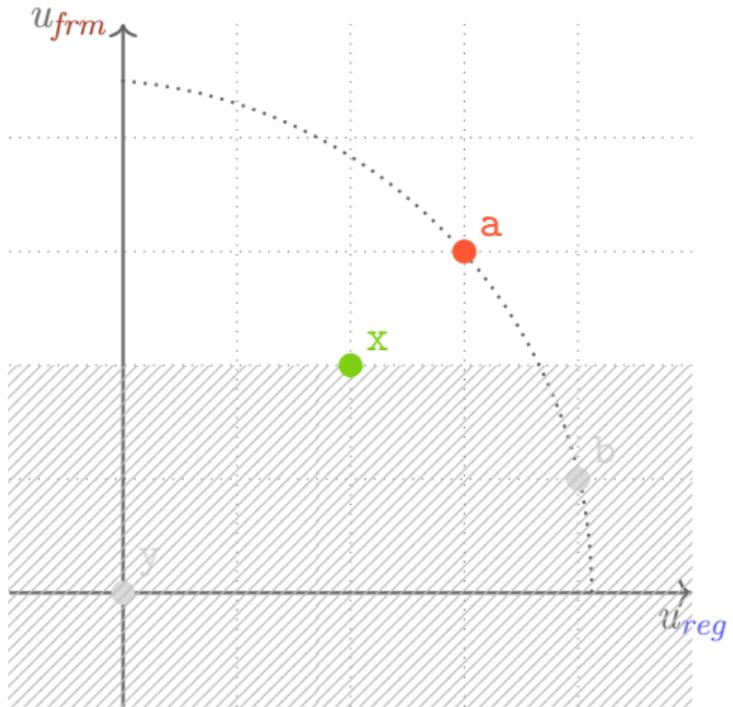
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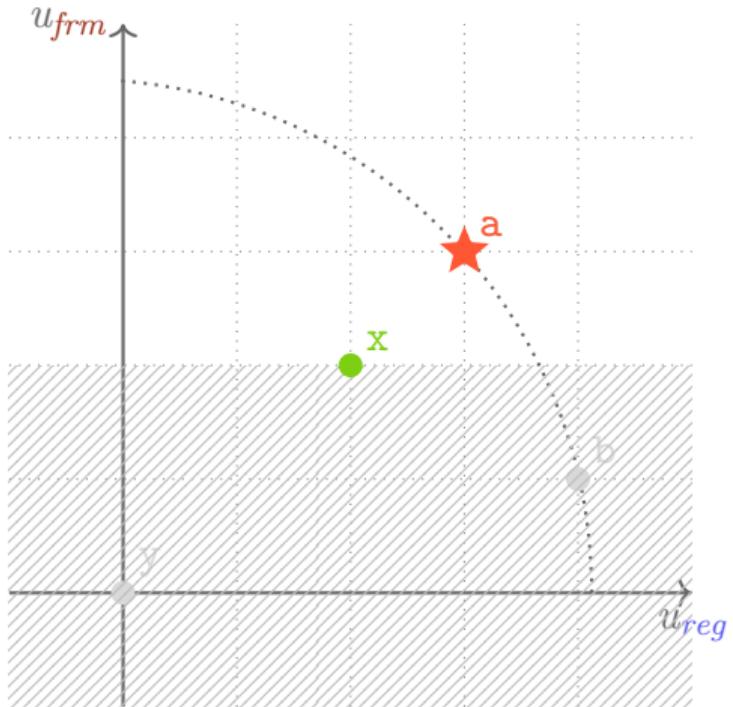
Example B



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Example B





Iterated Revelation Protocol



Iterated Revelation Protocol



INITIAL STEP — The **decision maker** announces $r_0 = r^\dagger$, and shortlists $a_0 \in r_0$.

ITERATIVE STEP — Given (r_0, \dots, r_{n-1}) distinct prior revelations, the **expert** reveals $r_n \in \mathcal{R}$.

- ◊ If $r_{n-1} \subsetneq r_n$, the **dm** shortlists $a_n \in r_n$, and the ITERATIVE STEP is repeated
- ◊ Otherwise, the protocol moves to the FINAL STEP

FINAL STEP — Given (r_0, \dots, r_n) distinct revelations, the **expert** chooses an action $a \in \{a_0, \dots, a_n\}$.

Importantly:

- ◊ This protocol can be explained / contracted to without having to express any specific actions/outcomes
- ◊ Specifically, the only contractual obligations in an IRP are actions that *have already been* revealed.



Strategies

Given the IRP, a **strategy**

- ◊ for the **dm** is a function from *sequences of revelations* to actions:

$$s : (r_0 \dots r_n) \mapsto a_n \in r_n$$

- ◊ for the **ex** is a function from *sequences of shortlisted actions* to revelations:

$$\sigma : (a_0 \dots a_{n-1}) \mapsto r_n \in \mathcal{R}$$

(and a choice out of the final shortlist)

Implementation



Let $a(s, \sigma)$ denote the action enacted by playing strategies s and σ .

Say that s **implements** the selection rule f if for all $r \in \mathcal{R}$

$$f(r) = a(s, \sigma) \quad \text{for some best response for type } r$$

and **fully implements** f if

$$f(r) = a(s, \sigma) \quad \text{for every best response for type } r$$

Theorem

The following are equivalent for a selection rule f

- (1) f is monotone (resp. strongly monotone)
- (2) there exists some s that implements f , (resp. fully implements)



Greedy Strategies & Efficiency



Each shortlist proposal in an IRP specifies:

- (1) The outcome should the game end
 - ◊ dm wants to maximize his own payoff
- (2) The implicit incentive constraint should the game continue
 - ◊ dm wants to minimize ex's payoff

In the examples, IRPs solved (1) ignoring (2)

Definition

Call a strategy s (for the decision maker) **mostly greedy** if for all $(r_0 \dots r_n)$, there is no $a \in r_n$ such that $V_e(s(r_0 \dots r_{n-1})) \leq V_e(a)$ and

$$u_d(a) > u_d(s(r_0 \dots r_n))$$

or such that

$$u_d(a) = u_d(s(r_0 \dots r_n)) \text{ and } u_e(a) > u_e(s(r_0 \dots r_n))$$

A mostly greedy strategy:

- ◊ maximizes the **dm**'s payoff myopically (subject to IC constraint)
- ◊ does not account for effect on future incentive constraints
- ◊ breaks ties in favor of the **expert** (hence only *mostly* greedy)

Theorem

Let s be mostly greedy. Then s implements the decision maker's preferred efficient selection rule, f^* .

- ◊ If f is any other monotone and efficient selection rule, then for all $r \supseteq r_d$

$$u_d(f^*(r)) \geq u_d(f(r))$$

Comparative Statics: Information Spillover

- ◊ Let \mathcal{R} and \mathcal{Q} be two different type spaces over the same set of actions:

$$\mathcal{R} \subseteq \mathcal{Q} \subseteq 2^{\mathcal{A}}$$

- ◊ Let $r^\dagger, r \in \mathcal{R}$ and $q^\dagger, q \in \mathcal{Q}$ be such that

$$r^\dagger = q^\dagger \subseteq r = q$$

Let $f^{\mathcal{R}}$ and $f^{\mathcal{Q}}$ denote the efficient, monotone selection rule induced by the mostly greedy strategy, then:

$$u_d(f^{\mathcal{Q}}(q)) \leq u_d(f^{\mathcal{R}}(r))$$

Comparative Statics: Information Spillover

- ◊ In the limit $\mathcal{R} = \{\mathcal{A} - r^\dagger\}$ (all actions reveal all other actions)
 - ◊ As if **dm** maximizes subject to individual rationality constraint
- ◊ In the limit $\mathcal{R} = 2^{\mathcal{A}}$ (all actions can be revealed independently)
 - ◊ As if **expert** maximizes subject to individual rationality constraint
 - ◊ This coincides with the expert preferred efficient selection rule
 - ◊ Corollary: efficient selection rule is unique



General Strategic Analysis



Definition

Call a strategy s **greedy** if for all $(r_0 \dots r_n)$, there is no $a \in r_n$ such that

$$V_e(s(r_0 \dots r_{n-1})) \leq V_e(a) \quad \text{and} \quad V_d(s(r_0 \dots r_n)) < V_d(a)$$

- ◊ There is no way to for the dm to increase his own payoff
- ◊ Generalization of mostly greedy strategy

Theorem

A selection rule f is implemented by a greedy s

if and only if

for all $r \in \mathcal{R}$, there is no other monotone selection rule f' such that

$$\inf_{r' \supseteq r} V_d(f(r')) < \inf_{r' \supseteq r} V_d(f'(r'))$$

Definition

Call a strategy s **locally rational** if for all $(r_0 \dots r_n)$, there is no $a \in r_n$ such that

$$V_e(s(r_0 \dots r_{n-1})) \leq V_e(a) < V_e(s(r_0 \dots r_n)) \quad \text{and} \quad V_d(s(r_0 \dots r_n)) < V_d(a)$$

- ◊ There is no way to simultaneously for the **dm** to
 - ◊ increase his own payoff
 - ◊ decrease the **expert**'s payoff

Theorem

A selection rule f is implemented by a locally rational s

if and only if

for all $r \in \mathcal{R}$, there is no other monotone selection rule f' such that

$$V_d(f(r')) \leq V_d(f'(r')) \quad \text{for all } r' \supseteq r,$$

$$V_d(f(r')) < V_d(f'(r')) \quad \text{for some } r' \supseteq r$$

- ◊ ‘if’ direction requires a richness condition on \mathcal{R}



Payoff Uncertainty



- ◊ The implementation above presupposes **dm** can anticipate **ex**'s acceptance / rejection
- ◊ What happens with private information:
 - ◊ Actions are state-dependent $u : \mathcal{A} \times \Omega \rightarrow \mathbb{R}$
 - ◊ assume **ex** knows the state, $\omega \in \Omega$
 - ◊ **dm** does not

Example C

- ◊ $\Omega = \{\omega_L, \omega_R\}$, **ex** knows the state, **dm** believes equally likely
 - ◊ Each action is therefore given by $(\langle x_{d,L}, x_{d,R} \rangle, \langle x_{e,L}, x_{e,R} \rangle)$.
- ◊ The **dm** is initially aware of one action:

$$x = (\langle 0, 0 \rangle, \langle 0, 0 \rangle)$$

- ◊ The **ex** is also aware of:

$$a_L = (\langle 3, -1 \rangle, \langle 3, -1 \rangle) \quad a_R = (\langle -1, 3 \rangle, \langle -1, 3 \rangle) \quad b = (\langle 2, 2 \rangle, \langle 2, 2 \rangle)$$

- ◊ The only revelation type is $\{a_L, a_R, b\}$.

Example C

$*x = \langle 0, 0 \rangle, \langle 0, 0 \rangle*$

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$$x = (\langle \mathbf{0}, \mathbf{0} \rangle, \langle \mathbf{0}, \mathbf{0} \rangle)$$

$$b = (\langle \mathbf{2}, \mathbf{2} \rangle, \langle \mathbf{2}, \mathbf{2} \rangle)$$

$$a_L = (\langle \mathbf{3}, -\mathbf{1} \rangle, \langle \mathbf{3}, -\mathbf{1} \rangle)$$

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Example C

- ◊ Preferences are completely aligned, but IRP does not allow delegation
- ◊ the protocol cannot use **ex**'s private info.
 - ◊ this creates inefficiency
- ◊ Instead, **dm** chooses a **set of actions** $p_1 \subseteq r$. After revelation, propose

$$p_1 = \{a_L, a_R\}$$

and let the **ex** choose.

- ◊ A **generalized IRP** allows the **dm** to choose a set of actions at each step:
 - ◊ At each $(r_0 \dots r_n), s(r_0 \dots r_n) \subseteq r_n$
- ◊ A **generalized selection rule** is a function $f: \Omega \times \mathcal{R} \rightarrow \mathcal{A}$
 - ◊ For each $r \in \mathcal{R}, w \in \Omega$, we have $f(w, t) \in t$



Theorem

The following are equivalent for a gen. selection rule f

- (1) f can be implemented by a gen. IRP
- (2) f is monotone: for all $\omega, r \in \Omega \times \mathcal{R}$

$$u_{\textcolor{red}{e}}(f(\omega', r'), \omega) \leq u_{\textcolor{red}{e}}(f(\omega, t), \omega)$$

for any other $\omega' \in \Omega$ and $r' \subseteq r$.



Thank You

