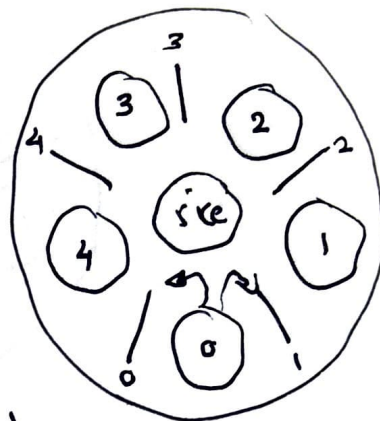


# II The Dining Philosophers problem

1

Five Philosophers who think and eat.

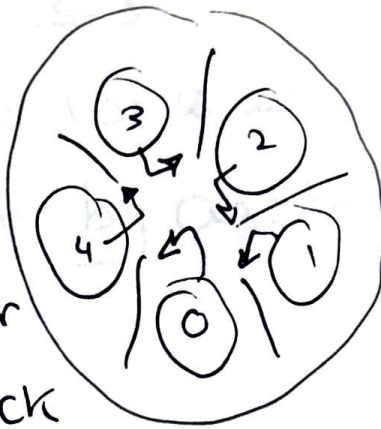
When a Philosopher gets hungry, they grab the 2 Chopsticks (clopset) and eat.



Semaphore chopstick[5]  
initialized to 1

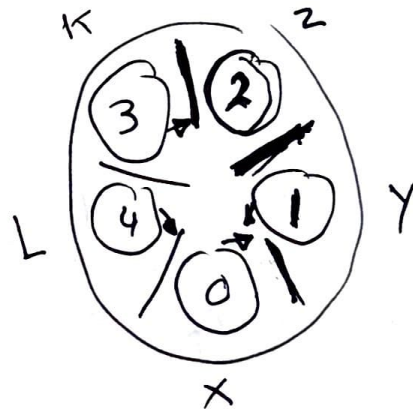
```
Pi: do { wait (chopstick[i]);  
        wait (chopstick[(i+1)%5]);  
        eat();  
        signal (chopstick[i]);  
        signal (chopstick[(i+1)%5]);  
        think();  
    } while (1);
```

Deadlock  $\Rightarrow$   
 When every Philosopher  
 grabs the chopstick  
 on their left first.



Soln

① Odd Philosophers grab left first  
 Even " " Right first



② Pick both  
 chopsticks at the  
 same time

$\Rightarrow$  using one mutex

```
wait(mutex);
grab both;
Signal(mutex);
eat();
```

# Deadlocks

3

A set of processes is in a deadlocked state if every process is waiting for an event that can only happen by another process in the set.

$P_1$	$P_2$
Wait(s)	Wait(q)
Wait(q)	Wait(s)
$\vdots$	$\vdots$

Dynamics of execution that determines if a deadlock happens or not.

## I System Model

Number of Processes ( $n$ )

Number of Resources (group by type).

A process requires a number of resources to finish (must be less than the total available).

[1] Request : If a resource is not available [4]  
process must wait

[2] Use

[3] Release

### Conditions for a deadlock

[1] Mutual Exclusion : only 1 process can use a resource at a time

[2] Hold and Wait : A process holds a resource while waiting for another

[3] No-preemption : A resource cannot be taken from a process.

[4] Circular Wait : A closed chain of waiting processes exists.

All 4 Conditions must be present for a deadlock to occur.



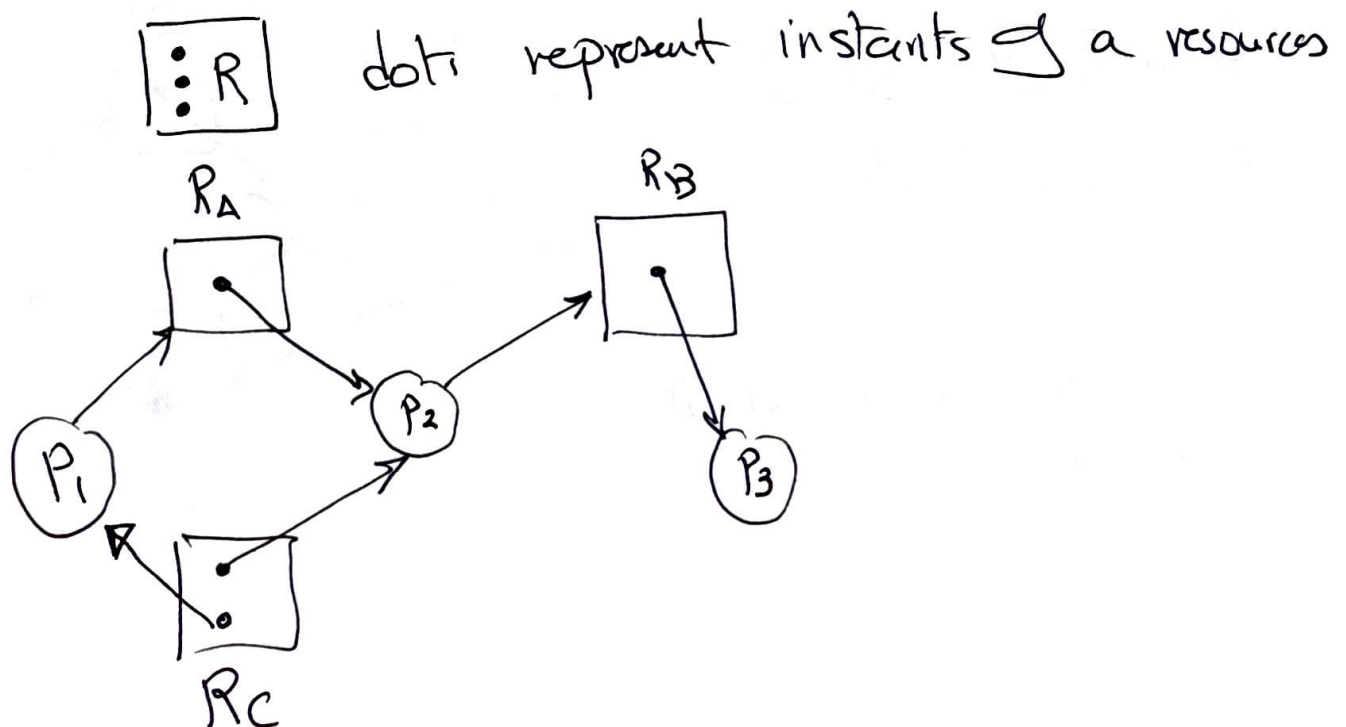
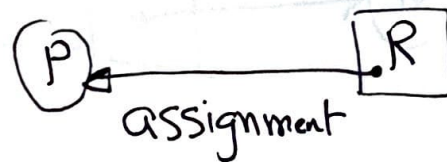
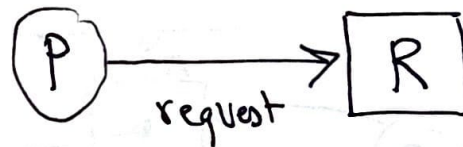
# Resource Allocation Graphs

Directed graph that shows the allocation of resources to processes.

Vertices (nodes) : processes  $\bigcirc$

resource  $\square$

Edges : requests and assignments.

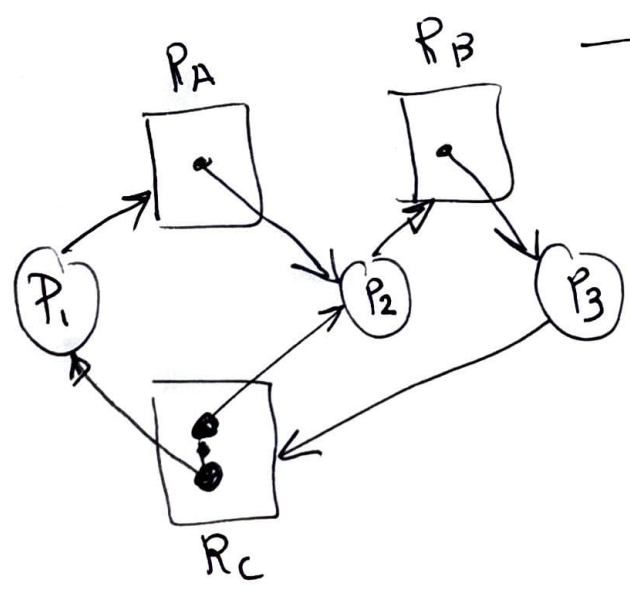
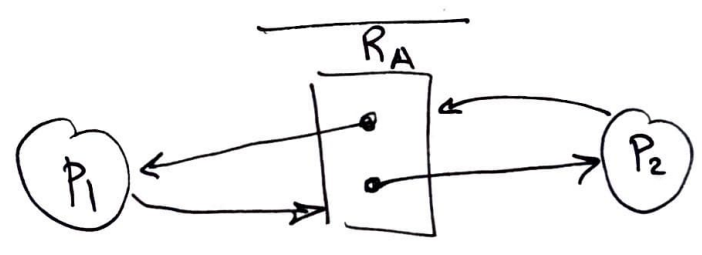


If graph has no cycle (no process is deadlocked)

else A deadlock may exist

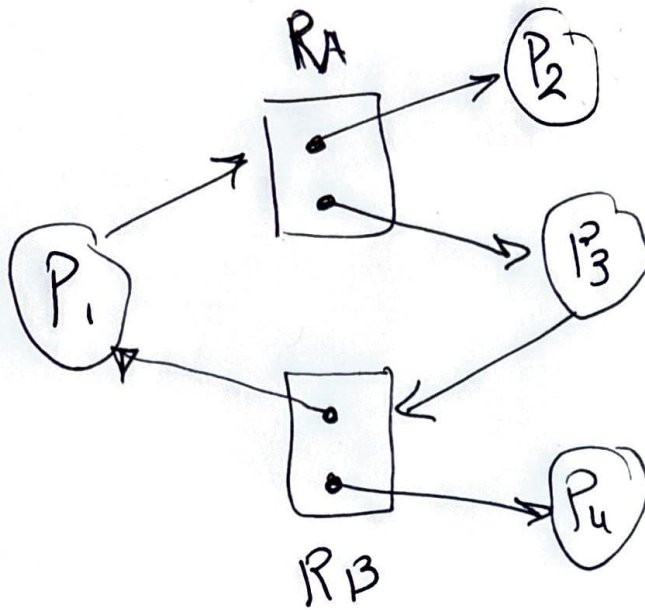
If each resource has exactly 1 instant  
 $\Rightarrow$  a cycle means a deadlock.

Otherwise  
 $\Rightarrow$  ^ May mean a deadlock.



$P_3 \ R_C \ P_1 \ R_A \ P_2 \ R_B \ P_3$   
 $R_B \ P_3 \ R_C \ P_2 \ R_B$

$P_1 \ P_2 \ P_3$  are deadlocked



$P_1 \quad RA \quad P_3 \quad RB \quad P_1$

Cycle

No deadlock.