Homework Series 1

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Some parsers:

1.1 ANTLR4

Written in Java

Target Language C#, Java, Python

Algorithms LL

 $Source\ https://theantlrguy.atlassian.net/wiki/display/ANTLR4/$

1.2 JFLAP

Written in Java

Target Language Java

Algorithms LL1 and LALR1

Source http://en.wikipedia.org/wiki/Compiler-compiler

1.3 Parsec

Written in Haskell

Target Language Haskell

Algorithms LL and Backtracking

Source Lennart

1.4 PLY

Written in Python

Target Language Python

Algorithms LALR1

Source Experience with this parser generator

1.5 Racc

Written in Ruby

Target Language Ruby

Algorithms LALR1

Source https://github.com/tenderlove/racc

1.6 Bison

Written in C

Target Language C, C++, Java

Algorithms LALR1, LR1, IELR1, GLR

Source http://www.gnu.org/software/bison/

1.7 Yapps

Written in Python

Target Language Python

Algorithms LL1

Source Experience with this parser generator

1.8 Citrus

Written in Ruby

Target Language Ruby

Algorithms PEG

Source https://github.com/mjackson/citrus

1.9 Jison

Written in Javascript

Target Language Javascript

Algorithms LALR1, LR0, SLR1, LR1, LL1

Source https://zaach.github.io/jison/docs/

1.10 Happy

Written in Haskell

Target Language Haskell

Algorithms Generalized LR

Source https://www.haskell.org/happy/#what

 $\mathbf{2}$

2.1

literal	int	long
binary	/0b[0-1]+/	/0b[0-1]+(l—L)/
octal	/0(0 [1-7][0-7]) + /	/0(0-[1-7][0-7])+(1-L)/
hexadecimal	/0x[0-9a-f] + /	/0x[0-9a-f] + (l L)/
decimal	/(0 [1-9][0-9]*)/	/(0 [1-9][0-9]*(l L))/

2.2

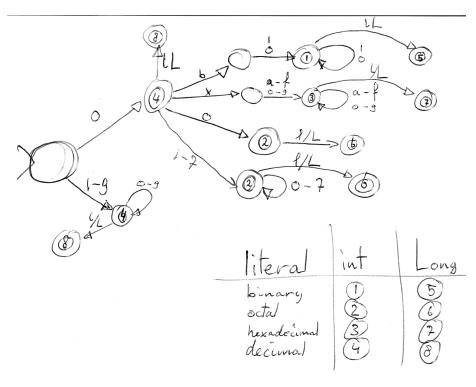


Figure 1: A DFA combining the token types

2.3

See the code in pp.s1378791.q1_2.Numbers.g4 and pp.s1378791.q1_2.NumbersTest

3

3.1

3.1.1 a[i+1] + b.field

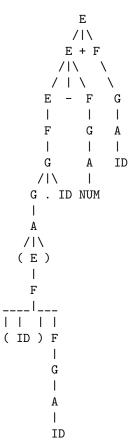
 ${\rm E\atop E+F}$

$$\begin{array}{l} F + F \\ G + F \\ G \mid E \mid + F \\ A \mid E \mid + F \\ ID \mid E + F \mid + F \\ ID \mid F + F \mid + F \\ ID \mid G + F \mid + F \\ ID \mid A + F \mid + F \\ ID \mid ID + F \mid + F \\ ID \mid ID + G \mid + F \\ ID \mid ID + A \mid + F \\ ID \mid ID + NUM \mid + F \\ ID \mid ID + NUM \mid + G \\ ID \mid ID + NUM \mid + G \\ ID \mid ID + NUM \mid + A \\ ID \mid ID + NUM \mid + A \\ ID \mid ID + NUM \mid + A \\ ID \mid ID + NUM \mid + A \\ ID \mid ID \mid ID + NUM \mid + A \\ ID \mid ID \mid ID + NUM \mid + A \\ ID \mid ID \mid ID + NUM \mid + A \\ ID \mid ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID + NUM \mid + ID \\ ID \mid ID +$$

3.1.2 ((Type) x).i - 10 + y

$$\begin{split} E & E + F \\ E - F + F \\ F - F + F \\ G - F + F \\ G . ID - F + F \\ A . ID - F + F \\ (E) . ID - F + F \\ (E) . ID - F + F \\ (F) . ID - F + F \\ (ID) F) . ID - F + F \end{split}$$

```
\begin{array}{l} (\;(\;\mathrm{ID}\;)\;\mathrm{G}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{F}\;+\;\mathrm{F}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{A}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{F}\;+\;\mathrm{F}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{F}\;+\;\mathrm{F}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{G}\;+\;\mathrm{F}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{A}\;+\;\mathrm{F}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{NUM}\;+\;\mathrm{F}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{NUM}\;+\;\mathrm{G}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{NUM}\;+\;\mathrm{A}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{NUM}\;+\;\mathrm{A}\\ (\;(\;\mathrm{ID}\;)\;\mathrm{ID}\;)\;.\;\mathrm{ID}\;\text{-}\;\mathrm{NUM}\;+\;\mathrm{ID} \end{array}
```



3.2

See pp.s1378791.q1_3

3.3

```
G -> A G'
G' -> "[" E "] G'
| "." ID G'
| epsilon
A -> "(" E ")"
| NUM
| ID
```

3.4

3.4.1 FIRST

	1	2	3	4
\overline{E}	Ø	(((, num, id
E^{\prime}	+, -, epsilon	+, -, epsilon	+, -, epsilon	+, -, epsilon
\mathbf{F}	(((, num, id	(, num $, id$
G	Ø	(, num, id	(, num, id	(, num, id
G'	[, ., epsilon	[, ., epsilon	[, ., epsilon	[, ., epsilon
A	(, num, id	(, num, id	(, num, id	(, num, id

3.4.2 FOLLOW

-	init	1	2
E	eof	eof],)	eof,],)
\mathbf{E}'	Ø	eof],)	eof,],)
F	Ø	eof],), +, -	eof,],), +, -
G	Ø	eof],), +, -	eof,],), +, -
G'	Ø	eof],), +, -	eof,],), +, -
A	Ø	a = a = a = a = a = a = a = a = a = a =	$\mathrm{eof},\],\),\ +,\ -,\ [,\ .$

3.4.3 FIRST+

Regel	first	follow	first+
$E \to FE'$	(, num, id	eof,],)	(, num, id
$E' \rightarrow "+"FE'$	+	eof,],)	+
$E' \rightarrow$ "-" FE'	-	eof,],)	-
$E' \rightarrow epsilon$	epsilon	eof,],)	eof,],)
$F \rightarrow$ "(" ID ")" F	(eof,],), +, -	(
$F \to G$	num, id	eof,],), +, -	(, num $, id$
$G \to AG'$	num, id	eof,],), +, -	(, num, id
$G' \rightarrow$ "[" E "] G'	[eof,],), +, -	[
$G' \rightarrow$ "." IDG'	•	eof,],), +, -	
$G' \to epsilon$	epsilon	eof,],), +, -	eof, +, -,],)
$A \rightarrow$ "("E")"	(eof,],), +, -, [,]	(
$A \to NUM$	num	eof,],), +, -, [, .	num
$A \to ID$	id	eof, [1, 1), +, -, [1, 1]	id

See lab files in pp.s1378791.q1_5