

G++ 2.91.57, cygnus\cygwin-b20\include\g++\stl_tree.h 完整列表

```
/*
 *
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 *
 */

/* NOTE: This is an internal header file, included by other STL headers.
 * You should not attempt to use it directly.
 */

#ifndef __SGI_STL_INTERNAL_TREE_H
#define __SGI_STL_INTERNAL_TREE_H
```

```
/*
本檔實作Red-black tree (紅-黑樹) class，用以實作 STL 關聯式容器（如set,
multiset, map, multimap）。所用之insertion 和deletion 演算法係以
Cormen, Leiserson 和 Rivest 所著之 Introduction to Algorithms
(MIT Press, 1990) 一書為基礎，唯以下兩點不同：
```

(1) header 不僅指向 root，也指向紅黑樹的最左節點，以便實作出常數時間之 begin()；並且也指向紅黑樹的最右節點，以便set 相關泛型演算法（如set_union 等等）有線性時間之表現。

(2) 當一個即將被刪除之節點擁有兩個子節點時，它的successor node is relinked into its place, rather than copied, 如此一來唯一失效 (invalidated) 的迭代器就只是那些referring to the deleted node.

```

*/

#include <stl_algobase.h>
#include <stl_alloc.h>
#include <stl_construct.h>
#include <stl_function.h>

__STL_BEGIN_NAMESPACE

typedef bool __rb_tree_color_type;
const __rb_tree_color_type __rb_tree_red = false; // 紅色為 0
const __rb_tree_color_type __rb_tree_black = true; // 黑色為 1

struct __rb_tree_node_base
{
    typedef __rb_tree_color_type color_type;
    typedef __rb_tree_node_base* base_ptr;

    color_type color; // 節點顏色，非紅即黑。
    base_ptr parent; // RB 樹的許多操作，必須知道父節點。
    base_ptr left; // 指向左節點。
    base_ptr right; // 指向右節點。

    static base_ptr minimum(base_ptr x)
    {
        while (x->left != 0) x = x->left; // 一直向左走，就會找到最小值，
        return x; // 這是二元搜尋樹的特性。
    }

    static base_ptr maximum(base_ptr x)
    {
        while (x->right != 0) x = x->right; // 一直向右走，就會找到最大值，
        return x; // 這是二元搜尋樹的特性。
    }
};

template <class Value>
struct __rb_tree_node : public __rb_tree_node_base
{
    typedef __rb_tree_node<Value>* link_type;
    Value value_field; // 節點實值
};

struct __rb_tree_base_iterator
{
    typedef __rb_tree_node_base::base_ptr base_ptr;
    typedef bidirectional_iterator_tag iterator_category;
    typedef ptrdiff_t difference_type;

```

```

base_ptr node; // 它用來與容器之間產生一個連結關係 (make a reference)

// 以下其實可實作於 operator++ 內，因為再無他處會呼叫此函式了。
void increment()
{
    if (node->right != 0) { // 如果有右子節點。狀況(1)
        node = node->right; // 就向右走
        while (node->left != 0) // 然後一直往左子樹走到底
            node = node->left; // 即是解答
    }
    else { // 沒有右子節點。狀況(2)
        base_ptr y = node->parent; // 找出父節點
        while (node == y->right) { // 如果現行節點本身是個右子節點，
            node = y; // 就一直上溯，直到「不為右子節點」止。
            y = y->parent;
        }
        if (node->right != y) // 「若此時的右子節點不等於此時的父節點」。
            node = y; // 狀況(3) 此時的父節點即為解答。
                        // 否則此時的node 為解答。狀況(4)
    }
    // 注意，以上判斷「若此時的右子節點不等於此時的父節點」，是為了應付一種
    // 特殊情況：我們欲尋找根節點的下一節點，而恰巧根節點無右子節點。
    // 當然，以上特殊作法必須配合 RB-tree 根節點與特殊之header 之間的
    // 特殊關係。
}

// 以下其實可實作於 operator-- 內，因為再無他處會呼叫此函式了。
void decrement()
{
    if (node->color == __rb_tree_red && // 如果是紅節點，且
        node->parent->parent == node) // 父節點的父節點等於自己，
        node = node->right; // 狀況(1) 右子節點即為解答。
    // 以上情況發生於node為header時 (亦即 node 為 end() 時)。
    // 注意，header 之右子節點即 mostright，指向整棵樹的 max 節點。
    else if (node->left != 0) { // 如果有左子節點。狀況(2)
        base_ptr y = node->left; // 令y指向左子節點
        while (y->right != 0) // 當y有右子節點時
            y = y->right; // 一直往右子節點走到底
        node = y; // 最後即為答案
    }
    else { // 既非根節點，亦無左子節點。
        base_ptr y = node->parent; // 狀況(3) 找出父節點
        while (node == y->left) { // 當現行節點身為左子節點
            node = y; // 一直交替往上走，直到現行節點
            y = y->parent; // 不為左子節點
        }
        node = y; // 此時之父節點即為答案
    }
}

```

```

};

template <class Value, class Ref, class Ptr>
struct __rb_tree_iterator : public __rb_tree_base_iterator
{
    typedef Value value_type;
    typedef Ref reference;
    typedef Ptr pointer;
    typedef __rb_tree_iterator<Value, Value&, Value*> iterator;
    typedef __rb_tree_iterator<Value, const Value&, const Value*> const_iterator;
    typedef __rb_tree_iterator<Value, Ref, Ptr> self;
    typedef __rb_tree_node<Value>* link_type;

    __rb_tree_iterator() {}
    __rb_tree_iterator(link_type x) { node = x; }
    __rb_tree_iterator(const iterator& it) { node = it.node; }

    reference operator*() const { return link_type(node)->value_field; }
#ifdef __SGI_STL_NO_ARROW_OPERATOR
    pointer operator->() const { return &(operator*()); }
#endif /* __SGI_STL_NO_ARROW_OPERATOR */

    self& operator++() { increment(); return *this; }
    self operator++(int) {
        self tmp = *this;
        increment();
        return tmp;
    }

    self& operator--() { decrement(); return *this; }
    self operator--(int) {
        self tmp = *this;
        decrement();
        return tmp;
    }
};

inline bool operator==(const __rb_tree_base_iterator& x,
                      const __rb_tree_base_iterator& y) {
    return x.node == y.node;
    // 兩個迭代器相等，意指其所指的節點相等。
}

inline bool operator!=(const __rb_tree_base_iterator& x,
                      const __rb_tree_base_iterator& y) {
    return x.node != y.node;
    // 兩個迭代器不等，意指其所指的節點不等。
}

```

```

#ifndef __STL_CLASS_PARTIAL_SPECIALIZATION

inline bidirectional_iterator_tag
iterator_category(const __rb_tree_base_iterator&) {
    return bidirectional_iterator_tag();
}

inline __rb_tree_base_iterator::difference_type*
distance_type(const __rb_tree_base_iterator&) {
    return (__rb_tree_base_iterator::difference_type*) 0;
}

template <class Value, class Ref, class Ptr>
inline Value* value_type(const __rb_tree_iterator<Value, Ref, Ptr>&) {
    return (Value*) 0;
}

#endif /* __STL_CLASS_PARTIAL_SPECIALIZATION */

// 以下都是全域函式：__rb_tree_rotate_left(), __rb_tree_rotate_right(),
// __rb_tree_rebalance(), __rb_tree_rebalance_for_erase()

// 新節點必為紅節點。如果安插處之父節點亦為紅節點，就違反紅黑樹規則，此時必須
// 做樹形旋轉（及顏色改變，在程式它處）。
inline void
__rb_tree_rotate_left(__rb_tree_node_base* x, __rb_tree_node_base*& root)
{
    // x 為旋轉點
    __rb_tree_node_base* y = x->right; // 令y 為旋轉點的右子節點
    x->right = y->left;
    if (y->left != 0)
        y->left->parent = x; // 別忘了回馬槍設定父節點
    y->parent = x->parent;

    // 令 y 完全頂替 x 的地位（必須將 x 對其父節點的關係完全接收過來）
    if (x == root) // x 為根節點
        root = y;
    else if (x == x->parent->left) // x 為其父節點的左子節點
        x->parent->left = y;
    else // x 為其父節點的右子節點
        x->parent->right = y;
    y->left = x;
    x->parent = y;
}

// 新節點必為紅節點。如果安插處之父節點亦為紅節點，就違反紅黑樹規則，此時必須
// 做樹形旋轉（及顏色改變，在程式它處）。
inline void
__rb_tree_rotate_right(__rb_tree_node_base* x, __rb_tree_node_base*& root)

```

```

{
    // x 為旋轉點
    __rb_tree_node_base* y = x->left; // y 為旋轉點的左子節點
    x->left = y->right;
    if (y->right != 0)
        y->right->parent = x; // 別忘了回馬槍設定父節點
    y->parent = x->parent;

    // 令 y 完全頂替 x 的地位（必須將 x 對其父節點的關係完全接收過來）
    if (x == root) // x 為根節點
        root = y;
    else if (x == x->parent->right) // x 為其父節點的右子節點
        x->parent->right = y;
    else // x 為其父節點的左子節點
        x->parent->left = y;
    y->right = x;
    x->parent = y;
}

// 重新令樹形平衡（改變顏色及旋轉樹形）
// 參數一為新增節點，參數二為 root
inline void
__rb_tree_rebalance(__rb_tree_node_base* x, __rb_tree_node_base*& root)
{
    x->color = __rb_tree_red; // 新節點必為紅
    while (x != root && x->parent->color == __rb_tree_red) { // 父節點為紅
        if (x->parent == x->parent->parent->left) { // 父節點為祖父節點之左子節點
            __rb_tree_node_base* y = x->parent->parent->right; // 令 y 為伯父節點
            if (y && y->color == __rb_tree_red) { // 伯父節點存在，且為紅
                x->parent->color = __rb_tree_black; // 更改父節點為黑
                y->color = __rb_tree_black; // 更改伯父節點為黑
                x->parent->parent->color = __rb_tree_red; // 更改祖父節點為紅
                x = x->parent->parent;
            }
        }
        else { // 無伯父節點，或伯父節點為黑
            if (x == x->parent->right) { // 如果新節點為父節點之右子節點
                x = x->parent;
                __rb_tree_rotate_left(x, root); // 第一參數為左旋點
            }
            x->parent->color = __rb_tree_black; // 改變顏色
            x->parent->parent->color = __rb_tree_red;
            __rb_tree_rotate_right(x->parent->parent, root); // 第一參數為右旋點
        }
    }
    else { // 父節點為祖父節點之右子節點
        __rb_tree_node_base* y = x->parent->parent->left; // 令 y 為伯父節點
        if (y && y->color == __rb_tree_red) { // 有伯父節點，且為紅
            x->parent->color = __rb_tree_black; // 更改父節點為黑
            y->color = __rb_tree_black; // 更改伯父節點為黑
        }
    }
}

```

```

        x->parent->parent->color = __rb_tree_red;    // 更改祖父節點為紅
        x = x->parent->parent;    // 準備繼續往上層檢查...
    }
    else {    // 無伯父節點，或伯父節點為黑
        if (x == x->parent->left) {    // 如果新節點為父節點之左子節點
            x = x->parent;
            __rb_tree_rotate_right(x, root);    // 第一參數為右旋點
        }
        x->parent->color = __rb_tree_black;    // 改變顏色
        x->parent->parent->color = __rb_tree_red;
        __rb_tree_rotate_left(x->parent->parent, root);    // 第一參數為左旋點
    }
}
} // while 結束
root->color = __rb_tree_black;    // 根節點永遠為黑
}

inline __rb_tree_node_base*
__rb_tree_rebalance_for_erase(__rb_tree_node_base* z,
                              __rb_tree_node_base*& root,
                              __rb_tree_node_base*& leftmost,
                              __rb_tree_node_base*& rightmost)
{
    __rb_tree_node_base* y = z;
    __rb_tree_node_base* x = 0;
    __rb_tree_node_base* x_parent = 0;
    if (y->left == 0)    // z has at most one non-null child. y == z.
        x = y->right;    // x might be null.
    else
        if (y->right == 0)    // z has exactly one non-null child. y == z.
            x = y->left;    // x is not null.
        else {    // z has two non-null children. Set y to
            // z's successor. x might be null.
            y = y->right;
            while (y->left != 0)
                y = y->left;
            x = y->right;
        }
    if (y != z) {    // relink y in place of z. y is z's successor
        z->left->parent = y;
        y->left = z->left;
        if (y != z->right) {
            x_parent = y->parent;
            if (x) x->parent = y->parent;
            y->parent->left = x;    // y must be a left child
            y->right = z->right;
            z->right->parent = y;
        }
        else
            x_parent = y;
    }
}

```

```

    if (root == z)
        root = y;
    else if (z->parent->left == z)
        z->parent->left = y;
    else
        z->parent->right = y;
    y->parent = z->parent;
    __STD::swap(y->color, z->color);
    y = z;
    // y now points to node to be actually deleted
}
else {
    // y == z
    x_parent = y->parent;
    if (x) x->parent = y->parent;
    if (root == z)
        root = x;
    else
        if (z->parent->left == z)
            z->parent->left = x;
        else
            z->parent->right = x;
    if (leftmost == z)
        if (z->right == 0) // z->left must be null also
            leftmost = z->parent;
    // makes leftmost == header if z == root
    else
        leftmost = __rb_tree_node_base::minimum(x);
    if (rightmost == z)
        if (z->left == 0) // z->right must be null also
            rightmost = z->parent;
    // makes rightmost == header if z == root
    else // x == z->left
        rightmost = __rb_tree_node_base::maximum(x);
}
if (y->color != __rb_tree_red) {
    while (x != root && (x == 0 || x->color == __rb_tree_black))
        if (x == x_parent->left) {
            __rb_tree_node_base* w = x_parent->right;
            if (w->color == __rb_tree_red) {
                w->color = __rb_tree_black;
                x_parent->color = __rb_tree_red;
                __rb_tree_rotate_left(x_parent, root);
                w = x_parent->right;
            }
            if ((w->left == 0 || w->left->color == __rb_tree_black) &&
                (w->right == 0 || w->right->color == __rb_tree_black)) {
                w->color = __rb_tree_red;
                x = x_parent;
                x_parent = x_parent->parent;
            }
        }
}

```



```

    } else {
        if (w->right == 0 || w->right->color == __rb_tree_black) {
            if (w->left) w->left->color = __rb_tree_black;
            w->color = __rb_tree_red;
            __rb_tree_rotate_right(w, root);
            w = x_parent->right;
        }
        w->color = x_parent->color;
        x_parent->color = __rb_tree_black;
        if (w->right) w->right->color = __rb_tree_black;
        __rb_tree_rotate_left(x_parent, root);
        break;
    }
} else { // same as above, with right <=> left.
    __rb_tree_node_base* w = x_parent->left;
    if (w->color == __rb_tree_red) {
        w->color = __rb_tree_black;
        x_parent->color = __rb_tree_red;
        __rb_tree_rotate_right(x_parent, root);
        w = x_parent->left;
    }
    if ((w->right == 0 || w->right->color == __rb_tree_black) &&
        (w->left == 0 || w->left->color == __rb_tree_black)) {
        w->color = __rb_tree_red;
        x = x_parent;
        x_parent = x_parent->parent;
    } else {
        if (w->left == 0 || w->left->color == __rb_tree_black) {
            if (w->right) w->right->color = __rb_tree_black;
            w->color = __rb_tree_red;
            __rb_tree_rotate_left(w, root);
            w = x_parent->left;
        }
        w->color = x_parent->color;
        x_parent->color = __rb_tree_black;
        if (w->left) w->left->color = __rb_tree_black;
        __rb_tree_rotate_right(x_parent, root);
        break;
    }
}
    if (x) x->color = __rb_tree_black;
}
return y;
}

template <class Key, class Value, class KeyOfValue, class Compare,
          class Alloc = alloc>
class rb_tree {
protected:

```

```

typedef void* void_pointer;
typedef __rb_tree_node_base* base_ptr;
typedef __rb_tree_node<Value> rb_tree_node;
typedef simple_alloc<rb_tree_node, Alloc> rb_tree_node_allocator;
typedef __rb_tree_color_type color_type;
public:
    // 注意，沒有定義 iterator (喔，不，定義在後面)
    typedef Key key_type;
    typedef Value value_type;
    typedef value_type* pointer;
    typedef const value_type* const_pointer;
    typedef value_type& reference;
    typedef const value_type& const_reference;
    typedef rb_tree_node* link_type;
    typedef size_t size_type;
    typedef ptrdiff_t difference_type;
protected:
    link_type get_node() { return rb_tree_node_allocator::allocate(); }
    void put_node(link_type p) { rb_tree_node_allocator::deallocate(p); }

    link_type create_node(const value_type& x) {
        link_type tmp = get_node(); // 配置空間
        __STL_TRY {
            construct(&tmp->value_field, x); // 建構內容
        }
        __STL_UNWIND(put_node(tmp));
        return tmp;
    }

    link_type clone_node(link_type x) { // 複製一個節點 (的值和色)
        link_type tmp = create_node(x->value_field);
        tmp->color = x->color;
        tmp->left = 0;
        tmp->right = 0;
        return tmp;
    }

    void destroy_node(link_type p) {
        destroy(&p->value_field); // 解構內容
        put_node(p); // 釋還記憶體
    }

protected:
    // RB-tree 只以三筆資料表現。
    size_type node_count; // 追蹤記錄樹的大小 (節點數量)
    link_type header;
    Compare key_compare; // 節點間的鍵值大小比較準則。應該會是個 function object。

    // 以下三個函式用來方便取得 header 的成員

```

```

link_type& root() const { return (link_type&) header->parent; }
link_type& leftmost() const { return (link_type&) header->left; }
link_type& rightmost() const { return (link_type&) header->right; }

// 以下六個函式用來方便取得節點 x 的成員
static link_type& left(link_type x) { return (link_type&)(x->left); }
static link_type& right(link_type x) { return (link_type&)(x->right); }
static link_type& parent(link_type x) { return (link_type&)(x->parent); }
static reference value(link_type x) { return x->value_field; }
static const Key& key(link_type x) { return KeyOfValue()(value(x)); }
static color_type& color(link_type x) { return (color_type&)(x->color); }

// 以下六個函式用來方便取得節點 x 的成員
static link_type& left(base_ptr x) { return (link_type&)(x->left); }
static link_type& right(base_ptr x) { return (link_type&)(x->right); }
static link_type& parent(base_ptr x) { return (link_type&)(x->parent); }
static reference value(base_ptr x) { return ((link_type)x)->value_field; }
static const Key& key(base_ptr x) { return KeyOfValue()(value(link_type(x))); }
static color_type& color(base_ptr x) { return (color_type&)(link_type(x)->color); }

// 求取極大值和極小值。node class 有實作此功能，交給它們完成即可。
static link_type minimum(link_type x) {
    return (link_type) __rb_tree_node_base::minimum(x);
}
static link_type maximum(link_type x) {
    return (link_type) __rb_tree_node_base::maximum(x);
}

public:
    typedef __rb_tree_iterator<value_type, reference, pointer> iterator;
    typedef __rb_tree_iterator<value_type, const_reference, const_pointer>
        const_iterator;

#ifdef __STL_CLASS_PARTIAL_SPECIALIZATION
    typedef reverse_iterator<const_iterator> const_reverse_iterator;
    typedef reverse_iterator<iterator> reverse_iterator;
#else /* __STL_CLASS_PARTIAL_SPECIALIZATION */
    typedef reverse_bidirectional_iterator<iterator, value_type, reference,
        difference_type>
        reverse_iterator;
    typedef reverse_bidirectional_iterator<const_iterator, value_type,
        const_reference, difference_type>
        const_reverse_iterator;
#endif /* __STL_CLASS_PARTIAL_SPECIALIZATION */
private:
    iterator __insert(base_ptr x, base_ptr y, const value_type& v);
    link_type __copy(link_type x, link_type p);
    void __erase(link_type x);
    void init() {

```

```

    header = get_node();    // 產生一個節點空間，令 header 指向它
    color(header) = __rb_tree_red; // 令 header 為紅色，用來區分 header
                                   // 和 root (在 iterator.operator++ 中)

    root() = 0;
    leftmost() = header;    // 令 header 的左子節點為自己。
    rightmost() = header;   // 令 header 的右子節點為自己。
}
public:
                                   // allocation/deallocation
    rb_tree(const Compare& comp = Compare())
        : node_count(0), key_compare(comp) { init(); }

    // 以另一個 rb_tree 物件 x 為初值
    rb_tree(const rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& x)
        : node_count(0), key_compare(x.key_compare)
    {
        header = get_node();    // 產生一個節點空間，令 header 指向它
        color(header) = __rb_tree_red;    // 令 header 為紅色
        if (x.root() == 0) {    // 如果 x 是個空白樹
            root() = 0;
            leftmost() = header; // 令 header 的左子節點為自己。
            rightmost() = header; // 令 header 的右子節點為自己。
        }
        else {    // x 不是一個空白樹
            __STL_TRY {
                root() = __copy(x.root(), header);    // ???
            }
            __STL_UNWIND(put_node(header));
            leftmost() = minimum(root()); // 令 header 的左子節點為最小節點
            rightmost() = maximum(root()); // 令 header 的右子節點為最大節點
        }
        node_count = x.node_count;
    }
    ~rb_tree() {
        clear();
        put_node(header);
    }
    rb_tree<Key, Value, KeyOfValue, Compare, Alloc>&
    operator=(const rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& x);

public:
                                   // accessors:
    Compare key_comp() const { return key_compare; }
    iterator begin() { return leftmost(); }    // RB 樹的起頭為最左（最小）節點處
    const_iterator begin() const { return leftmost(); }
    iterator end() { return header; }    // RB 樹的終點為 header 所指處
    const_iterator end() const { return header; }
    reverse_iterator rbegin() { return reverse_iterator(end()); }
    const_reverse_iterator rbegin() const {

```

```

        return const_reverse_iterator(end());
    }
    reverse_iterator rend() { return reverse_iterator(begin()); }
    const_reverse_iterator rend() const {
        return const_reverse_iterator(begin());
    }
    bool empty() const { return node_count == 0; }
    size_type size() const { return node_count; }
    size_type max_size() const { return size_type(-1); }

    void swap(rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& t) {
        // RB-tree 只以三個資料成員表現。所以互換兩個 RB-trees時，
        // 只需將這三個成員互換即可。
        __STD::swap(header, t.header);
        __STD::swap(node_count, t.node_count);
        __STD::swap(key_compare, t.key_compare);
    }

public:
        // insert/erase
        // 將 x 安插到 RB-tree 中（保持節點值獨一無二）。
        pair<iterator, bool> insert_unique(const value_type& x);
        // 將 x 安插到 RB-tree 中（允許節點值重複）。
        iterator insert_equal(const value_type& x);

        iterator insert_unique(iterator position, const value_type& x);
        iterator insert_equal(iterator position, const value_type& x);

#ifdef __STL_MEMBER_TEMPLATES
        template <class InputIterator>
        void insert_unique(InputIterator first, InputIterator last);
        template <class InputIterator>
        void insert_equal(InputIterator first, InputIterator last);
#else /* __STL_MEMBER_TEMPLATES */
        void insert_unique(const_iterator first, const_iterator last);
        void insert_unique(const value_type* first, const value_type* last);
        void insert_equal(const_iterator first, const_iterator last);
        void insert_equal(const value_type* first, const value_type* last);
#endif /* __STL_MEMBER_TEMPLATES */

        void erase(iterator position);
        size_type erase(const key_type& x);
        void erase(iterator first, iterator last);
        void erase(const key_type* first, const key_type* last);
        void clear() {
            if (node_count != 0) {
                __erase(root());
                leftmost() = header;
                root() = 0;
            }
        }

```

```

        rightmost() = header;
        node_count = 0;
    }
}

public:
    // 集合 (set) 的各種操作行為：
    iterator find(const key_type& x);
    const_iterator find(const key_type& x) const;
    size_type count(const key_type& x) const;
    iterator lower_bound(const key_type& x);
    const_iterator lower_bound(const key_type& x) const;
    iterator upper_bound(const key_type& x);
    const_iterator upper_bound(const key_type& x) const;
    pair<iterator, iterator> equal_range(const key_type& x);
    pair<const_iterator, const_iterator> equal_range(const key_type& x) const;

public:
    // Debugging.
    bool __rb_verify() const;
};

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
inline bool operator==(const rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& x,
                      const rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& y) {
    return x.size() == y.size() && equal(x.begin(), x.end(), y.begin());
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
inline bool operator<(const rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& x,
                    const rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& y) {
    return lexicographical_compare(x.begin(), x.end(), y.begin(), y.end());
}

#ifdef __STL_FUNCTION_TMPL_PARTIAL_ORDER

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
inline void swap(rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& x,
               rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& y) {
    x.swap(y);
}

#endif /* __STL_FUNCTION_TMPL_PARTIAL_ORDER */

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>&
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::
operator=(const rb_tree<Key, Value, KeyOfValue, Compare, Alloc>& x) {

```

```

    if (this != &x) {
        // Note that Key may be a constant type.
        clear();
        node_count = 0;
        key_compare = x.key_compare;
        if (x.root() == 0) {
            root() = 0;
            leftmost() = header;
            rightmost() = header;
        }
        else {
            root() = __copy(x.root(), header);
            leftmost() = minimum(root());
            rightmost() = maximum(root());
            node_count = x.node_count;
        }
    }
    return *this;
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::
__insert(base_ptr x_, base_ptr y_, const Value& v) {
    // 參數x_ 為新值安插點，參數y_ 為安插點之父節點，參數v 為新值。
    link_type x = (link_type) x_;
    link_type y = (link_type) y_;
    link_type z;

    // key_compare 是鍵值大小比較準則。應該會是個 function object。
    if (y == header || x != 0 || key_compare(KeyOfValue()(v), key(y))) {
        z = create_node(v); // 產生一個新節點
        left(y) = z; // 這使得當 y 即為 header時，leftmost() = z
        if (y == header) {
            root() = z;
            rightmost() = z;
        }
        else if (y == leftmost()) // 如果y為最左節點
            leftmost() = z; // 維護leftmost()，使它永遠指向最左節點
    }
    else {
        z = create_node(v); // 產生一個新節點
        right(y) = z; // 令新節點成為安插點之父節點 y 的右子節點
        if (y == rightmost())
            rightmost() = z; // 維護rightmost()，使它永遠指向最右節點
    }
    parent(z) = y; // 設定新節點之父節點
    left(z) = 0; // 設定新節點的左子節點
    right(z) = 0; // 設定新節點的右子節點
}

```

```

        // 新節點的顏色將在 __rb_tree_rebalance() 設定 (並調整)
__rb_tree_rebalance(z, header->parent); // 參數一為新增節點, 參數二為 root
++node_count; // 節點數累加
return iterator(z); // 傳回一個迭代器, 指向新增節點
}

// 安插新值: 節點鍵值允許重複。
// 注意, 傳回值是一個 RB-tree 迭代器, 指向新增節點
template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::insert_equal(const Value& v)
{
    link_type y = header;
    link_type x = root(); // 從根節點開始
    while (x != 0) { // 從根節點開始, 往下尋找適當的安插點
        y = x;
        x = key_compare(KeyOfValue()(v), key(x)) ? left(x) : right(x);
        // 以上, 遇「大」則往左, 遇「小於或等於」則往右
    }
    return __insert(x, y, v);
}

// 安插新值: 節點鍵值不允許重複, 若重複則安插無效。
// 注意, 傳回值是個pair, 第一元素是個 RB-tree 迭代器, 指向新增節點,
// 第二元素表示安插成功與否。
template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
pair<typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator, bool>
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::insert_unique(const Value& v)
{
    link_type y = header;
    link_type x = root(); // 從根節點開始
    bool comp = true;
    while (x != 0) { // 從根節點開始, 往下尋找適當的安插點
        y = x;
        comp = key_compare(KeyOfValue()(v), key(x)); // v 鍵值小於目前節點之鍵值?
        x = comp ? left(x) : right(x); // 遇「大」則往左, 遇「小於或等於」則往右
    }
    // 離開 while 迴圈之後, y 所指即安插點之父節點 (此時的它必為葉節點)

    iterator j = iterator(y); // 令迭代器 j 指向安插點之父節點 y
    if (comp) // 如果離開 while 迴圈時 comp 為真 (表示遇「大」, 將安插於左側)
        if (j == begin()) // 如果安插點之父節點為最左節點
            return pair<iterator, bool>(__insert(x, y, v), true);
            // 以上, x 為安插點, y 為安插點之父節點, v 為新值。
        else // 否則 (安插點之父節點不為最左節點)
            --j; // 調整 j, 回頭準備測試...
    if (key_compare(key(j.node), KeyOfValue()(v)))
        // 小於新值 (表示遇「小」, 將安插於右側)
        return pair<iterator, bool>(__insert(x, y, v), true);
}

```



```

// 進行至此，表示新值一定與樹中鍵值重複，那麼就不該插入新值。
return pair<iterator,bool>(j, false);
}

template <class Key, class Val, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Val, KeyOfValue, Compare, Alloc>::iterator
rb_tree<Key, Val, KeyOfValue, Compare, Alloc>::insert_unique(iterator position,
                                                             const Val& v) {
    if (position.node == header->left) // begin()
        if (size() > 0 && key_compare(KeyOfValue()(v), key(position.node)))
            return __insert(position.node, position.node, v);
    // first argument just needs to be non-null
    else
        return insert_unique(v).first;
    else if (position.node == header) // end()
        if (key_compare(key(rightmost()), KeyOfValue()(v)))
            return __insert(0, rightmost(), v);
        else
            return insert_unique(v).first;
    else {
        iterator before = position;
        --before;
        if (key_compare(key(before.node), KeyOfValue()(v))
            && key_compare(KeyOfValue()(v), key(position.node)))
            if (right(before.node) == 0)
                return __insert(0, before.node, v);
            else
                return __insert(position.node, position.node, v);
    // first argument just needs to be non-null
    else
        return insert_unique(v).first;
    }
}

template <class Key, class Val, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Val, KeyOfValue, Compare, Alloc>::iterator
rb_tree<Key, Val, KeyOfValue, Compare, Alloc>::insert_equal(iterator position,
                                                             const Val& v) {
    if (position.node == header->left) // begin()
        if (size() > 0 && key_compare(KeyOfValue()(v), key(position.node)))
            return __insert(position.node, position.node, v);
    // first argument just needs to be non-null
    else
        return insert_equal(v);
    else if (position.node == header) // end()
        if (!key_compare(KeyOfValue()(v), key(rightmost())))
            return __insert(0, rightmost(), v);

```

```

        else
            return insert_equal(v);
    else {
        iterator before = position;
        --before;
        if (!key_compare(KeyOfValue()(v), key(before.node))
            && !key_compare(key(position.node), KeyOfValue()(v)))
            if (right(before.node) == 0)
                return __insert(0, before.node, v);
            else
                return __insert(position.node, position.node, v);
        // first argument just needs to be non-null
    }
    else
        return insert_equal(v);
}

#ifdef __STL_MEMBER_TEMPLATES

template <class K, class V, class KoV, class Cmp, class Al> template<class II>
void rb_tree<K, V, KoV, Cmp, Al>::insert_equal(II first, II last) {
    for ( ; first != last; ++first)
        insert_equal(*first);
}

template <class K, class V, class KoV, class Cmp, class Al> template<class II>
void rb_tree<K, V, KoV, Cmp, Al>::insert_unique(II first, II last) {
    for ( ; first != last; ++first)
        insert_unique(*first);
}

#else /* __STL_MEMBER_TEMPLATES */

template <class K, class V, class KoV, class Cmp, class Al>
void
rb_tree<K, V, KoV, Cmp, Al>::insert_equal(const V* first, const V* last) {
    for ( ; first != last; ++first)
        insert_equal(*first);
}

template <class K, class V, class KoV, class Cmp, class Al>
void
rb_tree<K, V, KoV, Cmp, Al>::insert_equal(const_iterator first,
                                         const_iterator last) {
    for ( ; first != last; ++first)
        insert_equal(*first);
}

template <class K, class V, class KoV, class Cmp, class A>

```

```

void
rb_tree<K, V, KoV, Cmp, A>::insert_unique(const V* first, const V* last) {
    for ( ; first != last; ++first)
        insert_unique(*first);
}

template <class K, class V, class KoV, class Cmp, class A>
void
rb_tree<K, V, KoV, Cmp, A>::insert_unique(const_iterator first,
                                         const_iterator last) {
    for ( ; first != last; ++first)
        insert_unique(*first);
}

#endif /* __STL_MEMBER_TEMPLATES */

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
inline void
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::erase(iterator position) {
    link_type y = (link_type) __rb_tree_rebalance_for_erase(position.node,
                                                             header->parent,
                                                             header->left,
                                                             header->right);

    destroy_node(y);
    --node_count;
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::size_type
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::erase(const Key& x) {
    pair<iterator, iterator> p = equal_range(x);
    size_type n = 0;
    distance(p.first, p.second, n);
    erase(p.first, p.second);
    return n;
}

template <class K, class V, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<K, V, KeyOfValue, Compare, Alloc>::link_type
rb_tree<K, V, KeyOfValue, Compare, Alloc>::__copy(link_type x, link_type p) {
    // structural copy. x and p must be non-null.
    link_type top = clone_node(x);
    top->parent = p;

    __STL_TRY {
        if (x->right)
            top->right = __copy(right(x), top);
        p = top;
        x = left(x);
    }
}

```

```

        while (x != 0) {
            link_type y = clone_node(x);
            p->left = y;
            y->parent = p;
            if (x->right)
                y->right = __copy(right(x), y);
            p = y;
            x = left(x);
        }
    }
    __STL_UNWIND(__erase(top));

    return top;
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
void rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::__erase(link_type x) {
    // erase without rebalancing
    while (x != 0) {
        __erase(right(x));
        link_type y = left(x);
        destroy_node(x);
        x = y;
    }
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
void rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::erase(iterator first,
                                                             iterator last) {
    if (first == begin() && last == end())
        clear();
    else
        while (first != last) erase(first++);
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
void rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::erase(const Key* first,
                                                             const Key* last) {
    while (first != last) erase(*first++);
}

// 尋找 RB 樹中是否有鍵值為 k 的節點
template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::find(const Key& k) {
    link_type y = header;           // Last node which is not less than k.
    link_type x = root();           // Current node.

```

```

while (x != 0)
    // 以下，key_compare 是節點鍵值大小比較準則。應該會是個 function object。
    if (!key_compare(key(x), k))
        // 進行到這裡，表示 x 鍵值大於 k。遇到大值就向左走。
        y = x, x = left(x);    // 注意語法!
    else
        // 進行到這裡，表示 x 鍵值小於 k。遇到小值就向右走。
        x = right(x);

    iterator j = iterator(y);
    return (j == end() || key_compare(k, key(j.node))) ? end() : j;
}

// 尋找 RB 樹中是否有鍵值為 k 的節點
template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::const_iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::find(const Key& k) const {
    link_type y = header; /* Last node which is not less than k. */
    link_type x = root(); /* Current node. */

    while (x != 0) {
        // 以下，key_compare 是節點鍵值大小比較準則。應該會是個 function object。
        if (!key_compare(key(x), k))
            // 進行到這裡，表示 x 鍵值大於 k。遇到大值就向左走。
            y = x, x = left(x);    // 注意語法!
        else
            // 進行到這裡，表示 x 鍵值小於 k。遇到小值就向右走。
            x = right(x);
    }
    const_iterator j = const_iterator(y);
    return (j == end() || key_compare(k, key(j.node))) ? end() : j;
}

// 計算 RB 樹中鍵值為 k 的節點個數
template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::size_type
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::count(const Key& k) const {
    pair<const_iterator, const_iterator> p = equal_range(k);
    size_type n = 0;
    distance(p.first, p.second, n);
    return n;
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::lower_bound(const Key& k) {
    link_type y = header; /* Last node which is not less than k. */
    link_type x = root(); /* Current node. */

```

```

    while (x != 0)
        if (!key_compare(key(x), k))
            y = x, x = left(x);
        else
            x = right(x);

    return iterator(y);
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::const_iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::lower_bound(const Key& k) const {
    link_type y = header; /* Last node which is not less than k. */
    link_type x = root(); /* Current node. */

    while (x != 0)
        if (!key_compare(key(x), k))
            y = x, x = left(x);
        else
            x = right(x);

    return const_iterator(y);
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::upper_bound(const Key& k) {
    link_type y = header; /* Last node which is greater than k. */
    link_type x = root(); /* Current node. */

    while (x != 0)
        if (key_compare(k, key(x)))
            y = x, x = left(x);
        else
            x = right(x);

    return iterator(y);
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::const_iterator
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::upper_bound(const Key& k) const {
    link_type y = header; /* Last node which is greater than k. */
    link_type x = root(); /* Current node. */

    while (x != 0)
        if (key_compare(k, key(x)))
            y = x, x = left(x);
        else
            x = right(x);

```

```

        x = right(x);

    return const_iterator(y);
}

template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
inline pair<typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator,
           typename rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::iterator>
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::equal_range(const Key& k) {
    return pair<iterator, iterator>(lower_bound(k), upper_bound(k));
}

template <class Key, class Value, class KoV, class Compare, class Alloc>
inline pair<typename rb_tree<Key, Value, KoV, Compare, Alloc>::const_iterator,
           typename rb_tree<Key, Value, KoV, Compare, Alloc>::const_iterator>
rb_tree<Key, Value, KoV, Compare, Alloc>::equal_range(const Key& k) const {
    return pair<const_iterator, const_iterator>(lower_bound(k), upper_bound(k));
}

// 計算從 node 至 root 路徑中的黑節點數量。
inline int __black_count(__rb_tree_node_base* node, __rb_tree_node_base* root)
{
    if (node == 0)
        return 0;
    else {
        int bc = node->color == __rb_tree_black ? 1 : 0;
        if (node == root)
            return bc;
        else
            return bc + __black_count(node->parent, root); // 累加
    }
}

// 驗證己身這棵樹是否符合 RB 樹的條件
template <class Key, class Value, class KeyOfValue, class Compare, class Alloc>
bool
rb_tree<Key, Value, KeyOfValue, Compare, Alloc>::__rb_verify() const
{
    // 空樹，符合RB樹標準
    if (node_count == 0 || begin() == end())
        return node_count == 0 && begin() == end() &&
            header->left == header && header->right == header;

    // 最左（葉）節點至 root 路徑內的黑節點數
    int len = __black_count(leftmost(), root());
    // 以下走訪整個RB樹，針對每個節點（從最小到最大）...
    for (const_iterator it = begin(); it != end(); ++it) {
        link_type x = (link_type) it->node; // __rb_tree_base_iterator::node
        link_type L = left(x);             // 這是左子節點
    }
}

```

```

link_type R = right(x);      // 這是右子節點

if (x->color == __rb_tree_red)
    if ((L && L->color == __rb_tree_red) ||
        (R && R->color == __rb_tree_red))
        return false;  // 父子節點同為紅色，不符合 RB 樹的要求。

if (L && key_compare(key(x), key(L))) // 目前節點的鍵值小於左子節點鍵值
    return false;                    // 不符合二元搜尋樹的要求。
if (R && key_compare(key(R), key(x))) // 目前節點的鍵值大於右子節點鍵值
    return false;                    // 不符合二元搜尋樹的要求。

// 「葉節點至 root」路徑內的黑節點數，與「最左節點至 root」路徑內的黑節點數不同。
// 這不符合 RB 樹的要求。
if (!L && !R && __black_count(x, root()) != len)
    return false;
}

if (leftmost() != __rb_tree_node_base::minimum(root()))
    return false; // 最左節點不為最小節點，不符合二元搜尋樹的要求。
if (rightmost() != __rb_tree_node_base::maximum(root()))
    return false; // 最右節點不為最大節點，不符合二元搜尋樹的要求。

return true;
}

__STL_END_NAMESPACE

#endif /* __SGI_STL_INTERNAL_TREE_H */

// Local Variables:
// mode:C++
// End:

```