Combining Verifiers in Conditional Model Checking¹

Marie-Christine Jakobs

Joint work with Dirk Beyer, Thomas Lemberger, and Heike Wehrheim

LMU Munich, Germany





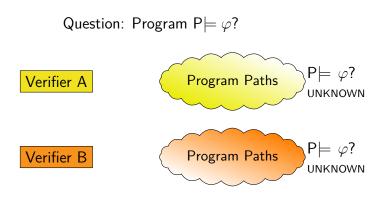


¹https://www.sosy-lab.org/research/reducer/

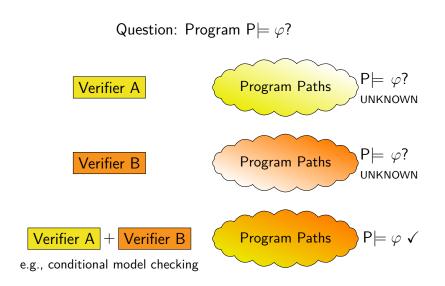
Many Verification Tools Available



Facing Hard Verification Tasks



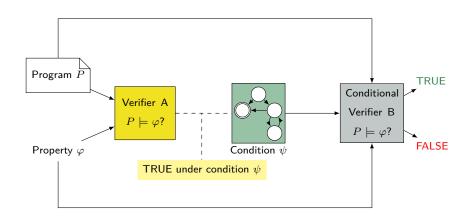
Facing Hard Verification Tasks



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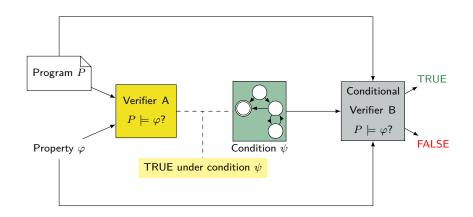
Conditional Model Checking

[Beyer/Henzinger/Keremoglu/Wendler FSE'12]



Conditional Model Checking

[Beyer/Henzinger/Keremoglu/Wendler FSE'12]

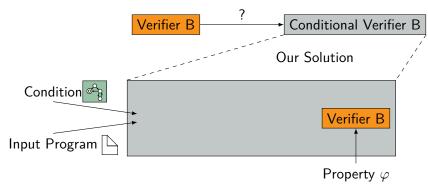


Problem: Often, verifiers are not conditional

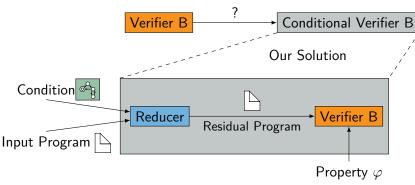
[Beyer/Jakobs/Lemberger/Wehrheim ICSE'18]

Verifier B ? Conditional Verifier B

[Beyer/Jakobs/Lemberger/Wehrheim ICSE'18]



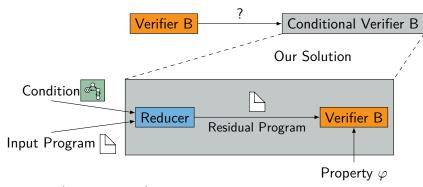
[Beyer/Jakobs/Lemberger/Wehrheim ICSE'18]



Reducer (preprocessor)

- ► Builds standard input (C program)
- Representing a subset of paths
- Contains at least all non-verified paths

[Beyer/Jakobs/Lemberger/Wehrheim ICSE'18]



Reducer (preprocessor)

- Builds standard input (C program)
- Representing a subset of paths
- Contains at least all non-verified paths
- + Verifier-unspecific approach + Many conditional verifiers

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Example Program and Condition

```
Program

0: if (notThursday)

1: discount=day%7;
else

2: discount=5;
3: assert(0<=discount<7);
4:

Program

OntThursday

OntThursday

Odiscount=day%7;

Odiscount=5;

Odiscount=6;

Odiscount=7);

Odiscount=7);

Odiscount=7);
```

Example Program and Condition

```
Program

0: if(notThursday)

1: discount=day%7;
else

2: discount=5;
3: assert(0<=discount<7);
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assert(0<=discount<7);

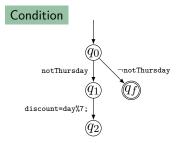
assert(0<=discount<7);

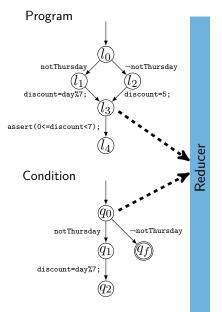
assert(0<=discount<7);
```

Verifier A only proofs else branch

Example Program and Condition

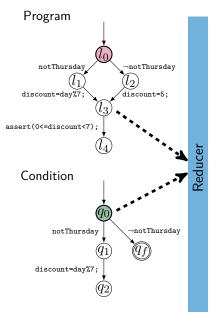
Verifier A only proofs else branch

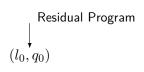




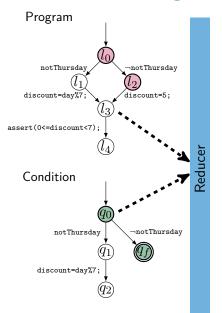
Residual Program

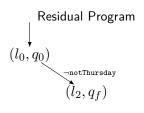
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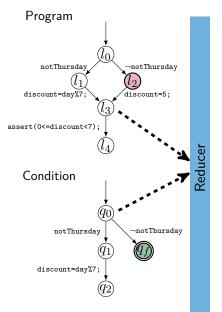


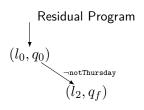
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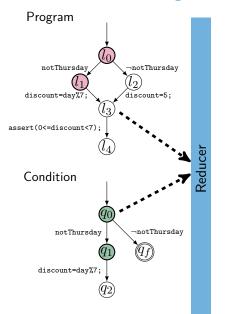


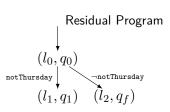
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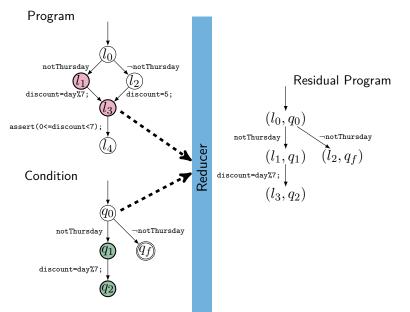


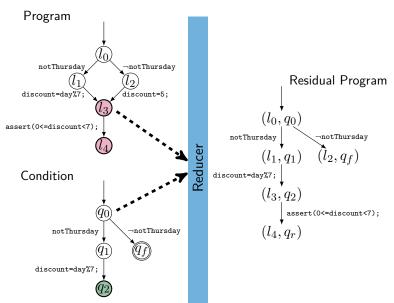
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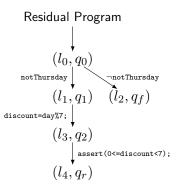


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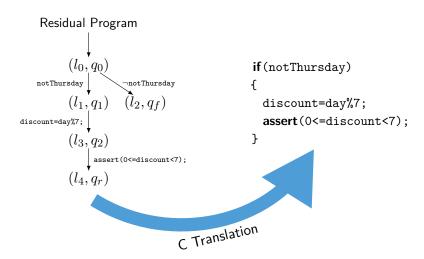




Reducer: C Transformation

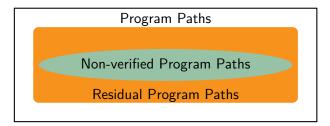


Reducer: C Transformation



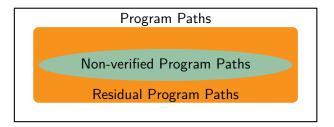
Reducer: Soundness

Residual Condition



Reducer: Soundness

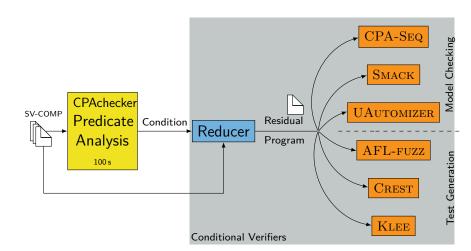
Residual Condition



Theorem

Presented reducer fulfills residual condition.

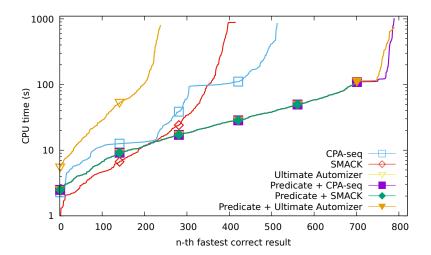
Evaluation Setup



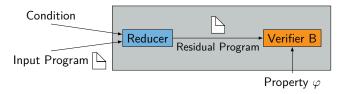
Small Extract of Results

		CPA-SEQ UAUT			PREDICATE		PREDICATE		
				UAUTOMIZER		+Reducer		+Reducer	
						+CPA-SEQ		+UAUTOMIZER	
Task	R	S	t(s)	S	t(s)	S	t(s)	S	t(s)
P15I01	Т	Х	910	X	900	1	120	1	130
flood4	Т	X	910	X	910	1	450	X	1100
newt3_6	F	X	950	X	490	X	910	1	260

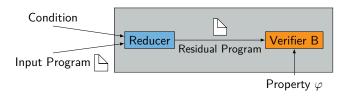
Effectiveness on Hard Tasks



Template-based conditional verifier construction

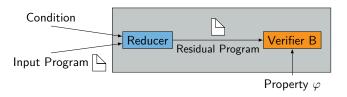


Template-based conditional verifier construction



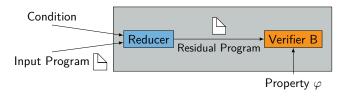
- One Reducer
 - Proven sound
 - Used in many conditional verifiers

Template-based conditional verifier construction



- One Reducer
 - Proven sound
 - Used in many conditional verifiers
- ▶ Effective on hard tasks for verifiers and test tools

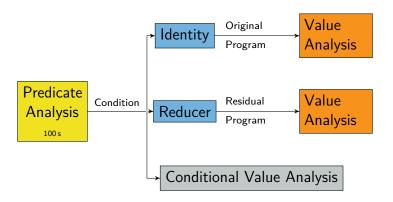
Template-based conditional verifier construction



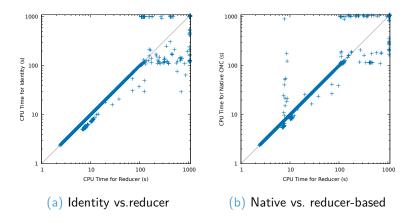
- One Reducer
 - Proven sound
 - Used in many conditional verifiers
- Effective on hard tasks for verifiers and test tools
- Future Work
 - More reducers
 - Using conditions from other tools

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Comparison Setup



Comparison Results



References I



D. Bever, T. A. Henzinger, M. E. Keremoglu, and P. Wendler. Conditional Model Checking: A Technique to Pass Information Between Verifiers. In Proc. FSE. ACM, 2012.



D. Beyer, M.-C. Jakobs, T. Lemberger, and H. Wehrheim. Reducer-Based Construction of Conditional Verifiers. In Proc. ICSE. ACM. 2018.