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Team Note of Powered by Zigui

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ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

0 Quotes and Prerequisites

```
evenharder : Mental Abuse To Humans
djkim0613 : 열심히 응원하겠습니다.
SoulTch : How much is this bus ticket?
* This template is brought from that of 'Deobureo Minkyu Party'
```

Run script

```
#!/bin/bash
g++ -fsanitize=undefined -std=c++14 -02 -o /tmp/pow $1.cpp
time /tmp/pow < $1.in
# export PATH=~:$PATH</pre>
```

Debug Code

```
#define setz(x) memset(x, 0, sizeof(x))
#define sz(x) ((int)(x).size())
#define rep(i, e) for (int i = 0, _##i = (e); i < _##i; i++)
#define repp(i, s, e) for (int i = (s), _##i = (e); i < _##i; i++)
#define repr(i, s, e) for (int i = (s)-1, _##i = (e); i \ge _{\#}i; i--)
#define repi(i, x) for (auto &i : (x))
// using namespace std;
using ll = long long;
using pii = pair<int, int>;
using pll = pair<11, 11>;
template<typename T>
ostream &operator<<(ostream &os, const vector<T>& v) {
    cout << "[":
    for (auto p : v) cout << p << ",";</pre>
    cout << "]":
    return os;
}
#ifndef SOULTCH
#define debug(...) 0
#define endl '\n'
#define debug(...) cout << " [-] ", _dbg(#__VA_ARGS__, __VA_ARGS__)</pre>
template<class TH> void _dbg(const char *sdbg, TH h){ cout << sdbg << '=' << h <<
endl; }
template<class TH, class... TA> void _dbg(const char *sdbg, TH h, TA... a) {
    while(*sdbg != ',') cout << *sdbg++;</pre>
    cout << '=' << (h) << ',';
    _dbg(sdbg+1, a...);
}
#endif
```

Reminders

Should be added.

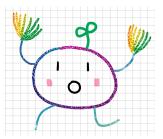


Figure 1: 풀다가 막힐 때는 이 그림을 봅시다. 아자아자 화이팅!

1 Math

1.1 Basic Mathematics

1.1.1 Trigonometry

- $\sin(\alpha \pm \beta) = \sin \alpha \cos \beta \pm \cos \alpha \sin \beta$
- $\cos(\alpha \pm \beta) = \cos \alpha \cos \beta \mp \sin \alpha \sin \beta$
- $\tan(\alpha \pm \beta) = \frac{\tan \alpha \pm \tan \beta}{1 \mp \tan \alpha \tan \beta}$
- $\sin 2\theta = 2\sin \theta \cos \theta$
- $\bullet \ c^2 = a^2 + b^2 2ab\cos\gamma$

1.1.2 Generating Function

- $\sum_{n} (pn+q)x^{n} = \frac{p}{1-x} + \frac{q}{(1-x)^{2}}$ (Arithmetic progression)
- $\sum_{n} (rx)^n = (1 rx)^{-1}$ (Geometric progression)
- $\sum_{n} {m \choose n} x^n = (1+x)^m$ (Binomial coefficient)
- $\sum_{n} {m+n-1 \choose n} x^n = (1-x)^{-m}$ (Multiset coefficient)

1.1.3 Calculus

- $\int_a^b f(x) dx = \frac{b-a}{6} \left[f(a) + 4f(\frac{a+b}{2}) + f(b) \right]$ (Simpson's Rule, for cubic poly)
- $\int u'v \ dx = uv \int uv' \ dx$ (Integration by parts)

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1.2 Number Theory

1.2.1 Lattice Points under Line

for(ll j=0;j<m;j++) {</pre>

```
// 0 \le x \le n, 0 \le y \le (a/c)x+(b/c)
11 calc(l1 a,l1 b,l1 c,l1 n){
    if(!n)return 0;
    11 tmp=a/c*n*(n-1)/2;
    tmp+=b/c*n:
    return tmp+calc(c,(a*n+b)%c,a%c,((a%c)*n+b%c)/c);
}
1.2.2 Shanks' Baby-step Giant-step
  Should be revised.
11 mexp(ll x, ll y, ll p) {
    if(!v) return 1:
    if (y \& 1) return x * mexp(x*x%p, y>>1, p) % p;
    return mexp(x*x\%p, y>>1, p);
}
vector<ll> get_factor(ll n) {
    vector<ll> v;
    for(ll i=2:i*i<=n:i++) {
        if(n \% i == 0) {
            v.push_back(i);
            while(n \% i == 0) n /= i:
        }
    }
    if(n > 1) v.push_back(n);
    return v:
}
ll get_primitive(ll n) {
    ll phi = n-1; // assume n is prime
    vector<ll> fact = get_factor(phi);
    for(11 x=2;x<=n;x++) {
        int ves = 1:
        for(ll y : fact) {
            yes &= (mexp(x, phi / y, n) != 1);
        if(yes) return x;
    }
    return -1;
// find x s.t. x^k \mod n = a \rightarrow (g^k)^y \mod n = a, where x = g^y
11 bsgs(ll k, ll a, ll n) {
    11 g = get_primitive(n);
    11 phi = n-1; // assume n is prime
    if(g == -1) return -1;
    ll m = ceil(sqrt(n) + 1e-9);
    vector<pl> prec(m);
```

```
prec[j] = {mexp(g, j * k % phi, n), j};
   }
   sort(prec.begin(), prec.end());
   ll cur = a, ncur = mexp(g, (phi - m) * k % phi, n);
   for(ll i=0:i<m:i++) {
        auto it = lower_bound(prec.begin(), prec.end(), pl(cur, 0));
        if(it->first == cur) {
           ll ans = mexp(g, (i*m + it->second) \% phi, n);
            assert(mexp(ans, k, n) == a);
            return ans:
        cur = cur * ncur % n;
   }
   return 0;
}
1.2.3 Extended Euclidean Algorithm
// ax + by = gcd(a,b). x, y?
pll ext_gcd(ll a,ll b) {
   if(b) {
        auto tmp = ext_gcd(b, a%b);
        return {tmp.second, tmp.first - (a/b) * tmp.second};
   }
    else return {1, 0};
// ax = gcd(a, m) mod m. x?
11 mod inv(ll a, ll m) {
   return (ext_gcd(a, m).first + m) % m;
1.2.4 Chinese Remainder Theorem
ll pos_rem(ll a, ll m) { // m > 0. a % m?
   11 \text{ res} = abs(a) \% m:
   return a > 0 ? res : (res ? m - res : 0);
// ax = c mod m, bx = d mod n. x?
11 solve(ll a, ll c, ll m, ll b, ll d, ll n) {
    a = pos_rem(a, m); c = pos_rem(c, m); // if a, c not in [0, m)
   b = pos_rem(b, n); d = pos_rem(d, n); // if b, d not in [0, n)
   11 g = gcd(a, gcd(c, m)); a \neq g, c \neq g, m \neq g;
        g = gcd(b, gcd(d, n)); b /= g, d /= g, n /= g;
   if(c % _gcd(a, m) || d % _gcd(b, n)) return inf;
   ll t1 = (mod_inv(a, m) * c) % m;
   11 t2 = (mod inv(b, n) * d) \% n:
    g = gcd(m, n);
   11 lc = m * n / g;
    if(abs(t1 - t2) % g) return inf;
    pl p = ext_gcd(m, n);
```

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```
11 q = (t1 * p.second * n/g + t2 * p.first * m/g);
    return pos_rem(q, lc);
}
1.2.5 Möbius Inversion Formula
  \forall n \in \mathbb{N} \ g(n) = \sum_{d \mid n} f(d) \implies f(n) = \sum_{d \mid n} \mu(d)g(n/d)
1.3 FFT
  FFT: (a_0, a_1, \dots, a_{n-1}) \mapsto (\sum_{i=0}^{n-1} a_0(\omega^0)^j, \sum_{i=0}^{n-1} a_1(\omega^1)^j, \dots, \sum_{i=0}^{n-1} a_{n-1}(\omega^{n-1})^j)
void fft(vector<base>& a, bool inv) {
    int n = a.size(), j = 0;
    vector<ll> roots(n/2);
    for(int i=1;i<n;i++) {</pre>
         int bit = (n >> 1):
         while(j >= bit) {
             j -= bit;
             bit >>= 1;
         }
         j += bit;
         if(i < j) swap(a[i], a[j]);</pre>
    }
    double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
    for(int i=0;i<n/2;i++) {
         roots[i] = base(cos(ang * i), sin(ang * i));
    }
    /* In NTT, let prr = primitive root. Then,
    int ang = mexp(prr, (mod - 1) / n);
    if(inv) ang = mexp(ang, mod - 2);
    for(int i=0: i<n/2: i++){
         roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
    also, make sure to apply modulus under here
    for(int i=2:i<=n:i<<=1) {</pre>
         int step = n / i;
         for(int j=0;j<n;j+=i) {</pre>
             for(int k=0;k<i/2;k++) {</pre>
                  ll u = a[j+k], v = a[j+k+i/2] * roots[step * k];
                  a[i+k] = u+v:
                  a[j+k+i/2] = u-v;
             }
         }
    if(inv) for(int i=0;i<n;i++) a[i] /= n;
}
```

```
void conv(vector<base>& x, vector<base>& y) {
   int n = 2; while (n < max(x.size(), y.size())) n <<= 1;
   n <<= 1:
   x.resize(n); y.resize(n);
   fft(x, false): fft(v, false):
   for(int i=0;i<n;i++) x[i] *= y[i];</pre>
   fft(x, true); // access (ll)round(x[i].real())
1.4 Miller-Rabin + Pollard-Rho
//Prove By Solving - https://www.acmicpc.net/problem/4149
namespace miller_rabin{
   lint mul(lint x, lint y, lint mod) { return (_int128) x * y % mod; }
   lint ipow(lint x, lint y, lint p){
       lint ret = 1, piv = x \% p;
        while(v){
           if(y&1) ret = mul(ret, piv, p);
           piv = mul(piv, piv, p);
           y >>= 1;
       return ret;
   }
   bool miller_rabin(lint x, lint a){
        if (x \% a == 0) return 0;
       lint d = x - 1:
        while(1){
           lint tmp = ipow(a, d, x);
           if(d&1) return (tmp != 1 && tmp != x-1);
           else if(tmp == x-1) return 0;
           d >>= 1;
       }
   7
   bool isprime(lint x){
       for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
           if(x == i) return 1:
           if(x > 40 && miller_rabin(x, i)) return 0;
        if(x \le 40) return 0:
        return 1;
   }
}
namespace pollard_rho{
   lint f(lint x, lint n, lint c){
        return (c + miller_rabin::mul(x, x, n)) % n;
   void rec(lint n, vector<lint> &v){
       if(n == 1) return;
        if(n \% 2 == 0){
           v.push_back(2);
```

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```
rec(n/2, v);
            return:
        }
        if(miller_rabin::isprime(n)){
            v.push back(n):
            return;
        }
       lint a, b, c;
        while(1){
            a = rand() \% (n-2) + 2:
           b = a:
            c = rand() \% 20 + 1;
            do{
                a = f(a, n, c);
                b = f(f(b, n, c), n, c);
            }while(gcd(abs(a-b), n) == 1);
            if(a != b) break;
        lint x = gcd(abs(a-b), n);
        rec(x, v);
        rec(n/x, v):
    }
    vector<lint> factorize(lint n){
        vector<lint> ret;
        rec(n, ret);
        sort(ret.begin(), ret.end());
        return ret:
   }
}:
```

1.5 Black Box Linear Algebra + Kitamasa

```
vector<int> berlekamp_massey(vector<int> x){
    vector<int> ls. cur:
    int lf, ld;
   for(int i=0: i<x.size(): i++){</pre>
       lint t = 0:
        for(int j=0; j<cur.size(); j++){</pre>
            t = (t + 111 * x[i-j-1] * cur[j]) \% mod;
        if((t - x[i]) \% mod == 0) continue;
        if(cur.empty()){
            cur.resize(i+1):
            lf = i:
            1d = (t - x[i]) \% mod;
            continue;
        }
        lint k = -(x[i] - t) * ipow(ld, mod - 2) % mod;
        vector<int> c(i-lf-1);
        c.push_back(k);
        for(auto &j : ls) c.push_back(-j * k % mod);
```

```
if(c.size() < cur.size()) c.resize(cur.size());</pre>
        for(int j=0; j<cur.size(); j++){</pre>
            c[i] = (c[i] + cur[i]) \% mod;
        if(i-lf+(int)ls.size()>=(int)cur.size()){
            tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) % mod);
        cur = c:
   for(auto &i : cur) i = (i % mod + mod) % mod:
   return cur:
int get_nth(vector<int> rec, vector<int> dp, lint n){
   int m = rec.size();
    vector<int> s(m), t(m);
   s[0] = 1:
    if(m != 1) t[1] = 1;
   else t[0] = rec[0];
   auto mul = [&rec](vector<int> v, vector<int> w){
        int m = v.size();
        vector<int> t(2 * m):
        for(int j=0; j<m; j++){</pre>
            for(int k=0; k<m; k++){</pre>
                t[j+k] += 111 * v[j] * w[k] % mod;
                if(t[j+k] >= mod) t[j+k] -= mod;
            }
        }
        for(int j=2*m-1; j>=m; j--){
            for(int k=1: k<=m: k++){
                t[j-k] += 111 * t[j] * rec[k-1] % mod;
                if(t[j-k] >= mod) t[j-k] -= mod;
           }
        t.resize(m):
        return t:
   };
   while(n){
        if(n \& 1) s = mul(s, t);
        t = mul(t, t);
        n >>= 1:
   lint ret = 0;
   for(int i=0; i<m; i++) ret += 111 * s[i] * dp[i] % mod;
   return ret % mod;
int guess_nth_term(vector<int> x, lint n){ // init with > 3k, 0(1^2 lg n)
    if(n < x.size()) return x[n];</pre>
    vector<int> v = berlekamp_massey(x);
   if(v.empty()) return 0;
   return get_nth(v, x, n);
```

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```
}
struct elem{int x, y, v;}; // A_(x, y) <- v, 0-based. no duplicate please..
vector<int> get min polv(int n. vector<elem> M){
    // smallest poly P such that A^i = sum_{j < i} {A^j \times P_j}</pre>
    vector<int> rnd1, rnd2:
    mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
    for(int i=0: i<n: i++){</pre>
        rnd1.push_back(randint(1, mod - 1));
        rnd2.push_back(randint(1, mod - 1));
    }
    vector<int> gobs;
    for(int i=0; i<2*n+2; i++){
        int tmp = 0:
        for(int j=0; j<n; j++){</pre>
            tmp += 111 * rnd2[j] * rnd1[j] % mod;
            if(tmp >= mod) tmp -= mod;
        }
        gobs.push_back(tmp);
        vector<int> nxt(n);
        for(auto &i : M){ // sparse matrix * vector
            nxt[i.x] += 111 * i.v * rnd1[i.y] % mod;
            if(nxt[i.x] >= mod) nxt[i.x] -= mod;
        }
        rnd1 = nxt:
    }
    auto sol = berlekamp_massey(gobs);
    reverse(sol.begin(), sol.end());
    return sol;
}
lint det(int n, vector<elem> M){
    vector<int> rnd:
    mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
    };
    for(int i=0; i<n; i++) rnd.push_back(randint(1, mod - 1));</pre>
    for(auto &i : M){
        i.v = 111 * i.v * rnd[i.v] % mod;
    auto sol = get_min_poly(n, M)[0];
    if(n \% 2 == 0) sol = mod - sol;
    for(auto &i : rnd) sol = 111 * sol * ipow(i, mod - 2) % mod:
    return sol;
}
```

2 Geometry

2.1 struct Point

```
const double eps = 1e-10:
template <class T>
struct point{
   typedef point P;
   T x, v;
   point(T x=0, T y=0) : x(x), y(y) {}
   bool operator< (P a) const {return fabs(x-a.x) > eps ? x<a.x : y<a.y;}</pre>
   bool operator == (P a) const {return max(fabs(x-a.x), fabs(y-a.y)) < eps;}
   P operator+ (P a) const {return P(x+a.x, v+a.v);}
   P operator- (P a) const {return P(x-a.x, y-a.y);}
   P operator- () const {return P(-x, -y);};
   T operator* (P a) const {return x*a.x + y*a.y;} // inner prod
   T operator/ (P a) const {return x*a.y - y*a.x;} // outer prod
   T dist2() const {return x*x + y*y;}
   double dist() const {return sqrt(double(dist2()));}
   P perp() const {return P(-y, x);}; // rotate 90 deg ccw
   P mult(T t) const {return P(x*t, y*t);}
   P unit() const {return P(x/dist(), v/dist());}
   P rotate(double a){
       return P(x*cos(a)-y*sin(a), x*sin(a)+y*cos(a));
   }
};
int sgn(double x) {return (x > eps) - (x < -eps);}
typedef point<double> P;
2.2 Distance, Intersection
2.2.1 Point-to-Line
double lineDist(P a, P b, P p) {
   return ((b-a)/(p-a))/(b-a).dist(); // a->b : left (+), right : (-);
2.2.2 Point-to-Segment
double segDist(P s, P e, P p) {
   if(s == e) return (p-s).dist(): // mind the eps
   double d = (e-s).dist2(), t = min(d, max(.0, (p-s)*(e-s)));
   return ((p-s).mult(d)-(e-s).mult(t)).dist() / d;
2.2.3 Line intersection
template<class P>
pair<int, P> lineInter(P a, P b, P c, P d){
   if((b-a)/(d-c) == 0) // parallel, mind the eps
```

return $\{-((b-a)/(c-a) == 0), a\}$;

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```
double oa = (d-c)/(a-c), ob = (d-c)/(b-c);
    return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
} // 1.0.-1(inf) : inter
2.2.4 Segment Intersection
bool onSegment(P s, P e, P p) {return segDist(s, e, p) < eps;}</pre>
template<class P> vector<P> segInter(P a, P b, P c, P d){
    double oa = (d-c)/(a-c), ob = (d-c)/(b-c),
            oc = (b-a)/(c-a), od = (b-a)/(d-a);
    if(sgn(oa)*sgn(ob) < 0 && sgn(oc)*sgn(od) < 0)
        return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
    if(onSegment(c, d, a)) S.insert(a);
    if(onSegment(c, d, b)) S.insert(b);
    if(onSegment(a, b, c)) S.insert(c);
    if(onSegment(a, b, d)) S.insert(d);
    return vector<P>(S.begin(), S.end()):
}
```

2.2.5 Circle-Line Intersection

Should be added.

2.3 Convex Hull

```
vector<pll> get_CV(vector<pll> V){
    sort(V.begin(), V.end());
    sort(V.begin() + 1, V.end(), [&](pll x, pll y){
        pll xx = x - V[0];
        pll yy = y - V[0];
        11 \text{ res} = xx / vv:
        if(res != 0) return res > 0;
        if(xx.first != yy.first) return xx.first < yy.first;</pre>
        return xx.second < yy.second;</pre>
    });
    vector<pll> ret;
    for(auto val : V){
        while(ret.size() > 1){
            pll xx = ret[ret.size() - 2] - val;
            pll yy = ret[ret.size() - 1] - val;
            if(xx / yy <= 0) ret.pop_back();</pre>
            else break:
        }
        ret.push_back(val);
    }
```

```
return ret;
2.4 Rotating Calipers
void rotating_calipers(vector<pll> CV){
   int pos = 0;
   for(int i = 0; i < CV.size(); i++) if(CV[pos] < CV[i]) pos = i;
   int ind1 = 0, ind2 = pos:
   11 dist = (CV[ind1] - CV[ind2]) * (CV[ind1] - CV[ind2]);
   auto get_v = [\&](int x) { return CV[(x + 1) \% CV.size()] - <math>CV[x];};
   for(int i = 0 ; i < CV.size() ; i++){</pre>
        pll v = get_v(i);
        while((-v) / get_v(pos) < 0) pos = (pos + 1) % CV.size();
        11 tmp_dist = (CV[pos] - CV[i]) * (CV[pos] - CV[i]);
        if(dist < tmp_dist) {</pre>
           dist = tmp_dist;
           ind1 = i; ind2 = pos;
   }
    printf("%lld %lld %lld %lld\n", CV[ind1].first, CV[ind1].second, CV[ind2].first,
   CV[ind2].second);
2.5 Sorting Points by Angle
// credit : http://koosaga.com/97
auto angle_sort = [&] (P &a, P &b){
   if((a < point(0, 0)) ^ (b < point(0, 0))) return b < a;</pre>
   if (a / b != 0) return a / b > 0;
   return a.dist2() < b.dist2(); // norm</pre>
}; // clockwise sort
2.6 Smallest Enclosing Circle
//Prove By Solving - https://www.acmicpc.net/problem/11930
int main(){
   scanf("%d", &N):
   for(int i = 1; i \le N; i++) scanf("%lf%lf", &A[i].x, &A[i].y, &A[i].z);
   int t = 70000:
   double rate = 1.0;
   point cur = (point)\{0, 0, 0\};
   for(int i = 1 : i \le t : i++){}
        int ind = 1;
       for(int j = 1; j \le N; j++) if((A[j] - cur) * (A[j] - cur) > (A[ind] -
        cur) * (A[ind] - cur)) ind = j;
        cur = cur + (A[ind] - cur) * rate;
```

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```
rate *= 0.99;
}
double r = 0;
for(int i = 1; i <= N; i++) r = max(r, (A[i] - cur) * (A[i] - cur));
cout << sqrt(r);
return 0;
} // Non-deterministic, deterministic O(n lg n) requires Voronoi diagram</pre>
```

2.7 Polygon Area

2.7.1 Polygon Area

```
double ans = 0; // ans : double area
for(int i=0;i<points.size();i++)
   ans += points[i] / points[(i+1 == points.size() ? 0 : i+1)];</pre>
```

3 Strings

3.1 Aho-Corasick Algorithm

```
namespace aho_corasick {
    const int MAXN = 100000, MAXC = 26;
    int trans[MAXN+1][MAXC];
    int fail[MAXN+1];
   bool term[MAXN+1];
    void build(const vector<string> &v) {
       setz(trans), setz(fail), setz(term);
        int cnode = 1;
       repi(s, v) {
           int p = 0;
           repi(j, s) {
               char c = i-'a':
               if (!trans[p][c]) trans[p][c] = cnode++;
                p = trans[p][c];
           }
            term[p] = true;
       }
        queue<int> q; rep(i, MAXC) if (trans[0][i]) q.push(trans[0][i]);
        while(!q.empty()) {
           int t = q.front(); q.pop();
           rep(i, MAXC) {
               if (trans[t][i]) {
                    int p = fail[t];
                    while(p and not trans[p][i]) p = fail[p];
                    p = trans[p][i];
                    fail[trans[t][i]] = p;
                    if (term[p]) term[trans[t][i]] = true;
                    q.push(trans[t][i]);
```

```
}
    bool query(string &t) {
        int p = 0;
        repi(i, t) {
            char c = i-'a';
            while(p and not trans[p][c]) p = fail[p];
            p = trans[p][c];
            if (term[p]) return true;
        }
        return false;
   }
}
3.2 Suffix Array
// str : abracadabra
// SA : 10 7 0 3 5 8 1 4 6 9 2
// LCP : 1 4 1 1 0 3 0 0 0 2
vector<int> make_sa(const string& s) {
    int n = s.length();
   int \lim = \max(128, n+1):
    vector\langle int \rangle sa(n), g(n+1), ng(n+1), cnt(lim), ind(lim+1);
   for(int i=0:i<n:i++) {
        sa[i] = i; g[i] = s[i];
   }
    g[n] = 0;
   for(int t=1;t<s.length();t<<=1)</pre>
        auto cmp = [&] (int a, int b) {
            return g[a] != g[b] ? g[a] < g[b] : g[a+t] < g[b+t];
        }:
        for(int i=0;i<n;i++) cnt[g[min(i+t, n)]]++;</pre>
        for(int i=1;i<lim;i++) cnt[i] += cnt[i-1];</pre>
        for(int i=n-1; i>=0; i--) ind[--cnt[g[min(i+t, n)]]] = i;
        for(int i=0;i<lim;i++) cnt[i] = 0;</pre>
        for(int i=0; i \le n; i++) cnt[g[i]]++; // same as cnt[g[ind[i]]]++
        for(int i=1;i<lim;i++) cnt[i] += cnt[i-1];</pre>
        for(int i=n-1;i>=0;i--) sa[--cnt[g[ind[i]]]] = ind[i];
        ng[sa[0]] = 1;
        for(int i=1;i<n;i++) {</pre>
            ng[sa[i]] = ng[sa[i-1]] + cmp(sa[i-1], sa[i]);
        g = ng;
        fill(cnt.begin(), cnt.end(), 0);
        fill(ind.begin(), ind.end(), 0);
```

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```
}
    return sa;
}
vector<int> make lcp(const string& s. const vector<int>& sa) {
    int n = s.length();
    vector<int> lcp(n-1), rank(n);
    for(int i=0;i<n;i++)</pre>
       rank[sa[i]] = i;
    int len = 0:
   for(int i=0;i<n;i++) {</pre>
       if(rank[i]) {
           int j = sa[rank[i]-1];
           int lc = n - max(i,j);
           while(len < lc && s[i+len] == s[j+len]) len++;</pre>
           lcp[rank[i]-1] = len;
       if(len) len--;
    }
    return lcp;
     Manacher's Algorithm
// 0-based
//s = #h#e#l#l#o#
// ret = 0 1 0 1 0 1 2 1 0 1 0
vector<int> manacher(const string& s) {
    int n = s.size(), r = -1, k = -1;
   vector<int> p(n);
   for (int i=0; i<n; i++) {
       if (i<=r) p[i] = min(r-i, p[2*k-i]);</pre>
       while (i-p[i]-1>=0 and i+p[i]+1<n and s[i-p[i]-1] == s[i+p[i]+1]) p[i]++;
       if (r < i+p[i]) r = i+p[i], k = i;
   }
    return p;
}
3.4 Z Algorithm
// O-based
//s = abcababca
// ret = 9 0 0 2 0 4 0 0 1
vector<int> z_algo(const string &s) {
    int 1 = 0, r = 0, N = sz(s):
    vector<int> Z(N);
   Z[0] = N;
                                                                                      // memset(finished, -1, sizeof(finished));
    repp(i, 1, N) {
                                                                                      int scc(int here){
       if (i > r) {
```

```
while(r < N \text{ and } s[r] == s[r-1]) r++;
           Z[i] = r-1+1;
       } else {
            int k = i-1;
           if (Z[k] < r-i+1) Z[i] = Z[k];
           else {
               1 = i;
               while(r < N \text{ and } s[r] == s[r-1]) r++;
               Z[i] = r-l+1;
           }
       }
   }
   return Z;
3.5 Lexicographically Smallest String Rotation
// rotate(v.begin(), v.begin()+min_rotation(v), v.end());
int min_rotation(string s) {
   int a=0, N=sz(s); s += s;
   rep(b,0,N) rep(i,0,N) {
       if (a+i == b \mid | s[a+i] < s[b+i]) \{b += max(0, i-1); break;\}
       if (s[a+i] > s[b+i]) \{ a = b : break : \}
   }
   return a;
4 Graph Theory
4.1 Strongly Connected Component
const int MAXN = 2e5 + 10; // > 2*N
int N. M:
int dfsn[MAXN], low[MAXN], finished[MAXN], cnt;
vector<int> ADJ[MAXN];
vector<vector<int>> G:
stack<int> S;
int f(int x){ // 0 1 2 3 4 5... -> f(1) f(-1) f(2) f(-2) f(3) f(-3)...
   return 2 * (abs(x) - 1) + (x < 0):
void add_edge(int x, int y){ // call by f(x), f(y)
   ADJ[x ^ 1].push_back(y);
   ADJ[y ^ 1].push_back(x);
```

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```
static vector<int> tmp;
S.push(here);
dfsn[here] = low[here] = ++cnt:
int &ret = low[here];
for(int there : ADJ[here]){
   if(dfsn[there] == 0) ret = min(ret, scc(there));
    else if(finished[there] == -1) ret = min(ret, dfsn[there]);
}
if(dfsn[here] == low[here]){
    while(1){
       int x = S.top(); S.pop();
       finished[x] = G.size();
       tmp.push_back(x);
       if(x == here) break;
   G.push_back(tmp);
    tmp.clear();
}
return ret;
```

4.1.1 2-SAT

}

- scc를 실행시켜 f(i) 와 f(-i)가 같은 component에 있다면, 모순.
- f(i) 와 f(-i) 중 finished 배열의 수가 작은 것이 참이다.
 - SCC numbering의 역순이 위상정렬이기에, $F \rightarrow T$ 를 유지하기 위함

4.2 Biconnected Component

```
// https://gist.github.com/koosaga/6f6fd50dd7067901f1b1
void dfs(int x, int p){
    dfn[x] = low[x] = ++piv;
    par[x] = p;
   for(int i=0; i<graph[x].size(); i++){</pre>
        int w = graph[x][i];
        if(w == p) continue;
        if(!dfn[w]){
            dfs(w. x):
            low[x] = min(low[x], low[w]);
        }
        else{
            low[x] = min(low[x], dfn[w]);
    }
}
void color(int x. int c){
    if(c > 0) bcc[x].push_back(c); // c == 0 : first component
```

```
vis[x] = 1;
   for(int i=0; i<graph[x].size(); i++){</pre>
       int w = graph[x][i];
       if(vis[w]) continue;
       if(dfn[x] <= low[w]){</pre>
           bcc[x].push_back(++cpiv);
           color(w, cpiv);
       }
       else{
           color(w, c);
   }
}
4.3 Euler Tour
struct Edge{
   int to, cnt; // to: 인접한 정점, cnt: 남은 사용 횟수
   Edge *dual; // dual: 역방향 간선을 가리키는 포인터
   Edge(): Edge(-1, 0){}
   Edge(int to1, int cnt1): to(to1), cnt(cnt1), dual(nullptr){}
void Eulerian(int curr){
   for(Edge *e: adj[curr]){
       if(e\rightarrow cnt > 0){
           e->cnt--:
           e->dual->cnt--;
           Eulerian(e->to); // dfs
   }
   cout << curr << '\n';
4.4 Heavy-Light Decomposition
int N. M:
vector<int> ADJ[MAXN];
int S[MAXN];
int hld_head[MAXN], color[MAXN], dfsn[MAXN], dcnt, hcnt;
int P[MAXN];
void dfs1(int here, int par){
 S[here] = 1; P[here] = par;
 for(int there : ADJ[here]) if(there != par) dfs1(there, here), S[here] +=
 S[there];
void dfs2(int here, int c){ // dfs reordering
 if(hld_head[c] == 0) hld_head[c] = here;
 dfsn[here] = ++dcnt: color[here] = c:
```

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```
sort(ADJ[here].begin(), ADJ[here].end(), [&](int x, int y){
   return S[x] > S[y];
});
int cnt = 0;
for(int there : ADJ[here]) if(there != P[here]){
   if(++cnt == 1) dfs2(there, c);
   else dfs2(there, ++hcnt);
}
```

4.5 Dominator Tree

Should be added.

4.6 Negative Cycle Detection

Should be added.

4.7 Tree Compress

Should be added.

4.8 Global Min Cut.

```
// Stoer-Wagner Algorithm, O(VE lg E)
int minimum_cut_phase(int n, int &s, int &t, vector<vector<int>> &adj, vector<int>
vis){
    vector<int> dist(n);
    int mincut = 1e9:
    while(true){
        int pos = -1, cur = -1e9;
       for(int i=0: i<n: i++){
            if(!vis[i] && dist[i] > cur){
                cur = dist[i];
                pos = i:
            }
        if(pos == -1) break:
        s = t;
        t = pos;
        mincut = cur;
       vis[pos] = 1;
       for(int i=0; i<n; i++){</pre>
            if(!vis[i]) dist[i] += adj[pos][i];
       }
   }
    return mincut; // optimal s-t cut here is, {t} and V \ {t}
int solve(int n, vector<vector<int>> adj){
    if(n <= 1) return 0;
```

```
vector<int> vis(n);
int ans = 1e9;
for(int i=0; i<n-1; i++){
    int s, t;
    ans = min(ans, minimum_cut_phase(n, s, t, adj, vis));
    vis[t] = 1;
    for(int j=0; j<n; j++){
        if(!vis[j]){
            adj[s][j] += adj[t][j];
            adj[j][s] += adj[j][t];
        }
    }
    adj[s][s] = 0;
}
return ans;
</pre>
```

5 Network Flow

5.1 Theorems

Max-flow Min-cut theorem : 정점 s에서 정점 t까지 흐를 수 있는 최대 유량(max-flow)은 정점 s와 정점 t를 분리하는 간선들의 가중치 합(min-cut)과 같다.

Vertex cover : 어떤 그래프의 정점의 집합 *S*에 대해 그래프의 모든 간선이 *S*의 원소 중 최소 하나와 연결되어 있을 때, *S*를 해당 그래프의 vertex cover라고 하며, minimum vertex cover는 최소 개수의 정점을 사용한 vertex cover이다.

Independent set : 어떤 그래프의 정점의 집합 S에 대해 S의 서로 다른 두 정점을 연결하는 간선이 없을 때, S를 해당 그래프의 independent set이라고 하며, maximum independent set은 최대 개수의 정점을 사용한 independent set이다.

Matching (independent edge set): 어떤 그래프의 간선의 집합 S에 대해 S의 서로 다른 두 간선이 공통된 정점을 가지지 않을 때, S를 해당 그래프의 matching이라고 하며, maximum matching은 최대 개수의 간선을 사용한 matching이다.

König's theorem : 이분 그래프의 maximum matching의 크기는 minimum vertex cover의 것과 같다. Dinic's Algorithm : 시간 복잡도 $O(V^2E)$, unit capacity에서는 $\min(V^{2/3}E, E^{3/2})$.

Circulation Problem : 새로운 source/sink s_n , t_n 를 만들어서 다음과 같이 간선을 추가하고 $maxflow(s_n \to t_n) = \sum l_i$ 인지 확인, 이후 $s \to t$ 로 maxflow

```
• s_n \to b (l), a \to t_n (l), a \to b (r-l), t \to s (\infty)
```

5.2 Dinic's Algorithm

```
const int INF = 1e9;
struct Dinic{
  int N;
  struct edge{
    int index, cap, rev;
    edge() : index(0), cap(0), rev(0) {}
    edge(int index, int cap, int rev) : index(index), cap(cap), rev(rev) {}
};

vector<vector<edge>> ADJ;
```

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```
vector<int> R, W;
Dinic() {}
Dinic(int N) : N(N){
   ADJ.resize(N): R.resize(N):
                                    W.resize(N):
void CE(int node1, int node2, int cap){
    ADJ[node1].push_back(edge(node2, cap, ADJ[node2].size()));
   ADJ[node2].push_back(edge(node1, 0, ADJ[node1].size() - 1));
}
bool bfs(int src, int sink){
   fill(R.begin(), R.end(), -1);
   R[src] = 0;
   queue<int> Q; Q.push(src);
   while(Q.size()){
       int here = Q.front(); Q.pop();
        for(auto e : ADJ[here]){
           if(e.cap > 0 \&\& R[e.index] == -1)
                R[e.index] = R[here] + 1, Q.push(e.index);
        }
   }
   return R[sink] != -1;
int dfs(int here, int sink, int f){
    if(here == sink) return f;
   for(int &i = W[here] : i < ADJ[here].size() : i++){</pre>
        auto &e = ADJ[here][i];
        if(e.cap > 0 && R[here] < R[e.index]){
           int res = dfs(e.index, sink, min(f, e.cap));
           if(res) {
                e.cap -= res:
                ADJ[e.index][e.rev].cap += res;
                return res;
           }
        }
   }
   return 0:
int solve(int src. int sink){
   int ret = 0;
    while(bfs(src, sink)){
       fill(W.begin(), W.end(), 0);
       int res:
        while((res = dfs(src, sink, INF))) ret += res;
   }
   return ret;
```

```
};
5.3 MCMF with SPFA
const int INF = 1e9;
struct MCMF {
   struct EDGE {
       int there, cap, cost, rev;
       EDGE() : there(0), cap(0), cost(0), rev(0) {}
       EDGE(int there, int cap, int cost, int rev) : there(there), cap(cap),
        cost(cost), rev(rev) {}
   };
   int N:
   vector<vector<EDGE>> ADJ:
   vector<int> R, INQ, C, I;
   MCMF() : N(0) {}
   MCMF(int N) : N(N) { ADJ.resize(N + 1); R.resize(N + 1); INQ.resize(N + 1);
   C.resize(N + 1); I.resize(N + 1); }
   void connect_edge(int i, int j, int cap, int cost) {
        ADJ[i].push_back(EDGE(j, cap, cost, ADJ[j].size()));
        ADJ[j].push_back(EDGE(i, 0, -cost, ADJ[i].size() - 1));
   }
   bool SPFA(int src, int sink) {
       queue<int> Q; Q.push(src);
       fill(R.begin(), R.end(), -1);
                                        R[src] = 0;
                                        C[src] = 0:
       fill(C.begin(), C.end(), -1);
       fill(INQ.begin(), INQ.end(), 0); INQ[src] = 1;
       while (0.size()) {
           int here = Q.front(); Q.pop();
           INQ[here] = 0;
           for (int i = 0: i < ADJ[here].size(): i++) {</pre>
               auto e = ADJ[here][i];
               if (e.cap > 0 && (C[e.there] == -1 || C[e.there] > C[here] +
               e.cost)) {
                   C[e.there] = C[here] + e.cost; R[e.there] = here;
                   I[e.there] = i:
                   if (!INQ[e.there]) INQ[e.there] = 1, Q.push(e.there);
           }
       if (C[sink] == -1) return false;
       return true:
   }
   pii mcmf(int src, int sink) {
       pii ret = { 0, 0 };
```

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```
while (SPFA(src, sink)) {
            int flow = INF, cost = 0;
            for (int here = sink; here != src; here = R[here]) flow = min(flow.
            ADJ[R[here]][I[here]].cap);
            for (int here = sink: here != src: here = R[here]) {
                auto &e = ADJ[R[here]][I[here]];
                cost += e.cost * flow;
                e.cap -= flow;
                ADJ[e.there][e.rev].cap += flow;
            }
            ret.first += flow, ret.second += cost;
        }
        return ret;
    }
};
5.4 Hungarian Method
namespace Hung {
    const int MX = 2000:
    // IMPORTANT : n <= m, 1-based
    using T = long double;
    T \max v = 1e200:
    T a[MX][MX], n, m;
    void init(int nn, int mm) { n = nn; m = mm; }
    void set_value(int x, int y, T val) { a[x][y] = val; }
    T solve(vector <int> &ans) {
        vectorT> v(m+1), u(n+1);
        vector<int> p (m+1), way (m+1);
        for (int i=1; i<=n; ++i) {
           p[0] = i;
            int i0 = 0:
            vector<T> minv (m+1, maxv);
            vector<char> used (m+1, false):
            do {
                used[i0] = true;
                T delta = maxv:
                int i0 = p[j0], j1;
                for (int j=1; j<=m; ++j) if (!used[j]) {</pre>
```

T cur = a[i0][j]-u[i0]-v[j];

minv[j] = cur, way[j] = j0;

delta = minv[j], j1 = j;

if (cur < minv[j]) {</pre>

if (minv[j] < delta) {</pre>

for (int j=0; j<=m; ++j) {

if (used[i]) {

}

}

}

```
u[p[i]] += delta, v[i] -= delta;
                    else {
                        minv[j] -= delta;
               }
                j0 = j1;
           } while (p[j0] != 0);
           do {
                int j1 = way[j0];
                p[j0] = p[j1];
                j0 = j1;
           } while (j0);
        ans.resize(n + 1);
       for(int j=1;j<=m;++j) {</pre>
           ans[p[j]] = j;
       return -v[0];
5.5 Hopcroft-Karp Algorithm
struct hopcroft_karp{
   int N:
   vector<vector<int>> ADJ;
   vector<int> L, rev, used;
   hopcroft_karp() {}
   hopcroft_karp(int N) : N(N) {
        ADJ.resize(N);
        L.resize(N), rev.resize(N, -1), used.resize(N, 0);
   }
   void CE(int here, int there){
        ADJ[here].push_back(there);
   }
   void bfs(){
        queue<int> Q:
       for(int i = 0 ; i < N ; i++) {</pre>
           if(used[i]) L[i] = -1;
            else L[i] = 0, Q.push(i);
       }
        while(0.size()){
           int here = Q.front(); Q.pop();
           for(int there : ADJ[here]){
                if(rev[there] != -1 && L[rev[there]] == -1) {
```

L[rev[there]] = L[here] + 1;

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```
Q.push(rev[there]);
                }
            }
        }
    }
    bool dfs(int here){
        for(int there : ADJ[here]){
            if(rev[there] == -1 || (L[here] < L[rev[there]] && dfs(rev[there]))){
                rev[there] = here;
                used[here] = 1;
                return true;
            }
        }
        return false;
    }
    int solve(){
        int ret = 0;
        while(1){
            bfs():
            int res = 0;
            for(int i = 0 ; i < N ; i++) {</pre>
                if(used[i]) continue;
                res += dfs(i);
            if(res == 0) break:
            ret += res;
        }
        return ret;
};
```

6 Optimization Tricks

6.1 Knuth Optimization

```
• Recurrence : D[i][j] = \min_{i < k < j} (D[i][k] + D[k][j]) + C[i][j]
```

- Quadrangle Inequality : $C[a][c] + C[b][d] \le C[a][d] + C[b][c], \ a \le b \le c \le d$
- $\bullet \mbox{ Monotonicity}: C[b][c] \leq C[a][d], \ a \leq b \leq c \leq d$
- • A[i][j] = (min. k s.t. D[i][j] is min.). Then $A[i][j-1] \leq A[i][j] \leq A[i+1][j]$

• $O(N^2)$ time complexity

```
// opt[i-1][i] = i
for(int d=2;d<=n;d++) {
   for(int i=1;i+d<=n+1;i++) {
     for(int k=opt[i][j-1], j=i+d; k<=opt[i+1][j]; k++) {
        int v = dp[i][k] + dp[k][j] + c[i][j];</pre>
```

```
if(dp[i][j] > v) dp[i][j] = v, opt[i][j] = k;
}
}
```

6.2 Divide and Conquer Optimization

- Recurrence : $D[t][i] = \min_{k < i} (D[t-1][k] + C[k][i])$
- Min index : $A[t][i] \le A[t][i+1]$ ($A[t][i] = (\min. k \ s.t. \ D[t][i] \ is min.))$
- Quadrangle Inequality : $C[a][c] + C[b][d] \le C[a][d] + C[b][c], \ a \le b \le c \le d$
- Able to Divide and Conquer base on calculating D[t][i]
- $O(TN \lg N)$ time complexity

```
// range of index : [1,r]
// range of dp : [s,e]
void dnc(int t, int 1, int r, int s, int e)
{
    if(s > e) return;
    int m = (s+e)/2;
    D[t][m] = 2e9;
    for(int k=1;k<m&&<=r;k++)
    {
        int tmp = D[t-1][k] + C[k][m];
        if(D[t][m] > tmp)
            D[t][m] = tmp, A[t][m] = k;
    }
    dnc(t, 1, A[t][m], s, D[t][m]-1);
    dnc(t, A[t][m], r, D[t][m]+1, e);
}
```

6.3 Convex Hull Trick

- Recurrence : $dp[i] = \min_{j < i} (dp[j] + a[i]b[j]), \ b[i-1] \le b[i]$
- Think as $dp[x = a[i]] = \min_{j < i} (b[j] \cdot x + dp[j])$
- Thus push lines and find minimum (by binary search)
- If $a[i] \le a[i+1]$ sweeping is possible
- Intersection of $y = a_i x + b_i$ and $y = a_{i+1} x + b_{i+1} : x = \frac{b_{i+1} b_i}{a_i a_{i+1}}$

6.4 Centroid Decomposition

```
// credit : https://gist.github.com/igorcarpanese/75162f3253bd230abd0f32f9950bf384
int dfs(int u, int p) {
    sub[u] = 1;
    for (auto v : tree[u])
```

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```
if (v != p) sub[u] += dfs(v, u);
  return sub[u] + 1;
}

// each tree has at most two centroids
int centroid(int u, int p, int r) { // r : root
  for (auto v : tree[u])
      if (v != p and sub[v] > sub[r]/2) return centroid(v, u);
  return u;
}
```

7 Data Structure

7.1 Persistent Segment Tree

```
const MAXN = 1e5 + 10:
struct node{
    node *1, *r;
    int cnt:
    node () {}
} pool[(1 << 17) * 17], *tree_head[MAXN];</pre>
int tcnt:
node* alloc(){
    memset(pool + tcnt, 0, sizeof(node));
    return pool + tcnt++;
}
node * init(int 1, int r){
    node *ret = alloc();
    if(1 != r) {
        int mid = (1 + r) / 2;
        ret->l = init(l, mid):
        ret->r = init(mid + 1, r);
    }
    return ret;
}
void update(node * here, node *par, int 1, int r, int val){
    if(1 == r) {
        here->cnt = par->cnt + 1;
        return:
    }
    int mid = (1 + r) / 2;
    if(val <= mid){</pre>
        here->l = alloc();
        here->r = par->r;
        update(here->1, par->1, 1, mid, val);
    }
```

```
else {
        here->1 = par->1;
        here->r = alloc():
        update(here->r, par->r, mid + 1, r, val);
   here->cnt = here->l->cnt + here->r->cnt;
int query(node *node1, node *node2, int 1, int r, int k){
   if(1 == r) return 1:
   int ccc = node1->l->cnt - node2->l->cnt;
   int mid = (1 + r) / 2;
   if(k <= ccc) return query(node1->1, node2->1, 1, mid, k);
    else return query(node1->r, node2->r, mid + 1, r, k - ccc);
7.2 Link-Cut Tree
struct node{
   node *pp, *p, *l, *r;
   int val;
   node() \{ p = 0, 1 = 0, r = 0; \}
   node(int val) : val(val) \{ p = 0, l = 0, r = 0; \}
};
void push(node *x){}
void pull(node *x){}
void rotate(node *x){
    if(!x->p) return;
    push(x->p); // if there's lazy stuff
   push(x);
   node *p = x->p;
   bool is left = (p->1 == x):
   node *b = (is_left ? x->r : x->1);
   x->p = p->p;
   if (x-p \&\& x-p-1 == p) x-p-1 = x;
    if (x-p \&\& x-p-r == p) x-p-r = x;
   if(is left){
        if(b) b \rightarrow p = p;
        p->1 = b;
        p->p = x;
        x->r = p;
   }
    else{
        if(b) b \rightarrow p = p;
        p->r = b;
        p->p = x;
        x->1 = p;
    pull(p); // if there's something to pull up
```

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```
pull(x);
    //if(!x->p) root = x; // IF YOU ARE SPLAY TREE
    if(p->pp){ // IF YOU ARE LINK CUT TREE
        x->pp = p->pp;
        p->pp = nullptr;
    }
}
void splay(node *x){
    while(x->p){
        node *p = x->p;
        node *g = p->p;
        if(g){
            if((p->1 == x) ^ (g->1 == p)) rotate(x);
            else rotate(p);
        }
        rotate(x);
    }
}
void access(node *x){
    splay(x);
    push(x);
    if(x->r){
       x->r->pp = x;
        x->r->p = nullptr;
        x->r = nullptr;
    }
    pull(x);
    while(x->pp){
       node *nxt = x->pp;
        splay(nxt);
        push(nxt);
        if(nxt->r){
            nxt->r->pp = nxt;
           nxt->r->p = nullptr;
            nxt->r = nullptr;
        }
        nxt->r = x;
        x->p = nxt;
        x->pp = nullptr;
        pull(nxt);
        splay(x);
    }
}
node *root(node *x){
    access(x):
    while(x->1){
        push(x);
        x = x -> 1:
    }
    access(x);
```

```
return x;
node *par(node *x){
    access(x);
   if(!x->1) return nullptr;
   push(x);
   x = x->1;
    while(x->r){
        push(x);
        x = x->r;
   }
   access(x);
    return x;
node *lca(node *s, node *t){
   access(s):
   access(t);
   splay(s);
   if(s->pp == nullptr) return s;
   return s->pp;
void link(node *par, node *son){
   access(par);
   access(son);
   //son->rev ^= 1; // remove if needed
   push(son);
   son->1 = par;
   par->p = son;
   pull(son);
void cut(node *p){
   access(p);
    push(p);
   if(p->1){
        p->1->p = nullptr;
        p->1 = nullptr;
   }
    pull(p);
7.3 Dynamic Convex Hull
// https://github.com/niklasb/contest-algos/blob/master/convex_hull/dynamic.cpp
const ll is_query = -(1LL<<62);</pre>
struct Line {
   ll m, b;
    mutable function<const Line*()> succ:
   bool operator<(const Line& rhs) const {</pre>
        if (rhs.b != is_query) return m < rhs.m;</pre>
        const Line* s = succ();
        if (!s) return 0;
```

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```
// if result is a/b, return as {a, b}
       11 x = rhs.m;
       return b - s -> b < (s -> m - m) * x;
                                                                                          11 a = 0, b = 1: // 1
   }
                                                                                          11 c = 1, d = 0; // r
};
                                                                                           int 1, r;
struct HullDvnamic : public multiset<Line> { // will maintain upper hull for maximum
                                                                                          bool chg = true:
    bool bad(iterator y) {
       auto z = next(y);
                                                                                          while(chg) {
       if (v == begin()) {
                                                                                               chg = false;
           if (z == end()) return 0;
            return y->m == z->m && y->b <= z->b;
                                                                                               // to left
       }
                                                                                               1 = 0, r = (N-d-1)/b+1;
       auto x = prev(y);
                                                                                               while(1 < r)  {
       if (z == end()) return y->m == x->m && y->b <= x->b;
                                                                                                   int mid = (1+r+1)/2;
       return (x-b - y-b)*(z-m - y-m) >= (y-b - z-b)*(y-m - x-m);
                                                                                                  if (test(a*mid+c, b*mid+d)) r = mid-1;
    }
                                                                                                   else 1 = mid;
    void insert line(ll m, ll b) {
       auto y = insert({ m, b });
       y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
                                                                                               c += a*1:
       if (bad(y)) { erase(y); return; }
                                                                                               d += b*1:
       while (next(y) != end() && bad(next(y))) erase(next(y));
                                                                                               chg |= (1 > 0);
        while (y != begin() && bad(prev(y))) erase(prev(y));
    }
                                                                                               // to right
    11 eval(ll x) {
                                                                                               1 = 0, r = (d?(N-b-1)/d+1:M);
       auto 1 = *lower_bound((Line) { x, is_query });
                                                                                               while(1 < r)  {
       return 1.m * x + 1.b;
                                                                                                   int mid = (1+r+1)/2;
    }
                                                                                                  if (test(a+mid*c, b+mid*d)) l = mid;
}:
                                                                                                   else r = mid-1:
                                                                                               }
7.4 Stern-Brocot Tree
                                                                                               a += c*1:
// __int128 is recommended
                                                                                               b += d*1;
bool test(11 a, 11 b) { // for testing directions, vary by prob
                                                                                               chg |= (1 > 0);
    // return true if (true value) >= a/b
                                                                                          }
   11 n = 0, m = 1:
                                                                                          return {a, b};
   rep(i, N) {
       if (n < m*A[i].fi) n = A[i].fi, m = 1;</pre>
                                                                                       7.5 Rope
       11 c = b*n+m*a. d = m*b:
       11 g = gcd(c, d);
                                                                                       #include <bits/stdc++.h>
       n = c/g;
                                                                                       #include <ext/rope>
                                                                                       using namespace std;
       m = d/g;
                                                                                       using namespace __gnu_cxx;
       if (n > m*A[i].se) return false;
   }
                                                                                       int main()
    return true;
}
                                                                                          ios::svnc with stdio(false):
                                                                                           cin.tie(0);
pair<11, 11> stern_brocot(11 M, 11 N) {
   // M : max value
                                                                                           crope rp; // rope<char>
    // N : max divisor
                                                                                           string s("Lorem-ipsum");
```

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```
int n = s.length();
rp.append(s.c_str()); // add element

int x = 3, y = 8; // split and merge below
rp = rp.substr(x, y-x) + rp.substr(0, x) + rp.substr(y, n);
cout << rp.at(0) << '\n'; // get element, 'e'
cout << rp << '\n'; // print, "em-ip|Lor|sum"
}</pre>
```

7.6 Policy Based Data Structure

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
using namespace std;
using namespace __gnu_pbds;
typedef tree<
int.
null_type,
less<int>.
rb_tree_tag,
tree_order_statistics_node_update >
ordered set:
// less<int> : not allow for duplicate
// less_equal<int> : allow for duplicate
// use upper_bound when you erase from set used less_equal
int N;
int main(void) {
    iostream::sync_with_stdio(false);
    cin.tie(nullptr);
    ordered_set X;
    X.insert(1);
    X.insert(2):
    X.insert(4):
    X.insert(8):
    X.insert(16);
    cout<<*X.find_by_order(1)<<endl; // 2</pre>
    cout<<*X.find by order(2)<<endl: // 4
    cout<<*X.find_by_order(4)<<endl; // 16</pre>
    cout<<(end(X)==X.find_by_order(6))<<endl; // true</pre>
    cout<<X.order_of_key(-5)<<endl; // 0</pre>
```

```
cout<<X.order_of_key(1)<<endl;  // 0
cout<<X.order_of_key(3)<<endl;  // 2
cout<<X.order_of_key(4)<<endl;  // 2
cout<<X.order_of_key(400)<<endl;  // 5</pre>
```

8 Miscellaneous

8.1 Misc Formulae and Algorithms

8.1.1 Faulhaber's Formula

$$T(n,k) = \sum_{i=1}^{n} i^{k} = \frac{(n+1)^{k+1} - 1^{k+1} - \sum_{j=0}^{k-1} {k+1 \choose j} T(n,j)}{{k+1 \choose k}}$$

Also use

$$(x+1)^d - x^d = 1 + {d \choose 1}x + {d \choose 2}x^2 + \dots + {d \choose d-1}x^{d-1}$$

to get each coef.

8.1.2 Maximum Clique

```
typedef long long 11;
11 G[40]; // O-index
int N. M:
int cur;
void get_clique(int R = 0, 11 P = (111 << N)-1, 11 X = 0){
   if((P|X) == 0){
        cur = max(cur, R);
        return;
   }
   int u = __builtin_ctzll(P|X);
   11 c = P\&^{G}[u]:
   while(c){
        int v = builtin ctzll(c):
        get_clique(R + 1, P&G[v], X&G[v]);
       P ^= 111 << v;
       X = 111 << v:
        c ^= 111 << v;
   }
```

8.1.3 De Brujin Sequence

```
// https://github.com/koosaga/DeobureoMinkyuParty/blob/master/teamnote.tex
// alphabet = [0, k - 1], substr length n, res starts with 0 (cyclic)
int res[10000000], aux[10000000]; // >= k^n, k*n
int de_bruijn(int k, int n, int lim) { // returns size (k^n)
    if(k == 1) {
        res[0] = 0;
```

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```
return 1;
}
for(int i = 0; i < k * n; i++) aux[i] = 0;
int sz = 0;
function<void(int, int)> db = [&](int t, int p) {
   if(sz > lim) return;
   if(t > n) {
        if(n \% p == 0)
           for(int i = 1; i <= p; i++)
                res[sz++] = aux[i]:
   }
    else {
        aux[t] = aux[t - p];
        db(t + 1, p);
        for(int i = aux[t - p] + 1; i < k; i++) {
            aux[t] = i:
            db(t + 1, t);
        }
   }
};
db(1, 1);
return sz;
```

8.2 Highly Composite Numbers, Large Prime

}

< 10	^k number	divisors	2	3	5	7	11	13	17	19:	23:	29:	31	37	
1	6	4	1	1											
2	60	12	2	1	1										
3	840	32	3	1	1	1									
4	7560	64	3	3	1	1									
5	83160	128	3	3	1	1	1								
6	720720	240	4	2	1	1	1	1							
7	8648640	448	6	3	1	1	1	1							
8	73513440	768	5	3	1	1	1	1	1						
9	735134400	1344	6	3	2	1	1	1	1						
10	6983776800	2304	5	3	2	1	1	1	1	1					
11	97772875200	4032	6	3	2	2	1	1	1	1					
12	963761198400	6720	6	4	2	1	1	1	1	1	1				
13	9316358251200	10752	6	3	2	1	1	1	1	1	1	1			
14	97821761637600	17280	5	4	2	2	1	1	1	1	1	1			
15	866421317361600	26880	6	4	2	1	1	1	1	1	1	1	1		
16	8086598962041600	41472	8	3	2	2	1	1	1	1	1	1	1		
17	74801040398884800	64512	6	3	2	2	1	1	1	1	1	1	1	1	
18	897612484786617600	103680	8	4	2	2	1	1	1	1	1	1	1	1	
	< 10^k prime	> 10^k prime					# of prime								
1	7	:	 11								4			-	
2	97	10	01								25				

```
1009
                                                          168
   4
                   9973
                                      10007
                                                          1229
                   99991
                                      100003
                                                         9592
    6
                 999983
                                    1000003
                                                        78498
                9999991
                                    10000019
                                                       664579
   8
                                                      5761455
               99999989
                                  100000007
   9
                                                     50847534
              99999937
                                 1000000007
                     < 10 k prime
                                                 > 10<sup>k</sup> prime
    10
                                                  10000000019
                       999999967
   11
                      9999999977
                                                 10000000003
   12
                     99999999989
                                                1000000000039
    13
                   999999999971
                                               1000000000037
   14
                   9999999999973
                                              100000000000031
   15
                 9999999999999
                                             100000000000037
    16
                                            100000000000000061
                99999999999937
   17
               999999999999997
                                           100000000000000003
    18
              999999999999999
                                          1000000000000000003
NTT Prime:
 469762049 = 7 \times 2^{26} + 1. Primitive root : 3.
 998244353 = 119 \times 2^{23} + 1. Primitive root: 3.
 985661441 = 235 \times 2^{22} + 1. Primitive root: 3.
 1012924417 = 483 \times 2^{21} + 1. Primitive root: 5.
Primes near 10^9: 10^9 + [7, 9, 21, 33, 87]
8.3 Fast Integer IO
// credit : https://github.com/koosaga/DeobureoMinkyuParty/blob/master/teamnote.tex
static char buf[1 << 19]; // size : any number geq than 1024
static int idx = 0:
static int bytes = 0;
static inline int _read() {
   if (!bytes || idx == bytes) {
        bytes = (int)fread(buf, sizeof(buf[0]), sizeof(buf), stdin);
        idx = 0:
   }
   return buf[idx++]:
static inline int _readInt() {
   int x = 0, s = 1;
   int c = _read();
   while (c \leq 32) c = _read();
   if (c == '-') s = -1, c = read():
   while (c > 32) x = 10 * x + (c - '0'), c = _read();
   if (s < 0) x = -x;
   return x;
```

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8.4 C++ Tips / Environments

```
#include <bits/stdc++.h> // magic header
using namespace std; // magic namespace
struct StupidGCCCantEvenCompileThisSimpleCode{
    pair<int, int> array[1000000];
}; // https://gcc.gnu.org/bugzilla/show_bug.cgi?id=68203
// how to use rand (in 2017)
mt19937 rng(0xdeadbeef);
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
int randint(int lb, int ub){ return uniform_int_distribution<int>(lb, ub)(rng); }
shuffle(permutation.begin(), permutation.end(), rng);
mt19937_64 _R(chrono::steady_clock::now().time_since_epoch().count()); // _R()
// comparator overload
auto cmp = [](seg a, seg b){return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
int x, y;
bool operator == (const point &p)const{ return x == p.x && y == p.y; }
};
struct hasher {
size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
};
unordered_map<point, int, hasher> hsh;
// c++ setprecision example
#include <iostream>
                        // std::cout, std::fixed
#include <iomanip>
                        // std::setprecision
int main () {
    double f = 3.14159:
    std::cout << std::setprecision(5) << f << '\n'; // 3.1416
    std::cout << std::setprecision(9) << f << '\n'; // 3.14159
    std::cout << std::fixed;</pre>
    std::cout << std::setprecision(5) << f << '\n'; // 3.14159
    std::cout << std::setprecision(9) << f << '\n'; // 3.141590000
    return 0;
}
```