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Team Note of Powered by Zigui

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ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

0 Quotes and Prerequisites

```
evenharder : Mental Abuse To Humans
djkim0613 : 열심히 응원하겠습니다.
SoulTch : How much is this bus ticket?
* This template is brought from that of 'Deobureo Minkyu Party'
```

Run script

```
#!/bin/bash
g++ -fsanitize=undefined -std=c++14 -02 -o /tmp/pow $1.cpp
time /tmp/pow < $1.in
# export PATH=~:$PATH</pre>
```

Debug Code

```
#define setz(x) memset(x, 0, sizeof(x))
#define sz(x) ((int)(x).size())
#define rep(i, e) for (int i = 0, _##i = (e); i < _##i; i++)
#define repp(i, s, e) for (int i = (s), _##i = (e); i < _##i; i++)
#define repr(i, s, e) for (int i = (s)-1, _##i = (e); i \ge _{\#}i; i--)
#define repi(i, x) for (auto &i : (x))
// using namespace std;
using ll = long long;
using pii = pair<int, int>;
using pll = pair<11, 11>;
template<typename T>
ostream &operator<<(ostream &os, const vector<T>& v) {
    cout << "[":
   for (auto p : v) cout << p << ",";
    cout << "]":
    return os;
}
#ifndef __SOULTCH
#define debug(...) 0
#define endl '\n'
#define debug(...) cout << " [-] ", _dbg(#__VA_ARGS__, __VA_ARGS__)</pre>
template<class TH> void _dbg(const char *sdbg, TH h){ cout << sdbg << '=' << h <<
endl; }
template<class TH, class... TA> void _dbg(const char *sdbg, TH h, TA... a) {
    while(*sdbg != ',') cout << *sdbg++;</pre>
    cout << '=' << (h) << ',';
    _dbg(sdbg+1, a...);
}
#endif
```

Reminders

Pre-submit	Wrong answer:		
예제 작성해보기 (최소, 최대) 메모리, overflow 분석하기 올바른 문제에 제출하기	코드 + debug output 출력 다중 테케 문제에서 초기화 확인하기 알고리즘이 제한조건을 전부 다루는지 확인하기 지문 다시 읽어보기 corner case 찾아보기 초기화 안 된 지역변수 찾아보기 초기화 안 된 지역변수 찾아보기 N, M, i, j 등 변수 확인하기 풀이 증명하기 STL 함수 다시 생각해보기 팀노트에서 그대로 가져온 변수값 다시 확인하기 이 목록 다시 읽어보기 알고리즘 팀원에게 설명하기 팀원이랑 코드 보기 잠깐 일어나서 생각 재정비하고 오기		
Runtime error:	Time limit exceeded: / Memory limit exceeded:		
코너 케이스 처리해보기	무한 루프 확인하기		
초기화 안 된 변수 찾기	알고리즘 시간 복잡도 확인하기		
out-of-range 확인하기	data copy 어느 정도 하는지 확인하기 (reference)		
팀노트에서 그대로 가져온 변수값 다시 확인하기	입출력 규모 생각하기 (scanf 고려해보기)		
assertion 넣어보기	vector, map 최소화하기		
무한 재귀 확인하기	팀원에게 알고리즘 물어보기		
null pointer 확인하기	최대 메모리 사용량 계산하기		
메모리 사용량 확인하기	다중 테케 문제에서 초기화하기		

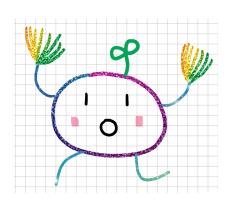




Figure 1: 풀다가 막힐 때는 이 그림을 봅시다. 아자아자 화이팅!

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1 Math

1.1 Basic Mathematics

1.1.1 Trigonometry

- $\sin(\alpha \pm \beta) = \sin \alpha \cos \beta \pm \cos \alpha \sin \beta$
- $\cos(\alpha \pm \beta) = \cos \alpha \cos \beta \mp \sin \alpha \sin \beta$
- $\tan(\alpha \pm \beta) = \frac{\tan \alpha \pm \tan \beta}{1 \mp \tan \alpha \tan \beta}$
- $\sin 2\theta = 2\sin\theta\cos\theta$
- $c^2 = a^2 + b^2 2ab\cos\gamma$

1.1.2 Generating Function

- $\sum_{n} (pn+q)x^{n} = \frac{p}{1-x} + \frac{q}{(1-x)^{2}}$ (Arithmetic progression)
- $\sum_{n} (rx)^n = (1 rx)^{-1}$ (Geometric progression)
- $\sum_{n} {m \choose n} x^n = (1+x)^m$ (Binomial coefficient)
- $\sum_{n} {m+n-1 \choose n} x^n = (1-x)^{-m}$ (Multiset coefficient)

1.1.3 Calculus

- $\int_a^b f(x) dx = \frac{b-a}{6} \left[f(a) + 4f(\frac{a+b}{2}) + f(b) \right]$ (Simpson's Rule, for cubic poly)
- $\int u'v \ dx = uv \int uv' \ dx$ (Integration by parts)

1.1.4 Combinatorics

- $\begin{bmatrix} n \\ k \end{bmatrix}$ (Stirling numbers of the first kind / n elem, k cycle)
- ${n+1 \brack k} = n {n \brack k} + {n \brack k-1}, {n \brack 0} = {0 \brack n} = {0 \brack 0} = 1$
- ${n \brace k}$ (Stirling numbers of the second kind / n elem, k unlabeled subset)
- ${n+1 \choose k} = k {n \choose k} + {n \choose k-1}, {n \choose 0} = {0 \choose n} = {0 \choose 0} = 1$

 B_n (Bell number, n elem, partition)

- $B_{n+1} = \sum_{k=0}^{n} {n \choose k} B_k, B_{n+1} = \sum_{k=0}^{n} {n \choose k}$
- $\bullet \ B_{p^m+n} \equiv mB_n + B_{n+1}$

1.2 Number Theory

1.2.1 Lattice Points under Line

```
divsum(to, c, k, m) = \sum_{i=0}^{\text{to}-1} \lfloor \frac{ki+c}{m} \rfloor, modsum(to, c, k, m) = \sum_{i=0}^{\text{to}-1} (ki+c)\%m
ull sumsq(ull to) { return to / 2 * ((to-1) | 1); }
ull divsum(ull to, ull c, ull k, ull m) {
    ull res = k / m * sumsq(to) + c / m * to;
    k %= m; c %= m;
    if (!k) return res;
    ull to2 = (to * k + c) / m:
    return res + (to - 1) * to2 - divsum(to2, m-1 - c, m, k);
11 modsum(ull to, 11 c, 11 k, 11 m) {
    c = ((c \% m) + m) \% m; k = ((k \% m) + m) \% m;
    return to * c + k * sumsq(to) - m * divsum(to, c, k, m):
1.2.2 Shanks' Baby-step Giant-step
11 mexp(ll x, ll y, ll p) {
    if(!y) return 1;
    if (y \& 1) return x * mexp(x*x%p, y>>1, p) % p;
    return mexp(x*x\%p, y>>1, p);
vector<ll> get_factor(ll n) {
    vector<11> v;
    for(ll i=2;i*i<=n;i++) {</pre>
        if(n \% i == 0) {
             v.push_back(i);
             while(n \% i == 0) n /= i;
    }
    if(n > 1) v.push_back(n);
    return v:
ll get_primitive(ll n) {
    ll phi = n-1; // assume n is prime
    vector<ll> fact = get_factor(phi);
    for(11 x=2;x<=n;x++) {
        int yes = 1;
        for(ll v : fact) {
             yes &= (mexp(x, phi / y, n) != 1);
         if(yes) return x;
    }
    return -1;
// find k s.t. g^k = b mod n
ll bsgs(ll g, ll b, ll n) {
```

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pl p = ext_gcd(m, n);

```
11 \text{ sq} = \text{ceil}(\text{sqrt}(n) + 1e-9);
    vector<pl> prec(sq);
    for(ll i=0:i<sa:i++) {</pre>
        prec[i] = {mexp(g, i, n), i};
    }
    sort(prec.begin(), prec.end());
    ll base = mexp(g, n-1-sq, n); // g^(-m), assume n : prime
    11 cur = b:
    for(int i=0;i<sq;i++) {</pre>
        auto it = lower_bound(prec.begin(), prec.end(), pl(cur, 0));
        if(it->first == cur) {
            return (i*sq + it->second) % (n-1);
        cur = cur * base % n;
    }
    return -1;
1.2.3 Extended Euclidean Algorithm
// ax + by = gcd(a,b). x, y?
pll ext_gcd(ll a,ll b) {
    if(b) {
        auto tmp = ext_gcd(b, a%b);
        return {tmp.second, tmp.first - (a/b) * tmp.second};
    }
    else return {1, 0};
}
// ax = gcd(a, m) mod m. x?
11 mod_inv(11 a, 11 m) {
    return (ext_gcd(a, m).first + m) % m;
}
1.2.4 Chinese Remainder Theorem
ll pos rem(ll a, ll m) { // m > 0, a % m?
    11 res = abs(a) % m;
    return a > 0 ? res : (res ? m - res : 0);
// ax = c mod m, bx = d mod n. x?
ll solve(ll a, ll c, ll m, ll b, ll d, ll n) {
    a = pos_rem(a, m); c = pos_rem(c, m); // if a, c not in [0, m)
    b = pos_rem(b, n); d = pos_rem(d, n); // if b, d not in [0, n)
    11 g = gcd(a, gcd(c, m)); a \neq g, c \neq g, m \neq g;
        g = gcd(b, gcd(d, n)); b /= g, d /= g, n /= g;
    if(c % _gcd(a, m) || d % _gcd(b, n)) return inf;
    ll t1 = (mod inv(a, m) * c) \% m:
    11 t2 = (mod_inv(b, n) * d) \% n;
    g = gcd(m, n);
    11 lc = m * n / g;
    if(abs(t1 - t2) % g) return inf;
```

```
11 q = (t1 * p.second * n/g + t2 * p.first * m/g);
    return pos_rem(q, lc);
1.2.5 Möbius Inversion Formula
 \forall n \in \mathbb{N} \ g(n) = \sum_{d \mid n} f(d) \implies f(n) = \sum_{d \mid n} \mu(d) g(n/d)
1.3 FFT
 FFT: (a_0, a_1, \dots, a_{n-1}) \mapsto (\sum_{j=0}^{n-1} a_j(\omega^0)^j, \sum_{j=0}^{n-1} a_j(\omega^1)^j, \dots, \sum_{j=0}^{n-1} a_j(\omega^{n-1})^j)
using base = complex<double>;
void fft(vector<base>& a. bool inv) {
    int n = a.size(), j = 0;
    vector<base> roots(n/2):
    for(int i=1;i<n;i++) {</pre>
        int bit = (n >> 1);
        while(j >= bit) {
             j -= bit;
             bit >>= 1;
        j += bit;
        if(i < j) swap(a[i], a[j]);</pre>
    double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
    for(int i=0;i<n/2;i++)
         roots[i] = base(cos(ang * i), sin(ang * i));
    /* In NTT, let prr = primitive root. Then,
    int prr = 3, ang = mexp(prr, (mod - 1) / n);
    if (inv) ang = mexp(ang, mod - 2);
    for(int i=0: i<n/2: i++){
         roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
    also, make sure to apply modulus under here
    and pre-compute mexp(n, mod-2)
    for(int i=2;i<=n;i<<=1) {</pre>
         int step = n / i:
        for(int j=0;j<n;j+=i) {</pre>
             for(int k=0:k<i/2:k++) {
                  base u = a[j+k], v = a[j+k+i/2] * roots[step * k];
                  a[j+k] = u+v;
                  a[j+k+i/2] = u-v;
    if(inv) for(int i=0;i<n;i++) a[i] /= n;
```

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```
void conv(vector<base>& x, vector<base>& y) {
    int n = 2; while(n < x.size()+y.size()) n <<= 1;
    x.resize(n); y.resize(n);
   fft(x, false): fft(v, false):
    for(int i=0;i<n;i++) x[i] *= y[i];</pre>
    fft(x, true); // access (ll)round(x[i].real())
}
1.4 Miller-Rabin + Pollard-Rho
//Prove By Solving - https://www.acmicpc.net/problem/4149
namespace miller_rabin{
    lint mul(lint x, lint y, lint mod){ return (__int128) x * y % mod; }
    lint ipow(lint x, lint y, lint p){
       lint ret = 1, piv = x \% p;
       while(v){
            if(y&1) ret = mul(ret, piv, p);
            piv = mul(piv, piv, p);
           y >>= 1;
       }
       return ret;
    }
    bool miller_rabin(lint x, lint a){
       if(x % a == 0) return 0;
       lint d = x - 1:
       while(1){
            lint tmp = ipow(a, d, x);
            if(d&1) return (tmp != 1 && tmp != x-1);
            else if(tmp == x-1) return 0;
            d >>= 1;
       }
    }
    bool isprime(lint x){
       for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
           if(x == i) return 1:
            if(x > 40 && miller_rabin(x, i)) return 0;
       if(x \le 40) return 0:
       return 1;
    }
}
namespace pollard_rho{
    lint f(lint x, lint n, lint c){
        return (c + miller_rabin::mul(x, x, n)) % n;
   }
    void rec(lint n, vector<lint> &v){
       if(n == 1) return;
       if(n \% 2 == 0){
            v.push_back(2);
```

```
rec(n/2, v);
           return:
       if(miller_rabin::isprime(n)){
            v.push_back(n);
           return;
       }
       lint a, b, c;
       while(1){
           a = rand() \% (n-2) + 2;
           b = a:
           c = rand() \% 20 + 1;
           do{
               a = f(a, n, c);
               b = f(f(b, n, c), n, c);
           }while(gcd(abs(a-b), n) == 1);
           if(a != b) break:
       lint x = gcd(abs(a-b), n);
       rec(x, v);
       rec(n/x, v):
   vector<lint> factorize(lint n){
        vector<lint> ret;
       rec(n, ret);
       sort(ret.begin(), ret.end());
       return ret:
   }
};
1.5 Black Box Linear Algebra + Kitamasa
const int mod = 998244353:
using lint = long long:
lint ipow(lint x, lint p){
 lint ret = 1, piv = x;
 while(p){
   if(p & 1) ret = ret * piv % mod;
   piv = piv * piv % mod;
   p >>= 1;
 return ret;
vector<int> berlekamp_massey(vector<int> x){
   vector<int> ls, cur;
   int lf, ld;
   for(int i=0: i<x.size(): i++){</pre>
       lint t = 0;
       for(int j=0; j<cur.size(); j++){</pre>
           t = (t + 111 * x[i-j-1] * cur[j]) \% mod;
```

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```
if((t - x[i]) % mod == 0) continue;
        if(cur.empty()){
            cur.resize(i+1):
            lf = i:
           1d = (t - x[i]) \% mod:
            continue;
       }
       lint k = -(x[i] - t) * ipow(ld, mod - 2) % mod:
        vector<int> c(i-lf-1);
       c.push_back(k);
       for(auto &j : ls) c.push_back(-j * k % mod);
        if(c.size() < cur.size()) c.resize(cur.size());</pre>
       for(int j=0; j<cur.size(); j++){</pre>
            c[i] = (c[i] + cur[i]) \% mod;
       }
        if(i-lf+(int)ls.size()>=(int)cur.size()){
            tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) \% mod);
       }
       cur = c:
   }
   for(auto &i : cur) i = (i % mod + mod) % mod:
    return cur;
int get_nth(vector<int> rec, vector<int> dp, lint n){
    int m = rec.size();
   vector<int> s(m), t(m):
   s[0] = 1:
   if(m != 1) t[1] = 1;
   else t[0] = rec[0]:
    auto mul = [&rec](vector<int> v, vector<int> w){
       int m = v.size();
       vector<int> t(2 * m);
       for(int j=0; j<m; j++){</pre>
            for(int k=0: k<m: k++){
                t[i+k] += 111 * v[j] * w[k] % mod;
                if(t[j+k] >= mod) t[j+k] -= mod;
           }
        for(int j=2*m-1; j>=m; j--){
           for(int k=1: k<=m: k++){
               t[j-k] += 111 * t[j] * rec[k-1] % mod;
                if(t[j-k] >= mod) t[j-k] -= mod;
            }
       }
       t.resize(m):
       return t;
   }:
    while(n){
       if(n \& 1) s = mul(s, t);
       t = mul(t, t):
```

```
n >>= 1;
   }
   lint ret = 0:
   for(int i=0; i<m; i++) ret += 111 * s[i] * dp[i] % mod;
   return ret % mod;
int guess_nth_term(vector<int> x, lint n){ // init with > 3k, 0(1^2 lg n)
    if (n < x.size()) return x[n]:
   vector<int> v = berlekamp_massey(x);
   if(v.empty()) return 0;
   return get_nth(v, x, n);
struct elem\{int x, y, v;\}; // A_(x, y) <- v, 0-based. no duplicate please...
vector<int> get_min_poly(int n, vector<elem> M){
   // smallest poly P such that A^i = sum_{j < i} {A^j \times P_j}
    vector<int> rnd1. rnd2:
    mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
   };
   for(int i=0: i<n: i++){
        rnd1.push_back(randint(1, mod - 1));
        rnd2.push_back(randint(1, mod - 1));
   }
   vector<int> gobs;
   for(int i=0: i<2*n+2: i++){
        int tmp = 0:
        for(int j=0; j<n; j++){</pre>
            tmp += 111 * rnd2[i] * rnd1[i] % mod:
            if(tmp >= mod) tmp -= mod;
        gobs.push_back(tmp);
        vector<int> nxt(n);
        for(auto &i : M){ // sparse matrix * vector
            nxt[i.x] += 111 * i.v * rnd1[i.v] % mod:
            if(nxt[i.x] >= mod) nxt[i.x] -= mod;
        rnd1 = nxt;
   auto sol = berlekamp massev(gobs):
   reverse(sol.begin(), sol.end());
   return sol:
lint det(int n, vector<elem> M){
   vector<int> rnd:
    mt19937 rng(0x14004);
   auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
   };
   for(int i=0; i<n; i++) rnd.push_back(randint(1, mod - 1));</pre>
```

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```
for(auto &i : M){
       i.v = 111 * i.v * rnd[i.v] % mod:
    }
    auto sol = get_min_poly(n, M)[0];
    if(n \% 2 == 0) sol = mod - sol:
    for(auto &i : rnd) sol = 111 * sol * ipow(i, mod - 2) % mod;
    return sol;
}
2 Geometry
     struct Point
const double eps = 1e-10:
template <class T>
struct point{
    typedef point P;
    T x, y;
    point(T x=0, T y=0) : x(x), y(y) {}
    bool operator< (P a) const {return x!=a.x ? x<a.x : y<a.y;}
    bool operator == (P a) const {return x == a.x && y == a.y;}
    P operator+ (P a) const {return P(x+a.x, y+a.y);}
    P operator- (P a) const {return P(x-a.x, y-a.y);}
    P operator- () const {return P(-x, -y);};
    T operator* (P a) const {return x*a.x + v*a.v:} // inner prod
    T operator/ (P a) const {return x*a.y - y*a.x;} // outer prod
    T dist2() const {return x*x + y*y;}
    double dist() const {return sqrt(double(dist2()));}
    P perp() const {return P(-v, x);}; // rotate 90 deg ccw
    P mult(T t) const {return P(x*t, y*t);}
    P unit() const {return P(x/dist(), y/dist());}
    P rotate(double a){
        return P(x*cos(a)-v*sin(a), x*sin(a)+v*cos(a)):
    }
};
int sgn(double x) {return (x > eps) - (x < -eps);}</pre>
typedef point<ll> P; // use point<double> at caution with fabs, eps
2.2 Distance, Intersection
2.2.1 Point-to-Line
double lineDist(P a, P b, P p) {
    return ((b-a)/(p-a))/(b-a).dist(); // a->b : left (+), right : (-);
}
2.2.2 Point-to-Segment
double segDist(P s, P e, P p) {
    if(s == e) return (p-s).dist(); // mind the eps
    double d = (e-s).dist2(), t = min(d, max(.0, (p-s)*(e-s)));
```

```
return ((p-s).mult(d)-(e-s).mult(t)).dist() / d;
2.2.3 Line intersection
template<class P>
pair<int, P> lineInter(P a, P b, P c, P d){
   if((b-a)/(d-c) == 0) // parallel, mind the eps
        return \{-((b-a)/(c-a) == 0), a\};
   double oa = (d-c)/(a-c), ob = (d-c)/(b-c);
   return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
} // 1.0.-1(inf) : inter
2.2.4 Segment Intersection
bool onSegment(P s, P e, P p) {return segDist(s, e, p) < eps;}</pre>
template < class P > vector < P > segInter(P a, P b, P c, P d) {
   double oa = (d-c)/(a-c), ob = (d-c)/(b-c),
            oc = (b-a)/(c-a), od = (b-a)/(d-a);
   if(sgn(oa)*sgn(ob) < 0 && sgn(oc)*sgn(od) < 0)
        return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
   set<P> S:
   if(onSegment(c, d, a)) S.insert(a);
   if(onSegment(c, d, b)) S.insert(b);
   if(onSegment(a, b, c)) S.insert(c);
   if(onSegment(a, b, d)) S.insert(d);
   return vector<P>(S.begin(), S.end()):
}
2.2.5 Circle-Line Intersection
vector<P> circLine(P A, P B, P O, double r){
   vector<P> v;
   P X = O-A, D = B-A;
   double rat = 1.0 * (X*D) / (D*D);
   double det = (X*D)*(X*D) - (D*D)*(X*X - r*r);
   if(det < 0) return {}:</pre>
   else if(det < eps) return {P(A + D.mult(rat))};</pre>
   det = sqrt(det);
   return {P(A + D.mult(rat + det/(D*D))).
           P(A + D.mult(rat - det/(D*D)));
2.2.6 Circle-Line Tangent
vector<P> circTangent(P A, P O, double r){
   if((0-A).dist2() < r*r + eps) return {}:
   double th = asin(r/(0-A).dist());
   return {A + (0-A).rotate(th). mult(cos(th)),
            A + (O-A).rotate(-th).mult(cos(th)):
```

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2.3 Convex Hull

```
vector<P> make_hull(vector<P>& v) {
    sort(v.begin(), v.end());
    int n = v.size(), sz = 0;
    vector<P> L. U:
    for(int i=0;i<n;i++) {</pre>
        while(sz >= 2 && (L[sz-1]-L[sz-2])/(v[i]-L[sz-1]) <= 0)
            L.pop_back(), sz--;
        L.push_back(v[i]); sz++;
    }
    sz = 0;
    for(int i=n-1:i>=0:i--) {
        while(sz >= 2 && (U[sz-1]-U[sz-2])/(v[i]-U[sz-1]) <= 0)
            U.pop_back(), sz--;
        U.push_back(v[i]); sz++;
   L.pop_back(), U.pop_back();
    L.insert(L.end(), U.begin(), U.end());
    return L;
}
```

2.4 Sorting Points by Angle

```
// credit : http://koosaga.com/97
auto angle_sort = [&] (P &a, P &b){
    if((a < P(0, 0)) ^ (b < P(0, 0))) return b < a;
    if(a / b != 0) return a / b > 0;
    return a.dist2() < b.dist2(); // norm
}; // counter-clockwise sort</pre>
```

2.5 Rotating Calipers

```
11 rotating_calipers(vector<P> v) {
    v = make_hull(v);
    ll mdist = 0;
    int j = 1, n = v.size();
    v.insert(v.end(), v.begin(), v.end());
    for(int i=0;i<n;i++) {
        while((v[i+1]-v[i])/(v[j+1]-v[j]) > 0) j++;
        for(int k=0;k<2;k++)
            mdist = max(mdist, (v[j+k]-v[i]).dist2());
    }
    return mdist;
}</pre>
```

2.6 Smallest Enclosing Circle

```
//Prove By Solving - https://www.acmicpc.net/problem/11930
int main(){
    scanf("%d", &N);
```

```
for(int i = 1; i \le N; i ++) scanf("%lf%lf%lf", &A[i].x, &A[i].y, &A[i].z);
   int t = 70000:
   double rate = 1.0;
    point cur = (point)\{0, 0, 0\};
   for(int i = 1 ; i <= t; i++){
        int ind = 1;
        for(int j = 1 ; j <= N ; j++)
        if( (A[i] - cur) * (A[i] - cur) >
            (A[ind] - cur) * (A[ind] - cur)) ind = j;
        cur = cur + (A[ind] - cur).mult(rate);
        rate *= 0.99:
   }
   double r = 0;
   for(int i = 1; i \le N; i++) r = max(r, (A[i] - cur) * (A[i] - cur));
   cout << setprecision(10) << fixed << sqrt(r);</pre>
    return 0;
} // Non-deterministic, deterministic O(n lg n) requires Voronoi diagram
2.7 Circumcircle
double cc_radius(P& A, P& B, P& C){
    return (B-A).dist() * (C-B).dist() * (A-C).dist() / fabs((A-B) / (B-C)) / 2;
P cc_center(P& A, P& B, P& C){
   P b = C-A, c = B-A:
    return A + (b.mult(c.dist2()) - c.mult(b.dist2())).perp().mult(0.5/(b/c));
2.8 Polygon Area
double ans = 0; // ans : double area
for(int i=0;i<points.size();i++)</pre>
   ans += points[i] / points[(i+1 == points.size() ? 0 : i+1)];
```

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3 Strings

3.1 Aho-Corasick Algorithm

```
// required macro : rep, repi, setz
namespace aho_corasick {
    const int MAXN = 200002, MAXC = 26;
    int go[MAXN+1][MAXC];
    int fail[MAXN+1];
    bool term[MAXN+1];
    void build(const vector<string> &v) {
        setz(go), setz(fail), setz(term);
        int cnode = 1;
        repi(s, v) {
            int p = 0;
                                // root is 0
            repi(j, s) {
                char c = j-'a'; // check starting alphabet
                if (!go[p][c]) go[p][c] = cnode++;
                p = go[p][c];
            }
            term[p] = true;
        }
        queue<int> q; rep(i, MAXC) if (go[0][i]) q.push(go[0][i]);
        while(!q.empty()) {
            int t = q.front(); q.pop();
            rep(i, MAXC) {
                if (go[t][i]) {
                    int p = fail[t];
                    while(p and not go[p][i]) p = fail[p];
                    p = go[p][i];
                    fail[go[t][i]] = p;
                    if (term[p]) term[go[t][i]] = true;
                    q.push(go[t][i]);
               }
           }
       }
    }
    int step(int p, int c) {
        if(go[p][c]) return go[p][c];
        if(p) return go[p][c] = step(fail[p], c); // if not root
        return go[p][c] = -1; // -1 is same as 0, but is visited
    }
    bool query(string &t) {
        int p = 0;
        repi(i, t) {
            char c = i-'a'; // check starting alphabet
```

```
p = max(0, step(p, c));
            /* if you need the original trie, use this or make length array
            while(p and not go[p][c]) p = fail[p];
            p = max(0, go[p][c]);
            */
            if (term[p]) return true;
       return false;
}
3.2 Suffix Array
// required macro : rep, repr, repp
// str : abracadabra
// SA : 10 7 0 3 5 8 1 4 6 9 2
                                    (0-based)
// LCP : 1 4 1 1 0 3 0 0 0 2
                                    (lcp[i] : lcp of sa[i], sa[i+1])
vector<int> make_sa(const string& s) {
    int n = s.length();
    int \lim = \max(128, n+1);
   vector<int> sa(n), g(n+1), ng(n+1), cnt(lim), ind(lim+1);
   rep(i, n) sa[i] = i, g[i] = s[i];
    g[n] = 0;
   for(int t=1;t<s.length();t<<=1)</pre>
        auto cmp = [&] (int a, int b) {
            return g[a] != g[b] ? g[a] < g[b] : g[a+t] < g[b+t];
        };
       rep (i, n)
                        cnt[g[min(i+t, n)]]++;
        repp(i, 1, lim) cnt[i] += cnt[i-1];
        repr(i, n, 0) ind[--cnt[g[min(i+t, n)]]] = i;
        rep (i, lim) cnt[i] = 0;
        rep (i, n)
                        cnt[g[i]]++; // same as cnt[g[ind[i]]]++
       repp(i, 1, lim) cnt[i] += cnt[i-1]:
        repr(i, n, 0) sa[--cnt[g[ind[i]]] = ind[i];
        ng[sa[0]] = 1:
        repp(i, 1, n) ng[sa[i]] = ng[sa[i-1]] + cmp(sa[i-1], sa[i]);
        g = ng;
       fill(cnt.begin(), cnt.end(), 0);
        fill(ind.begin(), ind.end(), 0);
   }
   return sa:
vector<int> make_lcp(const string& s, const vector<int>& sa) {
    int n = s.length(), len = 0;
   vector<int> lcp(n-1), rank(n);
   rep(i, n) rank[sa[i]] = i;
   rep(i, n) {
       if(rank[i]) {
            int j = sa[rank[i]-1];
```

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```
int lc = n - max(i,j);
            while(len < lc && s[i+len] == s[j+len]) len++;</pre>
           lcp[rank[i]-1] = len;
       if(len) len--:
    }
    return lcp;
}
     Lexicographically Smallest String Rotation
int min_rotation(string s) {
    int a=0, N=s.size(); s += s;
    rep(b,N) rep(i,N) {
       if (a+i == b \mid \mid s[a+i] < s[b+i]) \{b += max(0, i-1); break;\}
       if (s[a+i] > s[b+i]) \{ a = b; break; \}
    }
    return a; // rotate(v.begin(), v.begin()+min_rotation(v), v.end());
}
     Manacher's Algorithm
// 0-based
// s = # h # e # l # l # o #
// ret = 0 1 0 1 0 1 2 1 0 1 0
vector<int> manacher(const string& s) {
    int n = s.size(), r = -1, k = -1;
    vector<int> p(n);
   for (int i=0; i<n; i++) {
       if (i<=r) p[i] = min(r-i, p[2*k-i]);</pre>
       while (i-p[i]-1>=0 and i+p[i]+1<n and s[i-p[i]-1] == s[i+p[i]+1]) p[i]++;
       if (r < i+p[i]) r = i+p[i], k = i;
    }
    return p;
}
3.5 Z Algorithm
// 0-based
//s = abcababca
// ret = 9 0 0 2 0 4 0 0 1
vector<int> z_algo(const string &s) {
    int 1 = 0, r = 0, N = sz(s);
    vector<int> Z(N);
   Z[0] = N:
    repp(i, 1, N) {
       if (i > r) {
           l = r = i:
            while(r < N and s[r] == s[r-1]) r++;
```

```
r--;
    Z[i] = r-l+1;
} else {
    int k = i-l;
    if (Z[k] < r-i+1) Z[i] = Z[k];
    else {
        1 = i;
        while(r < N and s[r] == s[r-l]) r++;
        r--;
        Z[i] = r-l+1;
    }
}
return Z;
}</pre>
```

4 Graph Theory

4.1 Biconnected Component

```
// https://gist.github.com/koosaga/6f6fd50dd7067901f1b1
void dfs(int x, int p){
   dfn[x] = low[x] = ++piv;
   par[x] = p;
   for(int i=0; i<graph[x].size(); i++){</pre>
        int w = graph[x][i];
        if(w == p) continue;
        if(!dfn[w]){
            dfs(w, x);
            low[x] = min(low[x], low[w]);
        else low[x] = min(low[x], dfn[w]);
   }
}
void color(int x, int c){
    if(c > 0) bcc[x].push_back(c); // c == 0 : first component
   vis[x] = 1:
   for(int i=0; i<graph[x].size(); i++){</pre>
        int w = graph[x][i];
        if(vis[w]) continue;
        if(dfn[x] <= low[w]){</pre>
            bcc[x].push_back(++cpiv);
            color(w, cpiv);
        else color(w, c);
   }
```

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4.2 Strongly Connected Component

```
const int MAXN = 2e5 + 10; // > 2*N
int N. M:
int dfsn[MAXN], low[MAXN], finished[MAXN], cnt;
vector<int> ADJ[MAXN];
vector<vector<int>> G;
stack<int> S;
int f(int x){ // f(1) f(-1) f(2) f(-2) f(3) f(-3) ... -> 0 1 2 3 4 5 ...
    return 2 * (abs(x) - 1) + (x < 0);
}
void add_edge(int x, int y){ // call by f(x), f(y)
    ADJ[x ^ 1].push_back(y);
    ADJ[y ^ 1].push_back(x);
}
// memset(finished, -1, sizeof(finished));
int scc(int here){
    static vector<int> tmp;
    S.push(here);
    dfsn[here] = low[here] = ++cnt;
    int &ret = low[here]:
    for(int there : ADJ[here]){
        if(dfsn[there] == 0) ret = min(ret, scc(there));
        else if(finished[there] == -1) ret = min(ret, dfsn[there]);
    }
    if(dfsn[here] == low[here]){
        while(1){
            int x = S.top(); S.pop();
            finished[x] = G.size();
            tmp.push_back(x);
            if(x == here) break;
        G.push_back(tmp);
        tmp.clear();
    }
    return ret;
}
4.2.1 2-SAT
```

scc를 실행시켜 f(i) 와 f(-i)가 같은 component에 있다면, 모순. f(i) 와 f(-i) 중 finished 배열의 수가 작은 것이 참이다. (SCC numbering의 역순이 위상정렬이기에, $F \rightarrow T$ 를 유지하기 위함)

Euler Tour 4.3

```
struct Edge{
   int to, cnt; // to: 인접한 정점, cnt: 남은 사용 횟수
   Edge *dual; // dual: 역방향 간선을 가리키는 포인터
```

```
Edge(): Edge(-1, 0){}
    Edge(int to1, int cnt1): to(to1), cnt(cnt1), dual(nullptr){}
void Eulerian(int curr){
    for(Edge *e: adi[curr]){
        if(e->cnt > 0){
            e->cnt--;
            e->dual->cnt--;
            Eulerian(e->to); // dfs
    }
    cout << curr << '\n';
 4.4 Heavy-Light Decomposition
int par[16] [MAXN], head[MAXN], sz[MAXN], in[MAXN], out[MAXN], lv[MAXN];
 vector<int> v[MAXN];
int dfs0(int x, int p) {
    lv[x] = lv[p] + 1; // lv[0] = -1;
    sz[x] = 1:
    for(int& y : v[x]) {
        if(y == p) continue;
        par[0][y] = x;
        sz[x] += dfs0(y, x);
        if(sz[v[x][0]] < sz[y]) swap(y, v[x][0]);
    }
    return sz[x];
void dfs1(int x, int p) {
    static int cnt = 0:
    in[x] = ++cnt:
    for(int v : v[x]) {
        if(y == p) continue;
        head[y] = y == v[x][0] ? head[x] : y;
        dfs1(v, x);
    }
    out[x] = cnt;
int get_query(segtree& seg, int x, int p) {
    int ret = 0:
    while(head[x] != head[p]) {
        // query(in[head[x]], in[x]);
        x = par[0][head[x]];
    // query(in[p], in[x]);
    return ret;
// for(int i=1;i<=n;i++) if(!in[i]) head[i] = i, dfs0(i, 0), dfs1(i, 0);
```

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4.5 Dominator Tree

```
namespace Dtree {
    const int MAXN = 250001:
    vector<int> E[MAXN], RE[MAXN], rdom[MAXN];
    int S[MAXN], RS[MAXN], cs;
    int par[MAXN], val[MAXN];
    int sdom[MAXN], rp[MAXN];
    int dom[MAXN];
    int Find(int x, int c = 0) {
       if (par[x] == x) return c?-1:x;
       int p = Find(par[x], 1);
       if (p == -1) return c?par[x]:val[x];
       if (sdom[val[x]] > sdom[val[par[x]]]) val[x] = val[par[x]];
       par[x] = p;
       return c?p:val[x];
   }
   void Union(int x, int y) {
       par[x] = y;
   }
   void dfs(int x) {
       RS[S[x] = ++cs] = x;
       par[cs] = sdom[cs] = val[cs] = cs;
       for(int e : E[x]) {
           if (S[e] == 0) dfs(e), rp[S[e]] = S[x];
           RE[S[e]].pb(S[x]);
       }
   }
    int Do(int s, int *up) {
       dfs(s):
       for (int i = cs;i;i--) {
           for (int e : RE[i]) sdom[i] = min(sdom[i], sdom[Find(e)]);
           if (i > 1) rdom[sdom[i]].pb(i);
           for (int e:rdom[i]) {
                int p = Find(e):
               if (sdom[p] == i) dom[e] = i;
               else dom[e] = p;
           }
            if (i > 1) Union(i, rp[i]);
       for (int i = 2; i <= cs; i++) if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
        for (int i = 2; i <= cs; i++) {
           up[RS[i]] = RS[dom[i]]:
       }
       return cs;
   }
```

```
void addE(int x, int y) {E[x].pb(y);}
4.6 Global Min Cut
// Stoer-Wagner Algorithm, O(VE lg E)
int minimum_cut_phase(int n, int &s, int &t,
   vector<vector<int>> &adj, vector<int> vis){
   vector<int> dist(n):
   int mincut = 1e9;
   while(true){
        int pos = -1, cur = -1e9;
       for(int i=0; i<n; i++){
            if(!vis[i] && dist[i] > cur){
                cur = dist[i];
                pos = i;
           }
        if (pos == -1) break;
        s = t:
       t = pos;
        mincut = cur;
        vis[pos] = 1;
        for(int i=0; i<n; i++){</pre>
            if(!vis[i]) dist[i] += adj[pos][i];
   }
   return mincut; // optimal s-t cut here is, {t} and V \ {t}
int solve(int n, vector<vector<int>> adj){
   if(n <= 1) return 0;
   vector<int> vis(n):
   int ans = 1e9:
   for(int i=0; i<n-1; i++){</pre>
        int s, t;
        ans = min(ans, minimum_cut_phase(n, s, t, adj, vis));
        vis[t] = 1;
       for(int j=0; j<n; j++){</pre>
           if(!vis[j]){
                adj[s][j] += adj[t][j];
                adj[j][s] += adj[j][t];
           }
       }
        adj[s][s] = 0;
   }
   return ans;
```

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5 Network Flow

5.1 Theorems

Max-flow Min-cut theorem : 정점 s에서 정점 t까지 흐를 수 있는 최대 유량(max-flow)은 정점 s와 정점 t를 분리하는 간선들의 가중치 합(min-cut)과 같다.

Vertex cover: 어떤 그래프의 정점의 집합 S에 대해 그래프의 모든 간선이 S의 원소 중 최소 하나와 연결되어 있을 때, S를 해당 그래프의 vertex cover라고 하며, minimum vertex cover는 최소 개수의 정점을 사용한 vertex cover이다.

Independent set : 어떤 그래프의 정점의 집합 S에 대해 S의 서로 다른 두 정점을 연결하는 간선이 없을 때, S를 해당 그래프의 independent set이라고 하며, maximum independent set은 최대 개수의 정점을 사용한 independent set이다.

Matching (independent edge set) : 어떤 그래프의 간선의 집합 S에 대해 S의 서로 다른 두 간선이 공통된 정점을 가지지 않을 때, S를 해당 그래프의 matching이라고 하며, maximum matching은 최대 개수의 간선을 사용한 matching이다.

König's theorem : 이분 그래프의 maximum matching의 크기는 minimum vertex cover의 것과 같다. Dinic's Algorithm : 시간 복잡도 $O(V^2E)$, unit capacity에서는 $\min(V^{2/3}E, E^{3/2})$.

Circulation Problem : 새로운 source/sink s_n , t_n 를 만들어서 다음과 같이 간선을 추가하고 $maxflow(s_n \to t_n) = \sum l_i$ 인지 확인, 이후 $s \to t$ 로 maxflow

```
• s_n \to b (l), a \to t_n (l), a \to b (r-l), t \to s (\infty)
```

5.2 Dinic's Algorithm

```
const int INF = 1e9;
struct Dinic{
    int N:
    struct edge{
       int index, cap, rev;
        edge(): index(0), cap(0), rev(0) {}
        edge(int index, int cap, int rev) : index(index), cap(cap), rev(rev) {}
    };
    vector<vector<edge>> ADJ;
    vector<int> R. W:
   Dinic() {}
   Dinic(int N) : N(N){
        ADJ.resize(N); R.resize(N);
                                        W.resize(N);
    }
    void CE(int node1, int node2, int cap){
        ADJ[node1].push_back(edge(node2, cap, ADJ[node2].size()));
        ADJ[node2].push_back(edge(node1, 0, ADJ[node1].size() - 1));
   }
   bool bfs(int src, int sink){
       fill(R.begin(), R.end(), -1):
        R[src] = 0;
        queue<int> Q; Q.push(src);
        while(Q.size()){
            int here = Q.front(); Q.pop();
```

```
for(auto e : ADJ[here]){
                if(e.cap > 0 && R[e.index] == -1)
                    R[e.index] = R[here] + 1, Q.push(e.index);
        return R[sink] != -1;
   }
    int dfs(int here, int sink, int f){
       if(here == sink) return f:
       for(int &i = W[here] ; i < ADJ[here].size() ; i++){</pre>
            auto &e = ADJ[here][i];
           if(e.cap > 0 && R[here] < R[e.index]){</pre>
                int res = dfs(e.index, sink, min(f, e.cap));
                if(res) {
                    e.cap -= res:
                    ADJ[e.index][e.rev].cap += res;
                    return res;
           }
       }
       return 0;
   }
   int solve(int src, int sink){
       int ret = 0:
        while(bfs(src. sink)){
           fill(W.begin(), W.end(), 0);
           int res:
            while((res = dfs(src, sink, INF))) ret += res;
       }
       return ret;
   }
}:
5.3 MCMF with SPFA
template <typename T>
struct MCMF {
   const T INF = 1e9;
   struct edge {
       T there, cap, cost, rev;
        edge(T there=0, T cap=0, T cost=0, T rev=0)
            : there(there), cap(cap), cost(cost), rev(rev) {}
   };
   int N;
   vector<vector<edge>> ADJ:
   vector<T> R, INQ, C, I;
   MCMF(int N=0) : N(N), ADJ(N+1), R(N+1), INQ(N+1), C(N+1), I(N+1) {}
```

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```
void CE(int i, int j, T cap, T cost) {
        ADJ[i].push_back(edge(j, cap, cost, ADJ[j].size()));
        ADJ[i].push_back(edge(i, 0, -cost, ADJ[i].size() - 1));
    }
    bool SPFA(int src, int sink) {
        queue<int> Q;
                        Q.push(src);
       fill(R.begin(), R.end(), -1);
                                           R[src] = 0:
       fill(C.begin(), C.end(), INF);
                                           C[src] = 0;
       fill(INQ.begin(), INQ.end(), 0);
                                           INQ[src] = 1;
        while (Q.size()) {
            int here = Q.front();
                                   Q.pop();
            INO[here] = 0:
            for (int i = 0; i < ADJ[here].size(); i++) {</pre>
                auto e = ADJ[here][i];
                if ((C[e.there] == INF || C[e.there] > C[here] + e.cost) &&
                    e.cap > 0) {
                   C[e.there] = C[here] + e.cost;
                   R[e.there] = here:
                   I[e.there] = i;
                   if (!INQ[e.there]) INQ[e.there] = 1, Q.push(e.there);
               }
            }
       }
        if (C[sink] == INF) return false;
        return true:
    }
    pair<T. T> mcmf(int src, int sink) {
       pair<T, T> ret{};
       while (SPFA(src, sink)) {
            T flow = INF. cost = 0:
            for (int here = sink; here != src; here = R[here])
               flow = min(flow, ADJ[R[here]][I[here]].cap);
            for (int here = sink: here != src: here = R[here]) {
                auto &e = ADJ[R[here]][I[here]];
                cost += e.cost * flow;
               e.cap -= flow;
               ADJ[e.there][e.rev].cap += flow;
            }
            ret.first += flow, ret.second += cost;
       }
       return ret;
    }
};
5.4 Hungarian Method
```

```
namespace Hung {
    const int MX = 2000;
   // IMPORTANT : n <= m, 1-based
```

```
using T = long double;
    T \max v = 1e200:
    T a[MX][MX], n, m;
    void init(int nn, int mm) { n = nn; m = mm; }
    void set_value(int x, int y, T val) { a[x][y] = val; }
    T solve(vector <int> &ans) {
        vectorT> v(m+1), u(n+1);
        vector\langle int \rangle p (m+1), way (m+1);
        for (int i=1; i<=n; ++i) {
            p[0] = i;
            int j0 = 0;
            vector<T> minv (m+1, maxv);
            vector<char> used (m+1, false);
            do {
                used[j0] = true;
                T delta = maxv;
                int i0 = p[j0], j1;
                for (int j=1; j<=m; ++j) if (!used[j]) {</pre>
                    T cur = a[i0][j]-u[i0]-v[j];
                    if (cur < minv[j]) minv[j] = cur, wav[j] = j0;</pre>
                    if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
                for (int j=0; j<=m; ++j) {
                    if (used[j]) u[p[j]] += delta, v[j] -= delta;
                    else minv[i] -= delta:
                }
                i0 = j1;
            } while (p[j0] != 0);
            do {
                int j1 = way[j0];
                p[j0] = p[j1];
                j0 = j1;
            } while (j0);
        ans.resize(n + 1);
        for(int j=1; j<=m;++j) {</pre>
            ans[p[j]] = j;
        return -v[0];
    }
}
5.5 Hopcroft-Karp Algorithm
struct hopcroft karp{
    int N;
```

```
vector<vector<int>> ADJ;
vector<int> L, rev, used;
```

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```
hopcroft_karp() {}
hopcroft_karp(int N) : N(N) {
   ADJ.resize(N):
   L.resize(N), rev.resize(N, -1), used.resize(N, 0);
}
void CE(int here, int there){
    ADJ[here].push_back(there);
}
void bfs(){
    queue<int> Q;
   for(int i = 0; i < N; i++) {
        if(used[i]) L[i] = -1;
        else L[i] = 0, Q.push(i);
   }
    while(Q.size()){
        int here = Q.front(); Q.pop();
        for(int there : ADJ[here]){
            if(rev[there] != -1 && L[rev[there]] == -1) {
                L[rev[there]] = L[here] + 1;
                Q.push(rev[there]);
           }
        }
   }
}
bool dfs(int here){
   for(int there : ADJ[here]){
        if(rev[there] == -1 || (L[here] < L[rev[there]] && dfs(rev[there]))){</pre>
            rev[there] = here;
            used[here] = 1;
           return true:
        }
   }
   return false;
}
int solve(){
    int ret = 0;
    while(1){
        bfs():
        int res = 0;
        for(int i = 0; i < N; i++) {
           if(used[i]) continue;
           res += dfs(i);
        }
        if(res == 0) break;
        ret += res;
```

```
}
        return ret;
    }
};
6 Optimization Tricks
 An array P[i][j] is Monge array if \forall a < b < c < d, C[a][c] + C[b][d] < C[a][d] + C[b][c]
6.1 Knuth Optimization
 • Recurrence : D[i][j] = \min_{i < k < j} (D[i][k] + D[k][j]) + C[i][j]
 • Condition: \forall a < b < c < d, C[i][j] is Monge array and C[b][c] < C[a][d] (monotonic)
 • A[i][j] = (\min k \ s.t. \ D[i][j] \ \text{is min.}). Then A[i][j-1] \le A[i][j] \le A[i+1][j]
  • Time Complexity: O(N^3) \to O(N^2)
// opt[i-1][i] = i
for(int d=2;d<=n;d++) {</pre>
    for(int i=1;i+d<=n+1;i++) {</pre>
        for(int k=opt[i][j-1], j=i+d; k<=opt[i+1][j]; k++) {</pre>
             int v = dp[i][k] + dp[k][j] + c[i][j];
             if(dp[i][j] > v) dp[i][j] = v, opt[i][j] = k;
    }
}
6.2 Divide and Conquer Optimization
 • Recurrence : D[t][i] = \min_{k < i} (D[t-1][k] + C[k][i])
 • Condition: C[i][i] is Monge array.
 • A[t][i] \le A[t][i+1] (A[t][i] = (\min k \ s.t. \ D[t][i] \text{ is min.})) holds
  • Able to Divide and Conquer base on calculating D[t][i]
 • Time Complexity: O(TN^2) \to O(TN \lg N)
// range of index : [1,r]
// range of dp : [s,e]
void dnc(int t, int l, int r, int s, int e)
    if(s > e) return;
    int m = (s+e)/2;
    D[t][m] = 2e9:
    for(int k=1;k<m&k<=r;k++)</pre>
        int tmp = D[t-1][k] + C[k][m];
```

if(D[t][m] > tmp)

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```
D[t][m] = tmp, A[t][m] = k;
}
dnc(t, 1, A[t][m], s, D[t][m]-1);
dnc(t, A[t][m], r, D[t][m]+1, e);
}
```

6.3 Convex Hull Trick

- Recurrence : $dp[i] = \min_{j < i} (dp[j] + a[i]b[j])$
- Condition: None, but usually b[i] is monotonic. If not, refer to Section 7.3.
- Time Complexity: $O(N^2) \to O(N \lg N)$, O(N) if a[i] is monotonic by taking valid front.

```
using ll = long long;
struct line {
    ll a. b. i:
    line(ll a=0, ll b=0, ll i=0) : a(a), b(b), i(i) {}
   11 calc(ll x) { return a*x+b; }
    double inter(line 1) { return double(1.b - b) / (a - 1.a): }
};
bool check_back(line& p0, line& p1, line& p2) {
    return (p1.b - p0.b) * (p1.a - p2.a) >= (p0.a - p1.a) * (p2.b - p1.b);
    return p0.inter(p1) > p1.inter(p2); // double
bool check_front(ll x, line& p0, line& p1) {
    return x * (p0.a - p1.a) >= (p1.b - p0.b);
    return x > p0.inter(p1);
struct ConvexHull {
    int 1 = 0, r = 0:
    vector<line> v;
    void insert(line p) {
        while(r - 1 >= 2 && check back(v[r-2], v[r-1], p))
            v.pop_back(), r--;
        v.push back(p):
    }
    pair<11, 11> query(11 x) { // if x is increasing, else do binary search
        while(r - 1 >= 2 && check front(x, v[1], v[1+1]))
            1++:
        return pair<11, 11>{v[1].calc(x), v[1].i}:
    }
};
void calc() { // psuedo-code
    ConvexHull ch;
    11 val = 0, cnt = 0; // current dp and track
    for(int i=0:i<N:i++) {</pre>
        ch.insert(line{p[i], q[i]}); // j-dependent values (b[j], dp[j])
        pll res = ch.query(x[i]);
        val = res.first + c[i]; // i-dependent values
        cnt = res.second + 1;
```

```
return pll{val, cnt};
6.4 Centroid Decomposition
// credit : https://gist.github.com/igorcarpanese/75162f3253bd230abd0f32f9950bf384
int dfs(int u, int p) {
   sub[u] = 1:
   for (auto v : tree[u])
       if (v != p and !chk[v]) sub[u] += dfs(v, u);
   return sub[u] + 1:
} // calc subtree size, chk means that node was a centroid
// each tree has at most two centroids
int centroid(int u, int p, int r) { // r : root
   for (auto v : tree[u])
        if (v != p and !chk[v] and sub[v] > sub[r]/2) return centroid(v, u, r);
   return u;
6.5 Aliens Trick
 • Recurrence: dp[i][j] = Min_{k < j}(dp[i-1][k] + C[k+1][j])
 • Condition: dp[*][n] is convex (implied if dp[*][n] is Monge array)
11 calc(...) {
   // (...)
   k = min(k, (int)v.size());
   11 leff = some min value. reff = some big value: // check range and overflow
   int lind = v.size(), rind = 1; // range of k, decreases when eff increase
   while(reff - leff >= 2) {
        11 meff = (leff + reff) / 2:
        pll res = calc_with_lambda(v, meff); // {value, k'}
        if(res.second < rind) res.second = rind: // rind (<) lind
        if(res.second > lind) res.second = lind: // lind (>) rind
        if(res.second == k) return res.first - k * meff:
        else if(res.second > k) lind = res.second; leff = meff; // lind (>) rind
        else rind = res.second; reff = meff;
   // when minimizing, ans is max(calc with lambda().first - k * eff) for all eff
   // when maximizing, calculate min instead
   return max(calc_with_lambda(v, leff).first - k * leff,
               calc with lambda(v. reff).first - k * reff);
```

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7 Data Structure

7.1 Persistent Segment Tree

```
const MAXN = 1e5 + 10;
struct node{
   node *1, *r;
    int cnt;
   node () {}
} pool[(1 << 17) * 17], *tree_head[MAXN];</pre>
int tcnt:
node* alloc(){
    memset(pool + tcnt, 0, sizeof(node));
    return pool + tcnt++;
}
node * init(int 1, int r){
    node *ret = alloc();
   if(1 != r) {
       int mid = (1 + r) / 2;
        ret->l = init(l, mid);
       ret->r = init(mid + 1, r):
   }
    return ret;
}
void update(node * here, node *par, int 1, int r, int val){
    if(1 == r) {
       here->cnt = par->cnt + 1;
        return;
   }
    int mid = (1 + r) / 2:
    if(val <= mid){</pre>
       here->1 = alloc():
       here->r = par->r;
        update(here->1, par->1, 1, mid, val);
    }
    else {
        here->1 = par->1;
       here->r = alloc();
        update(here->r, par->r, mid + 1, r, val);
    }
    here->cnt = here->l->cnt + here->r->cnt;
}
int query(node *node1, node *node2, int 1, int r, int k){
    if(l == r) return 1;
    int ccc = node1->l->cnt - node2->l->cnt;
    int mid = (1 + r) / 2;
```

```
if(k <= ccc) return query(node1->1, node2->1, 1, mid, k);
    else return query(node1->r, node2->r, mid + 1, r, k - ccc);
7.2 Link-Cut Tree
struct node{
   node *pp, *p, *l, *r;
   int val;
   node(){p = 0, 1 = 0, r = 0;}
   node(int val) : val(val) { p = 0, 1 = 0, r = 0;}
};
void push(node *x){}
void pull(node *x){}
void rotate(node *x){
   if(!x->p) return;
   push(x->p); // if there's lazy stuff
   push(x);
   node *p = x->p;
   bool is_left = (p->l == x);
   node *b = (is_left ? x->r : x->1);
   x->p = p->p;
   if (x-p \&\& x-p-1 == p) x-p-1 = x;
   if(x->p \&\& x->p->r == p) x->p->r = x;
   if(is_left){
        if(b) b \rightarrow p = p;
        p->1 = b;
        p->p = x;
        x->r = p;
   }
    else{
        if(b) b \rightarrow p = p;
        p->r = b;
        p->p = x;
        x->1 = p;
   pull(p); // if there's something to pull up
   pull(x);
   //if(!x->p) root = x; // IF YOU ARE SPLAY TREE
   if(p->pp){ // IF YOU ARE LINK CUT TREE
        x->pp = p->pp;
        p->pp = nullptr;
   }
void splav(node *x){
   while(x->p){
        node *p = x->p;
        node *g = p->p;
        if(g){
```

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```
if((p->1 == x) ^ (g->1 == p)) rotate(x);
            else rotate(p);
        }
        rotate(x);
    }
}
void access(node *x){
    splay(x);
    push(x);
    if(x->r){
        x->r->pp = x;
        x->r->p = nullptr;
        x->r = nullptr;
    }
   pull(x);
    while(x->pp){
        node *nxt = x->pp;
        splay(nxt);
        push(nxt);
        if(nxt->r){
           nxt->r->pp = nxt;
            nxt->r->p = nullptr;
            nxt->r = nullptr;
        }
        nxt->r = x;
        x->p = nxt;
        x->pp = nullptr;
        pull(nxt);
        splay(x);
    }
}
node *root(node *x){
    access(x);
    while(x->1){
        push(x);
        x = x->1;
    }
    access(x);
    return x;
node *par(node *x){
    access(x);
    if(!x->1) return nullptr;
   push(x);
   x = x->1;
    while(x->r){}
        push(x);
        x = x->r;
    }
    access(x);
```

```
return x;
node *lca(node *s, node *t){
    access(s);
    access(t):
    splay(s);
    if(s->pp == nullptr) return s;
    return s->pp;
void link(node *par, node *son){
    access(par);
    access(son);
    //son->rev ^= 1; // remove if needed
    push(son);
    son->1 = par;
    par->p = son;
    pull(son);
void cut(node *p){
    access(p);
    push(p);
    if(p->1){
        p->1->p = nullptr;
        p->1 = nullptr;
    }
    pull(p);
7.3 Dynamic Convex Hull
// https://github.com/niklasb/contest-algos/blob/master/convex_hull/dynamic.cpp
const ll is_query = -(1LL<<62);</pre>
struct Line {
   ll m. b:
    mutable function<const Line*()> succ;
    bool operator<(const Line& rhs) const {</pre>
        if (rhs.b != is_query) return m < rhs.m;</pre>
        const Line* s = succ();
        if (!s) return 0:
        return b - s \rightarrow b < (s \rightarrow m - m) * rhs.m;
    }
};
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum
   bool bad(iterator y) {
        auto z = next(y);
        if (y == begin()) {
            if (z == end()) return 0:
            return y->m == z->m && y->b <= z->b;
        auto x = prev(y);
        if (z == end()) return y > m == x - m && y - b <= x - b;
```

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```
return (x-b - y-b)*(z-m - y-m) >= (y-b - z-b)*(y-m - x-m);
    }
    void insert_line(Line 1) {
        auto y = insert(1);
        v->succ = [=] { return next(v) == end() ? 0 : &*next(v): }:
        if (bad(y)) { erase(y); return; }
        while (next(y) != end() && bad(next(y))) erase(next(y));
        while (y != begin() && bad(prev(y))) erase(prev(y));
    }
   11 \text{ eval}(11 \text{ x}) 
        auto 1 = *lower_bound((Line) { x, is_query });
        return 1.m * x + 1.b;
    }
};
     Stern-Brocot Tree
// __int128 is recommended
bool test(11 a, 11 b) { // for testing directions, vary by prob
    // return true if (true value) >= a/b
   11 n = 0, m = 1;
    rep(i, N) {
        if (n < m*A[i].fi) n = A[i].fi, m = 1;</pre>
       11 c = b*n+m*a. d = m*b:
       11 g = gcd(c, d);
       n = c/g;
        m = d/g;
        if (n > m*A[i].se) return false;
   }
    return true;
}
pair<11, 11> stern_brocot(11 M, 11 N) {
   // M : max value
    // N : max divisor
   // if result is a/b, return as {a, b}
   11 a = 0, b = 1; // 1
   11 c = 1, d = 0; // r
    int 1, r;
    bool chg = true;
    while(chg) {
        chg = false;
        // to left
       1 = 0, r = (N-d-1)/b+1;
        while(1 < r)  {
            int mid = (1+r+1)/2;
```

```
if (test(a*mid+c, b*mid+d)) r = mid-1;
            else l = mid:
        }
        c += a*1:
        d += b*1;
        chg |= (1 > 0);
        // to right
        1 = 0, r = (d?(N-b-1)/d+1:M);
        while(l < r) {</pre>
            int mid = (1+r+1)/2;
            if (test(a+mid*c, b+mid*d)) l = mid;
            else r = mid-1;
        }
        a += c*1:
        b += d*1;
        chg |= (1 > 0);
   }
   return {a, b};
7.5 Rope
#include <bits/stdc++.h>
#include <ext/rope>
using namespace std;
using namespace __gnu_cxx;
int main()
   ios::svnc with stdio(false):
   cin.tie(0);
    crope rp; // rope<char>
    string s("Lorem-ipsum");
   int n = s.length():
   rp.append(s.c_str()); // add element
   int x = 3, y = 8; // split and merge below
   rp = rp.substr(x, y-x) + rp.substr(0, x) + rp.substr(y, n);
   cout << rp.at(0) << '\n'; // get element, 'e'</pre>
    cout << rp << '\n'; // print, "em-ip|Lor|sum"</pre>
7.6 Bitset
#include <bitset>
#include <iostream>
```

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```
using namespace std;
int main() {
    bitset<8> b1(13);
                                     // 00001101
    bitset<8> b2("10111"):
                                     // 00010111
    cout << b1.count() << endl;</pre>
                                     // 3
    cout << b1.test(6) << endl:</pre>
                                     // 0, since 2^{\circ}6-th bit is 0
                                     // set to 1, 1-fill if no param
    b1.set(6);
                                     // set to 0, 0-fill if no param
    b2.reset(2);
    // use 'flip' for flipping
    cout << "b1:" << b1 << endl:
                                             // b1:01001101
    cout << "b2:" << b2 << endl;
                                             // b2:00010011
    // use any, none, all (c++11) for bit checking
    // supported operators : &, |, ^, <<, >>, ~, ==, !=
    // these operators must match size (given to template)
    cout << "&: " << (b1 & b2) << endl;
                                             // &: 00000001
    cout << "^: " << (b1 ^ b2) << endl;
                                             // &: 01011110
    cout << "|: " << (b1 | b2) << endl:
                                            // |: 01011111
    cout << "~: " << (~b1) << endl;
                                             // ~: 10110010
                                             // <<:01101000
    cout << "<<:" << (b1 << 3) << endl;
    cout << b2.to_ulong() << endl;</pre>
                                             // 19, c++11 supports ullong
}
```

7.7 Policy Based Data Structure

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
using namespace std:
using namespace __gnu_pbds;
typedef tree<
int.
null_type,
less<int>,
rb_tree_tag,
tree_order_statistics_node_update >
ordered_set;
// less<int> : not allow for duplicate
// less_equal<int> : allow for duplicate
// use upper_bound when you erase from set used less_equal
int N;
int main(void) {
```

```
iostream::sync_with_stdio(false);
cin.tie(nullptr);
ordered_set X;
X.insert(1);
X.insert(2);
X.insert(4):
X.insert(8);
X.insert(16):
cout<<*X.find_by_order(1)<<endl; // 2</pre>
cout<<*X.find_by_order(2)<<endl; // 4</pre>
cout<<*X.find_by_order(4)<<endl; // 16</pre>
cout<<(end(X)==X.find_by_order(6))<<end1; // true</pre>
cout<<X.order_of_key(-5)<<endl; // 0</pre>
cout<<X.order_of_key(1)<<endl; // 0</pre>
cout<<X.order_of_key(3)<<endl; // 2</pre>
cout<<X.order_of_key(4)<<endl; // 2</pre>
cout<<X.order_of_key(400)<<endl; // 5</pre>
```

8 Miscellaneous

8.1 Misc Formulae and Algorithms

8.1.1 Faulhaber's Formula

$$T(n,k) = \sum_{i=1}^{n} i^{k} = \frac{(n+1)^{k+1} - 1^{k+1} - \sum_{j=0}^{k-1} {k+1 \choose j} T(n,j)}{{k+1 \choose k}}$$

Also use

$$(x+1)^d - x^d = 1 + {d \choose 1}x + {d \choose 2}x^2 + \dots + {d \choose d-1}x^{d-1}$$

to get each coef.

8.1.2 Maximum Clique

```
typedef long long ll;
ll G[40]; // 0-index
int N, M;
int cur;
void get_clique(int R = 0, ll P = (1ll<<N)-1, ll X = 0){
    if((P|X) == 0){
        cur = max(cur, R);
        return;
    }
    int u = __builtin_ctzll(P|X);
    ll c = P&~G[u];</pre>
```

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```
while(c){
        int v = __builtin_ctzll(c);
        get_clique(R + 1, P&G[v], X&G[v]);
       P ^= 111 << v;
       X = 111 << v:
       c ^= 111 << v;
   }
}
8.1.3 De Brujin Sequence
// https://github.com/koosaga/DeobureoMinkyuParty/blob/master/teamnote.tex
// alphabet = [0, k - 1], substr length n, res starts with 0 (cyclic)
int res[10000000], aux[10000000]; // >= k^n, k*n
int de_bruijn(int k, int n, int lim) { // returns size (k^n)
    if(k == 1) {
       res[0] = 0;
       return 1;
    }
    for(int i = 0; i < k * n; i++) aux[i] = 0;
    function<void(int, int)> db = [&](int t, int p) {
        if(sz > lim) return;
        if(t > n) {
            if(n \% p == 0)
               for(int i = 1; i <= p; i++)
                    res[sz++] = aux[i];
       }
        else {
            aux[t] = aux[t - p];
            db(t + 1, p);
            for(int i = aux[t - p] + 1; i < k; i++) {
               aux[t] = i;
                db(t + 1, t);
            }
       }
    };
    db(1, 1);
    return sz;
}
```

8.2 Highly Composite Numbers, Large Prime

< 10^k	`k number divisors		2 3 5 71113171923293137
1	6	4	1 1
2	60	12	2 1 1
3	840	32	3 1 1 1
4	7560	64	3 3 1 1
5	83160	128	3 3 1 1 1
6	720720	240	4 2 1 1 1 1

```
8648640
                                     448 6 3 1 1 1 1
   8
                  73513440
                                     768 5 3 1 1 1 1 1
   9
                 735134400
                                    1344 6 3 2 1 1 1 1
   10
                6983776800
                                    2304 5 3 2 1 1 1 1 1
   11
               97772875200
                                    4032 6 3 2 2 1 1 1 1
   12
              963761198400
                                    6720 6 4 2 1 1 1 1 1 1
   13
             9316358251200
                                   10752 6 3 2 1 1 1 1 1 1 1
   14
            97821761637600
                                   17280 5 4 2 2 1 1 1 1 1 1
   15
           866421317361600
                                   26880 6 4 2 1 1 1 1 1 1 1 1
   16
          8086598962041600
                                   41472 8 3 2 2 1 1 1 1 1 1 1
                                   64512 6 3 2 2 1 1 1 1 1 1 1 1
        74801040398884800
   18 897612484786617600
                                  103680 8 4 2 2 1 1 1 1 1 1 1 1
           < 10<sup>k</sup> prime
                              > 10<sup>k</sup> prime
                                                   # of prime
                                        11
                                                           4
                    97
                                        101
                                                           25
   3
                    997
                                       1009
                                                          168
                   9973
                                      10007
                                                         1229
   5
                                                         9592
                  99991
                                     100003
   6
                 999983
                                   1000003
                                                        78498
                9999991
                                   10000019
                                                      664579
   8
               99999989
                                                      5761455
                                  100000007
   9
              99999937
                                                     50847534
                                 100000007
                                                > 10<sup>k</sup> prime
                    < 10 k prime
   10
                      999999967
                                                 10000000019
   11
                     9999999977
                                                100000000003
   12
                    99999999999
                                               1000000000039
   13
                   999999999971
                                              1000000000037
   14
                  999999999973
                                             10000000000031
   15
                 99999999999989
                                            100000000000037
   16
                99999999999937
                                           10000000000000061
   17
                                          100000000000000003
               999999999999997
   18
              999999999999999
                                         1000000000000000003
NTT Prime:
 469762049 = 7 \times 2^{26} + 1. Primitive root : 3.
 998244353 = 119 \times 2^{23} + 1. Primitive root: 3.
 985661441 = 235 \times 2^{22} + 1. Primitive root: 3.
 1012924417 = 483 \times 2^{21} + 1. Primitive root: 5.
Primes near 10^9: 10^9 + [7, 9, 21, 33, 87]
```

8.3 Fast Integer IO

```
// credit : https://github.com/koosaga/DeobureoMinkyuParty/blob/master/teamnote.tex
static char buf[1 << 19]; // size : any number geq than 1024
static int idx = 0;
static int bytes = 0;
static inline int _read() {</pre>
```

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```
if (!bytes || idx == bytes) {
       bytes = (int)fread(buf, sizeof(buf[0]), sizeof(buf), stdin);
       idx = 0:
    }
    return buf[idx++]:
}
static inline int _readInt() {
    int x = 0, s = 1;
    int c = _read();
    while (c \le 32) c = read();
    if (c == '-') s = -1, c = _read();
    while (c > 32) x = 10 * x + (c - '0'), c = _read();
    if (s < 0) x = -x;
    return x;
}
8.4 C++ Tips / Environments
#include <bits/stdc++.h> // magic header
using namespace std;
                         // magic namespace
struct StupidGCCCantEvenCompileThisSimpleCode{
    pair<int, int> array[1000000];
}; // https://gcc.gnu.org/bugzilla/show_bug.cgi?id=68203
// how to use rand (in 2017)
mt19937 rng(0xdeadbeef);
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
int randint(int lb, int ub){ return uniform_int_distribution<int>(lb, ub)(rng); }
shuffle(permutation.begin(), permutation.end(), rng);
mt19937_64 _R(chrono::steady_clock::now().time_since_epoch().count()); // _R()
// comparator overload
auto cmp = [](seg a, seg b){return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
bool operator==(const point &p)const{ return x == p.x && y == p.y; }
};
struct hasher {
size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
};
unordered_map<point, int, hasher> hsh;
// c++ setprecision example
#include <iostream>
                       // std::cout, std::fixed
#include <iomanip>
                       // std::setprecision
```

```
int main () {
    double f = 3.14159;
    std::cout << std::setprecision(5) << f << '\n'; // 3.1416
    std::cout << std::setprecision(9) << f << '\n'; // 3.14159
    std::cout << std::fixed;
    std::cout << std::setprecision(5) << f << '\n'; // 3.14159
    std::cout << std::setprecision(9) << f << '\n'; // 3.141590000
}</pre>
```