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# Team Note of Powered by Zigui

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### ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

## 0 Quotes and Prerequisites

```
evenharder : Mental Abuse To Humans
djkim0613 : 열심히 응원하겠습니다.
SoulTch : How much is this bus ticket?
* This template is brought from that of 'Deobureo Minkyu Party'
```

### Run script

```
#!/bin/bash
g++ -fsanitize=undefined -std=c++14 -02 -o /tmp/pow $1.cpp
time /tmp/pow < $1.in
# export PATH=~:$PATH</pre>
```

### Debug Code

```
#define setz(x) memset(x, 0, sizeof(x))
#define sz(x) ((int)(x).size())
#define rep(i, e) for (int i = 0, _##i = (e); i < _##i; i++)
#define repp(i, s, e) for (int i = (s), _##i = (e); i < _##i; i++)
#define repr(i, s, e) for (int i = (s)-1, _##i = (e); i \ge _{\#}i; i--)
#define repi(i, x) for (auto &i : (x))
// using namespace std;
using ll = long long;
using pii = pair<int, int>;
using pll = pair<ll, ll>;
template<typename T>
ostream &operator<<(ostream &os, const vector<T>& v) {
    cout << "[":
   for (auto p : v) cout << p << ",";
    cout << "]":
    return os;
}
#ifndef __SOULTCH
#define debug(...) 0
#define endl '\n'
#define debug(...) cout << " [-] ", _dbg(#__VA_ARGS__, __VA_ARGS__)</pre>
template<class TH> void _dbg(const char *sdbg, TH h){ cout << sdbg << '=' << h <<
endl; }
template<class TH, class... TA> void _dbg(const char *sdbg, TH h, TA... a) {
    while(*sdbg != ',') cout << *sdbg++;</pre>
    cout << '=' << (h) << ',';
    _dbg(sdbg+1, a...);
}
#endif
```

#### Reminders

Pre-submit	Wrong answer:						
예제 작성해보기 (최소, 최대) 메모리, overflow 분석하기 올바른 문제에 제출하기	코드 + debug output 출력 다중 테케 문제에서 초기화 확인하기 알고리즘이 제한조건을 전부 다루는지 확인하기 지문 다시 읽어보기 corner case 찾아보기 초기화 안 된 지역변수 찾아보기 차기화 안 된 지역변수 찾아보기 자, M, i, j 등 변수 확인하기 풀이 증명하기 STL 함수 다시 생각해보기 팀노트에서 그대로 가져온 변수값 다시 확인하기 이 목록 다시 읽어보기 알고리즘 팀원에게 설명하기 팀원이랑 코드 보기 잠깐 일어나서 생각 재정비하고 오기						
Runtime error:	Time limit exceeded: / Memory limit exceeded:						
코너 케이스 처리해보기	무한 루프 확인하기						
초기화 안 된 변수 찾기	알고리즘 시간 복잡도 확인하기						
out-of-range 확인하기	data copy 어느 정도 하는지 확인하기 (reference)						
팀노트에서 그대로 가져온 변수값 다시 확인하기	입출력 규모 생각하기 (scanf 고려해보기)						
assertion 넣어보기	vector, map 최소화하기						
무한 재귀 확인하기	팀원에게 알고리즘 물어보기						
null pointer 확인하기	최대 메모리 사용량 계산하기						
메모리 사용량 확인하기	다중 테케 문제에서 초기화하기						

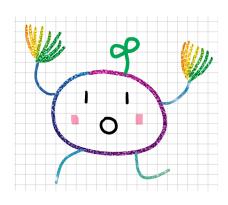




Figure 1: 풀다가 막힐 때는 이 그림을 봅시다. 아자아자 화이팅!

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#### Math 1

#### **Basic Mathematics**

#### 1.1.1 Trigonometry

- $\sin(\alpha \pm \beta) = \sin \alpha \cos \beta \pm \cos \alpha \sin \beta$
- $\cos(\alpha \pm \beta) = \cos \alpha \cos \beta \mp \sin \alpha \sin \beta$
- $\tan(\alpha \pm \beta) = \frac{\tan \alpha \pm \tan \beta}{1 \mp \tan \alpha \tan \beta}$
- $\sin 2\theta = 2\sin\theta\cos\theta$
- $c^2 = a^2 + b^2 2ab\cos\gamma$

#### 1.1.2 Generating Function

- $\sum_{n} (pn+q)x^{n} = \frac{p}{1-x} + \frac{q}{(1-x)^{2}}$  (Arithmetic progression)
- $\sum_{m} (rx)^{n} = (1 rx)^{-1}$  (Geometric progression)
- $\sum_{n} {m \choose n} x^n = (1+x)^m$  (Binomial coefficient)
- $\sum_{n} {m+n-1 \choose n} x^n = (1-x)^{-m}$  (Multiset coefficient)

#### 1.1.3 Calculus

- $\int_a^b f(x) dx = \frac{b-a}{6} \left[ f(a) + 4f(\frac{a+b}{2}) + f(b) \right]$  (Simpson's Rule, for cubic poly)
- $\int u'v \ dx = uv \int uv' \ dx$  (Integration by parts)

#### 1.1.4 Combinatorics

- $\binom{n}{k}$  (Stirling numbers of the first kind / n elem, k cycle)
- $\binom{n+1}{k} = n \binom{n}{k} + \binom{n}{k-1}, \binom{n}{0} = \binom{0}{n} = \binom{0}{0} = 1$
- $\binom{n}{k}$  (Stirling numbers of the second kind / n elem, k unlabeled subset)
- ${n+1 \brace k} = k \begin{Bmatrix} n \cr k \end{Bmatrix} + {n \brace k-1}, {n \brace 0} = {0 \brack n} = {0 \brack 0} = 1$

 $B_n$  (Bell number, n elem, partition)

- $B_{n+1} = \sum_{k=0}^{n} \binom{n}{k} B_k, B_{n+1} = \sum_{k=0}^{n} \binom{n}{k}$
- $\bullet \ B_{n^m+n} \equiv mB_n + B_{n+1}$

### 1.2 Number Theory

vector<ll> get\_factor(ll n) {

for(11 i=2;i\*i<=n;i++) {  $if(n \% i == 0) {$ v.push\_back(i);

if(n > 1) v.push\_back(n);

for(ll x=2;x<=n;x++) { int yes = 1;

for(ll v : fact) {

if(yes) return x;

ll sq = ceil(sqrt(n) + 1e-9);

sort(prec.begin(), prec.end());

prec[i] = {mexp(g, i, n), i};

// find k s.t. g^k = b mod n ll bsgs(ll g, ll b, ll n) {

vector<pl> prec(sq);

for(ll i=0;i<sq;i++) {</pre>

for(int i=0;i<sq;i++) {</pre>

11 phi = n-1; // assume n is prime vector<ll> fact = get\_factor(phi);

vector<ll> v;

}

}

return -1;

11 cur = b:

return v;

ll get\_primitive(ll n) {

#### 1.2.1 Lattice Points under Line

```
11 calc(ll a,ll b,ll c,ll n){
                                     if(!n)return 0;
                                     11 tmp=a/c*n*(n-1)/2;
                                       tmp+=b/c*n;
                                     return tmp+calc(c,(a*n+b)%c,a%c,((a%c)*n+b%c)/c);
 \frac{1}{y} = \frac{1}
1.2.2 Shanks' Baby-step Giant-step
 11 mexp(ll x, ll y, ll p) {
                                       if(!y) return 1;
                                       if (y \& 1) return x * mexp(x*x%p, y>>1, p) % p;
                                       return mexp(x*x\%p, y>>1, p);
```

while(n % i == 0) n /= i:

yes &= (mexp(x, phi / y, n) != 1);

ll base = mexp(g, n-1-sq, n); //  $g^(-m)$ , assume n : prime

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```
auto it = lower_bound(prec.begin(), prec.end(), pl(cur, 0));
        if(it->first == cur) {
            return (i*sq + it->second) % (n-1);
        cur = cur * base % n:
    }
    return -1;
}
1.2.3 Extended Euclidean Algorithm
// ax + by = gcd(a,b). x, y?
pll ext_gcd(ll a,ll b) {
    if(b) {
        auto tmp = ext_gcd(b, a%b);
        return {tmp.second, tmp.first - (a/b) * tmp.second};
    }
    else return {1, 0};
// ax = gcd(a, m) mod m. x?
11 mod inv(ll a, ll m) {
    return (ext_gcd(a, m).first + m) % m;
}
1.2.4 Chinese Remainder Theorem
ll pos rem(ll a, ll m) { // m > 0, a % m?
   ll res = abs(a) % m;
    return a > 0 ? res : (res ? m - res : 0);
// ax = c mod m, bx = d mod n. x?
ll solve(ll a, ll c, ll m, ll b, ll d, ll n) {
    a = pos_rem(a, m); c = pos_rem(c, m); // if a, c not in [0, m)
    b = pos_rem(b, n); d = pos_rem(d, n); // if b, d not in [0, n)
    11 g = gcd(a, gcd(c, m)); a \neq g, c \neq g, m \neq g;
        g = gcd(b, gcd(d, n)); b /= g, d /= g, n /= g;
    if(c % _gcd(a, m) || d % _gcd(b, n)) return inf;
    ll t1 = (mod_inv(a, m) * c) % m;
    11 t2 = (mod_inv(b, n) * d) \% n;
    g = gcd(m, n):
    11 lc = m * n / g;
    if(abs(t1 - t2) % g) return inf;
    pl p = ext_gcd(m, n);
    11 q = (t1 * p.second * n/g + t2 * p.first * m/g);
    return pos_rem(q, lc);
}
1.2.5 Möbius Inversion Formula
 \forall n \in \mathbb{N} \ g(n) = \sum_{d \mid n} f(d) \implies f(n) = \sum_{d \mid n} \mu(d)g(n/d)
```

#### 1.3 FFT

```
FFT: (a_0, a_1, \dots, a_{n-1}) \mapsto (\sum_{j=0}^{n-1} a_j(\omega^0)^j, \sum_{j=0}^{n-1} a_j(\omega^1)^j, \dots, \sum_{j=0}^{n-1} a_j(\omega^{n-1})^j)
using base = complex<double>;
void fft(vector<base>& a, bool inv) {
   int n = a.size(), j = 0;
   vector<base> roots(n/2);
   for(int i=1:i<n:i++) {
        int bit = (n >> 1);
        while(j >= bit) {
            j -= bit;
            bit >>= 1;
        j += bit;
        if(i < j) swap(a[i], a[j]);</pre>
   }
    double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
    for(int i=0;i<n/2;i++)
        roots[i] = base(cos(ang * i), sin(ang * i));
   /* In NTT, let prr = primitive root. Then,
    int prr = 3, ang = mexp(prr, (mod - 1) / n);
   if(inv) ang = mexp(ang, mod - 2);
   for(int i=0: i<n/2: i++){
        roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
    also, make sure to apply modulus under here
    and pre-compute mexp(n, mod-2)
   for(int i=2;i<=n;i<<=1) {
        int step = n / i;
        for(int i=0:i<n:i+=i) {</pre>
            for(int k=0; k<i/2; k++) {
                 base u = a[i+k], v = a[i+k+i/2] * roots[step * k]:
                 a[j+k] = u+v;
                 a[j+k+i/2] = u-v;
        }
    if(inv) for(int i=0;i<n;i++) a[i] /= n;
void conv(vector<base>& x, vector<base>& y) {
    int n = 2; while(n < x.size()+y.size()) n <<= 1;</pre>
   x.resize(n): v.resize(n):
   fft(x, false); fft(y, false);
   for(int i=0;i<n;i++) x[i] *= y[i];</pre>
   fft(x, true); // access (ll)round(x[i].real())
```

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#### 1.4 Miller-Rabin + Pollard-Rho

```
//Prove By Solving - https://www.acmicpc.net/problem/4149
namespace miller rabin{
    lint mul(lint x, lint y, lint mod){ return (__int128) x * y % mod; }
    lint ipow(lint x, lint y, lint p){
       lint ret = 1, piv = x \% p;
        while(y){
            if(y&1) ret = mul(ret, piv, p);
           piv = mul(piv, piv, p);
           y >>= 1;
        return ret;
    }
    bool miller_rabin(lint x, lint a){
        if(x \% a == 0) return 0;
        lint d = x - 1:
        while(1){
            lint tmp = ipow(a, d, x);
            if(d&1) return (tmp != 1 && tmp != x-1);
            else if(tmp == x-1) return 0;
            d >>= 1;
        }
    }
    bool isprime(lint x){
        for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
            if(x == i) return 1;
            if(x > 40 && miller_rabin(x, i)) return 0;
        if(x \le 40) return 0;
        return 1;
    }
}
namespace pollard_rho{
    lint f(lint x. lint n. lint c){
        return (c + miller_rabin::mul(x, x, n)) % n;
    void rec(lint n. vector<lint> &v){
        if(n == 1) return;
        if(n \% 2 == 0){
            v.push_back(2);
            rec(n/2, v);
            return;
       }
        if(miller_rabin::isprime(n)){
            v.push back(n):
            return;
        }
        lint a, b, c;
        while(1){
```

```
a = rand() \% (n-2) + 2;
            b = a:
            c = rand() \% 20 + 1:
                a = f(a, n, c):
                b = f(f(b, n, c), n, c);
            \frac{1}{2} while (\gcd(abs(a-b), n) == 1);
            if(a != b) break:
        lint x = gcd(abs(a-b), n);
        rec(x, v);
        rec(n/x, v);
   vector<lint> factorize(lint n){
        vector<lint> ret;
        rec(n. ret):
        sort(ret.begin(), ret.end());
        return ret;
   }
};
1.5 Black Box Linear Algebra + Kitamasa
vector<int> berlekamp_massey(vector<int> x){
    vector<int> ls, cur;
    int lf. ld:
   for(int i=0; i<x.size(); i++){</pre>
        lint t = 0:
        for(int j=0; j<cur.size(); j++){</pre>
            t = (t + 111 * x[i-j-1] * cur[j]) % mod;
        if((t - x[i]) \% mod == 0) continue;
        if(cur.empty()){
            cur.resize(i+1):
            lf = i;
            1d = (t - x[i]) \% mod:
            continue;
        lint k = -(x[i] - t) * ipow(ld, mod - 2) % mod;
        vector<int> c(i-lf-1);
        c.push_back(k);
        for(auto &j : ls) c.push_back(-j * k % mod);
        if(c.size() < cur.size()) c.resize(cur.size());</pre>
        for(int j=0; j<cur.size(); j++){</pre>
            c[j] = (c[j] + cur[j]) % mod;
        if(i-lf+(int)ls.size()>=(int)cur.size()){
            tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) \% mod);
        cur = c;
```

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```
for(auto &i : cur) i = (i % mod + mod) % mod;
    return cur:
}
int get_nth(vector<int> rec, vector<int> dp, lint n){
    int m = rec.size();
    vector<int> s(m), t(m);
    s[0] = 1;
    if(m != 1) t[1] = 1:
    else t[0] = rec[0];
    auto mul = [&rec](vector<int> v. vector<int> w){
        int m = v.size();
        vector\langle int \rangle t(2 * m);
        for(int j=0; j<m; j++){
            for(int k=0; k<m; k++){</pre>
                t[j+k] += 111 * v[j] * w[k] % mod;
                if(t[i+k] >= mod) t[i+k] -= mod:
            }
        }
        for(int j=2*m-1; j>=m; j--){
            for(int k=1; k<=m; k++){</pre>
                t[j-k] += 111 * t[j] * rec[k-1] % mod;
                if(t[j-k] >= mod) t[j-k] -= mod;
            }
        }
        t.resize(m);
        return t:
    }:
    while(n){
        if(n \& 1) s = mul(s, t):
        t = mul(t, t);
        n >>= 1;
    }
    lint ret = 0;
    for(int i=0: i<m: i++) ret += 111 * s[i] * dp[i] % mod:
    return ret % mod:
int guess_nth_term(vector<int> x, lint n){ // init with > 3k, 0(1^2 lg n)
    if(n < x.size()) return x[n];</pre>
    vector<int> v = berlekamp_massey(x);
    if(v.emptv()) return 0:
    return get_nth(v, x, n);
}
struct elem{int x, y, v;}; // A_(x, y) <- v, O-based. no duplicate please...
vector<int> get_min_poly(int n, vector<elem> M){
    // smallest poly P such that A^i = sum_{j} < i \ A^j \times P_{j}
    vector<int> rnd1, rnd2;
    mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
    };
```

```
for(int i=0; i<n; i++){</pre>
        rnd1.push_back(randint(1, mod - 1));
        rnd2.push_back(randint(1, mod - 1));
   vector<int> gobs:
   for(int i=0; i<2*n+2; i++){
        int tmp = 0;
        for(int j=0; j<n; j++){
            tmp += 111 * rnd2[j] * rnd1[j] % mod;
            if(tmp >= mod) tmp -= mod:
        gobs.push_back(tmp);
        vector<int> nxt(n):
        for(auto &i : M){ // sparse matrix * vector
            nxt[i.x] += 111 * i.v * rnd1[i.v] % mod;
            if(nxt[i.x] >= mod) nxt[i.x] -= mod;
        rnd1 = nxt;
   }
    auto sol = berlekamp_massey(gobs);
   reverse(sol.begin(), sol.end());
    return sol;
lint det(int n, vector<elem> M){
   vector<int> rnd;
   mt19937 rng(0x14004):
   auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
   }:
   for(int i=0; i<n; i++) rnd.push_back(randint(1, mod - 1));</pre>
    for(auto &i : M){
        i.v = 111 * i.v * rnd[i.y] % mod;
   auto sol = get_min_poly(n, M)[0];
   if (n \% 2 == 0) sol = mod - sol:
   for(auto &i : rnd) sol = 111 * sol * ipow(i, mod - 2) % mod;
   return sol:
2 Geometry
2.1 struct Point
const double eps = 1e-10;
template <class T>
struct point{
    typedef point P;
   T x, y;
    point(T x=0, T y=0) : x(x), y(y) {}
```

bool operator< (P a) const {return fabs(x-a.x) > eps ? x<a.x : y<a.y;}

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```
bool operator == (P a) const {return max(fabs(x-a.x), fabs(y-a.y)) < eps;}
                                                                                            if(sgn(oa)*sgn(ob) < 0 \&\& sgn(oc)*sgn(od) < 0)
    P operator+ (P a) const {return P(x+a.x, y+a.y);}
                                                                                                return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
    P operator- (P a) const {return P(x-a.x, v-a.v);}
                                                                                            set<P> S:
    P operator- () const {return P(-x, -y);};
                                                                                            if(onSegment(c, d, a)) S.insert(a);
    T operator* (P a) const {return x*a.x + v*a.v:} // inner prod
                                                                                            if(onSegment(c, d, b)) S.insert(b);
    T operator/ (P a) const {return x*a.y - y*a.x;} // outer prod
                                                                                            if(onSegment(a, b, c)) S.insert(c);
    T dist2() const {return x*x + y*y;}
                                                                                            if(onSegment(a, b, d)) S.insert(d);
    double dist() const {return sqrt(double(dist2()));}
                                                                                            return vector<P>(S.begin(), S.end());
    P perp() const {return P(-v, x);}; // rotate 90 deg ccw
    P mult(T t) const {return P(x*t, y*t);}
    P unit() const {return P(x/dist(), y/dist());}
                                                                                        2.2.5 Circle-Line Intersection
    P rotate(double a){
        return P(x*cos(a)-y*sin(a), x*sin(a)+y*cos(a));
                                                                                        vector<P> circLine(P A, P B, P O, double r){
    }
                                                                                            vector<P> v;
};
                                                                                            P X = O-A, D = B-A;
                                                                                            double rat = 1.0 * (X*D) / (D*D):
int sgn(double x) {return (x > eps) - (x < -eps);}</pre>
                                                                                            double det = (X*D)*(X*D) - (D*D)*(X*X - r*r):
typedef point<double> P;
                                                                                            if(det < 0) return {};</pre>
                                                                                            else if(det < eps) return {P(A + D.mult(rat))};</pre>
                                                                                            det = sqrt(det);
2.2 Distance, Intersection
                                                                                            return {P(A + D.mult(rat + det/(D*D))),
2.2.1 Point-to-Line
                                                                                                    P(A + D.mult(rat - det/(D*D)));
double lineDist(P a, P b, P p) {
    return ((b-a)/(p-a))/(b-a).dist(); // a->b : left (+), right : (-);
                                                                                        2.2.6 Circle-Line Tangent
}
                                                                                        vector<P> circTangent(P A, P O, double r){
2.2.2 Point-to-Segment
                                                                                            if((0-A).dist2() < r*r + eps) return {};
                                                                                            double th = asin(r/(0-A).dist());
double segDist(P s, P e, P p) {
                                                                                            return {A + (0-A).rotate(th). mult(cos(th)),
    if(s == e) return (p-s).dist(); // mind the eps
                                                                                                    A + (O-A).rotate(-th).mult(cos(th))};
    double d = (e-s).dist2(), t = min(d, max(.0, (p-s)*(e-s)));
    return ((p-s).mult(d)-(e-s).mult(t)).dist() / d;
}
                                                                                        2.3 Convex Hull
2.2.3 Line intersection
                                                                                        vector<pll> get CV(vector<pll> V){
                                                                                            sort(V.begin(), V.end());
template<class P>
                                                                                            sort(V.begin() + 1, V.end(), [&](pll x, pll y){
pair<int, P> lineInter(P a, P b, P c, P d){
                                                                                                pll xx = x - V[0]:
    if((b-a)/(d-c) == 0) // parallel, mind the eps
                                                                                                pll yy = y - V[0];
        return \{-((b-a)/(c-a) == 0), a\}:
                                                                                                11 \text{ res} = xx / yy;
    double oa = (d-c)/(a-c), ob = (d-c)/(b-c);
                                                                                                if(res != 0) return res > 0;
    return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
                                                                                                if(xx.first != yy.first) return xx.first < yy.first;</pre>
} // 1.0.-1(inf) : inter
                                                                                                return xx.second < yy.second;</pre>
                                                                                            });
2.2.4 Segment Intersection
bool onSegment(P s, P e, P p) {return segDist(s, e, p) < eps;}</pre>
                                                                                            vector<pll> ret:
                                                                                            for(auto val : V){
                                                                                                while(ret.size() > 1){
template<class P> vector<P> segInter(P a, P b, P c, P d){
    double oa = (d-c)/(a-c), ob = (d-c)/(b-c),
                                                                                                    pll xx = ret[ret.size() - 2] - val;
            oc = (b-a)/(c-a), od = (b-a)/(d-a);
                                                                                                    pll vy = ret[ret.size() - 1] - val;
```

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```
if(xx / yy <= 0) ret.pop_back();</pre>
            else break:
        ret.push back(val):
    }
    return ret:
}
      Sorting Points by Angle
// credit : http://koosaga.com/97
auto angle_sort = [&] (P &a, P &b){
    if ((a < P(0, 0)) \hat{b} < P(0, 0)) return b < a;
    if(a / b != 0) return a / b > 0;
    return a.dist2() < b.dist2(); // norm</pre>
}; // counter-clockwise sort
2.5 Rotating Calipers
void rotating_calipers(vector<pll> CV){
    int pos = 0:
    for(int i = 0 ; i < CV.size() ; i++) if(CV[pos] < CV[i]) pos = i;</pre>
    int ind1 = 0, ind2 = pos;
    11 dist = (CV[ind1] - CV[ind2]) * (CV[ind1] - CV[ind2]);
    auto get_v = [\&](int x) { return CV[(x + 1) \% CV.size()] - <math>CV[x];};
    for(int i = 0 ; i < CV.size() ; i++){</pre>
        pll v = get_v(i);
        while((-v) / get_v(pos) < 0) pos = (pos + 1) % CV.size();
        11 tmp_dist = (CV[pos] - CV[i]) * (CV[pos] - CV[i]);
        if(dist < tmp_dist) {</pre>
            dist = tmp_dist;
            ind1 = i; ind2 = pos;
       }
    }
    printf("%1ld %1ld %1ld %1ld\n", CV[ind1].first, CV[ind1].second, CV[ind2].first,
    CV[ind2].second);
}
2.6 Smallest Enclosing Circle
//Prove By Solving - https://www.acmicpc.net/problem/11930
int main(){
    scanf("%d", &N);
    for(int i = 1; i \le N; i++) scanf("%lf%lf", &A[i].x, &A[i].y, &A[i].z);
    int t = 70000;
```

```
double rate = 1.0;
    point cur = (point)\{0, 0, 0\};
   for(int i = 1; i <= t; i++){
        int ind = 1;
        for(int i = 1 : i \le N : i++)
        if( (A[j] - cur) * (A[j] - cur) >
            (A[ind] - cur) * (A[ind] - cur)) ind = j;
        cur = cur + (A[ind] - cur).mult(rate);
        rate *= 0.99;
   }
   double r = 0;
   for(int i = 1; i \le N; i \leftrightarrow r = max(r, (A[i] - cur) * (A[i] - cur));
   cout << setprecision(10) << fixed << sqrt(r);</pre>
} // Non-deterministic, deterministic O(n lg n) requires Voronoi diagram
2.7 Circumcircle
double cc radius(P& A. P& B. P& C){
    return (B-A).dist() * (C-B).dist() * (A-C).dist() / fabs((A-B) / (B-C)) / 2:
P cc_center(P& A, P& B, P& C){
   P b = C-A, c = B-A;
   return A + (b.mult(c.dist2()) - c.mult(b.dist2())).perp().mult(0.5/(b/c));
2.8 Polygon Area
double ans = 0; // ans : double area
for(int i=0;i<points.size();i++)</pre>
    ans += points[i] / points[(i+1 == points.size() ? 0 : i+1)];
```

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### 3 Strings

### 3.1 Aho-Corasick Algorithm

```
// required macro : rep, repi, setz
namespace aho_corasick {
    const int MAXN = 200002, MAXC = 26;
    int go[MAXN+1][MAXC];
    int fail[MAXN+1];
    bool term[MAXN+1];
    void build(const vector<string> &v) {
        setz(go), setz(fail), setz(term);
        int cnode = 1;
        repi(s, v) {
            int p = 0;
                                // root is 0
            repi(j, s) {
                char c = j-'a'; // check starting alphabet
                if (!go[p][c]) go[p][c] = cnode++;
                p = go[p][c];
            }
            term[p] = true;
        }
        queue<int> q; rep(i, MAXC) if (go[0][i]) q.push(go[0][i]);
        while(!q.empty()) {
            int t = q.front(); q.pop();
            rep(i, MAXC) {
                if (go[t][i]) {
                    int p = fail[t];
                    while(p and not go[p][i]) p = fail[p];
                    p = go[p][i];
                    fail[go[t][i]] = p;
                    if (term[p]) term[go[t][i]] = true;
                    q.push(go[t][i]);
               }
           }
       }
    }
    int step(int p, int c) {
        if(go[p][c]) return go[p][c];
        if(p) return go[p][c] = step(fail[p], c); // if not root
        return go[p][c] = -1; // -1 is same as 0, but is visited
    }
    bool query(string &t) {
        int p = 0;
        repi(i, t) {
            char c = i-'a'; // check starting alphabet
```

```
p = max(0, step(p, c));
            /* if you need the original trie, use this or make length array
            while(p and not go[p][c]) p = fail[p];
            p = max(0, go[p][c]);
            */
            if (term[p]) return true;
       return false;
}
3.2 Suffix Array
// required macro : rep, repr, repp
// str : abracadabra
// SA : 10 7 0 3 5 8 1 4 6 9 2
                                    (0-based)
// LCP : 1 4 1 1 0 3 0 0 0 2
                                    (lcp[i] : lcp of sa[i], sa[i+1])
vector<int> make_sa(const string& s) {
    int n = s.length();
    int \lim = \max(128, n+1);
   vector<int> sa(n), g(n+1), ng(n+1), cnt(lim), ind(lim+1);
   rep(i, n) sa[i] = i, g[i] = s[i];
    g[n] = 0;
   for(int t=1;t<s.length();t<<=1)</pre>
        auto cmp = [&] (int a, int b) {
            return g[a] != g[b] ? g[a] < g[b] : g[a+t] < g[b+t];
        };
       rep (i, n)
                        cnt[g[min(i+t, n)]]++;
        repp(i, 1, lim) cnt[i] += cnt[i-1];
        repr(i, n, 0) ind[--cnt[g[min(i+t, n)]]] = i;
        rep (i, lim) cnt[i] = 0;
        rep (i, n)
                        cnt[g[i]]++; // same as cnt[g[ind[i]]]++
       repp(i, 1, lim) cnt[i] += cnt[i-1]:
        repr(i, n, 0) sa[--cnt[g[ind[i]]] = ind[i];
        ng[sa[0]] = 1:
        repp(i, 1, n) ng[sa[i]] = ng[sa[i-1]] + cmp(sa[i-1], sa[i]);
        g = ng;
       fill(cnt.begin(), cnt.end(), 0);
        fill(ind.begin(), ind.end(), 0);
   }
   return sa:
vector<int> make_lcp(const string& s, const vector<int>& sa) {
    int n = s.length(), len = 0;
   vector<int> lcp(n-1), rank(n);
   rep(i, n) rank[sa[i]] = i;
   rep(i, n) {
       if(rank[i]) {
            int j = sa[rank[i]-1];
```

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```
int lc = n - max(i,j);
            while(len < lc && s[i+len] == s[j+len]) len++;</pre>
           lcp[rank[i]-1] = len;
       if(len) len--:
    }
    return lcp;
}
     Lexicographically Smallest String Rotation
int min_rotation(string s) {
    int a=0, N=s.size(); s += s;
    rep(b,N) rep(i,N) {
       if (a+i == b \mid \mid s[a+i] < s[b+i]) \{b += max(0, i-1); break;\}
       if (s[a+i] > s[b+i]) \{ a = b; break; \}
    }
    return a; // rotate(v.begin(), v.begin()+min_rotation(v), v.end());
}
     Manacher's Algorithm
// 0-based
// s = # h # e # l # l # o #
// ret = 0 1 0 1 0 1 2 1 0 1 0
vector<int> manacher(const string& s) {
    int n = s.size(), r = -1, k = -1;
    vector<int> p(n);
   for (int i=0; i<n; i++) {
       if (i<=r) p[i] = min(r-i, p[2*k-i]);</pre>
       while (i-p[i]-1>=0 and i+p[i]+1<n and s[i-p[i]-1] == s[i+p[i]+1]) p[i]++;
       if (r < i+p[i]) r = i+p[i], k = i;
    }
    return p;
}
3.5 Z Algorithm
// 0-based
//s = abcababca
// ret = 9 0 0 2 0 4 0 0 1
vector<int> z_algo(const string &s) {
    int 1 = 0, r = 0, N = sz(s);
    vector<int> Z(N);
   Z[0] = N:
    repp(i, 1, N) {
       if (i > r) {
           l = r = i:
            while(r < N and s[r] == s[r-1]) r++;
```

```
r--;
    Z[i] = r-l+1;
} else {
    int k = i-l;
    if (Z[k] < r-i+1) Z[i] = Z[k];
    else {
        1 = i;
        while(r < N and s[r] == s[r-l]) r++;
        r--;
        Z[i] = r-l+1;
    }
}
return Z;
}</pre>
```

## 4 Graph Theory

### 4.1 Biconnected Component

```
// https://gist.github.com/koosaga/6f6fd50dd7067901f1b1
void dfs(int x, int p){
   dfn[x] = low[x] = ++piv;
   par[x] = p;
   for(int i=0; i<graph[x].size(); i++){</pre>
        int w = graph[x][i];
        if(w == p) continue;
        if(!dfn[w]){
            dfs(w, x);
            low[x] = min(low[x], low[w]);
        else low[x] = min(low[x], dfn[w]);
   }
}
void color(int x, int c){
    if(c > 0) bcc[x].push_back(c); // c == 0 : first component
   vis[x] = 1:
   for(int i=0; i<graph[x].size(); i++){</pre>
        int w = graph[x][i];
        if(vis[w]) continue;
        if(dfn[x] <= low[w]){</pre>
            bcc[x].push_back(++cpiv);
            color(w, cpiv);
        else color(w, c);
   }
```

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#### 4.2 Strongly Connected Component

```
const int MAXN = 2e5 + 10; // > 2*N
int N. M:
int dfsn[MAXN], low[MAXN], finished[MAXN], cnt;
vector<int> ADJ[MAXN];
vector<vector<int>> G;
stack<int> S;
int f(int x){ // f(1) f(-1) f(2) f(-2) f(3) f(-3) ... -> 0 1 2 3 4 5 ...
    return 2 * (abs(x) - 1) + (x < 0);
}
void add_edge(int x, int y){ // call by f(x), f(y)
    ADJ[x ^ 1].push_back(y);
    ADJ[y ^ 1].push_back(x);
}
// memset(finished, -1, sizeof(finished));
int scc(int here){
    static vector<int> tmp;
    S.push(here);
    dfsn[here] = low[here] = ++cnt;
    int &ret = low[here]:
    for(int there : ADJ[here]){
        if(dfsn[there] == 0) ret = min(ret, scc(there));
        else if(finished[there] == -1) ret = min(ret, dfsn[there]);
    }
    if(dfsn[here] == low[here]){
        while(1){
            int x = S.top(); S.pop();
            finished[x] = G.size();
            tmp.push_back(x);
            if(x == here) break;
        G.push_back(tmp);
        tmp.clear();
    }
    return ret;
}
4.2.1 2-SAT
```

scc를 실행시켜 f(i) 와 f(-i)가 같은 component에 있다면, 모순. f(i) 와 f(-i) 중 finished 배열의 수가 작은 것이 참이다. (SCC numbering의 역순이 위상정렬이기에,  $F \to T$ 를 유지하기 위함)

#### 4.3 Euler Tour

```
struct Edge{
   int to, cnt; // to: 인접한 정점, cnt: 남은 사용 횟수
   Edge *dual; // dual: 역방향 간선을 가리키는 포인터
```

```
Edge(): Edge(-1, 0){}
    Edge(int to1, int cnt1): to(to1), cnt(cnt1), dual(nullptr){}
void Eulerian(int curr){
    for(Edge *e: adi[curr]){
        if(e->cnt > 0){
            e->cnt--;
            e->dual->cnt--;
            Eulerian(e->to); // dfs
    }
    cout << curr << '\n';
 4.4 Heavy-Light Decomposition
int par[16] [MAXN], head[MAXN], sz[MAXN], in[MAXN], out[MAXN], lv[MAXN];
 vector<int> v[MAXN];
int dfs0(int x, int p) {
    lv[x] = lv[p] + 1; // lv[0] = -1;
    sz[x] = 1:
    for(int& y : v[x]) {
        if(y == p) continue;
        par[0][y] = x;
        sz[x] += dfs0(y, x);
        if(sz[v[x][0]] < sz[y]) swap(y, v[x][0]);
    }
    return sz[x];
void dfs1(int x, int p) {
    static int cnt = 0:
    in[x] = ++cnt:
    for(int v : v[x]) {
        if(y == p) continue;
        head[y] = y == v[x][0] ? head[x] : y;
        dfs1(v, x);
    }
    out[x] = cnt;
int get_query(segtree& seg, int x, int p) {
    int ret = 0:
    while(head[x] != head[p]) {
        // query(in[head[x]], in[x]);
        x = par[0][head[x]];
    // query(in[p], in[x]);
    return ret;
// for(int i=1;i<=n;i++) if(!in[i]) head[i] = i, dfs0(i, 0), dfs1(i, 0);
```

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#### 4.5 Dominator Tree

```
namespace Dtree {
    const int MAXN = 250001:
    vector<int> E[MAXN], RE[MAXN], rdom[MAXN];
    int S[MAXN], RS[MAXN], cs;
    int par[MAXN], val[MAXN];
    int sdom[MAXN], rp[MAXN];
    int dom[MAXN];
    int Find(int x, int c = 0) {
       if (par[x] == x) return c?-1:x;
       int p = Find(par[x], 1);
       if (p == -1) return c?par[x]:val[x];
       if (sdom[val[x]] > sdom[val[par[x]]]) val[x] = val[par[x]];
       par[x] = p;
       return c?p:val[x];
   }
   void Union(int x, int y) {
       par[x] = y;
   }
   void dfs(int x) {
       RS[S[x] = ++cs] = x;
       par[cs] = sdom[cs] = val[cs] = cs;
       for(int e : E[x]) {
           if (S[e] == 0) dfs(e), rp[S[e]] = S[x];
           RE[S[e]].pb(S[x]);
       }
   }
    int Do(int s, int *up) {
       dfs(s):
       for (int i = cs;i;i--) {
           for (int e : RE[i]) sdom[i] = min(sdom[i], sdom[Find(e)]);
           if (i > 1) rdom[sdom[i]].pb(i);
           for (int e:rdom[i]) {
                int p = Find(e):
               if (sdom[p] == i) dom[e] = i;
               else dom[e] = p;
           }
            if (i > 1) Union(i, rp[i]);
       for (int i = 2; i <= cs; i++) if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
        for (int i = 2; i <= cs; i++) {
           up[RS[i]] = RS[dom[i]]:
       }
       return cs;
   }
```

```
void addE(int x, int y) {E[x].pb(y);}
4.6 Global Min Cut
// Stoer-Wagner Algorithm, O(VE lg E)
int minimum_cut_phase(int n, int &s, int &t,
   vector<vector<int>> &adj, vector<int> vis){
   vector<int> dist(n):
   int mincut = 1e9;
   while(true){
        int pos = -1, cur = -1e9;
       for(int i=0; i<n; i++){
            if(!vis[i] && dist[i] > cur){
                cur = dist[i];
                pos = i;
           }
        if (pos == -1) break;
        s = t:
       t = pos;
        mincut = cur;
        vis[pos] = 1;
        for(int i=0; i<n; i++){</pre>
            if(!vis[i]) dist[i] += adj[pos][i];
   }
   return mincut; // optimal s-t cut here is, {t} and V \ {t}
int solve(int n, vector<vector<int>> adj){
   if(n <= 1) return 0;
   vector<int> vis(n):
   int ans = 1e9:
   for(int i=0; i<n-1; i++){</pre>
        int s, t;
        ans = min(ans, minimum_cut_phase(n, s, t, adj, vis));
        vis[t] = 1;
       for(int j=0; j<n; j++){</pre>
           if(!vis[j]){
                adj[s][j] += adj[t][j];
                adj[j][s] += adj[j][t];
           }
       }
        adj[s][s] = 0;
   }
   return ans;
```

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#### 5 Network Flow

#### 5.1 Theorems

**Max-flow Min-cut theorem** : 정점 s에서 정점 t까지 흐를 수 있는 최대 유량(max-flow)은 정점 s와 정점 t를 분리하는 간선들의 가중치 합(min-cut)과 같다.

Vertex cover: 어떤 그래프의 정점의 집합 S에 대해 그래프의 모든 간선이 S의 원소 중 최소 하나와 연결되어 있을 때, S를 해당 그래프의 vertex cover라고 하며, minimum vertex cover는 최소 개수의 정점을 사용한 vertex cover이다.

Independent set : 어떤 그래프의 정점의 집합 S에 대해 S의 서로 다른 두 정점을 연결하는 간선이 없을 때, S를 해당 그래프의 independent set이라고 하며, maximum independent set은 최대 개수의 정점을 사용한 independent set이다.

Matching (independent edge set) : 어떤 그래프의 간선의 집합 S에 대해 S의 서로 다른 두 간선이 공통된 정점을 가지지 않을 때, S를 해당 그래프의 matching이라고 하며, maximum matching은 최대 개수의 간선을 사용한 matching이다.

König's theorem : 이분 그래프의 maximum matching의 크기는 minimum vertex cover의 것과 같다. Dinic's Algorithm : 시간 복잡도  $O(V^2E)$ , unit capacity에서는  $\min(V^{2/3}E, E^{3/2})$ .

**Circulation Problem** : 새로운 source/sink  $s_n$ ,  $t_n$ 를 만들어서 다음과 같이 간선을 추가하고  $maxflow(s_n \to t_n) = \sum l_i$ 인지 확인, 이후  $s \to t$ 로 maxflow

```
• s_n \to b (l), a \to t_n (l), a \to b (r-l), t \to s (\infty)
```

### 5.2 Dinic's Algorithm

```
const int INF = 1e9;
struct Dinic{
    int N:
    struct edge{
       int index, cap, rev;
        edge(): index(0), cap(0), rev(0) {}
        edge(int index, int cap, int rev) : index(index), cap(cap), rev(rev) {}
    };
    vector<vector<edge>> ADJ;
    vector<int> R. W:
   Dinic() {}
   Dinic(int N) : N(N){
        ADJ.resize(N); R.resize(N);
                                        W.resize(N);
    }
    void CE(int node1, int node2, int cap){
        ADJ[node1].push_back(edge(node2, cap, ADJ[node2].size()));
        ADJ[node2].push_back(edge(node1, 0, ADJ[node1].size() - 1));
   }
   bool bfs(int src, int sink){
       fill(R.begin(), R.end(), -1):
        R[src] = 0;
        queue<int> Q; Q.push(src);
        while(Q.size()){
            int here = Q.front(); Q.pop();
```

```
for(auto e : ADJ[here]){
                if(e.cap > 0 && R[e.index] == -1)
                    R[e.index] = R[here] + 1, Q.push(e.index);
        return R[sink] != -1;
   }
    int dfs(int here, int sink, int f){
        if(here == sink) return f:
        for(int &i = W[here] ; i < ADJ[here].size() ; i++){</pre>
            auto &e = ADJ[here][i];
            if(e.cap > 0 && R[here] < R[e.index]){</pre>
                int res = dfs(e.index, sink, min(f, e.cap));
                if(res) {
                    e.cap -= res:
                    ADJ[e.index][e.rev].cap += res;
                    return res;
            }
        }
        return 0;
   }
   int solve(int src, int sink){
        int ret = 0:
        while(bfs(src. sink)){
            fill(W.begin(), W.end(), 0);
            while((res = dfs(src, sink, INF))) ret += res;
       }
        return ret;
   }
};
5.3 MCMF with SPFA
const int INF = 1e9;
struct MCMF {
   struct edge {
        int there, cap, cost, rev;
        edge(): there(0), cap(0), cost(0), rev(0) {}
        edge(int there, int cap, int cost, int rev) : there(there), cap(cap),
            cost(cost), rev(rev) {}
   };
    int N;
   vector<vector<edge>> ADJ;
   vector<int> R, INQ, C, I;
```

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```
MCMF() : N(0) {}
    MCMF(int N) : N(N) { ADJ.resize(N + 1); R.resize(N + 1); INQ.resize(N + 1);
        C.resize(N + 1): I.resize(N + 1): }
    void CE(int i, int i, int cap, int cost) {
        ADJ[i].push_back(edge(j, cap, cost, ADJ[j].size()));
        ADJ[i].push_back(edge(i, 0, -cost, ADJ[i].size() - 1));
    }
    bool SPFA(int src, int sink) {
        queue<int> Q; Q.push(src);
        fill(R.begin(), R.end(), -1);
                                            R[src] = 0;
        fill(C.begin(), C.end(), -1);
                                            C[src] = 0;
        fill(INQ.begin(), INQ.end(), 0);
                                            INQ[src] = 1;
        while (Q.size()) {
            int here = Q.front(); Q.pop();
            INO[here] = 0:
            for (int i = 0; i < ADJ[here].size(); i++) {</pre>
                auto e = ADJ[here][i];
                if ((C[e.there] == -1 || C[e.there] > C[here] + e.cost)
                    && e.cap > 0) {
                    C[e.there] = C[here] + e.cost:
                    R[e.there] = here;
                    I[e.there] = i;
                    if (!INQ[e.there]) INQ[e.there] = 1, Q.push(e.there);
               }
            }
        }
        if (C[sink] == -1) return false;
        return true:
    }
    pii mcmf(int src, int sink) {
       pii ret = { 0, 0 };
        while (SPFA(src. sink)) {
            int flow = INF. cost = 0:
            for (int here = sink; here != src; here = R[here])
                flow = min(flow, ADJ[R[here]][I[here]].cap);
            for (int here = sink; here != src; here = R[here]) {
                auto &e = ADJ[R[here]][I[here]];
                cost += e.cost * flow:
                e.cap -= flow;
                ADJ[e.there][e.rev].cap += flow;
            ret.first += flow, ret.second += cost;
        }
        return ret;
    }
};
```

### 5.4 Hungarian Method

```
namespace Hung {
    const int MX = 2000;
    // IMPORTANT : n <= m, 1-based
   using T = long double;
   T \max v = 1e200;
   T a[MX][MX], n, m:
    void init(int nn, int mm) { n = nn; m = mm; }
    void set_value(int x, int y, T val) { a[x][y] = val; }
   T solve(vector <int> &ans) {
        vector<T> v(m+1), u(n+1):
        vector\langle int \rangle p (m+1), way (m+1);
        for (int i=1; i<=n; ++i) {
            p[0] = i;
            int j0 = 0;
            vector<T> minv (m+1, maxv);
            vector<char> used (m+1, false):
            do {
                used[j0] = true;
                T delta = maxv:
                int i0 = p[j0], j1;
                for (int j=1; j<=m; ++j) if (!used[j]) {</pre>
                    T cur = a[i0][j]-u[i0]-v[j];
                    if (cur < minv[j]) minv[j] = cur, wav[j] = j0;</pre>
                     if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
                for (int j=0; j<=m; ++j) {
                     if (used[j]) u[p[j]] += delta, v[j] -= delta;
                     else minv[j] -= delta;
                }
                j0 = j1;
            } while (p[j0] != 0);
            do {
                int j1 = wav[j0];
                p[j0] = p[j1];
                j0 = j1;
            } while (j0);
        ans.resize(n + 1);
        for(int j=1; j<=m; ++j) {</pre>
            ans[p[j]] = j;
        return -v[0];
   }
```

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### 5.5 Hopcroft-Karp Algorithm

```
struct hopcroft_karp{
    int N:
    vector<vector<int>> ADJ;
    vector<int> L, rev, used;
    hopcroft_karp() {}
   hopcroft_karp(int N) : N(N) {
       ADJ.resize(N);
       L.resize(N), rev.resize(N, -1), used.resize(N, 0);
   }
    void CE(int here, int there){
        ADJ[here].push_back(there);
   }
    void bfs(){
        queue<int> Q;
       for(int i = 0 ; i < N ; i++) {</pre>
            if(used[i]) L[i] = -1;
            else L[i] = 0, Q.push(i);
       }
        while(Q.size()){
            int here = Q.front(); Q.pop();
            for(int there : ADJ[here]){
                if(rev[there] != -1 && L[rev[there]] == -1) {
                    L[rev[there]] = L[here] + 1;
                    Q.push(rev[there]);
               }
            }
       }
   }
   bool dfs(int here){
       for(int there : ADJ[here]){
            if(rev[there] == -1 || (L[here] < L[rev[there]] && dfs(rev[there]))){
                rev[there] = here:
                used[here] = 1;
                return true;
            }
       }
       return false;
   }
    int solve(){
        int ret = 0;
        while(1){
            bfs();
            int res = 0;
```

```
for(int i = 0 ; i < N ; i++) {</pre>
                  if(used[i]) continue;
                  res += dfs(i):
             if(res == 0) break:
             ret += res;
        return ret;
};
    Optimization Tricks
6.1 Knuth Optimization
  • Recurrence : D[i][j] = \min_{i < k < j} (D[i][k] + D[k][j]) + C[i][j]
  • Quadrangle Inequality : C[a][c] + C[b][d] \le C[a][d] + C[b][c], \ a \le b \le c \le d
  • Monotonicity: C[b][c] < C[a][d], a < b < c < d
  • A[i|[j] = (\min k \ s.t. \ D[i][j] \ \text{is min.}). Then A[i][j-1] \le A[i][j] \le A[i+1][j]
  • O(N^2) time complexity
// opt[i-1][i] = i
for(int d=2:d<=n:d++) {</pre>
    for(int i=1;i+d<=n+1;i++) {</pre>
        for(int k=opt[i][j-1], j=i+d; k<=opt[i+1][j]; k++) {</pre>
             int v = dp[i][k] + dp[k][j] + c[i][j];
             if(dp[i][j] > v) dp[i][j] = v, opt[i][j] = k;
    }
}
6.2 Divide and Conquer Optimization
  • Recurrence : D[t][i] = \min_{k < i} (D[t-1][k] + C[k][i])
  • Min index : A[t][i] < A[t][i+1] (A[t][i] = (\min. k \ s.t. \ D[t][i] \text{ is min.}))
  • Quadrangle Inequality: C[a][c] + C[b][d] < C[a][d] + C[b][c], a < b < c < d
  • Able to Divide and Conquer base on calculating D[t][i]
  • O(TN \lg N) time complexity
// range of index : [1,r]
// range of dp : [s.e]
```

void dnc(int t, int l, int r, int s, int e)

if(s > e) return; int m = (s+e)/2; Korea University - Powered by Zigui Page 16 of 21

```
D[t][m] = 2e9;
    for(int k=1;k<m&k<=r;k++)</pre>
        int tmp = D[t-1][k] + C[k][m];
        if(D[t][m] > tmp)
            D[t][m] = tmp, A[t][m] = k;
    }
    dnc(t, 1, A[t][m], s, D[t][m]-1);
    dnc(t, A[t][m], r, D[t][m]+1, e);
}
```

#### Convex Hull Trick

- Recurrence:  $dp[i] = \min_{i < i} (dp[j] + a[i]b[j]), b[i-1] < b[i]$
- Think as  $dp[x = a[i]] = \min_{i < i} (b[i] \cdot x + dp[i])$
- Thus push lines and find minimum (by binary search)
- If a[i] < a[i+1] sweeping is possible
- Intersection of  $y = a_i x + b_i$  and  $y = a_{i+1} x + b_{i+1} : x = \frac{b_{i+1} b_i}{a_i a_{i+1}}$

### 6.4 Centroid Decomposition

```
// credit : https://gist.github.com/igorcarpanese/75162f3253bd230abd0f32f9950bf384
int dfs(int u, int p) {
    sub[u] = 1;
    for (auto v : tree[u])
        if (v != p \text{ and } !\text{chk}[v]) \text{ sub}[u] += dfs(v, u);
    return sub[u] + 1;
} // calc subtree size. chk means that node was a centroid
// each tree has at most two centroids
int centroid(int u, int p, int r) { // r : root
    for (auto v : tree[u])
        if (v != p and !chk[v] and sub[v] > sub[r]/2) return centroid(v, u, r);
    return u;
}
```

#### **Data Structure**

## 7.1 Persistent Segment Tree

```
const MAXN = 1e5 + 10;
struct node{
    node *1. *r:
    int cnt;
    node () {}
} pool[(1 << 17) * 17], *tree_head[MAXN];</pre>
```

```
int tcnt;
node* alloc(){
   memset(pool + tcnt, 0, sizeof(node));
   return pool + tcnt++;
node * init(int 1, int r){
   node *ret = alloc();
   if(1 != r) {
       int mid = (1 + r) / 2;
       ret->l = init(l, mid);
       ret->r = init(mid + 1, r);
   }
   return ret;
void update(node * here, node *par, int 1, int r, int val){
   if(1 == r) {
       here->cnt = par->cnt + 1;
       return;
   }
   int mid = (1 + r) / 2;
   if(val <= mid){</pre>
       here->1 = alloc();
       here->r = par->r:
       update(here->1, par->1, 1, mid, val):
   }
   else {
       here->1 = par->1;
       here->r = alloc();
       update(here->r, par->r, mid + 1, r, val);
   here->cnt = here->l->cnt + here->r->cnt:
int query(node *node1, node *node2, int 1, int r, int k){
   if(1 == r) return 1;
   int ccc = node1->l->cnt - node2->l->cnt;
   int mid = (1 + r) / 2:
   if(k <= ccc) return query(node1->1, node2->1, 1, mid, k);
   else return query(node1->r, node2->r, mid + 1, r, k - ccc);
7.2 Link-Cut Tree
struct node{
   node *pp, *p, *l, *r;
   int val;
   node(){p = 0, 1 = 0, r = 0;}
```

```
node(int val) : val(val) { p = 0, 1 = 0, r = 0;}
```

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```
};
void push(node *x){}
void pull(node *x){}
void rotate(node *x){
    if(!x->p) return;
    push(x->p); // if there's lazy stuff
    push(x);
    node *p = x->p;
    bool is_left = (p->1 == x);
    node *b = (is_left ? x->r : x->1);
    x->p = p->p;
    if (x-p \&\& x-p-1 == p) x-p-1 = x;
    if(x->p && x->p->r == p) x->p->r = x;
    if(is_left){
       if(b) b \rightarrow p = p;
        p->1 = b;
       p->p = x;
       x->r = p;
    }
    else{
        if(b) b->p = p;
        p->r = b;
       p->p = x;
       x->1 = p;
    }
   pull(p); // if there's something to pull up
   pull(x);
    //if(!x->p) root = x; // IF YOU ARE SPLAY TREE
    if(p->pp){ // IF YOU ARE LINK CUT TREE
        x->pp = p->pp;
        p->pp = nullptr;
    }
}
void splay(node *x){
    while(x->p){
        node *p = x->p;
        node *g = p->p;
        if(g){
            if((p->1 == x) ^ (g->1 == p)) rotate(x);
            else rotate(p);
        }
        rotate(x);
    }
}
void access(node *x){
    splay(x);
    push(x);
    if(x->r){
```

```
x->r->pp = x;
       x->r->p = nullptr;
       x->r = nullptr;
   pull(x);
   while(x->pp){
        node *nxt = x->pp;
        splay(nxt);
        push(nxt);
       if(nxt->r){
           nxt->r->pp = nxt;
            nxt->r->p = nullptr;
            nxt->r = nullptr;
       }
       nxt->r = x;
       x->p = nxt;
       x->pp = nullptr;
       pull(nxt);
        splay(x);
   }
node *root(node *x){
   access(x);
   while(x->1){
        push(x);
       x = x -> 1;
   }
   access(x);
   return x;
node *par(node *x){
   access(x);
   if(!x->1) return nullptr;
   push(x);
   x = x->1;
   while(x->r){
        push(x);
       x = x->r;
   }
   access(x):
   return x;
node *lca(node *s, node *t){
   access(s);
   access(t);
    splay(s);
   if(s->pp == nullptr) return s;
   return s->pp;
void link(node *par, node *son){
```

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auto l = \*lower\_bound((Line) { x, is\_query });

```
access(son);
                                                                                                 return 1.m * x + 1.b:
                                                                                             }
    //son->rev ^= 1: // remove if needed
    push(son);
                                                                                         };
    son->1 = par;
    par->p = son;
                                                                                          7.4 Stern-Brocot Tree
    pull(son);
                                                                                         // __int128 is recommended
                                                                                         bool test(11 a, 11 b) { // for testing directions, vary by prob
void cut(node *p){
                                                                                             // return true if (true value) >= a/b
    access(p);
    push(p);
                                                                                             11 n = 0, m = 1;
    if(p->1){
                                                                                             rep(i, N) {
        p->1->p = nullptr;
        p->1 = nullptr;
                                                                                                 if (n < m*A[i].fi) n = A[i].fi, m = 1;
    }
    pull(p);
                                                                                                 11 c = b*n+m*a, d = m*b;
}
                                                                                                 11 g = gcd(c, d);
                                                                                                 n = c/g;
                                                                                                 m = d/g;
     Dynamic Convex Hull
// https://github.com/niklasb/contest-algos/blob/master/convex_hull/dynamic.cpp
                                                                                                  if (n > m*A[i].se) return false;
const ll is_query = -(1LL<<62);</pre>
                                                                                             }
struct Line {
                                                                                             return true:
    11 m. b:
    mutable function<const Line*()> succ;
    bool operator<(const Line& rhs) const {</pre>
                                                                                         pair<11, 11> stern brocot(11 M, 11 N) {
                                                                                             // M : max value
        if (rhs.b != is_query) return m < rhs.m;</pre>
                                                                                             // N : max divisor
        const Line* s = succ():
        if (!s) return 0;
                                                                                             // if result is a/b, return as {a, b}
        return b - s -> b < (s -> m - m) * rhs.m;
                                                                                             11 a = 0, b = 1; // 1
    }
                                                                                             11 c = 1, d = 0; // r
};
                                                                                             int 1, r;
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum
                                                                                             bool chg = true;
    bool bad(iterator v) {
        auto z = next(y);
                                                                                             while(chg) {
        if (y == begin()) {
                                                                                                  chg = false;
            if (z == end()) return 0;
            return y->m == z->m && y->b <= z->b;
                                                                                                  // to left
                                                                                                 1 = 0, r = (N-d-1)/b+1:
        auto x = prev(y);
                                                                                                  while(1 < r) {
        if (z == end()) return y \rightarrow m == x \rightarrow m && y \rightarrow b <= x \rightarrow b;
                                                                                                      int mid = (1+r+1)/2:
        return (x-b - y-b)*(z-m - y-m) >= (y-b - z-b)*(y-m - x-m);
                                                                                                      if (test(a*mid+c, b*mid+d)) r = mid-1;
    }
                                                                                                      else l = mid:
    void insert_line(Line 1) {
        auto y = insert(1);
        y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
                                                                                                  c += a*1;
        if (bad(v)) { erase(v): return: }
                                                                                                  d += b*1:
        while (next(y) != end() && bad(next(y))) erase(next(y));
                                                                                                  chg |= (1 > 0);
        while (y != begin() && bad(prev(y))) erase(prev(y));
    }
                                                                                                  // to right
    11 eval(ll x) {
                                                                                                  1 = 0, r = (d?(N-b-1)/d+1:M);
```

access(par);

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X.insert(16);

```
while(1 < r) {</pre>
            int mid = (1+r+1)/2:
            if (test(a+mid*c, b+mid*d)) l = mid;
            else r = mid-1;
        }
        a += c*1;
        b += d*l:
        chg |= (1 > 0);
    }
    return {a, b};
}
7.5 Rope
#include <bits/stdc++.h>
#include <ext/rope>
using namespace std;
using namespace __gnu_cxx;
int main()
    ios::sync_with_stdio(false);
    cin.tie(0);
    crope rp; // rope<char>
    string s("Lorem-ipsum");
    int n = s.length():
    rp.append(s.c_str()); // add element
    int x = 3, y = 8; // split and merge below
    rp = rp.substr(x, y-x) + rp.substr(0, x) + rp.substr(y, n);
    cout << rp.at(0) << '\n'; // get element, 'e'</pre>
    cout << rp << '\n'; // print, "em-ip|Lor|sum"</pre>
}
7.6 Bitset
#include <bitset>
#include <iostream>
using namespace std;
int main() {
    bitset<8> b1(13);
                                     // 00001101
    bitset<8> b2("10111");
                                     // 00010111
    cout << b1.count() << endl;</pre>
    cout << b1.test(6) << endl;</pre>
                                     // 0, since 2^6-th bit is 0
    b1.set(6):
                                     // set to 1, 1-fill if no param
    b2.reset(2);
                                     // set to 0, 0-fill if no param
```

```
// use 'flip' for flipping
    cout << "b1:" << b1 << endl:
                                            // b1:01001101
    cout << "b2:" << b2 << endl;
                                            // b2:00010011
    // use anv. none, all (c++11) for bit checking
    // supported operators : &, |, ^, <<, >>, ~, ==, !=
    // these operators must match size (given to template)
    cout << "&: " << (b1 & b2) << endl;
                                            // &: 00000001
    cout << "^: " << (b1 ^ b2) << endl:
                                            // &: 01011110
    cout << "|: " << (b1 | b2) << endl;
                                            // |: 01011111
    cout << "~: " << (~b1) << endl;
                                            // ~: 10110010
    cout << "<<:" << (b1 << 3) << endl:
                                            // <<:01101000
    cout << b2.to_ulong() << endl;</pre>
                                            // 19, c++11 supports ullong
7.7 Policy Based Data Structure
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
using namespace std;
using namespace __gnu_pbds;
typedef tree<
int.
null_type,
less<int>.
rb_tree_tag,
tree_order_statistics_node_update >
ordered set:
// less<int> : not allow for duplicate
// less_equal<int> : allow for duplicate
// use upper_bound when you erase from set used less_equal
int N:
int main(void) {
    iostream::sync_with_stdio(false);
    cin.tie(nullptr);
    ordered_set X;
    X.insert(1):
    X.insert(2);
    X.insert(4);
    X.insert(8):
```

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```
cout<<*X.find_by_order(1)<<endl; // 2
cout<<*X.find_by_order(2)<<endl; // 4
cout<<*X.find_by_order(4)<<endl; // 16
cout<<(end(X)==X.find_by_order(6))<<endl; // true

cout<<X.order_of_key(-5)<<endl; // 0
cout<<X.order_of_key(1)<<endl; // 0
cout<<X.order_of_key(3)<<endl; // 2
cout<<X.order_of_key(4)<<endl; // 2
cout<<X.order_of_key(4)<<endl; // 5
}</pre>
```

#### 8 Miscellaneous

#### 8.1 Misc Formulae and Algorithms

#### 8.1.1 Faulhaber's Formula

$$T(n,k) = \sum_{i=1}^{n} i^{k} = \frac{(n+1)^{k+1} - 1^{k+1} - \sum_{j=0}^{k-1} {k+1 \choose j} T(n,j)}{{k+1 \choose k}}$$

Also use

$$(x+1)^d - x^d = 1 + {d \choose 1}x + {d \choose 2}x^2 + \dots + {d \choose d-1}x^{d-1}$$

to get each coef.

#### 8.1.2 Maximum Clique

```
typedef long long 11;
11 G[40]: // 0-index
int N, M;
int cur:
void get_clique(int R = 0, 11 P = (111 << N)-1, 11 X = 0){
    if((P|X) == 0){
        cur = max(cur, R);
        return;
    }
    int u = __builtin_ctzll(P|X);
   11 c = P\&^{G}[u];
    while(c){
        int v = __builtin_ctzll(c);
        get_clique(R + 1, P&G[v], X&G[v]);
        P ^= 111 << v:
       X = 111 << v;
        c ^= 111 << v;
    }
}
```

#### 8.1.3 De Brujin Sequence

```
// https://github.com/koosaga/DeobureoMinkyuParty/blob/master/teamnote.tex
// alphabet = [0, k - 1], substr length n, res starts with 0 (cyclic)
int res[10000000], aux[10000000]; // >= k^n, k*n
int de_bruijn(int k, int n, int lim) { // returns size (k^n)
   if(k == 1) {
        res[0] = 0;
        return 1;
   }
   for(int i = 0; i < k * n; i++) aux[i] = 0;
   int sz = 0:
   function<void(int, int)> db = [&](int t, int p) {
        if(sz > lim) return;
        if(t > n) {
           if(n \% p == 0)
                for(int i = 1; i <= p; i++)
                    res[sz++] = aux[i]:
       }
        else {
            aux[t] = aux[t - p];
           db(t + 1, p);
           for(int i = aux[t - p] + 1; i < k; i++) {
                aux[t] = i;
                db(t + 1, t);
           }
   };
   db(1, 1);
    return sz;
```

#### 8.2 Highly Composite Numbers, Large Prime

< 10^k	number	divisors	2 3 5 71113171923293137
1	6	4	1 1
2	60	12	2 1 1
3	840	32	3 1 1 1
4	7560	64	3 3 1 1
5	83160	128	3 3 1 1 1
6	720720	240	4 2 1 1 1 1
7	8648640	448	6 3 1 1 1 1
8	73513440	768	5 3 1 1 1 1 1
9	735134400	1344	6 3 2 1 1 1 1
10	6983776800	2304	5 3 2 1 1 1 1 1
11	97772875200	4032	6 3 2 2 1 1 1 1
12	963761198400	6720	6 4 2 1 1 1 1 1 1
13	9316358251200	10752	6 3 2 1 1 1 1 1 1 1
14	97821761637600	17280	5 4 2 2 1 1 1 1 1 1
15	866421317361600	26880	6 4 2 1 1 1 1 1 1 1 1

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16	8086598962041600	41472	8 3	2 2 1 1	1 1	1	1 1			
17	74801040398884800	64512	6 3	2 2 1 1	1 1	1	1 1	1		
18	897612484786617600	103680	8 4	2 2 1 1	1 1	1	1 1	1		
	< 10^k prime	· 10^k pri	ne	# of	pri	me				
1	7		 11			4		•		
2	97	10				25				
3	997	100			1	68				
4	9973	1000	07		12	29				
5	99991	10000	)3		95	92				
6	999983	100000	)3		784	98				
7	9999991	100000	19	6	645	79				
8	9999989	1000000	)7	57	7614	55				
9	99999937	100000000	)7	508	3475	34				
	< 10^k prim	ne		> 10^k	pri	me				
10	99999999			100000	0000	19		•		
11	999999999			100000						
12	9999999999			1000000						
13	99999999999			.00000000						
14	999999999999			00000000						
15	999999999999	39	100	00000000	0000	37				
16	9999999999999	37	1000	00000000	0000	61				
17	999999999999999	7	10000	00000000	0000	03				
18	99999999999999	39 :	100000	00000000	0000	03				
99824 98566 10129	rime: $2049 = 7 \times 2^{26} + 1$ . Primiti $4353 = 119 \times 2^{23} + 1$ . Prim $1441 = 235 \times 2^{22} + 1$ . Prim $24417 = 483 \times 2^{21} + 1$ . Primear $10^9 : 10^9 + [7, 9, 21, 33, 20]$	itive root: nitive root: mitive root:	3. 3.							
8.3 I	Fast Integer IO									
static static static static	<pre>it : https://github.com/ char buf[1 &lt;&lt; 19]; // si int idx = 0; int bytes = 0; inline int _read() { (!bytes    idx == bytes) bytes = (int)fread(buf,</pre>	ze : any 1	number	geq tha	an 1	024			eamno	te.tex

idx = 0;

return buf[idx++]:

int x = 0, s = 1;

int c = \_read();

static inline int \_readInt() {

}

```
while (c \leq 32) c = _read();
    if (c == '-') s = -1, c = _read();
    while (c > 32) x = 10 * x + (c - '0'), c = _read();
   if (s < 0) x = -x;
   return x:
8.4 C++ Tips / Environments
#include <bits/stdc++.h> // magic header
using namespace std;
                      // magic namespace
struct StupidGCCCantEvenCompileThisSimpleCode{
    pair<int, int> array[1000000];
}; // https://gcc.gnu.org/bugzilla/show_bug.cgi?id=68203
// how to use rand (in 2017)
mt19937 rng(0xdeadbeef);
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
int randint(int lb, int ub){ return uniform_int_distribution<int>(lb, ub)(rng); }
shuffle(permutation.begin(), permutation.end(), rng);
mt19937_64 _R(chrono::steady_clock::now().time_since_epoch().count()); // _R()
// comparator overload
auto cmp = [](seg a, seg b){return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
int x, y;
bool operator==(const point &p)const{ return x == p.x && y == p.y; }
struct hasher {
size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
unordered_map<point, int, hasher> hsh;
// c++ setprecision example
#include <iostream>
                        // std::cout, std::fixed
#include <iomanip>
                       // std::setprecision
int main () {
    double f = 3.14159;
    std::cout << std::setprecision(5) << f << '\n'; // 3.1416
    std::cout << std::setprecision(9) << f << '\n'; // 3.14159
    std::cout << std::fixed;</pre>
   std::cout << std::setprecision(5) << f << '\n'; // 3.14159
    std::cout << std::setprecision(9) << f << '\n'; // 3.141590000
```