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Team Note of Powered by Zigui

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ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

0 Quotes and Prerequisites

```
evenharder : Mental Abuse To Humans
djkim0613 : 열심히 응원하겠습니다.
SoulTch : How much is this bus ticket?
* This template is brought from that of 'Deobureo Minkyu Party'
```

Run script

```
#!/bin/bash
g++ -fsanitize=undefined -std=c++14 -02 -o /tmp/pow $1.cpp
time /tmp/pow < $1.in
# export PATH=~:$PATH</pre>
```

Debug Code

```
#define setz(x) memset(x, 0, sizeof(x))
#define sz(x) ((int)(x).size())
#define rep(i, e) for (int i = 0, _##i = (e); i < _##i; i++)
#define repp(i, s, e) for (int i = (s), _##i = (e); i < _##i; i++)
#define repr(i, s, e) for (int i = (s)-1, _##i = (e); i \ge _{\#}i; i--)
#define repi(i, x) for (auto &i : (x))
// using namespace std;
using ll = long long;
using pii = pair<int, int>;
using pll = pair<11, 11>;
template<typename T>
ostream &operator<<(ostream &os, const vector<T>& v) {
    cout << "[":
    for (auto p : v) cout << p << ",";</pre>
    cout << "]":
    return os;
}
#ifndef SOULTCH
#define debug(...) 0
#define endl '\n'
#define debug(...) cout << " [-] ", _dbg(#__VA_ARGS__, __VA_ARGS__)</pre>
template<class TH> void _dbg(const char *sdbg, TH h){ cout << sdbg << '=' << h <<
endl; }
template<class TH, class... TA> void _dbg(const char *sdbg, TH h, TA... a) {
    while(*sdbg != ',') cout << *sdbg++;</pre>
    cout << '=' << (h) << ',';
    _dbg(sdbg+1, a...);
}
#endif
```

Reminders

Should be added.

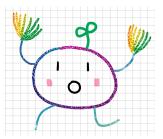


Figure 1: 풀다가 막힐 때는 이 그림을 봅시다. 아자아자 화이팅!

1 Math

1.1 Basic Mathematics

1.1.1 Trigonometry

- $\sin(\alpha \pm \beta) = \sin \alpha \cos \beta \pm \cos \alpha \sin \beta$
- $\cos(\alpha \pm \beta) = \cos \alpha \cos \beta \mp \sin \alpha \sin \beta$
- $\tan(\alpha \pm \beta) = \frac{\tan \alpha \pm \tan \beta}{1 \mp \tan \alpha \tan \beta}$
- $\sin 2\theta = 2\sin \theta \cos \theta$
- $\bullet \ c^2 = a^2 + b^2 2ab\cos\gamma$

1.1.2 Generating Function

- $\sum_{n} (pn+q)x^{n} = \frac{p}{1-x} + \frac{q}{(1-x)^{2}}$ (Arithmetic progression)
- $\sum_{n} (rx)^n = (1 rx)^{-1}$ (Geometric progression)
- $\sum_{n} {m \choose n} x^n = (1+x)^m$ (Binomial coefficient)
- $\sum_{n} {m+n-1 \choose n} x^n = (1-x)^{-m}$ (Multiset coefficient)

1.1.3 Calculus

- $\int_a^b f(x) dx = \frac{b-a}{6} \left[f(a) + 4f(\frac{a+b}{2}) + f(b) \right]$ (Simpson's Rule, for cubic poly)
- $\int u'v \, dx = uv \int uv' \, dx$ (Integration by parts)

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1.2 Number Theory

1.2.1 Lattice Points under Line

```
// 0 \le x \le n, 0 \le y \le (a/c)x+(b/c)
11 calc(11 a.11 b.11 c.11 n){
    if(!n)return 0:
    11 tmp=a/c*n*(n-1)/2;
    tmp+=b/c*n:
    return tmp+calc(c,(a*n+b)%c,a%c,((a%c)*n+b%c)/c);
}
1.2.2 Shanks' Baby-step Giant-step
11 mexp(ll x, ll y, ll p) {
    if(!y) return 1;
    if(y & 1) return x * mexp(x*x%p, y>>1, p) % p;
    return mexp(x*x\%p, y>>1, p);
vector<ll> get_factor(ll n) {
    vector<ll> v:
    for(ll i=2;i*i<=n;i++) {
        if(n \% i == 0) {
            v.push back(i):
            while(n \% i == 0) n /= i;
        }
    }
    if(n > 1) v.push_back(n);
    return v:
ll get_primitive(ll n) {
    11 phi = n-1; // assume n is prime
    vector<ll> fact = get_factor(phi);
    for(11 x=2;x<=n;x++) {
        int ves = 1:
        for(ll y : fact) {
            yes &= (mexp(x, phi / y, n) != 1);
        }
        if(yes) return x;
    }
    return -1;
// find x s.t. x^k \mod n = a \rightarrow (g^k)^y \mod n = a, where x = g^y
11 bsgs(ll k, ll a, ll n) {
    11 g = get_primitive(n);
    11 phi = n-1; // assume n is prime
    if (g == -1) return -1;
    ll m = ceil(sqrt(n) + 1e-9):
    vector<pl> prec(m);
    for(ll j=0;j<m;j++) {</pre>
        prec[j] = {mexp(g, j * k % phi, n), j};
```

```
sort(prec.begin(), prec.end());
   ll cur = a, ncur = mexp(g, (phi - m) * k % phi, n);
   for(ll i=0:i<m:i++) {</pre>
        auto it = lower_bound(prec.begin(), prec.end(), pl(cur, 0));
        if(it->first == cur) {
           ll ans = mexp(g, (i*m + it->second) \% phi, n);
            assert(mexp(ans, k, n) == a);
            return ans:
        cur = cur * ncur % n:
   }
   return 0;
1.2.3 Extended Euclidean Algorithm
// ax + by = gcd(a,b). x, y?
pll ext_gcd(ll a,ll b) {
   if(b) {
        auto tmp = ext_gcd(b, a%b);
       return {tmp.second, tmp.first - (a/b) * tmp.second};
   }
   else return {1, 0}:
// ax = gcd(a, m) mod m, x?
11 mod_inv(ll a, ll m) {
   return (ext_gcd(a, m).first + m) % m;
1.2.4 Chinese Remainder Theorem
ll pos_rem(ll a, ll m) { // m > 0. a % m?
   11 \text{ res} = abs(a) \% m:
    return a > 0 ? res : (res ? m - res : 0);
// ax = c mod m, bx = d mod n. x?
11 solve(ll a, ll c, ll m, ll b, ll d, ll n) {
   a = pos_rem(a, m); c = pos_rem(c, m); // if a, c not in [0, m)
   b = pos_rem(b, n); d = pos_rem(d, n); // if b, d not in [0, n)
   11 g = gcd(a, gcd(c, m)); a \neq g, c \neq g, m \neq g;
        g = gcd(b, gcd(d, n)); b /= g, d /= g, n /= g;
    if(c % _gcd(a, m) || d % _gcd(b, n)) return inf;
   ll t1 = (mod_inv(a, m) * c) % m;
   11 t2 = (mod_inv(b, n) * d) \% n;
    g = gcd(m, n);
   11 \ 1c = m * n / g:
   if(abs(t1 - t2) % g) return inf;
    pl p = ext_gcd(m, n);
   11 q = (t1 * p.second * n/g + t2 * p.first * m/g);
    return pos_rem(q, lc);
```

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```
}
1.2.5 Möbius Inversion Formula
 \forall n \in \mathbb{N} \ g(n) = \sum_{d|n} f(d) \implies f(n) = \sum_{d|n} \mu(d)g(n/d)
1.3 FFT
  FFT: (a_0, a_1, \dots, a_{n-1}) \mapsto (\sum_{j=0}^{n-1} a_0(\omega^0)^j, \sum_{j=0}^{n-1} a_1(\omega^1)^j, \dots, \sum_{j=0}^{n-1} a_{n-1}(\omega^{n-1})^j)
void fft(vector<base>& a, bool inv) {
    int n = a.size(), j = 0;
    vector<ll> roots(n/2);
    for(int i=1;i<n;i++) {</pre>
         int bit = (n >> 1);
         while(j >= bit) {
             i -= bit:
             bit >>= 1;
        }
         j += bit;
        if(i < j) swap(a[i], a[i]);</pre>
    }
    double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
    for(int i=0;i<n/2;i++) {</pre>
         roots[i] = base(cos(ang * i), sin(ang * i));
    }
    /* In NTT, let prr = primitive root. Then,
    int ang = mexp(prr, (mod - 1) / n);
    if(inv) ang = mexp(ang, mod - 2);
    for(int i=0; i<n/2; i++){
         roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
    }
    also, make sure to apply modulus under here
    for(int i=2;i<=n;i<<=1) {</pre>
         int step = n / i;
         for(int i=0:i<n:i+=i) {
             for(int k=0; k<i/2; k++) {
                  ll u = a[j+k], v = a[j+k+i/2] * roots[step * k];
                  a[j+k] = u+v;
                  a[j+k+i/2] = u-v;
             }
        }
    if(inv) for(int i=0:i<n:i++) a[i] /= n:
}
void conv(vector<base>& x, vector<base>& y) {
    int n = 2; while (n < max(x.size(), y.size())) n <<= 1;
```

```
n <<= 1;
   x.resize(n); y.resize(n);
   fft(x, false); fft(y, false);
   for(int i=0;i<n;i++) x[i] *= y[i];</pre>
   fft(x, true): // access (ll)round(x[i].real())
1.4 Miller-Rabin + Pollard-Rho
//Prove By Solving - https://www.acmicpc.net/problem/4149
namespace miller_rabin{
   lint mul(lint x, lint y, lint mod){ return (__int128) x * y % mod; }
   lint ipow(lint x, lint y, lint p){
       lint ret = 1, piv = x \% p;
        while(y){
           if(y&1) ret = mul(ret, piv, p);
           piv = mul(piv, piv, p);
           y >>= 1;
       return ret;
   bool miller_rabin(lint x, lint a){
        if(x \% a == 0) return 0:
       lint d = x - 1;
        while(1){
           lint tmp = ipow(a, d, x);
           if(d&1) return (tmp != 1 && tmp != x-1);
           else if(tmp == x-1) return 0;
           d >>= 1;
   7
   bool isprime(lint x){
       for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
           if(x == i) return 1:
            if (x > 40 \&\& miller_rabin(x, i)) return 0;
        if(x <= 40) return 0;
       return 1;
   }
}
namespace pollard_rho{
   lint f(lint x, lint n, lint c){
       return (c + miller_rabin::mul(x, x, n)) % n;
   }
   void rec(lint n, vector<lint> &v){
        if(n == 1) return:
       if(n \% 2 == 0){
           v.push_back(2);
           rec(n/2, v);
           return;
```

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```
if(miller_rabin::isprime(n)){
            v.push_back(n);
            return;
        }
        lint a, b, c;
        while(1){
            a = rand() \% (n-2) + 2:
            b = a;
            c = rand() \% 20 + 1;
            do{
                a = f(a, n, c);
                b = f(f(b, n, c), n, c);
            \frac{1}{2} while (\gcd(abs(a-b), n) == 1);
            if(a != b) break;
        lint x = gcd(abs(a-b), n);
        rec(x, v);
        rec(n/x, v);
    }
    vector<lint> factorize(lint n){
        vector<lint> ret;
        rec(n, ret);
        sort(ret.begin(), ret.end());
        return ret;
    }
}:
```

1.5 Black Box Linear Algebra + Kitamasa

```
vector<int> berlekamp_massey(vector<int> x){
    vector<int> ls, cur;
    int lf. ld:
    for(int i=0: i<x.size(): i++){</pre>
        lint t = 0;
        for(int i=0: i<cur.size(): i++){</pre>
            t = (t + 111 * x[i-j-1] * cur[j]) \% mod;
        if((t - x[i]) % mod == 0) continue:
        if(cur.empty()){
            cur.resize(i+1);
            lf = i:
            1d = (t - x[i]) \% mod;
            continue:
        }
        lint k = -(x[i] - t) * ipow(ld, mod - 2) % mod;
        vector<int> c(i-lf-1):
        c.push_back(k);
        for(auto &j : ls) c.push_back(-j * k % mod);
        if(c.size() < cur.size()) c.resize(cur.size());</pre>
        for(int j=0; j<cur.size(); j++){</pre>
```

```
c[i] = (c[i] + cur[i]) \% mod;
        if(i-lf+(int)ls.size()>=(int)cur.size()){
            tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) % mod);
        cur = c;
   }
   for(auto &i : cur) i = (i % mod + mod) % mod:
    return cur;
int get_nth(vector<int> rec, vector<int> dp, lint n){
    int m = rec.size();
   vector<int> s(m), t(m):
   s[0] = 1;
   if(m != 1) t[1] = 1;
   else t[0] = rec[0]:
    auto mul = [&rec](vector<int> v, vector<int> w){
        int m = v.size();
        vector < int > t(2 * m):
        for(int j=0; j<m; j++){</pre>
            for(int k=0: k < m: k++){
                t[j+k] += 111 * v[j] * w[k] % mod;
                if(t[j+k] >= mod) t[j+k] -= mod;
            }
        for(int j=2*m-1; j>=m; j--){
            for(int k=1: k<=m: k++){
                t[j-k] += 111 * t[j] * rec[k-1] % mod;
                if(t[i-k] \ge mod) t[i-k] -= mod:
        t.resize(m):
        return t;
   }:
    while(n){
        if(n \& 1) s = mul(s, t);
        t = mul(t, t);
        n >>= 1;
   lint ret = 0;
   for(int i=0; i<m; i++) ret += 111 * s[i] * dp[i] % mod;
   return ret % mod:
int guess_nth_term(vector<int> x, lint n){ // init with > 3k, 0(1^2 lg n)
    if(n < x.size()) return x[n];</pre>
   vector<int> v = berlekamp_massey(x);
   if(v.empty()) return 0;
   return get_nth(v, x, n);
struct elem{int x, y, v;}; // A_(x, y) <- v, 0-based. no duplicate please..
```

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```
vector<int> get_min_poly(int n, vector<elem> M){
    // smallest poly P such that A^i = sum_{j} < i \ A^j \times P_{j}
    vector<int> rnd1, rnd2:
    mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
    };
    for(int i=0: i<n: i++){
        rnd1.push_back(randint(1, mod - 1));
        rnd2.push_back(randint(1, mod - 1));
    }
    vector<int> gobs;
    for(int i=0; i<2*n+2; i++){</pre>
        int tmp = 0;
        for(int j=0; j<n; j++){</pre>
            tmp += 111 * rnd2[i] * rnd1[i] % mod:
            if(tmp >= mod) tmp -= mod;
        gobs.push_back(tmp);
        vector<int> nxt(n);
        for(auto &i : M){ // sparse matrix * vector
            nxt[i.x] += 111 * i.v * rnd1[i.v] % mod;
            if(nxt[i.x] >= mod) nxt[i.x] -= mod:
       }
        rnd1 = nxt;
    auto sol = berlekamp_massey(gobs);
    reverse(sol.begin(), sol.end());
    return sol:
lint det(int n, vector<elem> M){
    vector<int> rnd:
    mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){
        return uniform int distribution<int>(lb, ub)(rng):
    };
    for(int i=0; i<n; i++) rnd.push_back(randint(1, mod - 1));</pre>
    for(auto &i : M){
       i.v = 111 * i.v * rnd[i.v] % mod;
    }
    auto sol = get_min_poly(n, M)[0];
    if(n \% 2 == 0) sol = mod - sol;
    for(auto &i : rnd) sol = 111 * sol * ipow(i, mod - 2) % mod;
    return sol;
}
```

2 Geometry

2.1 struct Point

```
const double eps = 1e-10:
template <class T>
struct point{
    typedef point P;
   Тх, у;
   point(T x=0, T y=0) : x(x), y(y) {}
   bool operator< (P a) const {return fabs(x-a.x) > eps ? x<a.x : y<a.y;}</pre>
   bool operator == (P a) const {return max(fabs(x-a.x), fabs(y-a.y)) < eps;}
   P operator+ (P a) const {return P(x+a.x, y+a.y);}
   P operator- (P a) const {return P(x-a.x, y-a.y);}
   P operator- () const {return P(-x, -y);};
   T operator* (P a) const {return x*a.x + y*a.y;} // inner prod
   T operator/ (P a) const {return x*a.y - y*a.x;} // outer prod
   T dist2() const {return x*x + y*y;}
   double dist() const {return sqrt(double(dist2()));}
   P perp() const {return P(-y, x);}; // rotate 90 deg ccw
   P mult(T t) const {return P(x*t, y*t);}
   P unit() const {return P(x/dist(), v/dist());}
   P rotate(double a){
       return P(x*cos(a)-y*sin(a), x*sin(a)+y*cos(a));
   }
};
int sgn(double x) {return (x > eps) - (x < -eps);}
typedef point <double > P;
2.2 Distance, Intersection
2.2.1 Point-to-Line
double lineDist(P a, P b, P p) {
   return ((b-a)/(p-a))/(b-a).dist(); // a->b : left (+), right : (-);
2.2.2 Point-to-Segment
double segDist(P s, P e, P p) {
   if(s == e) return (p-s).dist(): // mind the eps
   double d = (e-s).dist2(), t = min(d, max(.0, (p-s)*(e-s)));
   return ((p-s).mult(d)-(e-s).mult(t)).dist() / d;
2.2.3 Line intersection
template<class P>
pair<int, P> lineInter(P a, P b, P c, P d){
   if((b-a)/(d-c) == 0) // parallel, mind the eps
```

return $\{-((b-a)/(c-a) == 0), a\}$;

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```
double oa = (d-c)/(a-c), ob = (d-c)/(b-c);
    return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
} // 1.0.-1(inf) : inter
2.2.4 Segment Intersection
bool onSegment(P s, P e, P p) {return segDist(s, e, p) < eps;}</pre>
template<class P> vector<P> segInter(P a, P b, P c, P d){
    double oa = (d-c)/(a-c), ob = (d-c)/(b-c),
            oc = (b-a)/(c-a), od = (b-a)/(d-a);
    if(sgn(oa)*sgn(ob) < 0 && sgn(oc)*sgn(od) < 0)
        return {(a.mult(ob) - b.mult(oa)).mult(1/(ob-oa))};
    if(onSegment(c, d, a)) S.insert(a);
    if(onSegment(c, d, b)) S.insert(b);
    if(onSegment(a, b, c)) S.insert(c);
    if(onSegment(a, b, d)) S.insert(d);
    return vector<P>(S.begin(), S.end()):
}
```

2.2.5 Circle-Line Intersection

Should be added.

2.3 Convex Hull

```
vector<pll> get_CV(vector<pll> V){
    sort(V.begin(), V.end());
    sort(V.begin() + 1, V.end(), [&](pll x, pll y){
        pll xx = x - V[0];
        pll yy = y - V[0];
        11 \text{ res} = xx / vv:
        if(res != 0) return res > 0;
        if(xx.first != yy.first) return xx.first < yy.first;</pre>
        return xx.second < yy.second;</pre>
    });
    vector<pll> ret;
    for(auto val : V){
        while(ret.size() > 1){
            pll xx = ret[ret.size() - 2] - val;
            pll yy = ret[ret.size() - 1] - val;
            if(xx / yy <= 0) ret.pop_back();</pre>
            else break:
        }
        ret.push_back(val);
    }
```

```
return ret;
2.4 Rotating Calipers
void rotating_calipers(vector<pll> CV){
   int pos = 0;
   for(int i = 0; i < CV.size(); i++) if(CV[pos] < CV[i]) pos = i;
   int ind1 = 0, ind2 = pos:
   11 dist = (CV[ind1] - CV[ind2]) * (CV[ind1] - CV[ind2]);
   auto get_v = [\&](int x) { return CV[(x + 1) \% CV.size()] - <math>CV[x];};
   for(int i = 0 ; i < CV.size() ; i++){</pre>
        pll v = get_v(i);
        while((-v) / get_v(pos) < 0) pos = (pos + 1) % CV.size();
        11 tmp_dist = (CV[pos] - CV[i]) * (CV[pos] - CV[i]);
        if(dist < tmp_dist) {</pre>
           dist = tmp_dist;
           ind1 = i; ind2 = pos;
   }
    printf("%lld %lld %lld %lld\n", CV[ind1].first, CV[ind1].second, CV[ind2].first,
   CV[ind2].second);
2.5 Sorting Points by Angle
// credit : http://koosaga.com/97
auto angle_sort = [&] (P &a, P &b){
   if((a < point(0, 0)) ^ (b < point(0, 0))) return b < a;</pre>
   if (a / b != 0) return a / b > 0;
   return a.dist2() < b.dist2(); // norm</pre>
}; // clockwise sort
2.6 Smallest Enclosing Circle
//Prove By Solving - https://www.acmicpc.net/problem/11930
int main(){
   scanf("%d", &N):
   for(int i = 1; i \le N; i++) scanf("%lf%lf", &A[i].x, &A[i].y, &A[i].z);
   int t = 70000:
   double rate = 1.0;
   point cur = (point)\{0, 0, 0\};
   for(int i = 1 : i \le t : i++){}
        int ind = 1;
       for(int j = 1; j \le N; j++) if((A[j] - cur) * (A[j] - cur) > (A[ind] -
        cur) * (A[ind] - cur)) ind = j;
        cur = cur + (A[ind] - cur) * rate;
```

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```
rate *= 0.99;
}
double r = 0;
for(int i = 1; i <= N; i++) r = max(r, (A[i] - cur) * (A[i] - cur));
cout << sqrt(r);
return 0;
} // Non-deterministic, deterministic O(n lg n) requires Voronoi diagram</pre>
```

2.7 Polygon Area

2.7.1 Polygon Area

```
double ans = 0; // ans : double area
for(int i=0;i<points.size();i++)
   ans += points[i] / points[(i+1 == points.size() ? 0 : i+1)];</pre>
```

3 Strings

3.1 Aho-Corasick Algorithm

```
namespace aho_corasick {
    const int MAXN = 100000, MAXC = 26;
    int trans[MAXN+1][MAXC];
    int fail[MAXN+1];
   bool term[MAXN+1];
    void build(const vector<string> &v) {
       setz(trans), setz(fail), setz(term);
        int cnode = 1;
       repi(s, v) {
           int p = 0;
           repi(j, s) {
               char c = i-'a':
               if (!trans[p][c]) trans[p][c] = cnode++;
                p = trans[p][c];
           }
            term[p] = true;
       }
        queue<int> q; rep(i, MAXC) if (trans[0][i]) q.push(trans[0][i]);
        while(!q.empty()) {
           int t = q.front(); q.pop();
           rep(i, MAXC) {
               if (trans[t][i]) {
                    int p = fail[t];
                    while(p and not trans[p][i]) p = fail[p];
                    p = trans[p][i];
                    fail[trans[t][i]] = p;
                    if (term[p]) term[trans[t][i]] = true;
                    q.push(trans[t][i]);
```

```
}
    bool query(string &t) {
        int p = 0;
        repi(i, t) {
            char c = i-'a';
            while(p and not trans[p][c]) p = fail[p];
            p = trans[p][c];
            if (term[p]) return true;
        }
        return false;
   }
}
3.2 Suffix Array
// str : abracadabra
// SA : 10 7 0 3 5 8 1 4 6 9 2
// LCP : 1 4 1 1 0 3 0 0 0 2
vector<int> make_sa(const string& s) {
    int n = s.length();
   int \lim = \max(128, n+1):
    vector\langle int \rangle sa(n), g(n+1), ng(n+1), cnt(lim), ind(lim+1);
   for(int i=0:i<n:i++) {
        sa[i] = i; g[i] = s[i];
   }
    g[n] = 0;
   for(int t=1;t<s.length();t<<=1)</pre>
        auto cmp = [&] (int a, int b) {
            return g[a] != g[b] ? g[a] < g[b] : g[a+t] < g[b+t];
        }:
        for(int i=0;i<n;i++) cnt[g[min(i+t, n)]]++;</pre>
        for(int i=1;i<lim;i++) cnt[i] += cnt[i-1];</pre>
        for(int i=n-1; i>=0; i--) ind[--cnt[g[min(i+t, n)]]] = i;
        for(int i=0;i<lim;i++) cnt[i] = 0;</pre>
        for(int i=0; i \le n; i++) cnt[g[i]]++; // same as cnt[g[ind[i]]]++
        for(int i=1;i<lim;i++) cnt[i] += cnt[i-1];</pre>
        for(int i=n-1;i>=0;i--) sa[--cnt[g[ind[i]]]] = ind[i];
        ng[sa[0]] = 1;
        for(int i=1;i<n;i++) {</pre>
            ng[sa[i]] = ng[sa[i-1]] + cmp(sa[i-1], sa[i]);
        g = ng;
        fill(cnt.begin(), cnt.end(), 0);
        fill(ind.begin(), ind.end(), 0);
```

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```
}
    return sa;
}
vector<int> make lcp(const string& s. const vector<int>& sa) {
    int n = s.length();
    vector<int> lcp(n-1), rank(n);
    for(int i=0;i<n;i++)</pre>
       rank[sa[i]] = i;
    int len = 0:
   for(int i=0;i<n;i++) {</pre>
       if(rank[i]) {
           int j = sa[rank[i]-1];
           int lc = n - max(i,j);
           while(len < lc && s[i+len] == s[j+len]) len++;</pre>
           lcp[rank[i]-1] = len;
       if(len) len--;
    }
    return lcp;
     Manacher's Algorithm
// 0-based
//s = #h#e#l#l#o#
// ret = 0 1 0 1 0 1 2 1 0 1 0
vector<int> manacher(const string& s) {
    int n = s.size(), r = -1, k = -1;
   vector<int> p(n);
   for (int i=0; i<n; i++) {
       if (i<=r) p[i] = min(r-i, p[2*k-i]);</pre>
       while (i-p[i]-1>=0 and i+p[i]+1<n and s[i-p[i]-1] == s[i+p[i]+1]) p[i]++;
       if (r < i+p[i]) r = i+p[i], k = i;
   }
    return p;
}
3.4 Z Algorithm
// O-based
//s = abcababca
// ret = 9 0 0 2 0 4 0 0 1
vector<int> z_algo(const string &s) {
    int 1 = 0, r = 0, N = sz(s):
    vector<int> Z(N);
   Z[0] = N;
    repp(i, 1, N) {
                                                                                      int scc(int here){
       if (i > r) {
```

```
while(r < N \text{ and } s[r] == s[r-1]) r++;
            Z[i] = r-1+1;
       } else {
            int k = i-1;
            if (Z[k] < r-i+1) Z[i] = Z[k];
            else {
               1 = i;
                while(r < N \text{ and } s[r] == s[r-1]) r++;
               Z[i] = r-l+1;
           }
        }
   }
   return Z;
3.5 Lexicographically Smallest String Rotation
// rotate(v.begin(), v.begin()+min_rotation(v), v.end());
int min_rotation(string s) {
   int a=0, N=sz(s); s += s;
   rep(b,0,N) rep(i,0,N) {
       if (a+i == b \mid | s[a+i] < s[b+i]) \{b += max(0, i-1); break;\}
        if (s[a+i] > s[b+i]) \{ a = b : break : \}
   }
   return a;
4 Graph Theory
4.1 Strongly Connected Component
const int MAXN = 2e5 + 10; // > 2*N
int N. M:
int dfsn[MAXN], low[MAXN], finished[MAXN], cnt;
vector<int> ADJ[MAXN];
vector<vector<int>> G:
stack<int> S;
int f(int x){ // 0 1 2 3 4 5... -> f(1) f(-1) f(2) f(-2) f(3) f(-3)...
   return 2 * (abs(x) - 1) + (x < 0):
void add_edge(int x, int y){ // call by f(x), f(y)
   ADJ[x ^ 1].push_back(y);
    ADJ[y ^ 1].push_back(x);
// memset(finished, -1, sizeof(finished));
```

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```
static vector<int> tmp;
    S.push(here);
    dfsn[here] = low[here] = ++cnt:
    int &ret = low[here];
    for(int there : ADJ[here]){
       if(dfsn[there] == 0) ret = min(ret, scc(there));
        else if(finished[there] == -1) ret = min(ret, dfsn[there]);
    }
    if(dfsn[here] == low[here]){
        while(1){
            int x = S.top(); S.pop();
           finished[x] = G.size();
            tmp.push_back(x);
            if(x == here) break;
       G.push_back(tmp);
        tmp.clear();
    }
    return ret;
}
```

4.1.1 2-SAT

- scc를 실행시켜 f(i) 와 f(-i)가 같은 component에 있다면, 모순.
- f(i) 와 f(-i) 중 finished 배열의 수가 작은 것이 참이다.
 - SCC numbering의 역순이 위상정렬이기에, $F \rightarrow T$ 를 유지하기 위함

4.2 Biconnected Component

```
// https://gist.github.com/koosaga/6f6fd50dd7067901f1b1
void dfs(int x, int p){
    dfn[x] = low[x] = ++piv;
    par[x] = p;
   for(int i=0; i<graph[x].size(); i++){</pre>
        int w = graph[x][i];
        if(w == p) continue;
        if(!dfn[w]){
            dfs(w. x):
            low[x] = min(low[x], low[w]);
        }
        else{
            low[x] = min(low[x], dfn[w]);
    }
}
void color(int x. int c){
    if(c > 0) bcc[x].push_back(c); // c == 0 : first component
```

```
vis[x] = 1;
   for(int i=0; i<graph[x].size(); i++){</pre>
       int w = graph[x][i];
       if(vis[w]) continue;
       if(dfn[x] <= low[w]){</pre>
           bcc[x].push_back(++cpiv);
           color(w, cpiv);
       }
       else{
           color(w, c);
   }
}
4.3 Euler Tour
struct Edge{
   int to, cnt; // to: 인접한 정점, cnt: 남은 사용 횟수
   Edge *dual; // dual: 역방향 간선을 가리키는 포인터
   Edge(): Edge(-1, 0){}
   Edge(int to1, int cnt1): to(to1), cnt(cnt1), dual(nullptr){}
void Eulerian(int curr){
   for(Edge *e: adj[curr]){
       if(e\rightarrow cnt > 0){
           e->cnt--:
           e->dual->cnt--;
           Eulerian(e->to); // dfs
   }
   cout << curr << '\n';
4.4 Heavy-Light Decomposition
int N. M:
vector<int> ADJ[MAXN];
int S[MAXN];
int hld_head[MAXN], color[MAXN], dfsn[MAXN], dcnt, hcnt;
int P[MAXN];
void dfs1(int here, int par){
 S[here] = 1; P[here] = par;
 for(int there : ADJ[here]) if(there != par) dfs1(there, here), S[here] +=
 S[there];
void dfs2(int here, int c){ // dfs reordering
 if(hld_head[c] == 0) hld_head[c] = here;
 dfsn[here] = ++dcnt: color[here] = c:
```

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```
sort(ADJ[here].begin(), ADJ[here].end(), [&](int x, int y){
   return S[x] > S[y];
 }):
  int cnt = 0:
  for(int there : ADJ[here]) if(there != P[here]){
   if(++cnt == 1) dfs2(there, c);
    else dfs2(there, ++hcnt):
}
4.5 Dominator Tree
namespace Dtree {
    const int MAXN = 250001;
    vector<int> E[MAXN], RE[MAXN], rdom[MAXN];
    int S[MAXN], RS[MAXN], cs;
    int par[MAXN], val[MAXN];
    int sdom[MAXN], rp[MAXN];
    int dom[MAXN];
    int Find(int x, int c = 0) {
       if (par[x] == x) return c?-1:x;
       int p = Find(par[x], 1);
       if (p == -1) return c?par[x]:val[x];
       if (sdom[val[x]] > sdom[val[par[x]]]) val[x] = val[par[x]];
       par[x] = p;
       return c?p:val[x];
    }
    void Union(int x, int y) {
       par[x] = y;
    }
    void dfs(int x) {
       RS[S[x] = ++cs] = x:
       par[cs] = sdom[cs] = val[cs] = cs;
       for(int e : E[x]) {
            if (S[e] == 0) dfs(e), rp[S[e]] = S[x];
            RE[S[e]].pb(S[x]);
       }
    }
    int Do(int s, int *up) {
       dfs(s);
       for (int i = cs;i;i--) {
            for (int e : RE[i]) sdom[i] = min(sdom[i], sdom[Find(e)]);
           if (i > 1) rdom[sdom[i]].pb(i);
            for (int e:rdom[i]) {
               int p = Find(e);
               if (sdom[p] == i) dom[e] = i;
```

```
else dom[e] = p;
           }
            if (i > 1) Union(i, rp[i]);
        for (int i = 2: i \le cs: i++) if (sdom[i] != dom[i]) dom[i] = dom[dom[i]]:
       for (int i = 2; i <= cs; i++) {
            up[RS[i]] = RS[dom[i]];
       return cs;
   }
   void addE(int x, int y) {E[x].pb(y);}
4.6 Global Min Cut
// Stoer-Wagner Algorithm, O(VE lg E)
int minimum_cut_phase(int n, int &s, int &t, vector<vector<int>> &adj, vector<int>
vis){
   vector<int> dist(n);
   int mincut = 1e9;
   while(true){
        int pos = -1, cur = -1e9:
       for(int i=0; i<n; i++){</pre>
            if(!vis[i] && dist[i] > cur){
                cur = dist[i]:
                pos = i;
           }
        if(pos == -1) break;
        s = t:
       t = pos;
       mincut = cur;
        vis[pos] = 1:
       for(int i=0; i<n; i++){
           if(!vis[i]) dist[i] += adi[pos][i]:
   return mincut: // optimal s-t cut here is, {t} and V \ {t}
int solve(int n, vector<vector<int>> adj){
   if(n <= 1) return 0;
   vector<int> vis(n):
   int ans = 1e9:
   for(int i=0; i<n-1; i++){
       int s, t;
        ans = min(ans, minimum_cut_phase(n, s, t, adj, vis));
       vis[t] = 1;
       for(int j=0; j<n; j++){</pre>
           if(!vis[i]){
                adi[s][i] += adi[t][i];
```

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```
adj[j][s] += adj[j][t];
}
adj[s][s] = 0;
}
return ans;
}
```

5 Network Flow

5.1 Theorems

Max-flow Min-cut theorem : 정점 s에서 정점 t까지 흐를 수 있는 최대 유량(max-flow)은 정점 s와 정점 t를 분리하는 간선들의 가중치 합(min-cut)과 같다.

Vertex cover: 어떤 그래프의 정점의 집합 S에 대해 그래프의 모든 간선이 S의 원소 중 최소 하나와 연결되어 있을 때, S를 해당 그래프의 vertex cover라고 하며, minimum vertex cover는 최소 개수의 정점을 사용한 vertex cover이다.

Independent set : 어떤 그래프의 정점의 집합 S에 대해 S의 서로 다른 두 정점을 연결하는 간선이 없을 때, S를 해당 그래프의 independent set이라고 하며, maximum independent set은 최대 개수의 정점을 사용한 independent set이다.

Matching (independent edge set) : 어떤 그래프의 간선의 집합 S에 대해 S의 서로 다른 두 간선이 공통된 정점을 가지지 않을 때, S를 해당 그래프의 matching이라고 하며, maximum matching은 최대 개수의 간선을 사용한 matching이다.

König's theorem : 이분 그래프의 maximum matching의 크기는 minimum vertex cover의 것과 같다. Dinic's Algorithm : 시간 복잡도 $O(V^2E)$, unit capacity에서는 $\min(V^{2/3}E, E^{3/2})$.

Circulation Problem : 새로운 source/sink s_n , t_n 를 만들어서 다음과 같이 간선을 추가하고 $maxflow(s_n \to t_n) = \sum l_i$ 인지 확인, 이후 $s \to t$ 로 maxflow

• $s_n \to b$ (l), $a \to t_n$ (l), $a \to b$ (r-l), $t \to s$ (∞)

5.2 Dinic's Algorithm

```
const int INF = 1e9:
struct Dinic{
    int N;
    struct edge{
       int index, cap, rev;
        edge(): index(0), cap(0), rev(0) {}
        edge(int index, int cap, int rev) : index(index), cap(cap), rev(rev) {}
   };
    vector<vector<edge>> ADJ;
    vector<int> R, W;
   Dinic() {}
   Dinic(int N) : N(N){
        ADJ.resize(N): R.resize(N):
                                        W.resize(N):
    }
    void CE(int node1, int node2, int cap){
        ADJ[node1].push_back(edge(node2, cap, ADJ[node2].size()));
```

```
ADJ[node2].push_back(edge(node1, 0, ADJ[node1].size() - 1));
   }
   bool bfs(int src, int sink){
       fill(R.begin(), R.end(), -1):
       R[src] = 0;
       queue<int> Q; Q.push(src);
        while(Q.size()){
            int here = Q.front(); Q.pop();
           for(auto e : ADJ[here]){
                if(e.cap > 0 && R[e.index] == -1)
                   R[e.index] = R[here] + 1, Q.push(e.index);
           }
       }
        return R[sink] != -1;
   }
   int dfs(int here, int sink, int f){
        if(here == sink) return f:
       for(int &i = W[here] ; i < ADJ[here].size() ; i++){</pre>
            auto &e = ADJ[here][i]:
           if(e.cap > 0 && R[here] < R[e.index]){
                int res = dfs(e.index, sink, min(f, e.cap));
                if(res) {
                    e.cap -= res;
                    ADJ[e.index][e.rev].cap += res;
                    return res:
           }
        }
        return 0;
   }
   int solve(int src, int sink){
       int ret = 0:
        while(bfs(src, sink)){
           fill(W.begin(), W.end(), 0);
            int res;
            while((res = dfs(src, sink, INF))) ret += res;
       }
        return ret;
   }
};
5.3 MCMF with SPFA
const int INF = 1e9:
struct MCMF {
   struct edge {
       int there, cap, cost, rev;
```

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```
edge(): there(0), cap(0), cost(0), rev(0) {}
    edge(int there, int cap, int cost, int rev) : there(there), cap(cap),
   cost(cost). rev(rev) {}
};
int N;
vector<vector<edge>> ADJ;
vector<int> R, INQ, C, I;
MCMF(): N(0) \{ \}
MCMF(int N) : N(N) { ADJ.resize(N + 1); R.resize(N + 1); INQ.resize(N + 1);
C.resize(N + 1); I.resize(N + 1); }
void CE(int i, int j, int cap, int cost) {
    ADJ[i].push_back(edge(j, cap, cost, ADJ[j].size()));
    ADJ[j].push_back(edge(i, 0, -cost, ADJ[i].size() - 1));
}
bool SPFA(int src, int sink) {
    queue<int> Q;
                    Q.push(src);
   fill(R.begin(), R.end(), -1);
                                    R[src] = 0;
   fill(C.begin(), C.end(), -1);
                                    C[src] = 0:
   fill(INQ.begin(), INQ.end(), 0); INQ[src] = 1;
    while (Q.size()) {
        int here = Q.front(); Q.pop();
        INQ[here] = 0;
        for (int i = 0: i < ADJ[here].size(): i++) {</pre>
           auto e = ADJ[here][i]:
           if (e.cap > 0 && (C[e.there] == -1 || C[e.there] > C[here] +
           e.cost)) {
                C[e.there] = C[here] + e.cost; R[e.there] = here;
                I[e.there] = i;
                if (!INQ[e.there]) INQ[e.there] = 1, Q.push(e.there);
           }
        }
    if (C[sink] == -1) return false;
   return true:
}
pii mcmf(int src, int sink) {
   pii ret = { 0, 0 };
   while (SPFA(src, sink)) {
        int flow = INF. cost = 0:
        for (int here = sink; here != src; here = R[here]) flow = min(flow,
        ADJ[R[here]][I[here]].cap);
        for (int here = sink; here != src; here = R[here]) {
           auto &e = ADJ[R[here]][I[here]];
           cost += e.cost * flow:
           e.cap -= flow;
            ADJ[e.there][e.rev].cap += flow;
```

```
ret.first += flow, ret.second += cost;
        return ret;
   }
};
5.4 Hungarian Method
namespace Hung {
    const int MX = 2000;
    // IMPORTANT : n <= m, 1-based
    using T = long double;
   T \max v = 1e200:
   T a[MX][MX], n, m;
   void init(int nn, int mm) { n = nn; m = mm; }
    void set_value(int x, int y, T val) { a[x][y] = val; }
   T solve(vector <int> &ans) {
        vectorT> v(m+1), u(n+1);
        vector\langle int \rangle p (m+1), way (m+1);
        for (int i=1: i<=n: ++i) {
            p[0] = i;
            int j0 = 0;
            vector<T> minv (m+1, maxv):
            vector<char> used (m+1, false);
            do {
                used[j0] = true;
                T delta = maxv:
                int i0 = p[j0], j1;
                for (int j=1; j<=m; ++j) if (!used[j]) {</pre>
                    T cur = a[i0][j]-u[i0]-v[j];
                    if (cur < minv[i]) {
                        minv[j] = cur, way[j] = j0;
                    }
                    if (minv[j] < delta) {</pre>
                        delta = minv[j], j1 = j;
                    }
                for (int j=0; j<=m; ++j) {</pre>
                    if (used[i]) {
                        u[p[j]] += delta, v[j] -= delta;
                    }
                    else {
                        minv[j] -= delta;
                }
                j0 = j1;
            } while (p[j0] != 0);
            do {
```

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```
p[j0] = p[j1];
               j0 = j1;
            } while (j0);
        ans.resize(n + 1);
        for(int j=1;j<=m;++j) {</pre>
            ans[p[j]] = j;
       return -v[0];
    }
}
      Hopcroft-Karp Algorithm
struct hopcroft_karp{
    int N:
    vector<vector<int>> ADJ;
    vector<int> L, rev, used;
    hopcroft_karp() {}
    hopcroft_karp(int N) : N(N) {
       ADJ.resize(N):
       L.resize(N), rev.resize(N, -1), used.resize(N, 0);
    }
    void CE(int here, int there){
        ADJ[here].push_back(there);
    }
    void bfs(){
        queue<int> Q;
       for(int i = 0; i < N; i++) {
           if(used[i]) L[i] = -1:
            else L[i] = 0, Q.push(i);
       }
        while(Q.size()){
            int here = Q.front(); Q.pop();
            for(int there : ADJ[here]){
                if(rev[there] != -1 && L[rev[there]] == -1) {
                    L[rev[there]] = L[here] + 1;
                    Q.push(rev[there]);
               }
            }
       }
    }
    bool dfs(int here){
       for(int there : ADJ[here]){
            if(rev[there] == -1 || (L[here] < L[rev[there]] && dfs(rev[there]))){
```

int j1 = wav[j0];

```
rev[there] = here;
                  used[here] = 1;
                  return true;
         return false;
    }
    int solve(){
         int ret = 0;
         while(1){
             bfs();
             int res = 0;
             for(int i = 0 ; i < N ; i++) {
                  if(used[i]) continue;
                  res += dfs(i):
             if(res == 0) break;
             ret += res:
        return ret;
    }
};
    Optimization Tricks
6.1 Knuth Optimization
  • Recurrence : D[i][j] = \min_{i < k < j} (D[i][k] + D[k][j]) + C[i][j]
  • Quadrangle Inequality : C[a][c] + C[b][d] \le C[a][d] + C[b][c], \ a \le b \le c \le d
  • Monotonicity: C[b][c] \le C[a][d], \ a \le b \le c \le d
  • A[i][j] = (\min k \ s.t. \ D[i][j] \ \text{is min.}). Then A[i][j-1] \le A[i][j] \le A[i+1][j]
  • O(N^2) time complexity
// opt[i-1][i] = i
for(int d=2;d<=n;d++) {</pre>
    for(int i=1:i+d<=n+1:i++) {
        for(int k=opt[i][j-1], j=i+d; k<=opt[i+1][j]; k++) {</pre>
             int v = dp[i][k] + dp[k][j] + c[i][j];
             if(dp[i][j] > v) dp[i][j] = v, opt[i][j] = k;
    }
```

6.2 Divide and Conquer Optimization

• Recurrence : $D[t][i] = \min_{k < i} (D[t-1][k] + C[k][i])$

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- Quadrangle Inequality : $C[a][c] + C[b][d] \le C[a][d] + C[b][c], \ a \le b \le c \le d$
- Able to Divide and Conquer base on calculating D[t][i]
- $O(TN \lg N)$ time complexity

```
// range of index : [1,r]
// range of dp : [s,e]
void dnc(int t, int 1, int r, int s, int e)
{
    if(s > e) return;
    int m = (s+e)/2;
    D[t][m] = 2e9;
    for(int k=1;k<m&&k<=r;k++)
    {
        int tmp = D[t-1][k] + C[k][m];
        if(D[t][m] > tmp)
            D[t][m] = tmp, A[t][m] = k;
    }
    dnc(t, 1, A[t][m], s, D[t][m]-1);
    dnc(t, A[t][m], r, D[t][m]+1, e);
}
```

6.3 Convex Hull Trick

- Recurrence: $dp[i] = \min_{j < i} (dp[j] + a[i]b[j]), b[i-1] \le b[i]$
- Think as $dp[x = a[i]] = \min_{i \le i} (b[i] \cdot x + dp[i])$
- Thus push lines and find minimum (by binary search)
- If a[i] < a[i+1] sweeping is possible
- Intersection of $y=a_ix+b_i$ and $y=a_{i+1}x+b_{i+1}$: $x=\frac{b_{i+1}-b_i}{a_i-a_{i+1}}$

6.4 Centroid Decomposition

```
// credit : https://gist.github.com/igorcarpanese/75162f3253bd230abd0f32f9950bf384
int dfs(int u, int p) {
    sub[u] = 1;
    for (auto v : tree[u])
        if (v != p) sub[u] += dfs(v, u);
    return sub[u] + 1;
}

// each tree has at most two centroids
int centroid(int u, int p, int r) { // r : root
    for (auto v : tree[u])
        if (v != p and sub[v] > sub[r]/2) return centroid(v, u);
    return u;
}
```

7 Data Structure

7.1 Persistent Segment Tree

```
const MAXN = 1e5 + 10;
struct node{
   node *1, *r;
    int cnt;
   node () {}
} pool[(1 << 17) * 17], *tree_head[MAXN];</pre>
int tcnt:
node* alloc(){
    memset(pool + tcnt, 0, sizeof(node));
   return pool + tcnt++;
node * init(int 1, int r){
    node *ret = alloc();
   if(1 != r) {
        int mid = (1 + r) / 2;
        ret->l = init(l, mid);
        ret->r = init(mid + 1, r):
   }
   return ret;
void update(node * here, node *par, int 1, int r, int val){
    if(1 == r) {
        here->cnt = par->cnt + 1;
        return;
   }
    int mid = (1 + r) / 2;
    if(val <= mid){</pre>
        here->1 = alloc():
        here->r = par->r;
        update(here->1, par->1, 1, mid, val);
   }
   else {
        here > 1 = par > 1;
        here->r = alloc();
        update(here->r, par->r, mid + 1, r, val);
   }
   here->cnt = here->l->cnt + here->r->cnt;
int query(node *node1, node *node2, int 1, int r, int k){
    if(l == r) return 1;
    int ccc = node1->l->cnt - node2->l->cnt;
    int mid = (1 + r) / 2;
```

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```
if(k <= ccc) return query(node1->1, node2->1, 1, mid, k);
    else return query(node1->r, node2->r, mid + 1, r, k - ccc);
}
7.2 Link-Cut Tree
struct node{
   node *pp, *p, *l, *r;
    int val;
   node(){p = 0, 1 = 0, r = 0;}
    node(int val) : val(val) { p = 0, 1 = 0, r = 0;}
};
void push(node *x){}
void pull(node *x){}
void rotate(node *x){
    if(!x->p) return;
    push(x->p); // if there's lazy stuff
   push(x);
    node *p = x->p;
    bool is_left = (p->1 == x);
   node *b = (is_left ? x->r : x->1);
    x->p = p->p;
    if (x-p \&\& x-p-1 == p) x-p-1 = x;
    if(x-p \&\& x-p-r == p) x-p-r = x;
    if(is_left){
        if(b) b \rightarrow p = p;
        p->1 = b;
       p->p = x;
       x->r = p;
    }
    else{
        if(b) b->p = p;
        p->r = b;
       p->p = x;
       x->1 = p;
    pull(p); // if there's something to pull up
   pull(x);
    //if(!x->p) root = x; // IF YOU ARE SPLAY TREE
    if(p->pp){ // IF YOU ARE LINK CUT TREE
        x->pp = p->pp;
        p->pp = nullptr;
    }
}
void splay(node *x){
    while(x->p){
       node *p = x->p;
        node *g = p->p;
        if(g){
```

```
if((p->l == x) ^ (g->l == p)) rotate(x);
            else rotate(p);
        }
        rotate(x);
   }
void access(node *x){
    splay(x);
    push(x);
   if(x->r){
        x->r->pp = x;
        x->r->p = nullptr;
        x->r = nullptr;
   }
   pull(x);
   while(x->pp){
        node *nxt = x->pp;
        splay(nxt);
        push(nxt);
        if(nxt->r){
           nxt->r->pp = nxt;
            nxt->r->p = nullptr;
            nxt->r = nullptr;
        }
        nxt->r = x;
        x->p = nxt;
        x->pp = nullptr;
        pull(nxt);
        splay(x);
   }
node *root(node *x){
   access(x);
    while (x->1) {
        push(x);
        x = x->1;
   }
   access(x);
   return x;
node *par(node *x){
   access(x);
    if(!x->1) return nullptr;
   push(x);
   x = x->1;
   while(x->r){
        push(x);
        x = x->r;
   }
   access(x);
```

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```
if (z == end()) return y > m == x - m && y - b <= x - b;
    return x;
}
                                                                                                return (x->b - y->b)*(z->m - y->m) >= (y->b - z->b)*(y->m - x->m);
node *lca(node *s. node *t){
                                                                                           }
    access(s):
                                                                                           void insert_line(ll m, ll b) {
    access(t):
                                                                                                auto v = insert({ m, b }):
    splay(s);
                                                                                                y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
    if(s->pp == nullptr) return s;
                                                                                                if (bad(y)) { erase(y); return; }
    return s->pp;
                                                                                                while (next(y) != end() && bad(next(y))) erase(next(y));
                                                                                                while (y != begin() && bad(prev(y))) erase(prev(y));
void link(node *par, node *son){
                                                                                           }
    access(par);
                                                                                           11 eval(ll x) {
    access(son);
                                                                                                auto 1 = *lower_bound((Line) { x, is_query });
    //son->rev ^= 1; // remove if needed
                                                                                                return 1.m * x + 1.b:
                                                                                           }
    push(son);
    son->1 = par;
                                                                                        };
    par->p = son;
    pull(son);
                                                                                        7.4 Stern-Brocot Tree
                                                                                       // __int128 is recommended
void cut(node *p){
    access(p);
                                                                                       bool test(11 a, 11 b) { // for testing directions, vary by prob
                                                                                           // return true if (true value) >= a/b
    push(p);
    if(p->1){
                                                                                           11 n = 0, m = 1;
       p->l->p = nullptr;
       p->1 = nullptr;
                                                                                           rep(i, N) {
                                                                                                if (n < m*A[i].fi) n = A[i].fi, m = 1;
    }
    pull(p);
}
                                                                                                11 c = b*n+m*a, d = m*b;
                                                                                                11 g = gcd(c, d);
                                                                                                n = c/g;
     Dynamic Convex Hull
                                                                                                m = d/g;
// https://github.com/niklasb/contest-algos/blob/master/convex_hull/dynamic.cpp
const ll is_query = -(1LL<<62);</pre>
                                                                                                if (n > m*A[i].se) return false;
                                                                                           }
struct Line {
   11 m. b:
                                                                                           return true:
    mutable function<const Line*()> succ;
                                                                                        }
   bool operator<(const Line& rhs) const {</pre>
        if (rhs.b != is_query) return m < rhs.m;</pre>
                                                                                       pair<11, 11> stern_brocot(11 M, 11 N) {
                                                                                           // M : max value
        const Line* s = succ();
        if (!s) return 0:
                                                                                           // N : max divisor
                                                                                           // if result is a/b, return as {a, b}
       11 x = rhs.m;
        return b - s->b < (s->m - m) * x:
                                                                                           11 a = 0, b = 1; // 1
   }
                                                                                           11 c = 1, d = 0; // r
};
                                                                                            int 1. r:
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum
                                                                                           bool chg = true;
    bool bad(iterator y) {
        auto z = next(y);
                                                                                            while(chg) {
        if (v == begin()) {
                                                                                                chg = false:
           if (z == end()) return 0;
            return y->m == z->m && y->b <= z->b;
                                                                                                // to left
        }
                                                                                                1 = 0, r = (N-d-1)/b+1:
                                                                                                while(1 < r)  {
        auto x = prev(y);
```

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```
int mid = (1+r+1)/2;
            if (test(a*mid+c, b*mid+d)) r = mid-1;
            else 1 = mid:
        c += a*1;
        d += b*1;
        chg |= (1 > 0);
       // to right
        1 = 0, r = (d?(N-b-1)/d+1:M);
        while(1 < r)  {
            int mid = (1+r+1)/2;
            if (test(a+mid*c, b+mid*d)) l = mid;
            else r = mid-1;
        }
        a += c*l:
        b += d*1:
        chg |= (1 > 0);
    }
    return {a, b};
}
```

7.5 Rope

```
#include <bits/stdc++.h>
#include <ext/rope>
using namespace std;
using namespace __gnu_cxx;

int main()
{
   ios::sync_with_stdio(false);
   cin.tie(0);

   crope rp; // rope<char>
   string s("Lorem-ipsum");
   int n = s.length();
   rp.append(s.c_str()); // add element

   int x = 3, y = 8; // split and merge below
   rp = rp.substr(x, y-x) + rp.substr(0, x) + rp.substr(y, n);
   cout << rp.at(0) << '\n'; // get element, 'e'
   cout << rp << '\n'; // print, "em-ip|Lor|sum"
}</pre>
```

7.6 Bitset

```
#include <bitset>
#include <iostream>
using namespace std;
int main() {
   bitset<8> b1(13);
                                    // 00001101
   bitset<8> b2("10111");
                                    // 00010111
                                    // 3
   cout << b1.count() << endl;</pre>
   cout << b1.test(6) << endl:</pre>
                                    // 0, since 2^6-th bit is 0
   b1.set(6):
                                    // set to 1, 1-fill if no param
   b2.reset(2);
                                    // set to 0, 0-fill if no param
   // use 'flip' for flipping
   cout << "b1:" << b1 << endl;
                                            // b1:01001101
   cout << "b2:" << b2 << endl:
                                            // b2:00010011
   // use any, none, all (c++11) for bit checking
   // supported operators : &, |, ^, <<, >>, ~, ==, !=
   // these operators must match size (given to template)
   cout << "&: " << (b1 & b2) << endl;
                                            // &: 00000001
   cout << "^: " << (b1 ^ b2) << endl;
                                            // &: 01011110
   cout << "|: " << (b1 | b2) << endl;
                                            // |: 01011111
   cout << "~: " << (~b1) << endl:
                                            // ~: 10110010
                                            // <<:01101000
   cout << "<<:" << (b1 << 3) << endl;
   cout << b2.to_ulong() << endl;</pre>
                                            // 19, c++11 supports ullong
```

7.7 Policy Based Data Structure

```
#include <bits/stdc++.h>
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
using namespace std;
using namespace __gnu_pbds;
typedef tree<
int.
null_type,
less<int>,
rb_tree_tag,
tree_order_statistics_node_update >
ordered set:
// less<int> : not allow for duplicate
// less_equal<int> : allow for duplicate
// use upper_bound when you erase from set used less_equal
```

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```
int N;
int main(void) {
    iostream::svnc with stdio(false):
    cin.tie(nullptr);
    ordered_set X;
    X.insert(1);
    X.insert(2);
    X.insert(4);
    X.insert(8):
    X.insert(16);
    cout<<*X.find_by_order(1)<<endl; // 2</pre>
    cout<<*X.find_by_order(2)<<endl; // 4</pre>
    cout<<*X.find_by_order(4)<<endl; // 16</pre>
    cout<<(end(X)==X.find_by_order(6))<<endl; // true</pre>
    cout<<X.order of kev(-5)<<endl: // 0
    cout<<X.order_of_key(1)<<endl; // 0</pre>
    cout<<X.order_of_key(3)<<endl; // 2</pre>
    cout<<X.order_of_key(4)<<endl; // 2</pre>
    cout<<X.order_of_key(400)<<endl; // 5</pre>
}
```

8 Miscellaneous

8.1 Misc Formulae and Algorithms

8.1.1 Faulhaber's Formula

$$T(n,k) = \sum_{i=1}^{n} i^{k} = \frac{(n+1)^{k+1} - 1^{k+1} - \sum_{j=0}^{k-1} {k+1 \choose j} T(n,j)}{{k+1 \choose k}}$$

Also use

$$(x+1)^d - x^d = 1 + {d \choose 1}x + {d \choose 2}x^2 + \dots + {d \choose d-1}x^{d-1}$$

to get each coef.

8.1.2 Maximum Clique

```
typedef long long ll;
ll G[40]; // 0-index
int N, M;
int cur;
void get_clique(int R = 0, ll P = (1ll<<N)-1, ll X = 0){
   if((P|X) == 0){
      cur = max(cur, R);
}</pre>
```

```
}
   int u = __builtin_ctzll(P|X);
   11 c = P\&^{G}[u];
   while(c){
        int v = __builtin_ctzll(c);
        get_clique(R + 1, P&G[v], X&G[v]);
       P ^= 111 << v:
       X = 111 << v;
        c ^= 111 << v:
   }
}
8.1.3 De Brujin Sequence
// https://github.com/koosaga/DeobureoMinkyuParty/blob/master/teamnote.tex
// alphabet = [0, k - 1], substr length n, res starts with 0 (cyclic)
int res[10000000], aux[10000000]; // >= k^n, k*n
int de_bruijn(int k, int n, int lim) { // returns size (k^n)
   if(k == 1) {
       res[0] = 0;
       return 1;
   }
   for(int i = 0: i < k * n: i++) aux[i] = 0:
   int sz = 0;
   function<void(int, int)> db = [&](int t, int p) {
       if(sz > lim) return;
       if(t > n) {
           if(n \% p == 0)
               for(int i = 1; i <= p; i++)
                   res[sz++] = aux[i]:
       }
        else {
           aux[t] = aux[t - p];
           db(t + 1, p);
           for(int i = aux[t - p] + 1; i < k; i++) {
                aux[t] = i:
                db(t + 1, t);
           }
       }
   };
   db(1, 1);
    return sz;
8.2 Highly Composite Numbers, Large Prime
    < 10<sup>k</sup>
                              divisors 2 3 5 71113171923293137
                                     4 1 1
```

12 2 1 1

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3	840	32	3 1 1 1
4	7560	64	3 3 1 1
5	83160	128	3 3 1 1 1
6	720720	240	4 2 1 1 1 1
7	8648640	448	6 3 1 1 1 1
8	73513440	768	5 3 1 1 1 1 1
9	735134400	1344	6 3 2 1 1 1 1
10	6983776800	2304	5 3 2 1 1 1 1 1
11	97772875200	4032	6 3 2 2 1 1 1 1
12	963761198400	6720	6 4 2 1 1 1 1 1 1
13	9316358251200	10752	6 3 2 1 1 1 1 1 1 1
14	97821761637600	17280	5 4 2 2 1 1 1 1 1 1
15	866421317361600	26880	6 4 2 1 1 1 1 1 1 1 1
16	8086598962041600	41472	8 3 2 2 1 1 1 1 1 1 1
17	74801040398884800	64512	6 3 2 2 1 1 1 1 1 1 1 1
18	897612484786617600	103680	8 4 2 2 1 1 1 1 1 1 1 1

	< 10^k prime	> 10^k prime	<pre># of prime</pre>
1	7	11	4
2	97	101	25
3	997	1009	168
4	9973	10007	1229
5	99991	100003	9592
6	999983	1000003	78498
7	9999991	10000019	664579
8	99999989	10000007	5761455
9	99999937	100000007	50847534

	< 10^k prime	> 10 ^k prime
10	999999967	1000000019
11	9999999977	10000000003
12	99999999999	100000000039
13	999999999971	1000000000037
14	9999999999973	10000000000031
15	99999999999999	100000000000037
16	99999999999937	10000000000000061
17	999999999999997	100000000000000003
18	9999999999999989	1000000000000000003

NTT Prime:

```
\begin{aligned} &469762049=7\times 2^{26}+1. \text{ Primitive root}: 3.\\ &998244353=119\times 2^{23}+1. \text{ Primitive root}: 3.\\ &985661441=235\times 2^{22}+1. \text{ Primitive root}: 3.\\ &1012924417=483\times 2^{21}+1. \text{ Primitive root}: 5. \end{aligned}
```

Primes near 10^9 : $10^9 + [7, 9, 21, 33, 87]$

8.3 Fast Integer IO

```
// credit : https://github.com/koosaga/DeobureoMinkyuParty/blob/master/teamnote.tex
static char buf[1 << 19]; // size : any number geq than 1024
static int idx = 0:
static int bytes = 0;
static inline int _read() {
    if (!bytes || idx == bytes) {
        bytes = (int)fread(buf, sizeof(buf[0]), sizeof(buf), stdin);
       idx = 0:
   }
   return buf[idx++]:
static inline int _readInt() {
   int x = 0, s = 1:
   int c = _read();
   while (c \le 32) c = read();
   if (c == '-') s = -1, c = _read();
   while (c > 32) x = 10 * x + (c - '0'), c = _read():
   if (s < 0) x = -x;
   return x:
8.4 C++ Tips / Environments
#include <bits/stdc++.h> // magic header
using namespace std; // magic namespace
struct StupidGCCCantEvenCompileThisSimpleCode{
    pair<int, int> array[1000000];
}; // https://gcc.gnu.org/bugzilla/show_bug.cgi?id=68203
// how to use rand (in 2017)
mt19937 rng(0xdeadbeef):
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
int randint(int lb. int ub){ return uniform int distribution<int>(lb. ub)(rng): }
shuffle(permutation.begin(), permutation.end(), rng);
mt19937_64 _R(chrono::steady_clock::now().time_since_epoch().count()); // _R()
// comparator overload
auto cmp = [](seg a, seg b){return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
int x, y;
bool operator==(const point &p)const{ return x == p.x && y == p.y; }
}:
struct hasher {
```

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```
size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
};
unordered_map<point, int, hasher> hsh;
// c++ setprecision example
                       // std::cout, std::fixed
#include <iostream>
#include <iomanip>
                        // std::setprecision
int main () {
   double f =3.14159;
    std::cout << std::setprecision(5) << f << '\n'; // 3.1416
    std::cout << std::setprecision(9) << f << '\n'; // 3.14159
    std::cout << std::fixed;</pre>
    std::cout << std::setprecision(5) << f << '\n'; // 3.14159
    std::cout << std::setprecision(9) << f << '\n'; // 3.141590000
    return 0;
}
```