Single chain completeness and some related properties

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joint work with

Stefano Aguzzoli

$\mathsf{FL}_\mathsf{ew}\text{-}\mathsf{algebras}$ and $\mathsf{MTL}\text{-}\mathsf{algebras}$

An FL_ew algebra is an algebra $\langle A, *, \to, \wedge, \vee, 0, 1 \rangle$. such that:

- **4** $\langle A, \wedge, \vee, 0, 1 \rangle$ is a bounded lattice with minimum 0 and maximum 1.
- (A, *, 1) is a commutative monoid.
- **(a)** $\langle *, \rightarrow \rangle$ forms a *residuated pair*: $z * x \le y$ iff $z \le x \to y$ for all $x, y, z \in A$. In particular, it holds that $x \to y = \max\{z \in A : z * x \le y\}$.
- 4 An MTL-algebra is an FL_{ew}-algebra satisfying

(Prelinearity)
$$(x \to y) \lor (y \to x) = 1.$$

- The class of MTL-algebras forms a variety, called MTL. The logic corresponding to MTL-algebras is called MTL.
- An axiomatic extension of MTL is a logic obtained by adding other axioms to it.
- Every axiomatic extension of MTL is algebraizable in the sense of [Blok and Pigozzi, 1989], and hence every subvariety of MTL induces a logic.

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- L enjoys the single chain completeness (SCC) if there is an L-chain such that L is complete w.r.t. it.
- L enjoys the finite strong single chain completeness (FSSCC) if there is an L-chain such that L is finitely strongly complete w.r.t. it.
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Theorem ([Montagna, 2011])

Let L be an axiomatic extension of MTL. If L enjoys the FSSCC, then the SSCC holds.

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Problem ([Montagna, 2011])

Does the SCC implies the SSCC, in general?

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- L has the deductive Maksimova's variable separation property (DMVP) if, for all sets of formulas $\Gamma \cup \{\varphi\}$ and $\Sigma \cup \{\psi\}$ that have no variables in common, $\Gamma, \Sigma \vdash_{L} \varphi \lor \psi$ implies $\Gamma \vdash_{L} \varphi$ or $\Sigma \vdash_{L} \psi$.

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Problem ([Galatos et al., 2007])

Find examples of substructural logics with DMVP, but without MVP.

The equivalence between HC and SCC

An FL_{ew}-algebra is said to be well connected whenever for every pair of elements x, y, if $x \sqcup y = 1$, then x = 1 or y = 1.

Theorem ([Galatos et al., 2007, Theorem 5.28])

Let L be an axiomatic extension of FL_{ew}. The following are equivalent:

- L has the Halldén completeness.
- 2 L is complete w.r.t. a well-connected FL_{ew}-algebra.
- L is meet irreducible (in the lattice of axiomatic extensions of FL_{ew}).

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Theorem

Let L be a logic over MTL. The following are equivalent:

- 1 L has the Halldén completeness.
- 2 L is complete w.r.t. an MTL-chain, that is L enjoys the SCC.
- **1** L is meet irreducible (in the lattice of axiomatic extensions of MTL).

Theorem ([Kihara, 2006, Theorem 6.9])

The following conditions are equivalent for every axiomatic extension L of MTL:

- L has the DMVP.
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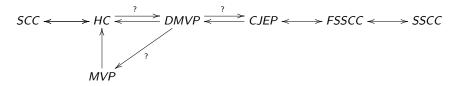
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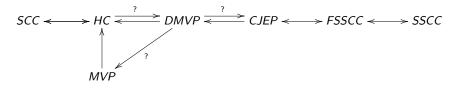
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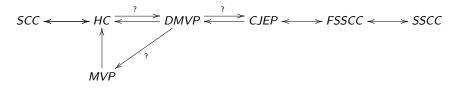
Corollary

Every axiomatic extension of MTL enjoying the SSCC has also the DMVP.



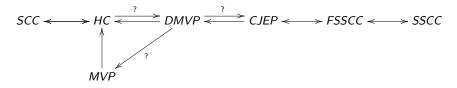


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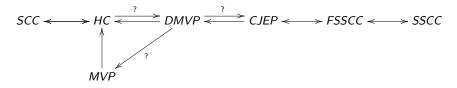
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- Does exist an axiomatic extension of MTL having the DVMP, but for which the MVP fails?
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The problem of the equivalence between DMVP and SSCC

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Definition

Let us call Q the MTL-chain (WNM-chain) $\langle \{0, a, b, c, 1\}, *, \Rightarrow, \min, \max, 0, 1 \rangle$, with 0 < a < b < c < 1 and such that:

X	$\sim x$
0	1
а	С
Ь	а
С	а
1	0

$$x*y = \begin{cases} 0 & \text{if } x \le \sim y \\ \min\{x,y\} & \text{otherwise.} \end{cases}$$

Let us call \mathbb{Q} the variety generated by \mathcal{Q} , and let \mathbb{Q} be the corresponding logic.

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- Notice that the three element Gödel-chain G_3 belongs to \mathbb{Q} . Just take the quotient of \mathcal{Q} over the filter $\{c,1\}$.
- Observe that $G_3 \not\hookrightarrow \mathcal{Q}$.
- From these fact it is easy to check that there is no Q-chain $\mathcal C$ such that $\mathcal Q \hookrightarrow \mathcal C$ and $\mathbf G_3 \hookrightarrow \mathcal C$. Hence the DMVP fails.



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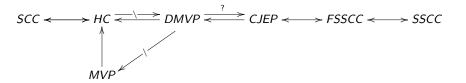
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Theorem

Let A be a non-trivial MTL-chain with less than five elements, and let $\mathbb{L} = \mathbf{V}(A)$. Then L enjoys the SSCC.

A general picture



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Extensions of MTL expanded with Δ

For every axiomatic extension L of MTL, we denote with L_{Δ} its expansion with an operator Δ satisfying the following axioms:

$$\Delta(\varphi) \vee \neg \Delta(\varphi).$$

$$(\Delta 2)$$
 $\Delta(\varphi \lor \psi) \to (\Delta(\varphi) \lor \Delta(\psi)).$

$$\Delta(\varphi) \to \varphi$$
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An MTL $_{\Delta}$ -chain is an MTL-chain expanded with an operation δ , interpreting Δ , such that, for every element x, $\delta(x)=1$ if x=1, whilst $\delta(x)=0$ if x<1.

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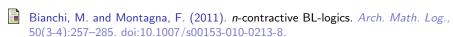
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- Do MTL and IMTL enjoy the SCC?
- How about the first-order case? Are there some logical properties characterizing the SCC and the SSCC? Note that the SCC for an axiomatic extension L of MTL do not necessarily implies the SCC for L∀. As shown in [Montagna, 2011] BL is a counterexample.

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APPENDIX

Semihoops and hoops

Definition

A semihoop is a structure $\mathcal{A} = \langle A, *, \sqcap, \Rightarrow, 1 \rangle$ such that $\langle A, \sqcap, 1 \rangle$ is an inf-semilattice with upper bound 1, * is a binary operation on A with unit 1, and \Rightarrow is a binary operation such that:

- x < y iff $x \Rightarrow y = 1$,
- \bullet $(x * y) \Rightarrow z = x \Rightarrow (y \Rightarrow z).$

A bounded semihoop is a semihoop with a minimum element; conversely, an unbounded hoop is a hoop without minimum.

- A hoop is a semihoop satisfying $x * (x \Rightarrow y) = y * (y \Rightarrow x)$.
- A Wajsberg hoop is a hoop satisfying $x \Rightarrow (x \Rightarrow y) = y \Rightarrow (y \Rightarrow x)$.



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Ordinal Sums

- Let $\langle I, \leq \rangle$ be a totally ordered set with minimum 0. For all $i \in I$, let \mathcal{A}_i be a totally ordered such that for $i \neq j$, $A_i \cap A_j = \{1\}$, and assume that \mathcal{A}_0 is bounded.
- Then $\bigoplus_{i \in I} A_i$ (the *ordinal sum* of the family $(A_i)_{i \in I}$) is the structure whose base set is $\bigcup_{i \in I} A_i$, whose bottom is the minimum of A_0 , whose top is 1, and whose operations are

$$x \to y = \begin{cases} x \to^{\mathcal{A}_i} y & \text{if } x, y \in A_i \\ y & \text{if } \exists i > j(x \in A_i \text{ and } y \in A_j) \\ 1 & \text{if } \exists i < j(x \in A_i \setminus \{1\} \text{ and } y \in A_j) \end{cases}$$

$$x * y = \begin{cases} x *^{\mathcal{A}_i} y & \text{if } x, y \in A_i \\ x & \text{if } \exists i < j(x \in A_i \setminus \{1\}, y \in A_j) \\ y & \text{if } \exists i < j(y \in A_i \setminus \{1\}, x \in A_j) \end{cases}$$

• As a consequence, if $x \in A_i \setminus \{1\}$, $y \in A_j$ and i < j then x < y.

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n-contractive logics

Definition

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Theorem ([Bianchi and Montagna, 2011])

Every n-contractive BL-chain is isomorphic to an ordinal sum of finite MV-chains, each of them having at most n elements.

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Proof.

We already know that the SSCC implies the DMVP, so assume that L has the DMVP. Let A, B be two L-chains: we show that the CJEP (and hence the SSCC) holds.

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Axiomatization of MTL

The basic connective are $\{\land,\&,\rightarrow,\bot\}$ (formulas built inductively: a theory is a set of formulas). Useful derived connectives are the following ones:

$$\begin{array}{ll} \text{(negation)} & \neg \varphi \stackrel{\text{def}}{=} \varphi \to \bot \\ \\ \text{(disjunction)} & \varphi \lor \psi \stackrel{\text{def}}{=} ((\varphi \to \psi) \to \psi) \land ((\psi \to \varphi) \to \varphi) \\ \\ \text{(top)} & \top \stackrel{\text{def}}{=} \neg \bot \end{array}$$

MTL can be axiomatized by using these axioms and modus ponens: $\frac{\varphi-\varphi\to\psi}{\psi}$.

(A1)
$$(\varphi \to \psi) \to ((\psi \to \chi) \to (\varphi \to \chi))$$

(A2)
$$(\varphi \& \psi) \to \varphi$$

(A3)
$$(\varphi \& \psi) \to (\psi \& \varphi)$$

(A4)
$$(\varphi \wedge \psi) \rightarrow \varphi$$

(A5)
$$(\varphi \wedge \psi) \to (\psi \wedge \varphi)$$

(A6)
$$(\varphi \& (\varphi \to \psi)) \to (\psi \land \varphi)$$

(A7a)
$$(\varphi \to (\psi \to \chi)) \to ((\varphi \& \psi) \to \chi)$$

(A7b)
$$((\varphi \& \psi) \to \chi) \to (\varphi \to (\psi \to \chi))$$

(A8)
$$((\varphi \to \psi) \to \chi) \to (((\psi \to \varphi) \to \chi) \to \chi)$$

$$(A9) \qquad \qquad \bot \to \varphi$$

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Sketch of the proof.

If L has the SSCC, then the DMVP holds. Assume that L has the DMVP, and let ${\cal A}$ is an L chain.

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If L has the SSCC, then the DMVP holds. Assume that L has the DMVP, and let A is an L chain. If A has a last component, then it is subdirectly irreducible, since every component is subdirectly irreducible by hypothesis.

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