Stochastic optimization for large scale optimal transport

Project report

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This report will study the matter of applying stochastic optimization techniques to solve optimal transport problems in the discrete and semi-discrete settings. In the discrete setting, the standard solver of the regularized OT problem is the Sinkhorn-Knopp algorithm which has a general computational complexity of $O(n^2)$. The problem is notoriously hard to solve, and the complexity of the Sinkhorn-Knopp algorithm is too high for a very large scale setting. Stochastic algorithms can be used to cope with that problem and compute solutions of the regularized OT problem in linear time.

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1 Introduction

Optimal Transport (OT) is well known for its many applications in various domains. It has recently had major successes in Computer Vision or Natural Language Processing

1.1 Optimal transport: problem formulations

[2]

1.1.1 Entropic regularization of OT

We consider two measures $\mu \in \mathcal{M}^1_+(\mathcal{X})$ and $\nu \in \mathcal{M}^1_+(\mathcal{Y})$ defined on metric spaces \mathcal{X} and \mathcal{Y} . The cost of moving a unit of mass from $x \in \mathcal{X}$ and $y \in \mathcal{Y}$ is defined by the continuous function $c \in \mathcal{C}(\mathcal{X}, \mathcal{Y})$, and written c(x, y). We also define the set of joint probability measures on $\mathcal{X} \times \mathcal{Y}$

$$\Pi(\mu,\nu) \triangleq \{\pi \in \mathcal{M}_{+}(\mathcal{X} \times \mathcal{Y}); \forall (A,B) \subset \mathcal{X} \times \mathcal{Y}, \pi(A,\mathcal{Y}) = \mu(A), \pi(\mathcal{X},B) = \nu(B)\}$$

The entropic regularized version of the OT problem [1] can be written as a single convex optimization problem in the following form: $\forall (\mu, \nu) \in \mathcal{M}^1_+(\mathcal{X}) \times \mathcal{M}^1_+(\mathcal{Y})$,

$$W_{\varepsilon}(\mu,\nu) \triangleq \min_{\pi \in \Pi(\mu,\nu)} \int_{\mathcal{X},\mathcal{V}} c(x,y) d\pi(x,y) + \varepsilon KL(\pi||\mu \otimes \nu)$$
 (\mathcal{P}_{\varepsilon})

With $\mathrm{KL}(\pi||\mu\otimes\nu)$ corresponding to the Kullback-Leibler divergence between joint probabilities π and $\mu\otimes\nu$, defined by $\mathrm{KL}(\pi||\xi)\triangleq\int_{\mathcal{X},\mathcal{Y}}\left(\log(\frac{d\pi}{d\xi}(x,y)-1\right)d\xi(x,y)$.

For $\varepsilon > 0$, the above problem is strongly convex. $(\mathcal{P}_{\varepsilon})$ is usually called the primal form of the regularized OT problem, by opposition to the dual and semi-dual form that will be studied further.

Sinkhorn for discrete OT In the discrete setting $\mu = \sum_i^n \delta_{x_i} \mu_i$ and $\nu = \sum_j^m \delta_{x_j} \nu_j$, the sums are finite and the cost is $\mathbf{C} \in \mathbb{R}^{n \times m}$. The structure of the KL divergence gives the optimal solution $\mathbf{P}_{\varepsilon} \in \Pi(\mu, \nu)$ a convenient structure that makes it possible solving the problem using Sinkhorn's algorithm [1]. There indeed exist two scaling variables $\mathbf{u}_{\varepsilon} \in \mathbb{R}^n$ and $\mathbf{v}_{\varepsilon} \in \mathbb{R}^m$ such that

$$\mathbf{P}_{\varepsilon} = \operatorname{diag}(\mathbf{u}_{\varepsilon}) \mathbf{K}_{\varepsilon} \operatorname{diag}(\mathbf{v}_{\varepsilon})$$

Where $(\mathbf{K}_{\varepsilon})_{i,j} = \exp(-\mathbf{C}_{i,j}/\varepsilon)$ [3]. Those scaling variables can be computed iteratively with the following update at step ℓ ,

$$\mathbf{u}_{\varepsilon}^{\ell+1} \triangleq \frac{\mu}{\mathbf{K}_{\varepsilon} \mathbf{v}_{\varepsilon}^{\ell}} \quad \text{and} \quad \mathbf{v}_{\varepsilon}^{\ell+1} \triangleq \frac{\nu}{\mathbf{K}_{\varepsilon}^{T} \mathbf{u}_{\varepsilon}^{\ell+1}}$$
 (1)

Because each step of the algorithm relies on a vector-matrix computation, the overall complexity of the algorithm is O(nm) in the most general configuration.

The algorithm can be used in a large scale setting by using hardware (multiple Wasserstein distances can be computed in parallel on a GPU [slomp2011gpu]) and in some other specific cases where the kernel **K** is separable or can be expressed as a convolution [3]. In the general case however, the complexity of Sinkhorn's algorithm can be prohibitively large for a large scale problem.

1.1.2 Dual formulation of OT

1.1.3 Semi-dual formulation of OT

2 Stochastic optimization for large scale optimal transport

- 2.1 Discrete Optimal Transport
- 2.2 Semi-discrete Optimal Transport
- 3 Conclusion and Perspectives

References

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