

# Binary Search Tree (IV): Traversals

CSCD 300 – Data Structures

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## Goal: learn and implement various BST traversal procedures.

```
class BST{
    BST_Node root;    //the root of the BST

    BST(){ root = null;} /* the constructor */

    void InOrder_Traversal(BST_Node subtree_root){...} /* Inorder traversal
        and print all the nodes in the subtree rooted at "subtree_root". */

    void PreOrder_Traversal(BST_Node subtree_root){...} /* Preorder traversal
        and print all the nodes in the subtree rooted at "subtree_root". */

    void PostOrder_Traversal(BST_Node subtree_root){...} /* Postorder traversal
        and print all the nodes in the subtree rooted at "subtree_root". */

    void LevelOrder_Traversal(BST_Node subtree_root){...} /* Level-order traversal
        and print all the nodes in the subtree rooted at "subtree_root". */

    /* Other methods will follow here */
}
```

# Outline

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# Introduction

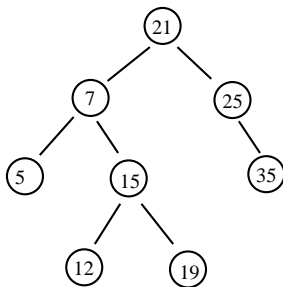
Tree traversal is to visit (print) all the nodes in the tree in some order.

- in-order traversal: traverse the left subtree first, then traverse the root, then traverse the right subtree of the root. Each subtree is also in-order traversed.
- pre-order traversal: traverse the root first, then traverse the left subtree, then traverse the right subtree of the root. Each subtree is also pre-order traversed.
- post-order traversal: traverse the left subtree first, then traverse the right subtree of the root, then traverse the root. Each subtree is also post-order traversed.
- Level-order traversal: traverse all the tree nodes from the top level to the bottom level. In each particular level, traverse the tree nodes from the left to the right.

## Theorem

*The time complexity of all the BST traversal procedures that we will present next is  $O(n)$ , where  $n$  is the number of nodes in the tree, because every node is visited for no more than a constant number of times.*

# Examples



- Inorder traversal: 5, 7, 12, 15, 19, 21, 25, 35
- Preorder traversal: 21, 7, 5, 15, 12, 19, 25, 35
- Postorder traversal: 5, 12, 19, 15, 7, 35, 25, 21
- Level-order traversal: 21, 7, 25, 5, 15, 35, 12, 19

# Inorder traversal

Idea: use recursion.

```
void BST::InOrder_Traversal(BST_Node subtree_root){  
    if(subtree_root != null)  
        InOrder_Traversal(subtree_root.left);  
    print(subtree_root.key);  
    InOrder_Traversal(subtree_root.right);  
}
```

Function call `InOrder_Traversal(root)`  
will in-order traverse the whole tree.

Inorder traversal prints the keys in the BST in ascending order

# Preorder traversal

Idea: use recursion.

```
void BST::PreOrder_Traversal(BST_Node subtree_root){  
    if(subtree_root != null)  
        print(subtree_root.key);  
        PreOrder_Traversal(subtree_root.left);  
        PreOrder_Traversal(subtree_root.right);  
}
```

Function call `PreOrder_Traversal(root)`  
will pre-order traverse the whole tree.

# Postorder traversal

Idea: use recursion.

```
void BST::PostOrder_Traversal(BST_Node subtree_root){  
    if(subtree_root != null)  
        PostOrder_Traversal(subtree_root.left);  
        PostOrder_Traversal(subtree_root.right);  
        print(subtree_root.key);  
}
```

Function call `PostOrder_Traversal(root)`  
will post-order traverse the whole tree.



## Level-order traversal

Idea: use a FIFO queue.

```
void BST::LevelOrder_Traversal(BST_Node subtree_root){
    Q = new FIFO; //Create a new FIFO queue of tree node type.

    Q.enqueue(subtree_root);

    while(Q.size > 0){
        BST_Node node = Q.dequeue();
        print(node.key);
        if(node.left != null)
            Q.enqueue(node.left);
        if(node.right != null)
            Q.enqueue(node.right);
    }
}
```

Function call `LevelOrder_Traversal(root)`  
will level-order traverse the whole tree.