

Sorting (II): Selection Sort

CSCD 300 – Data Structures

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Goal

We will learn the mechanism of the Selection Sort algorithm and then analyze its time complexity in the best as well as in the worst case.

Outline

- 1 Selection sort
- 2 The time complexity
- 3 Question

Selection sort

Basic idea

The algorithm repeatedly scans the **unsorted portion** of the sequence, and put the **smallest** number of the **unsorted portion** at the **first** location of the **unsorted portion**.

- The first pass is to scan $A[0 \dots n - 1]$, meaning the unsorted portion is the whole sequence.
- The second pass is to scan $A[1 \dots n - 1]$, meaning the unsorted portion is $A[1 \dots n - 1]$.
- ...
- The i th pass is to scan $A[i - 1 \dots n - 1]$, meaning the unsorted portion is $A[i - 1 \dots n - 1]$.
- ...

An example

The first pass: 4 3 2 ① \longrightarrow 1 3 2 4

The second pass: 1 3 ② 4 \longrightarrow 1 2 3 4

The third pass: 1 2 ③ 4 \longrightarrow 1 2 3 4

The fourth pass: 1 2 3 ④ \longrightarrow 1 2 3 4

unsorted portion ○ the smallest number in the unsorted portion

Pseudocode

SelectionSort(A)

Input: An array $A[0 \dots n - 1]$ of n numbers

Output: The sorted A .

```
for  $i = 0 \dots n - 1$  do
    min_index = i
    for  $j = i \dots n - 1$  do
        if  $A[j] < A[\text{min\_index}]$  then
            min_index = j;
    exchange( $A[i], A[\text{min\_index}]$ )
```

The time complexity

No matter what sequence of data is given:

- The first pass scans n elements.
- The second pass scans $n - 1$ elements.
- ...
- The i th pass scans $n - i + 1$ elements.
- ...

So the total time cost of the Selection sort in any case is:

$$n + (n - 1) + (n - 2) + \dots + 2 + 1 = \frac{n(n + 1)}{2} = O(n^2)$$

Question

How do you use Selection sort if the data sequence is saved in a singly linked list ?