

SUBIECTE

Nr 2:

- ① Conz. sunt "simile" fără memorie, cu simbolurile echiprobaabile. Avem matricea de zigomot.

a) $P(Y|X) = \begin{bmatrix} 5/8 & 1/4 & ? \\ ? & 1/4 & 5/8 \end{bmatrix}$.

$$\sum p(x_i, y_j) = 1 \Rightarrow \frac{5}{8} + \frac{1}{4} + ? = 1 \Rightarrow ? = 1/8 = g_1 = g_2. \rightarrow$$

⇒ $P(Y|X) = \begin{bmatrix} 5/8 & 1/4 & 1/8 \\ 1/8 & 1/4 & 5/8 \end{bmatrix}$

canal uniform după intrare: 3 permuteții pe linie.

- b) Calculati $H(Y)$, $H(Y|X)$, $H(X,Y)$, $H(X|Y)$, $I(X,Y)$, c.

$$P(X|Y) = P(Y|X) \cdot P(X)$$

$$P(X,Y) = \begin{bmatrix} 1/2 & 0 \\ 0 & 1/2 \end{bmatrix} \begin{bmatrix} 5/8 & 1/4 & 1/8 \\ 1/8 & 1/4 & 5/8 \end{bmatrix} = \begin{bmatrix} 5/16 & 1/8 & 1/16 \\ 1/16 & 1/8 & 5/16 \end{bmatrix}.$$

$$p(x_1) = \sum_{j=1}^3 p(x_i, y_j) = \quad p(y_1) = \sum_{i=1}^2 p(x_i, y_j)$$

$$p(x_2) = \sum_{j=1}^3 p(x_i, y_j) = \quad p(y_2) = \dots$$

$$p(y_3) = \dots$$

$$H(Y) = - \sum_{i,j} p(x_i, y_j) \cdot \log p(x_i, y_j)$$

$$H(X|Y) = - \sum_i p(x_i, y_i) \log p(x_i, y_i).$$

$$H(X|Y) = H(X, Y) - H(Y), \quad I(X, Y) = H(X) + H(Y) - H(X, Y).$$

$$c = \log m + \sum_{i=1}^m p(x_i, y_i) \log p(x_i, y_i)$$

c) $\ell = ?$ $P_e = -10 \text{ dB}$, $[S/Z]_{\text{dB}} = 20$, $\alpha = 3 \text{ dB/km}$, $Z_0 = 40 \Omega$

$$\boxed{10 \log \frac{P_{\text{out}}}{P_e} = -\alpha \cdot \ell} \Rightarrow 10 \log P_{\text{out}} - 10 \log P_e = -\alpha \ell \Leftrightarrow$$

$$\ell = \frac{10 \log P_{\text{out}} - 10 \log P_e}{-\alpha} = \frac{[P_e]_{\text{dBm}} - [P_{\text{out}}]_{\text{dBm}}}{\alpha}$$

$$[S/Z]_{\text{dB}} = 10 \log \frac{S}{Z} \Rightarrow 10 \log S - 10 \log Z = [S/Z]_{\text{dB}}$$

$$\Rightarrow [S]_{\text{dBm}} - [Z]_{\text{dBm}} = [S/Z]_{\text{dB}}$$

$$\ell = \underbrace{[P_e]_{\text{dBm}} - \left[\frac{S}{Z} \right]_{\text{dB}} - [Z_0]_{\text{dBm}}}_{\alpha} = \frac{-10 - 20 + 40}{3} = \frac{10}{3} \text{ km.}$$

②. Semnal modulat amplitudine:

$$S_{\text{MA}} = (7 + \sin 400\pi t - 20 \cos 100\pi t) \cos 10^4 \pi t$$

a) Expresii temporale ale puterilor si semnalelor modulate.

$$S_p = ?, S_m = ?$$

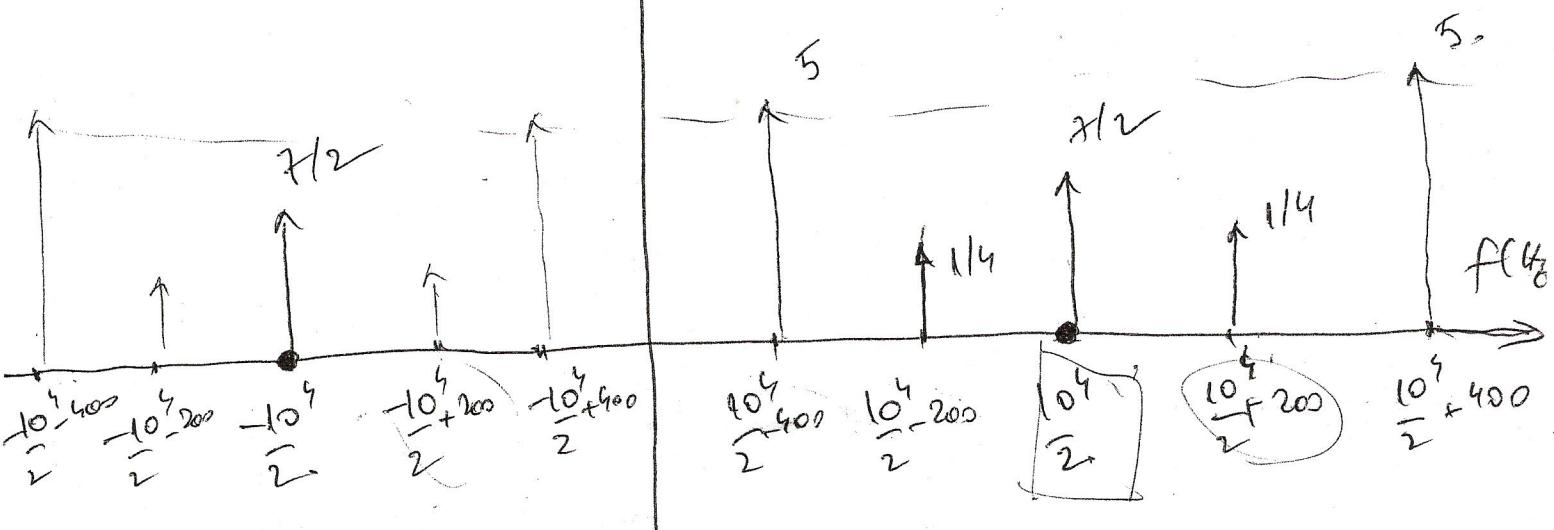
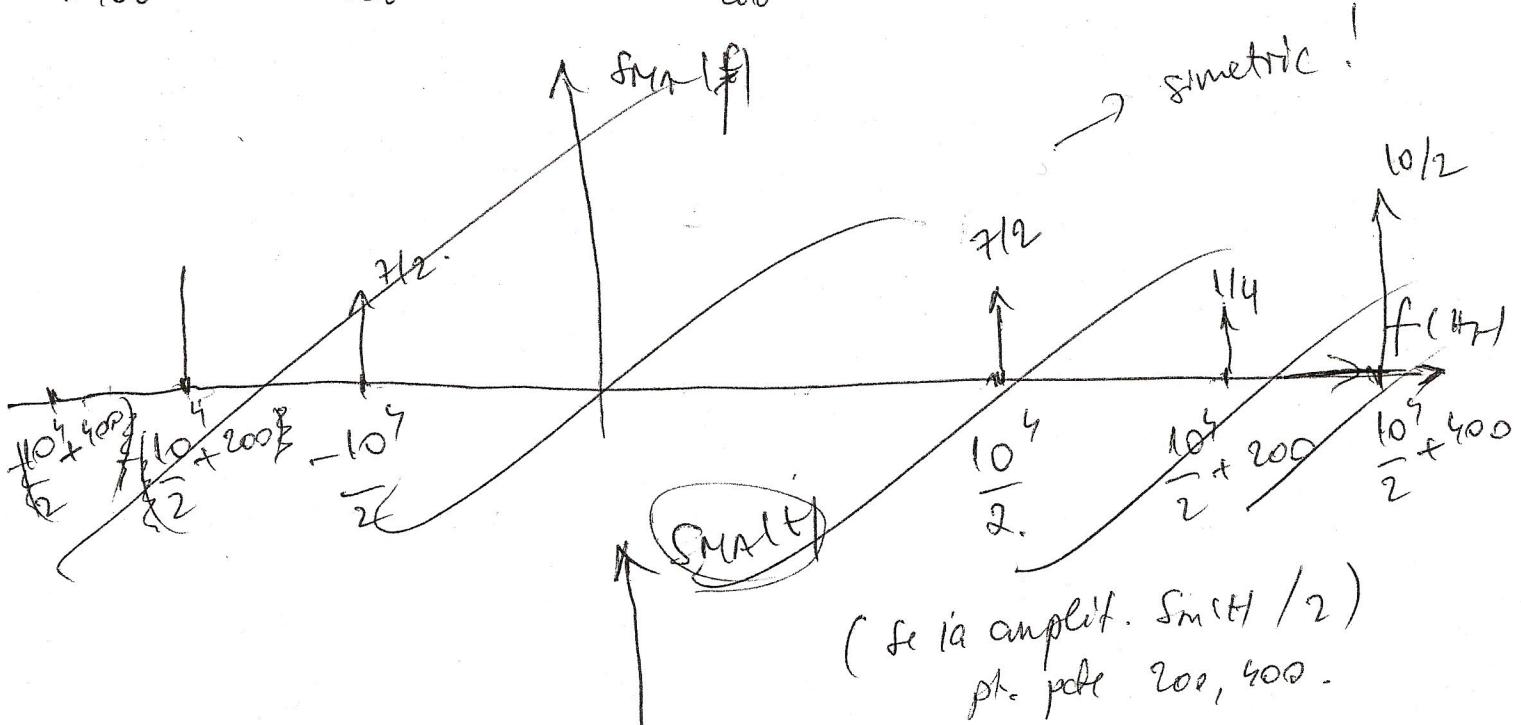
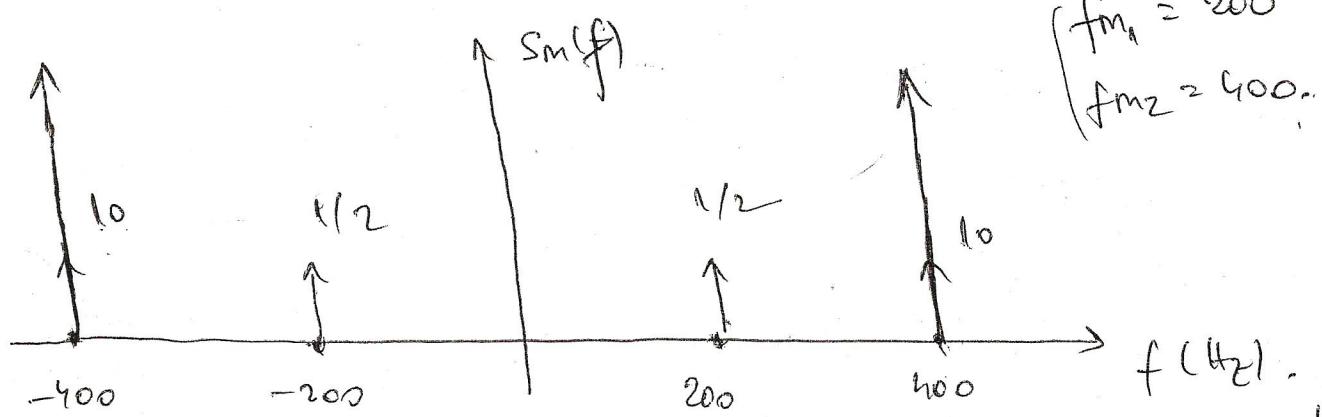
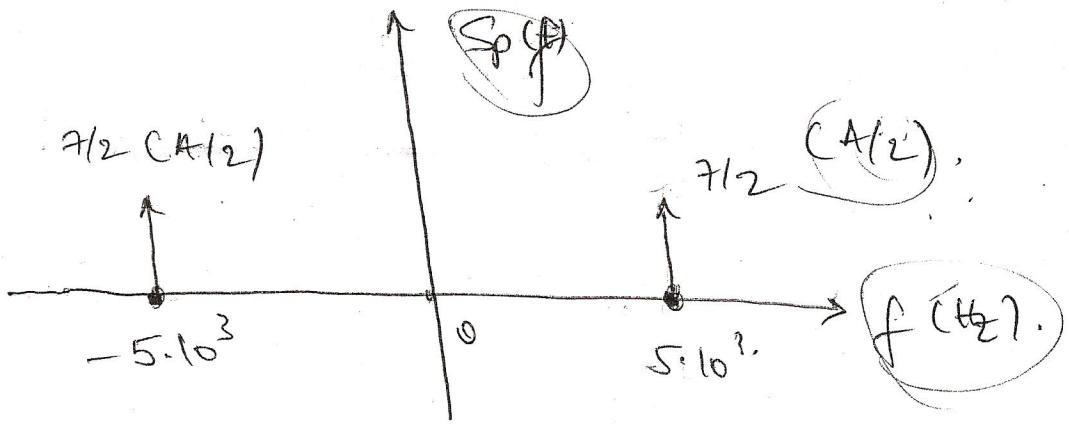
$$\boxed{S_{\text{MA}}(t) = (A_p + s_m(t)) \cos 2\pi f_p t} \Rightarrow$$

$$\Rightarrow \begin{cases} A_p = 7 \\ f_p = \frac{10^4}{2} = 5 \cdot 10^3 \text{ Hz} \end{cases}$$

$$\boxed{S_p(t) = A_p \cos(2\pi f_p t)} \Rightarrow S_p(t) = 7 \cos 10^4 \pi t$$

$$S_m(t) = \sin 400\pi t - 20 \cos 100\pi t$$

b) Spectrele semnalelor: s_m , S_p , S_{MA} .



c) Calculate index of modulation or efficiency transmission

$$m = ? , \eta = ?$$

$$m = \frac{A(t)_{\max} - A(t)_{\min}}{A(t)_{\max} + A(t)_{\min}}$$

$$\text{SMA LM} = A(t) \cos \omega_0 t$$

$$A(t) = 7 + \sin 400\pi t + \frac{1}{f_0} \cos 400\pi t$$

$$A(t) = 7 + \sin 400\pi t + \frac{1}{f_0} \cos 400\pi t = 0 \Rightarrow$$

$$\Rightarrow \cos 400\pi t = -\sin 400\pi t \Rightarrow \tan 400\pi t = -1 \quad (\text{at } f_0 = 2 \text{ Hz})$$

$$\Rightarrow \frac{1}{f_0} = \sin 400\pi t \quad \left(\frac{1}{f_0} \right)$$

$$\sin 400\pi t \neq 0 \quad \pm 1. \quad (\text{cond } \cos \neq 0).$$

$$A(t) = 7 + \sin 400\pi t + \frac{1}{f_0} \cos 400\pi t \quad (m = 1 - 2 \sin^2 x)$$

$$A(t) = 7 + 1 - 2 \cdot \frac{1}{f_0} = 7 + 1 - 2 \cdot \frac{1}{4} = 7 + 1 - 0.5 = 7.5$$

$$A(t) = 7 + \frac{1}{f_0} + 2 \cdot \frac{\sin 400\pi t}{f_0} = 7 + 0.125 + 1.99 \approx 9.125$$

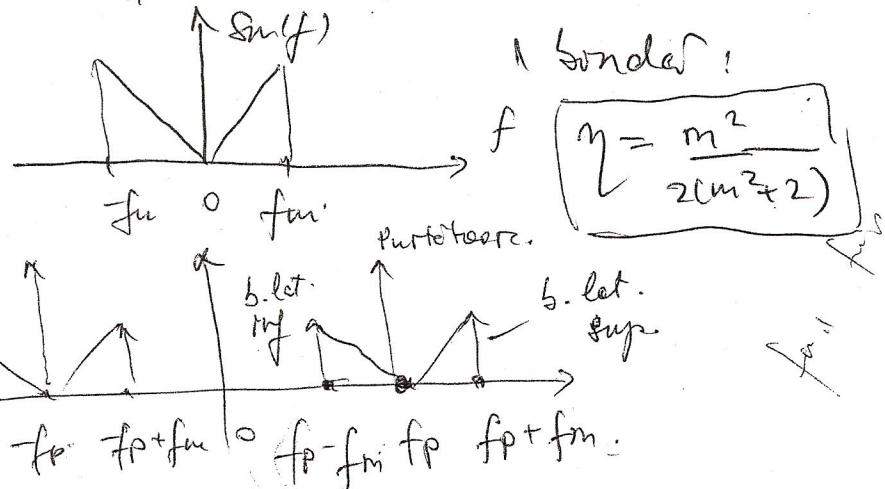
$$A(t)_{\min} = 7 + 0.125 = 7.125$$

$$m = \frac{4.125}{7.125} = 0.577 \approx 57.7\%$$

$$\eta = \frac{m^2}{m^2 + 2}$$

$$\eta = \frac{0.14}{4.14} = 0.03 = 3\%$$

(avem 2 bazei laterale).



- d) Propuneti o solutie de demodulare coerenta cu preciparea caracteristicilor blocajelor fct.

③

Conciliatia text:

ACCES_NELIMITAT_LA_ACESTE_TESTE_SI_LA_ALTELE_CARE

A : 7/51 N : 11/51 T : 6/51

C : 4/51 L : 6/51 Spatiu: 9/51

E : 10/51 I : 3/51

S : 14/51 M : 11/51

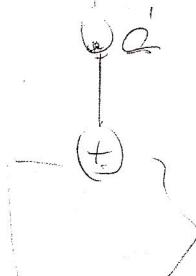
- a) probabilitatile de apariție:

$$p(CA) = 0.13 \quad p(CN) = 0.049$$

$$p(C) = 0.22 \quad p(sp) = 0.17$$

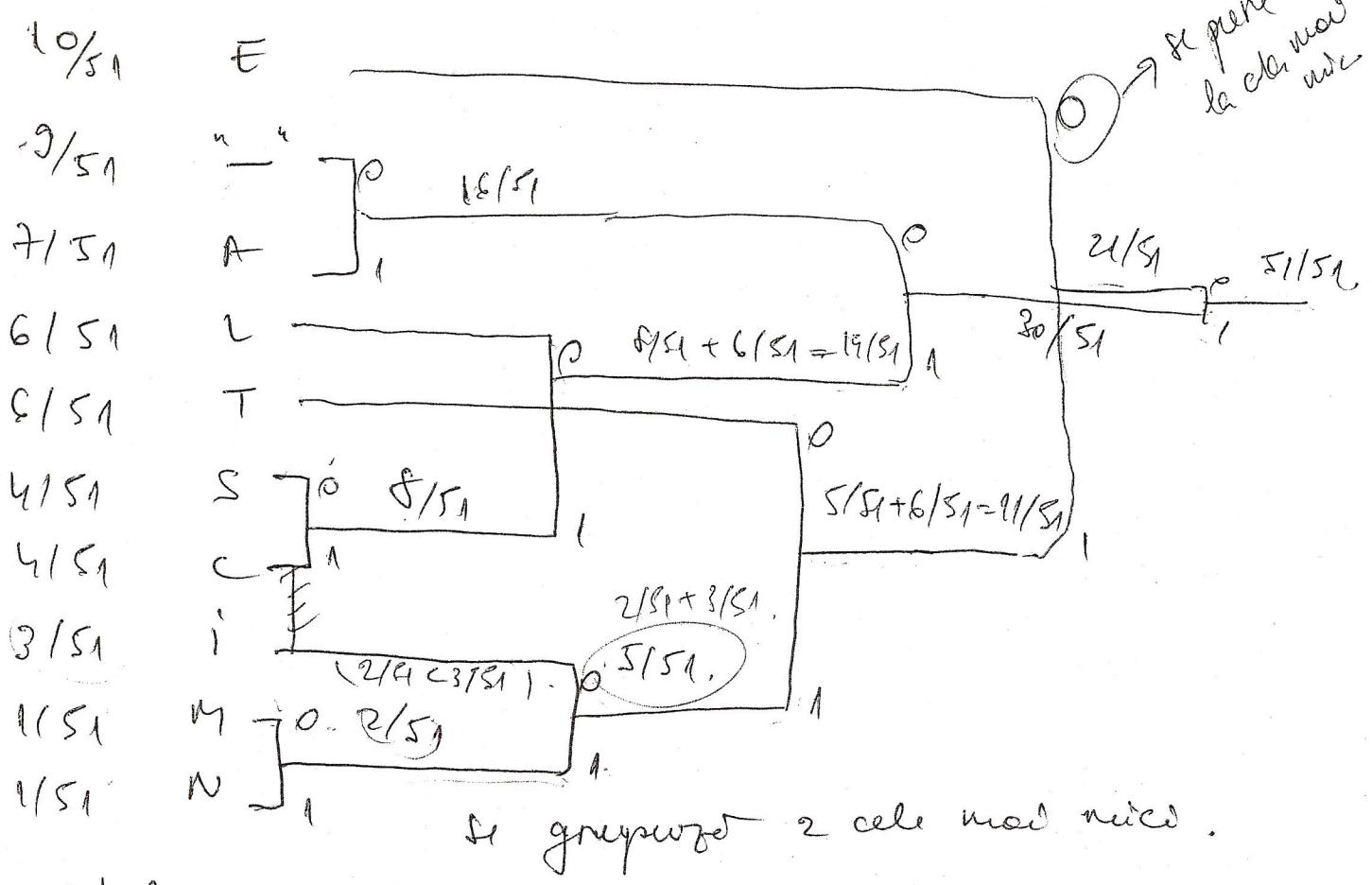
$$p(CE) = 0.19 \quad p(M) =$$

$$p(S) = 0.07$$



- b) Codificare cu Huffman:

Ludre. probabilitate ordonata crescator:



Tabel cod.

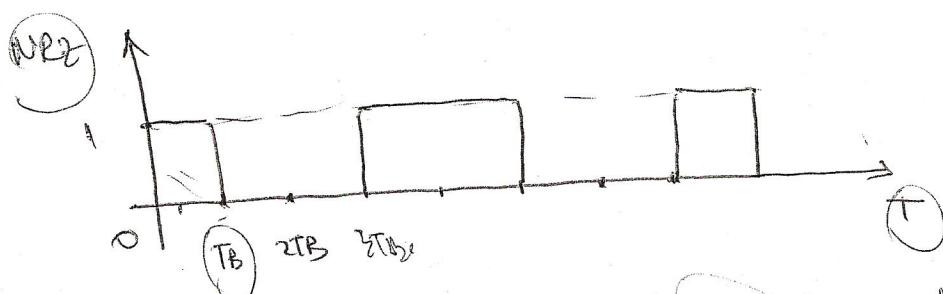
E	00
-	100
A	101
L	110
T	010
S	1110
C	1111
I	0110
M	01110
N	01111

(2)

Impuls negativat transmis în sistem transm. & terminalor
bitare în banda de bază la foarte voleme de interac. esențiale.

(k)	-3	-2	-1	0	1	2	3
Pr(k)	0	0	0.1	0.9	0.1	0	0

- a) semnalul emis în linie, la intrare în codor se aplică
10011001.



- b) proiectează egualizator cu 3 etape ce forțează reacția în 0 la
 $K = \pm 1, \pm 2, \pm 3$ și (in 1 la $K=0$) $2N+1=3 \Rightarrow (N=1)$

$$\begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} p_r(0) & p_r(-1) & p_r(-2) \\ p_r(1) & p_r(0) & p_r(-1) \\ p_r(2) & p_r(1) & p_r(0) \end{bmatrix} \cdot \begin{bmatrix} c_1 \\ c_0 \\ c_{-1} \end{bmatrix} \Rightarrow$$

$$\begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0.9 & 0.1 & 0 \\ 0.1 & 0.9 & 0.1 \\ 0 & 0.1 & 0.9 \end{bmatrix} \cdot \begin{bmatrix} c_1 \\ c_0 \\ c_{-1} \end{bmatrix} \Rightarrow$$

$$\begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} = \frac{1}{9} \begin{bmatrix} 9 & 1 & 0 \\ 1 & 9 & 1 \\ 0 & 1 & 9 \end{bmatrix} \begin{bmatrix} c_1 \\ c_0 \\ c_{-1} \end{bmatrix}$$

Inversarea matricei.

$$1 \quad 9AT = 9A$$

$$A^{-1} = \frac{A^T}{\Delta}$$

$$A^* = \begin{bmatrix} 80 & -9 & 1 \\ 9 & 81 & -1 \\ 1 & -9 & 81 \end{bmatrix} \Rightarrow A^{-1} = \frac{A^*}{711} \cdot \frac{1}{10}$$

$$A^{-1} \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} = \frac{1}{10} \begin{pmatrix} c_{-1} \\ c_0 \\ c_1 \end{pmatrix} \Rightarrow 10 A^{-1} \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} = \begin{pmatrix} c_{-1} \\ c_0 \\ c_1 \end{pmatrix} \Rightarrow$$

$$\Rightarrow \begin{pmatrix} c_{-1} \\ c_0 \\ c_1 \end{pmatrix} = \frac{10}{711} \begin{pmatrix} -9 \\ 8 \\ -9 \end{pmatrix} = \begin{pmatrix} -0.12 \\ 1.13 \\ -0.12 \end{pmatrix}$$

$$\text{peg}(k) = \sum_{n=-N}^N c_n \cdot \text{pr}(k-n), \Rightarrow$$

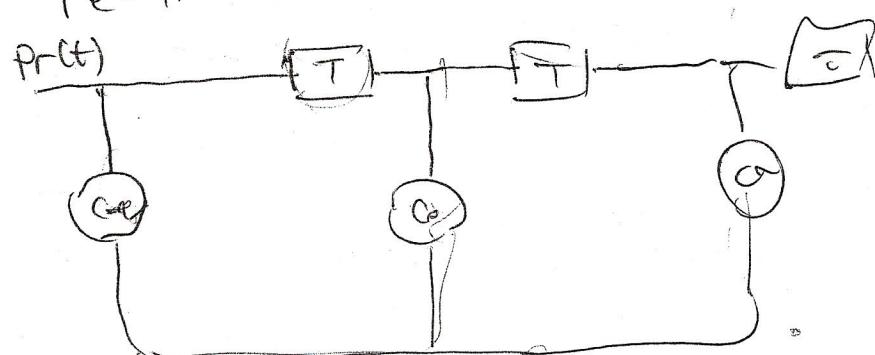
$$\Rightarrow \text{peg}(k) = \sum_{n=-1}^1 c_n \text{pr}(k-n),$$

$$\underline{k=0} \quad \text{peg}(0) = c_{-1} \text{pr}(1) + c_0 \text{pr}(0) + c_1 \text{pr}(-1) = \\ -0.12 \cdot 0.1 + 1.13 \cdot 0.9 - 0.12 \cdot 0.1 \\ = -0.024 + 0.817 = 0.993 \approx 1. \checkmark$$

$$\underline{k=1} \quad \text{peg}(1) = c_{-1} \text{pr}(2) + c_0 \text{pr}(1) + c_1 \text{pr}(0) = \\ -0.12 \cdot 0.1 + 1.13 \cdot 0.9 - 0.12 \cdot 0.1 \\ = 0.$$

$$\underline{k=2} \quad \text{peg}(2) = 0, \quad \underline{k=3} \quad \text{peg}(3) = 0, \quad \checkmark$$

c) Calculate $\text{pe}(k)$.



6. Un canal are matricea de probabilitate:

$$P(X|Y) = \begin{bmatrix} 1/2 & 1/3 & 2/1 \\ 2/2 & 2/3 & 1/4 \end{bmatrix}$$

$$\text{a)} \quad \frac{1}{2} + \frac{1}{3} + \frac{2}{1} = 1 \Rightarrow 5 + 6g_1 = 6 \Rightarrow g_1 = \frac{1}{6}.$$

$$\frac{1}{2} + \frac{2}{3} + \frac{1}{4} = 1 \Rightarrow 1 + 12g_2 = 1 \Rightarrow g_2 = \frac{1}{12}.$$

$$\Rightarrow P(X|Y) = \begin{bmatrix} 1/2 & 1/3 & 1/6 \\ 1/2 & 2/3 & 1/4 \end{bmatrix}$$

b) $H(X, Y)$ pt. simboluri echiprobaabile la intrare

$$P(X|Y) \cdot p(y) \Rightarrow \overline{P(X|Y)} = \begin{bmatrix} \frac{1}{2} \cdot \frac{1}{2} & \frac{1}{2} \cdot \frac{1}{3} & \frac{1}{2} \cdot \frac{1}{6} \\ \frac{1}{2} \cdot \frac{1}{2} & \frac{1}{2} \cdot \frac{2}{3} & \frac{1}{2} \cdot \frac{1}{4} \end{bmatrix}$$

$$\Rightarrow \overline{P(X|Y)} = \begin{bmatrix} \frac{1}{2} \cdot \frac{1}{2} & \frac{1}{2} \cdot \frac{1}{3} & \frac{1}{2} \cdot \frac{1}{6} \\ \frac{1}{2} \cdot \frac{1}{12} & \frac{1}{2} \cdot \frac{2}{3} & \frac{1}{2} \cdot \frac{1}{4} \end{bmatrix} = \begin{bmatrix} 1/4 & 1/6 & 1/12 \\ 1/24 & 2/18 & 1/8 \end{bmatrix}.$$

$$P(X, Y) = P(X|Y) \cdot P(Y).$$

$P(Y) = 1/2$: pt. ambele evenimente

$$\Rightarrow H(X, Y) = \left[\frac{1}{5} \cdot \log_2 \frac{1}{5} + \frac{1}{6} \cdot \log_2 \frac{1}{6} + \frac{1}{12} \cdot \log_2 \frac{1}{12} + \frac{1}{24} \cdot \log_2 \frac{1}{24} + \right. \\ \left. + \frac{1}{3} \cdot \log_2 \frac{1}{3} + \frac{1}{8} \cdot \log_2 \frac{1}{8} \right].$$

7. Pt. un canal diferit:

$$P(X|Y) = \begin{bmatrix} 0.1 & 0.25 \\ 0.2 & 0 \\ 0.3 & 0.15 \end{bmatrix}$$

$$\text{a)} \quad H(X), H(Y), H(X|Y), H(Y|X),$$

$$\text{b)} \quad H(X, Y), \text{ c)} \quad I(X, Y) \text{ transinformata.}$$

Og! Proprietatea $P(X, Y)$.

- suma elementelor pe linie:

$$\sum_{j=1}^m p(x_i, y_j) = p(x_i) \text{ cu } \sum_{i=1}^n p(x_i) = 1.$$

deci $p(x_1) = 0.35$, $p(x_2) = 0.2$, $p(x_3) = 0.45$.

$$H(X) = -[p(x_1) \cdot \log_2 p(x_1) + p(x_2) \cdot \log_2 p(x_2) + p(x_3) \cdot \log_2 p(x_3)]$$
$$= -[0.35 \log_2 0.35 + 0.2 \cdot \log_2 0.2 + 0.45 \cdot \log_2 0.45].$$

$$H(Y) = -[p(y_1) \cdot \log_2 p(y_1) + p(y_2) \cdot \log_2 p(y_2)].$$

- suma elementelor pe coloana:

$$\sum_{i=1}^n p(x_i, y_j) = p(y_j) \text{ cu } \sum_{j=1}^m p(y_j) = 1.$$

$$p(y_1) = 0.6 \text{ și } p(y_2) = 0.4.$$

$H(X, Y)$ se obtine din $P(X, Y)$.

Relație: $H(Y|X) = H(X, Y) - H(X) \Rightarrow H(Y|X) = \dots$ b.d.

$$H(X|Y) = H(X, Y) - H(Y) \Rightarrow H(X|Y) = \dots$$
 b.d.

$$I(X, Y) = H(X) + H(Y) - H(X, Y) = \dots$$
 b.d.

⑧ Pt canal gaussian cu $B = 1 \text{ MHz}$ și $S/I = 30 \text{ dB}$.

Formula: $\left[\frac{S}{I}\right]_{\text{dB}} = 10 \log_{10} \frac{S}{I} \Rightarrow \frac{S}{I} = 10^{\frac{S/I}{10}}$

$$\left[\frac{S}{I}\right]_{\text{dB}} = 30 \text{ dB} \Rightarrow \boxed{S/I = 10^3}$$

Calculați capacitatea și timpul în care se transmite pe canal 1 milion de caractere ASCII (8 biti)?

$$C = B \log_2 \left(1 + \frac{S}{N} \right) = 1 \cdot \log_2 (1+10^3) \approx 9.97 \text{ Mbit/s.}$$

10^6 de bit/s.

$$I = \frac{C \cdot T}{B - \text{trupel}} \Rightarrow I = \frac{10^6 \cdot 8}{9.97 \cdot 10^3} = 802,5 \text{ bit}$$

9. $B = 4 \text{ kHz}$

$f_E = 2,5 f_N$ (f_N : frecv. Nyquist) frecv. de esantionare

- frecvare esantion cuantizat pe 256 niveluri

- esantionare indep.

a) viteza de transmitere a info.

respre sursele?

b) eroarea de transmitere zero? pe canal gaussian cu banda

de 50 kHz si rap. $S/N = 23 \text{ dB}$.

c) B necesar transmitere fără eroare doar $S/N = 10 \text{ dB}$?

a) Criteriu Nyquist:

Pe canal cu 3 limitări se pot transmite cel mult 2 simboluri

(simboluri)/s pt. frecvare hertz de lungime de B .

$$f_E = 2,5 f_N \Rightarrow 2 \text{ simboluri/s} \xrightarrow{\text{---}} 2,5 \text{ Hz} \\ \times \text{ simbol/s} \xrightarrow{\text{---}} 4 \cdot 10^3 \text{ Hz} \text{ (B),}$$

$$x = 2 \cdot \frac{4 \cdot 10^3}{2,5} = \frac{8000}{2,5} = 3200 \text{ simbol/s.}$$

b) $C =$

Pentru
PUI.

Ac

$$⑩. \quad l = 10 \log_{10} \left(\frac{P_{out}}{P_{in}} \right) \rightarrow 10 \log_{10} \left(\frac{P_{out}}{P_{in}} \right) = -A \cdot l.$$

$$\underbrace{[P_{out}]_{dBm} + [P_{in}]_{dBm}}_A = l$$

$$[Net]_{dBm} = -40 dBm$$

$$[P_{in}]_{dBm} = -10 dBm.$$

$$[S/N] = 10 dB. \Rightarrow [S]_{dBm} = [N]_{dBm}$$

$$\boxed{l = \left(\frac{S}{N} \right) dB + [N]_{dBm} - [P_{in}]_{dBm}} \rightarrow -A.$$

$$\Rightarrow l = -\frac{1}{3} \cdot (10 - 40 + 10) = +20 \approx 7 km.$$

Pt. marge suplimentare: +6 dB.

$$l' = \frac{1}{3} (10 + 6 - 40 + 10) = \frac{14}{3} \approx 5 km.$$

marge LM 5 dB

b) lungimea medie a avonturiei de cod.

$$\overline{n}_{\text{med}} = \sum_{i=1}^{10} p_i \cdot n_i, \quad (n_i) - \text{lunghimea } i.$$

$$= \frac{10}{51} \cdot 2 + \frac{9}{51} \cdot 3 + \frac{7}{51} \cdot 3 + \frac{6}{51} \cdot 3 + \frac{6 \cdot 3}{51} + \frac{4 \cdot 8}{51} + \frac{3 \cdot 4}{51} +$$

$$+ \frac{2 \cdot 5}{51} = \dots$$

d) eficiență în raport cu ASE extinsă (abilitate),

$$m = \frac{\overline{n}_{\text{min}}}{\overline{n}} = \frac{H}{\overline{n} \log 2}, \quad H = \sum_{i=1}^{10} p_i \log p_i.$$

$g=2$. (0 și 1 reprezintă).

e) Varianta de cod de corecție în cazul unei singure erori.

NR 1

③ a) $x^4 + x + 1$ ca polinom generator pt cod ciclic (15, 11).
 $n=15$ sub total, $m=11$, $k=4$ efectiv folosit, 4 control, partitate.

$$\left\{ \begin{array}{l} n=15 \\ k=4 \\ m=11 \end{array} \right.$$

$$\begin{smallmatrix} n-1 \\ x \end{smallmatrix}$$

Reglea: $x^4 + x + 1$ divide $x^{15} - 1$ in modulo 2.

$$\begin{array}{r|l} x^5 - 1 & x^4 + x^3 + 1 \\ \hline -x^{12} - x^{11} & x^{11} + x^8 + x^7 + x^5 + x^4 + x^3 - x^2 - 3x - 3 \\ -x^{12} - x^{11} & \\ \hline x^{12} + x^9 + x^8 & \end{array}$$

$$\begin{array}{r|l} x^9 + x^8 + x^7 & -x^4 + x^9 + x^8 - 1 \\ \hline -x^9 - x^8 - x^7 & x^4 + x^8 + x^7 \\ -x^9 - x^8 - x^7 & x^9 + 2x^8 + x^7 - 1 \\ -x^9 - x^8 - x^7 & \end{array}$$

$$2x^4 + x^2 - x^6 - x^5 - 1$$

$$\begin{array}{r|l} 2x^4 - 2x^5 - 2x^4 & x^4 - x^6 - 3x^5 - 2x^4 - 1 \\ \hline x^4 - x^6 - 3x^5 - 2x^4 - 1 & x^4 - x^6 - x^3 \\ \hline x^4 - x^6 - x^3 & \end{array}$$

$$-x^4 - 3x^5 - 3x^4 - x^3 - 1$$

$$\begin{array}{r|l} x^6 + x^2 + x^1 & -3x^4 - 3x^5 + x^2 - 1 \\ \hline -3x^4 - 3x^5 + x^2 - 1 & \end{array}$$

$$3x^5 + 3x^2 + 3x$$

$$-3x^4 + 4x^2 + 3x - 1$$

$$3x^5 + 3x + 3$$

$$\begin{array}{r|l} 4x^2 + 6x + 2 & \\ \hline 0 & \\ 0 & \\ 0 & \end{array}$$

$$\cancel{x} + x = 0$$

$$x^{15} - 1 = (x^4 + x^3 + 1)(x^{11} + x^8 + x^7 + x^5 + x^3 + x^2 + x + 1)$$

divide!

b) Perechii G, H pt. codul cu polinom. gen x^4+x+1 ,

$$g(x) = x^4 + x + 1 \Rightarrow g_0 x^4 + \dots + g_k \quad (c)$$

$$G_{(m, n)} = \begin{bmatrix} 0 & 0 & \dots & 0 & g_0 & \dots & g_k \\ 0 & 0 & \dots & g_0 & \dots & g_k & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots \\ 0 & g_0 & \dots & g_k & 0 & \dots & 0 \\ g_0 & \dots & g_k & 0 & \dots & 0 & \end{bmatrix}$$

$$G = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix} | \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

H determinat de polinomul (c)

$$h(x) = x^6 + x^5 + x^4 + x^3 + \cancel{x^2} + x^1 + x + 1 =$$

$$x^{n-1} = g(x) \cdot h(x) \quad = h_m x^m + \dots + h_0$$

$$H_{(k, m)} = \begin{bmatrix} h_0 & h_1 & \dots & h_m & 0 & \dots & 0 \\ 0 & h_0 & h_1 & \dots & h_m & 0 & \dots & 0 \\ \vdots & \vdots \\ 0 & 0 & \dots & 0 & h_0 & h_1 & \dots & h_m \end{bmatrix}$$

$$H(4,12) = \left[\begin{array}{ccccccc} & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \end{array} \right].$$

c) Construcción: $u_0 = \underbrace{\text{0000000000}}_{Q_1 Q_2} \underbrace{\text{0000000000}}_{Q_3 Q_4}$

$$u_1 = \underbrace{\text{100111111111}}_{Q_2 Q_3} \in x^4 + x + 1.$$

$$H(4,15) = \boxed{\begin{array}{cccccccccc} 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 1 \end{array}}$$

$$u_0 = \begin{bmatrix} a_1 \\ \vdots \\ a_5 \end{bmatrix}$$

$u_0 \cdot H(4,15) = 0$ docé apartirne codeñiu.

$$H(4,15)u_0 = 0$$

$$\cancel{a_1} + \cancel{a_4} + \cancel{a_5} = 0$$

$$\underline{a_1} + \underline{a_2} + \underline{a_3} + \underline{a_4} + \underline{a_6} + \underline{a_7} + \underline{a_9} + \underline{a_{12}} = 0_{11}$$

$$\underline{a_2} + \underline{a_3} + \underline{a_4} + \underline{a_5} + \underline{a_7} + \underline{a_9} + \underline{a_{10}} + \underline{a_{13}} = 0$$

$$\underline{a_3} + \underline{a_4} + \underline{a_5} + \underline{a_6} + \underline{a_8} + \underline{a_{10}} + \underline{a_{11}} + \underline{a_{14}} = 0_{11}$$

no par. pt. u_1 , (1 fund)

u_0
apartirne

u_1 ^{num}
apartirne

d) Redundanță codulară

$$P = 1 - \eta \text{ Deficiență}$$

$$\eta = \frac{m}{n} = \frac{11}{15}$$

$$2^k \approx n$$

$$k = \log n \quad m = \underline{n} - \log n$$

e) Bif. control. pr. ciclic $(15, 11)$ și cod. de $d_H = 5$

$$\text{cod. impar} \Rightarrow 2r + 1 = 5 \Rightarrow$$

$$\Rightarrow r = 2 \text{ detect. 2 erori corectate } r-1 = 1 \text{ eroare}$$

$$\text{corector de lungime } 2 : r = 4$$

$$g \cdot r = \text{grad}(g(x))$$

$$(15, 11) \quad n = 15, m = 11, k = 4.$$

$$g(x) = x^4 + x + 1, \quad \text{grad}(g(x)) = 4$$

$$\text{grad}(g(x)) = 4 = 2^2 \text{ cifre test.}$$

$$d_H = 5 \rightarrow \text{DET: 2 erori}$$

$$\text{COR: 1 eroare}$$

$$2r + 1 = \text{det}$$

$$r = 1 \text{ - corect.}$$

$$2g \text{ este cel mult } 2^4 \text{ erori corectate}$$

$$2^4 = 16 \Rightarrow r = 4$$

$$2^{r-1} = n = 15$$

$$2^{r-1}, \quad \begin{cases} \text{Corect 1 eroare} \\ \text{Detect. 2 erori} \end{cases}$$

$$\text{pachete. } k \leq 4$$

Cod.

$$(2^{r-1}, 2^r - 2^{r-1})$$

$$(15, 10)$$

$$k = n - m = 4 \geq d_H$$

$$\Rightarrow 5 \geq d_H = 5 ?$$

Cod. corector de lungime:

$$2^{r-1} = 15 \quad \begin{cases} \text{Corect: 1 er.} \\ \text{Detect: 2 er.} \end{cases}$$

$$2^{r-1} = 15 \quad \begin{cases} \text{Corect: 1 er.} \\ \text{Detect: 2 er.} \end{cases}$$

$$\text{Pachete: } 2^r \cdot r \Rightarrow 4 \cdot 10$$

Codul not: $(2^{r_1}, 2^{r_2 \cdot r - 1}) = (15, 11)$

Pachete de lungime ≥ 4 . \Rightarrow fereastra de control. min 2g.

\times $k = n - m \geq d$, pachete lungime $\underline{d=4}$ $15 - 11 = \underline{4 \geq 4}$ ✓

① Surse binare fără memorie, simbolen' echiprobaabile:

$$P(Y|X) = \frac{\begin{bmatrix} 5/8 & 1/4 & ? \\ ? & 1/4 & 5/8 \end{bmatrix}}{\begin{bmatrix} 5/8 & 1/4 & 1/4 \\ 1/4 & 5/8 & 1/4 \end{bmatrix}} = 1. (\log 3 = 1.6, \log 5 = 2.3).$$

a) $\sum_j p(x_j, y_j) = 1 \Rightarrow \frac{5}{8} + \frac{1}{4} + \frac{1}{2} = 1 \Rightarrow 7/8 \cdot 2_1 = 8 \Rightarrow 2_1 = \frac{1}{7}$

$$\sum_i p(x_i, y_i) = 1 \Rightarrow 2_2 + \frac{1}{4} + \frac{5}{8} = 1 \Rightarrow 2_2 = \frac{1}{8}.$$

$$P(Y|X) = \begin{bmatrix} 5/8 & 1/4 & 1/8 \\ 1/8 & 1/4 & 5/8 \end{bmatrix} \quad \text{log } p(x_i)$$

- uniformitatea de intrare

b). calc. $H(Y)$, $H(Y|X)$, $H(X,Y)$, $H(X|Y)$, $I(X,Y)$, c

~~$P(Y|X) = P(Y|X) = \begin{bmatrix} 5/8 & 1/4 & 1/8 \\ 1/8 & 1/4 & 5/8 \end{bmatrix}.$~~

$$P(X,Y) = P(X) \cdot P(Y|X) = \begin{bmatrix} 1/2 & 0 \\ 0 & 1/2 \end{bmatrix} \begin{bmatrix} 5/8 & 1/4 & 1/8 \\ 1/8 & 1/4 & 5/8 \end{bmatrix} =$$

$$= \begin{bmatrix} \frac{5}{16} & \frac{1}{8} & \frac{1}{16} \\ \frac{1}{16} & \frac{1}{8} & \frac{5}{16} \end{bmatrix} = \underline{P(X|Y)}.$$

$$H(Y) = \sum_{i=1}^3 p_{yi} \log_2 p_{yi} = \left[\frac{6}{16} \cdot \log_2 \frac{6}{16} + \frac{2}{8} \cdot \log_2 \frac{2}{8} \right].$$

$$H(X) = \sum_{i=1}^2 p_{xi} \log_2 p_{xi} = C = \log m + \sum_{j=1}^m p(y_j|x_i) \cdot \log$$

$$H(X|Y) = \sum_{i=1}^2 \sum_{j=1}^3 p(x_i, y_j)$$

$$m = mn \cdot y - c' \\ C = \log 3 + (p(y_1|x_1) \cdot \log p(y_1|x_1) + p(y_2|x_2) \cdot \log p(y_2|x_2))$$

c)

(4)

$$\begin{array}{r}
 \frac{24}{128} \cdot 2 = \frac{27}{64} \\
 \frac{24}{64} \cdot 2 = \frac{27}{32} \\
 \frac{24}{32} \cdot 2 = \frac{27}{16} = 1 \text{ } \boxed{1} \\
 \frac{11}{16} \cdot 2 = \frac{11}{8} \\
 \frac{3}{8} \cdot 2 = \frac{3}{9} \\
 \frac{3}{4} \cdot 2 = \frac{3}{2} \\
 \frac{1}{2} \cdot 2 = 1 \text{ } \boxed{1}
 \end{array}$$

$\{10/51$	T	$\} \{ 32/51 \}$	000
$9/51$	-		001
$7/51$	A		010
$6/51$	L		011
<hr/>			
29	$6/51$	T	100
	$4/51$	S	101
	$4/51$	C	110
	$3/51$	I	1110
	$1/51$	M	11110
	$1/51$	N	11111
		$29/51$	-