

IMPORTANT: Explain your answer briefly. Do not just write a short answer or fill the assembly code.

1. [10pts] The following code fragment has a potential vulnerability.

```
void* copy_elements(void *ele_src[], int ele_cnt, size_t ele_size) {
    /*
     * Allocate buffer for ele_cnt objects, each of ele_size bytes
     * and copy from locations designated by ele_src
     */
    void *result = malloc(ele_cnt * ele_size);
    if (result == NULL)
        /* malloc failed */
        return NULL;
    void *next = result;
    int i;
    for (i = 0; i < ele_cnt; i++) {
        /* Copy object i to destination */
        memcpy(next, ele_src[i], ele_size);
        /* Move pointer to next memory region */
        next += ele_size;
    }
    return result;
}
```

- A. [4pts] Write possible values for ele_cnt and ele_size to crash the application/system. Explain the reason.

- B. [6pts] Write extra C codes to add before the malloc call to prevent the crash.

2. [10pts][floating point data representation]

- A. [4pts] Assume variables `x`, `f`, and `d` are of type `int`, `float`, and `double`, respectively. (Neither `f` nor `d` equals to +infinity, -infinity, or NaN). For each of the following expressions, either argue that it is always true or give a counterexample if it is not.

A-1) `x == (int) (double) x`

A-2) `f == -(-f)`

A-3) `1.0/2 == 1/2.0`

A-4) `d*d >= 0.0`

A-5) `(f+d) - f == d`

- B. [3pts] Write the rounded binary numbers for the following values. They should be rounded to nearest 1/4 (2 bits right of binary point, and must use “round-to-even” rule).

Explain the advantage of such “round-to-even” rule, compared to round-down or round-up.

10.00011 => _____

10.00110 => _____

10.11100 => _____

10.10100 => _____

- C. [3pts] Explain how a floating point compare instruction (fcmp) can be implemented for the IEEE fp format. How will it be different from the integer compare instruction (cmp)?

3. [8pts] Answer the two questions for the following C function and the corresponding assembly code.

```
long rfun(unsigned long x) {  
  
    if ( _____ )  
        return _____;  
  
    unsigned long nx = _____;  
    long rv = rfun(nx);  
    return _____;  
}
```

```
rfun:  
    pushq    %rbx  
    moveq    %rdi, %rbx  
    movl    $0, %eax  
    testq   %rdi, %rdi  
    je      .L2  
    shrq    $2, %rdi  
    call    rfun  
    addq    %rbx, %rax  
.L2:  
    popq    %rbx  
    ret
```

A. [3pts] What value in the C code does rfun store in %rbx?

B. [5pts] Fill in the missing expressions in the C code.

5. [10pts] Find the minimum number of operations to implement the given functions.

- Points will not be given if functions are implemented more than your minimum number of operations.
- You should justify your solution by implementing the functions.
- Assignment operator '=' is legal. But it will not be counted as an operator.

(1) addOK [5pts]

Description	Determine if can compute x+y without overflow
Examples	addOK(0x80000000,0x80000000) = 0 addOK(0x80000000,0x70000000) = 1
Legal Ops	! ~ & ^ + << >>

The minimum number of operations to implement this function is _____.

```
int addOK(int, int)
{
}
```

6. [10pts] Solve the following problems with the given assembly code.

C code
void phase(char *input) { int i; int numbers[6]; read_six_numbers(input, numbers); for(i = 1; i < 6; i++) { if ([]) explode_bomb(); } }

Assembly
Dump of assembler code for function phase: 0x000055555555522a <+0>: push %rbp 0x000055555555522b <+1>: push %rbx 0x000055555555522c <+2>: sub \$0x28,%rsp 0x0000555555555230 <+6>: mov %fs:0x28,%rax 0x0000555555555239 <+15>: mov %rax,0x18(%rsp) 0x000055555555523e <+20>: xor %eax,%eax 0x0000555555555240 <+22>: mov %rsp,%rbp 0x0000555555555243 <+25>: mov %rsp,%rsi 0x0000555555555246 <+28>: callq 0x555555555820 <read_six_numbers> 0x000055555555524b <+33>: mov %rsp,%rbx 0x000055555555524e <+36>: add \$0x14,%rbp 0x0000555555555252 <+40>: jmp 0x55555555525d <phase+51> 0x0000555555555254 <+42>: add \$0x4,%rbx 0x0000555555555258 <+46>: cmp %rbp,%rbx 0x000055555555525b <+49>: je 0x55555555526f <phase+69> 0x000055555555525d <+51>: mov (%rbx),%eax 0x000055555555525f <+53>: lea 0x2(%rax,%rax,2),%eax 0x0000555555555263 <+57>: cmp %eax,0x4(%rbx) 0x0000555555555266 <+60>: je 0x555555555254 <phase+42> 0x0000555555555268 <+62>: callq 0x5555555557fa <explode_bomb> 0x000055555555526d <+67>: jmp 0x555555555254 <phase+42> 0x000055555555526f <+69>: mov 0x18(%rsp),%rax 0x0000555555555274 <+74>: xor %fs:0x28,%rax 0x000055555555527d <+83>: jne 0x555555555286 <phase+92> 0x000055555555527f <+85>: add \$0x28,%rsp 0x0000555555555283 <+89>: pop %rbx 0x0000555555555284 <+90>: pop %rbp 0x0000555555555285 <+91>: retq 0x0000555555555286 <+92>: callq 0x555555554ed8 End of assembler dump.

A. Fill the C statement in the blank, based on the corresponding assembly codes. [5pts]

B. What is the solution if the first number of the solution is “1”? [5pts]

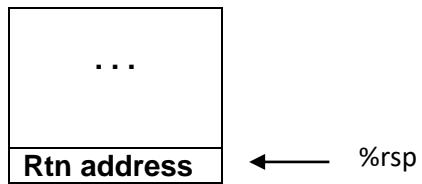
7. [7pts] Solve the following problems with the given C code and assembly code

C code
long mul(long *p,long val) { long x=*p; long y=x*val; *p=y; return x; } long call_mul(long x) { long v1=10; long v2=mul(&v1,300); return x+v2; }

mul: movq (%rdi), %rax imul %rax, %rsi movq %rsi, (%rdi) ret call_mul: subq \$16, %rsp movq \$1000, 8(%rsp) [] [] call mul ← addq 8(%rsp), %rax addq \$16, %rsp ret

- A. [3pts] Fill the assembly code in the blank(Assembly) [3pts]

B. [4pts] Assume that the program starts with the call_mul function. Draw the stack frame when the program **finishes executing “call mul” instructions.**



Initial Stack Structure