Gathered notes from:

2 Functor

 $\bullet$  Haskell Programming from First Principles [1]

todo

# 1 Monoid

 ${\rm todo}$ 

3 Applicative

4 Monad

todo todo

## 5 Arrow operator as Functor, Applicative, 5.3 Monad Monad

```
let F = (->) r = r ->
```

### 5.1 Functor

```
((->) r) = (r ->) as a functor
(r ->) * expects a type as argument
instances of (->) r * as a type class
examples of other functors: [] *, Maybe *
functor as a type constructor
```

### fmap:

```
<$>:: (a -> b) -> F a -> F b
<$>:: (a -> b) -> ((->) r) a -> ((->) r) b
<$>:: (a -> b) -> (r -> a) -> (r -> b)
<$>:: (a -> b) -> (r -> a) -> r -> b
composition operator:
(.) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c
therefore,
```

### 5.1.1 example

<\$> = (.) where F = (->) r

```
(+) <$> (*2)
(+) . (*2)
\x -> (+) ((*2) x)
\x -> (+) (x*2)
\x -> x*2 :: a -> a
(\x -> (+) (x*2)) :: a -> (a -> a)
(\x -> (+) (x*2)) :: a -> a -> a
(\x -> ((x*2)+) :: a -> a -> a
(\x -> (\y -> (x*2) + y) :: a -> a -> a
```

## 5.2 Applicative

```
<*>:: F (a -> b) -> F a -> F b
<*>:: ((->) r) (a -> b) -> ((->) r) a -> ((->) r) b
<*>:: (r -> a -> b) -> (r -> a) -> (r -> b)
h \leftrightarrow h = \r \rightarrow g r (h r)
where g :: r \rightarrow a \rightarrow b
where \hat{h} :: r \rightarrow a
pure :: a -> F a
pure :: a -> (->) r a
pure :: a -> r -> a
pure = const
```

```
5.2.1 example
(+) <$> (*2) <*> (+10)
(+) . (*2) <*> (+10)
(\x -> (+) (x*2)) <*> (\x -> x + 10)
\x -> (+) (x*2) (x+10)
types:
\x -> :: (->) r
(+) (x*2) :: a -> b where x is fixed
\x -> (+) (x*2) :: ((->) r) a -> b
\x -> x + 10 :: r -> a = ((->) r) a
(+) (x*2) (x+10) :: b where x is fixed
\x -> (+) (x*2) (x+10) :: ((->) r) b
```

<\*> :: (((->) r) a -> b) -> (((->) r) a) -> (((->) r) b)

thus types are as expected for applicative

return :: a -> F a

```
return :: a -> (->) r a = a -> r -> a
const :: a -> r -> a
return = const
bind:
>>= :: F a -> (a -> F b) -> F b
>>= :: (->) r a -> (a -> (->) r b) -> (->) r b
>>= :: (r -> a) -> (a -> r -> b) -> r -> b
g \gg h = \r \rightarrow h (g r) r
where g :: (r \rightarrow a)
where h :: (a \rightarrow r \rightarrow b)
g >>= h = \r \rightarrow (\x y z \rightarrow x z y) h (g r) r
g >>= h = \r \rightarrow (flip h) r (g r)
 g >>= h = (flip h) <*> g
```

# References

 $[1]\,$  Allen & Moronuki. Haskell programming from first principles, 2016.