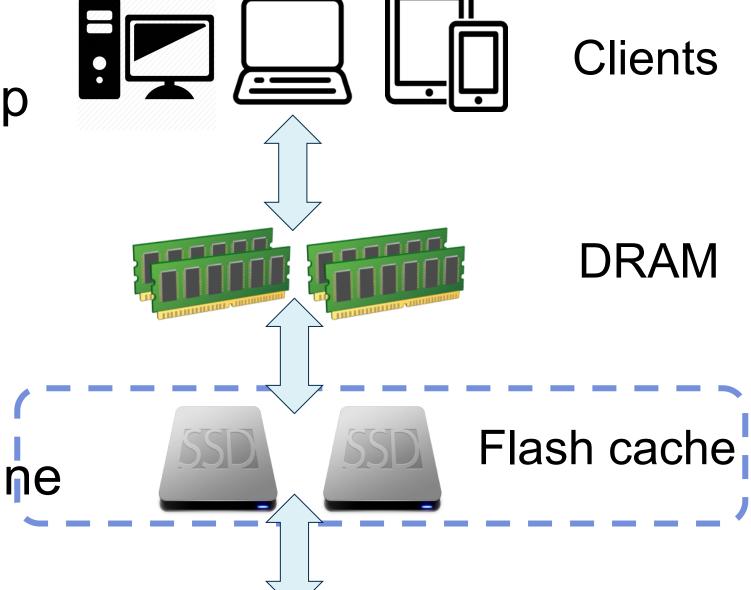
# Erasing Belady's Limitations: In Search of Flash Cache Offline Optimality

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## Motivation

- Using SSDs for an HDD cache
  - SSD cache btw. DRAM & HDD
    - Goal: to balance performance against endurance
  - Nitro [ATC'14], CacheDedup [FAST'16]
  - RIPQ [FAST'15]
- Pannier [Middleware'15]
- Are we doing well?
  - Comparison against an offline "best case"
  - But what is the offline optimal for flash cache?



**HDDs** 

## Background: Offline Optimality

- Belady's MIN: A simple offline caching alg. when the next access is known
  - Evicts the item that will be used furthest in the future
- Yields the optimal read hit ratio (RHR)
- Ignorant of minimizing erasures/block/day (EPBPD)
- MIN is not able to provide the optimal erasures in the context of flash caching
  - MIN inserts items that won't actually be read

## Offline Flash Caching

- Objectives
  - Minimize erasures s.t. maximal RHR
    - Never insert items if it does not increase RHR
  - Other objectives possible as discussed in the paper
- True optimal vs. Heuristic
  - Complexity of true optimal?
  - Approximation is the focus of this work

## 1) On a read miss, check if the

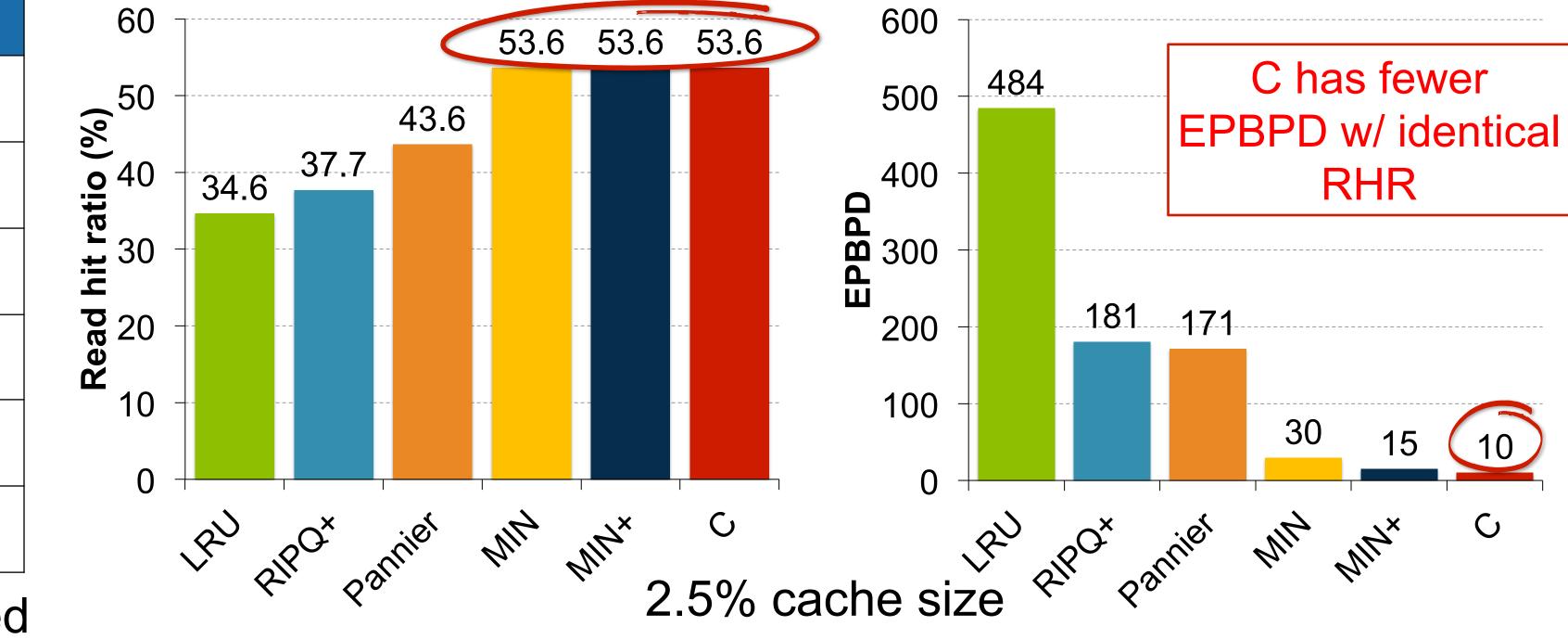
- 2 Pack a block into in-RAM write buffer and insert it to MIN PQ\*
- 3 Disperse the write buffer into containers and write them to the flash cache
- 4 On a read hit, read the block from the flash cache
- 5 At garbage collection, copyforward (CF) the valid block to the write buffer by checking if it will be reread before evicted \*PQ=Priority queue

#### Block Valid Container block will be read before evicted state: Block Block Index lookup Evict-pending Block 5 **Container PQ** Earliest to die Latest to die blocks MIN PQ Write buffer Next ref. furthest Next ref. nearest 3 **RAM** Flash

## Comparing Algorithms

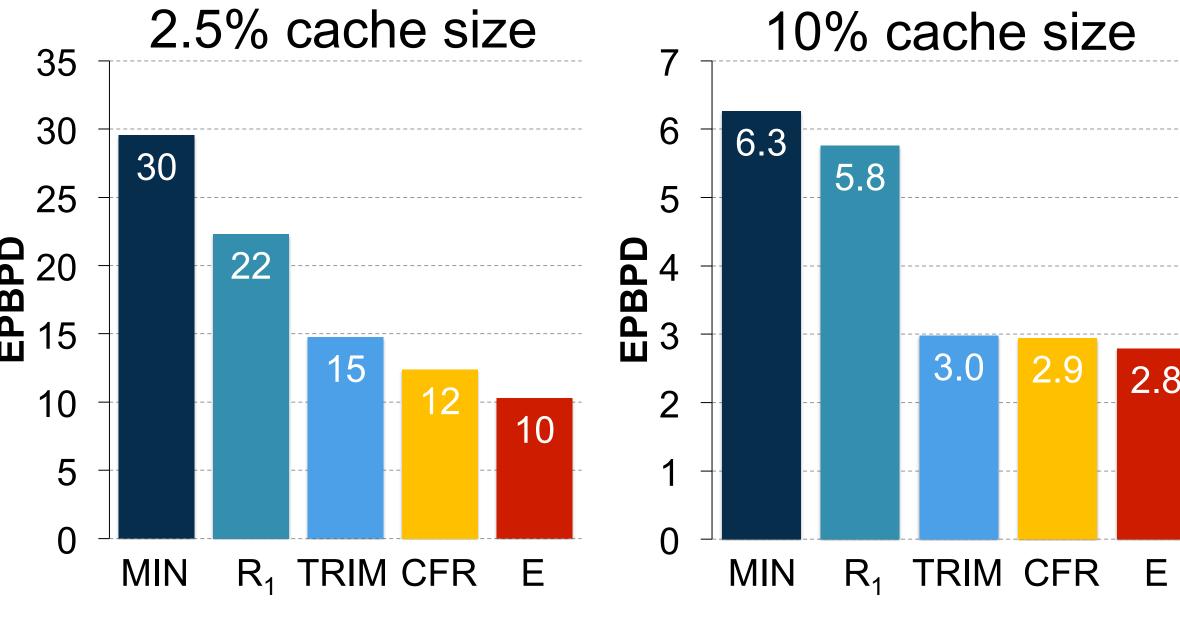
Container-optimized Heuristic

Policy	Description	<b>O</b> *	C+
LRU	Least recently used	X	X
RIPQ+	Static web content	X	X
Pannier	Handles divergent containers	X	1
MIN	Don't insert data w/ furthest next ref	<b>✓</b>	X
MIN+	Don't insert data evicted w/o read	<b>✓</b>	<b>✓</b>
C	Our container-optimized heuristic	<b>✓</b>	<b>/</b>
	*O=Offline. +C=Container-	optim	ize



## **Evaluating Optimization Techniques**

- R<sub>1</sub>: Omit insertions w/ no reread
- TRIM: Notify GC to omit dead blocks
- CFR: Avoid wasted ☐ 15 CF blocks
- E: Segregate blocks by evict timestamp



## Conclusion

- Important to have a baseline for the offline optimal considering both RHR and endurance
- Additional optimizations may be possible to move our 2 40 heuristic to the true optimal

