

Chapter 2

68HC12 Assembly Programming

Sample Program

Directive : Tells loader where to put program

ORG \$4000

Label	Opcode	Operand	Comment
main:	LDAA	\$800	; A = m[\$800]
	ADDA	\$801	; A = A + m[\$801]
	ADDA	\$802	; A = A + m[\$802]
	STAA	\$805	; m[\$805] = A

END

Directive : Tells assembler where program finished

Assembler Directives

- Commands used by development tools when building program

Example: END

- Ends a program to be processed by an assembler
- Any statement following END directive is ignored

Example: ORG

- Assembler uses **location counter** to keep track of memory
location where next machine code byte should be placed.
- Sets a new value for the location counter

Data Declarations

Write assembly program to implement the following C code.

```
byte aa, bb, cc, dd;

dd = aa + bb - cc;
```

Data Declaration Directives

- Used to declare variables.
- Often preceded by **org** directive.
- Many variations: Choose to suit the job
 - Two sets for byte and word

```
org $800
```

```
aa rmb 1
```

```
bb rmb 1
```

```
cd rmb 1
```

```
dd rmb 1
```

Directive :
Reserve
Memory Byte

```
org $4000
```

```
ldaa aa ; ldaa $800
```

```
adda bb ; adda $801
```

```
suba cc ; suba $802
```

```
staa dd ; staa $803
```

```
end
```

Directives to Simply Declare a Variable

- Simply allocate space.

	<u>Byte</u>	<u>Word</u>
Define Storage	<code>ds num</code>	<code>ds.w num</code>
Reserve Memory	<code>rmb num</code>	<code>rmw num</code>

Examples :

`buffer ds 100 ; Reserves 100 bytes at current location`

`dbuf ds.w 20 ; Reserve 40 bytes at current location.`

Directives to Declare and Initialize a Variable

	<u>Byte</u>	<u>Word</u>
Define byte	DB value	DW value
Define Constant Byte	DC.B value	DC.W value
Form Constant Byte	FCB value	FDB value

Examples :

```

        org $800
variable    dw    43
array      db    $11,$22,$33,$44
num        db    num-array

```

fcc (form constant character)

- Used to define a string of characters (a message).
- First (and last) character used as the delimiter.
- Last character must be the same as the first character.
- Delimiter must not appear in the string.
- Space character cannot be used as the delimiter.
- Each character is represented by its ASCII code.
- For example,

```
msg    fcc    "Please enter 1, 2 or 3:"
```

fill (fill memory)

- Allows the user to fill a certain number of memory locations with a given value.
- Syntax: **var_name fill value, count**
- For example,

```
space_line   fill    $20, 40
```


equ (equate)

- Assigns a value to a label.
- Using this directive makes our program more readable.

-Examples:

```
NUM_ELEMENTS equ 100  
DEFAULT equ 0  
array fill DEFAULT, NUM_ELEMENTS
```

Multiple Precision Arithmetic

- HC12: limited to 16-bit operations.
- Problem may involve larger-size data.

Example 2.7 Write a program to add two 4-byte numbers stored at \$800-\$803 and \$804-\$807, and save the sum at \$810-\$813.

	org	\$800	x =	\$ 88 33 9922
x	rmw	2	y =	\$ 11 88 8855
y	rmw	2		
	org	\$810	z =	\$ 99 BC 2177
z	rmb	4		

Multiple-Precision Solution:

Addition starts from the LSB and proceeds toward MSB.

```
        org    $4000
ldd     x+2    ; add and save the two LSBs
add     y+2    ;
std     z+2    ;

ldaa    x+1    ; add and save the second MSB
adca    y+1    ;
staa    z+1    ;

ldaa    x      ; add and save the MSB
adca    y      ;
staa    z      ;

        end
```

Assembling and running programs

(Demo)

Multiplication and Division

Table 2.1 Summary of 68HC12 multiply and divide instructions

Mnemonic	Function	Operation
EMUL EMULS MUL	unsigned 16 by 16 multiply signed 16 by 16 multiply unsigned 8 by 8 multiply	$(D) \times (Y) \rightarrow Y:D$ $(D) \times (Y) \rightarrow Y:D$ $(A) \times (B) \rightarrow A:B$
EDIV EDIVS	unsigned 32 by 16 divide signed 32 by 16 divide	$(Y:D) \div (X)$ quotient $\rightarrow Y$ remainder $\rightarrow D$ $(Y:D) \div (X)$ quotient $\rightarrow Y$ remainder $\rightarrow D$
IDIV IDIVS	unsigned 16 by 16 integer divide signed 16 by 16 integer divide	$(D) \div (X) \rightarrow X$ remainder $\rightarrow D$ $(D) \div (X) \rightarrow X$ remainder $\rightarrow D$

Example 2.10 Write instruction sequence to multiply 16-bit numbers stored at \$800-\$801 and \$802-\$803. Store product at \$900-\$903.

Solution:

ldd	\$800
ldy	\$802
emul	
sty	\$900
std	\$902

Example 2.11 Write instruction sequence to divide 16-bit number stored at \$820-\$821 into the 16-bit number stored at \$805-\$806. Store the quotient/remainder at \$900 and \$902, respectively.

Solution:

ldd	\$805	
ldx	\$820	
idiv		
stx	\$900	; store the quotient
std	\$902	; store remainder

Example 2.13 Write a program to convert a 16-bit number stored at \$800-\$801 into its ASCII representation as a decimal number.

Solution:

- Binary number converted to decimal ASCII: repeated division by 10.
- Largest 16-bit binary number = 65535 (five decimal digits).
- First division by 10: least significant digit; second division by 10: second least significant digit, etc.

```

        org      $800
data dc.w      12345      ; data to be tested
        org      $900
result        ds.b 5      ; reserve bytes to store the result

        org      $1000
ldd  data
ldy  #result
ldx  #10
ldiv
addb #$30      ; convert the digit into ASCII code
stab 4,Y      ; save the least significant digit
xgdx
ldx  #10

```

```
idiv
addb #$30
stab 3,Y      ; save second to least significant digit
xgdx
ldx  #10
idiv
addb #$30
stab 2,Y      ; save middle digit
xgdx
ldx  #10
idiv
addb #$30
stab 1,Y      ; save second most significant digit
xgdx
addb #$30
stab 0,Y      ; save most significant digit
end
```


Program Loops - Branches

Four **types** of branch instructions:

- **Unary (unconditional) branch:** always execute

BRA \$1000

- **Simple branches:** branch taken when specific bit of CCR in a specific status

BCC \$1000

- **Unsigned branches:** branches taken when comparison or test of **unsigned** numbers results in specific combination of CCR bits

BHI \$1000

- **Signed branches:** branches taken when comparison or test of **signed** quantities results in a specific combination of CCR bits

BGT \$1000

Example : Conditional Branches

Suppose A contains \$7F:

Unsigned Scenario

SUBA #\$80

BHI Bigger

Signed Scenario

SUBA #\$80

BGT Bigger

In each scenario, is the jump taken? Why?

Programmer MUST know how binary values are to be interpreted!
(e.g. value in AX above)

Branch Instructions use **Relative Mode** to encode the target


```

1:          =00001000          org $1000
2:      1000 B6 0800          ldaa $800
3:      1003 20 FB          bra  $1000
4:      1005                  END
    
```

8-bit relative offset

Let's execute-by-hand :

Fetch:	PC = \$1000	Instruction Register = ????????
Execute:	PC = \$1003	Instruction Register = B6 0800
Fetch:	PC = \$1003	Instruction Register = B6 0800
Execute:	PC = \$1005	Instruction be be executed = 20 FB
Fetch:	PC = \$1000	Instruction to be executed = 20 FB



Relative Mode

- Used only by branch instructions.
- Used even when programmer use labels for the target.
 - Assembler calculates actual branch offset (distance) from the instruction that follows the branch instruction.

minus . . .

bmi minus

- Two **categories** of Branches

- 1. Short Branches:** in the range of -128 ~ +127 bytes

- 8-bit opcode and a signed 8-bit offset.

- 2. Long Branches:** in the range of -32768 ~ +32767.

- 8-bit opcode and a signed 16-bit offset. Range of long relative mode is -32768 ~ +32767.

Compare and Test Instructions

- Condition flags need to be set up before simple and conditional branch instruction are to be executed.

Table 2.4 Summary of compare and test instructions

Mnemonic	Compare instructions	
	Function	Operation
CBA	Compare A to B	(A) - (B)
CMPA	Compare A to memory	(A) - (M)
CMPB	Compare B to memory	(B) - (M)
CPD	Compare D to memory	(D) - (M:M+1)
CPS	Compare SP to memory	(SP) - (M:M+1)
CPX	Compare X to memory	(X) - (M:M+1)
CPY	Compare Y to memory	(Y) - (M:M+1)
Test instructions		
	Mnemonic Function	Operation
TST	Test memory for zero or minus	(M) - \$00
TSTA	Test A for zero or minus	(A) - \$00
TSTB	Test B for zero or minus	(B) - \$00

Loop Primitive Instructions

- 68HC12 provides a group of instructions that either decrement or increment a loop count to determine if the looping should be continued.
- These are SHORT branches. *Counter* = A, B, D, X, Y, S

DBNE counter, rel (or **DBEQ**)

- Decrement counter and branch if != 0

IBNE counter, rel (or **IBEQ**)

- Increment counter and branch if != 0

TBEQ counter, rel (or **TBNE**)

- Test counter and branch if = 0
- where counter = A,B,D,X,Y or SP

Example Write a program to compute the sum of the first n whole numbers ($n < 255$), leaving the result in a word variable called “series”.

Solution:

```
N      equ 20
      org $800
series rmb 2
```

```
      org $4000
      ldab #N
      ldx  #0
```

next:

```
      abx
      dbne b, next
      stx  series
      bra  *
```

```
int series = 0;
for (byte i = N; i != 0; i-- ) {
    series += i;
}
```

```
; b = i = 0
; x = series = 0
```

```
; x (series) += i
; i--; while b != 0
; update series var.
; end of program
```

Example Write a code fragment showing how to implement the following pseudocode

```
boolean done = FALSE;
while ( ! done ) {
    ...
    ; at some point
    done = TRUE;
}
```

Solution:

```
TRUE    equ 1
FALSE   equ 0
```

```
        org $4000
        ldab #FALSE      ; b = done
notDone: TBNE b, end
        ; ...
        ; At some point,
        ldab #TRUE
        bra  notDone
end:     bra    *
```


Special Conditional Branch Instructions

```
[ <label> ]    BRCLR   (opr) (msk) (rel)  [ <comment> ]
[ <label> ]    BRSET   (opr) (msk) (rel)  [ <comment> ]
```

where

opr: memory location to check. Specified using direct/indexed addressing mode.

msk: 8-bit mask. Specifies which bits to check (those whose bit positions are 1s in the mask).

rel: branch offset (relative mode).

```
Example :      loop          inc    count
                                   ...
                                   brset  $66,$e0,loop
                                   ...
```

Question: When will the branch be taken ?

Example Write a program to test whether a variable is divisible by 4, leaving the boolean result in accumulator A.

Solution: A number divisible by 4 would have the least significant two bits equal 0s.

```
FALSE equ    0
TRUE  equ    1
      org     $800
variable db $19

      org     $4000
      brclr   variable,$03,yes ; check bits 1 and 0
      ldaa    #FALSE
      bra     continue
yes ldaa    #TRUE
continue

      . . .
      end
```

Example A robot has four motors, each of which can be off, on in forward direction, or on in reverse direction. Status of these motors are written by the robot into a status word, called “motors” in the following bitmask formation.

7	6	5	4	3	2	1	0
Motor1	Motor2	Motor3	Motor4				

where the two bits for each motor are set according

01	forward
10	reverse
11	off

Write a code fragment that waits while motor1 is off before continuing on.

Solution:

```

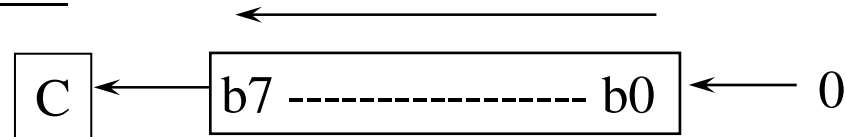
org      $800
motors rmb      1
org      $4000
brset    motors,$C0,*    ; check bits 7 and 6
. . .
end

```

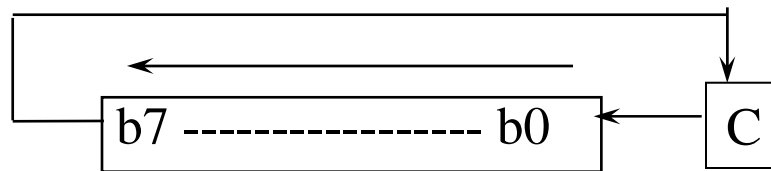
Shift and Rotate Instructions

Shift versus Rotate

Shift



Rotate



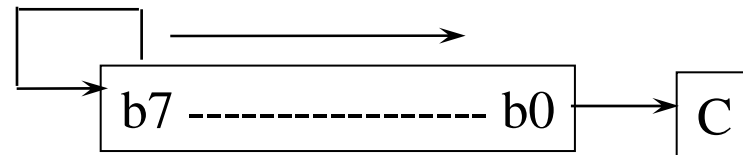
Arithmetic versus Logical Shift

- Different only when shifting right.

Logical : Zero is injected



Arithmetic : MSB is injected.
Why ?



Shift Instructions

- Apply to a memory location, accumulators A, B and D.
 - Memory operand specified using extended or indexed addressing modes.

8-bit arithmetic shift left instructions (**LOGICAL**: A → L):

ASL	opr	ASR	opr
ASLA		ASRA	
ASLB		ASRB	

16-bit arithmetic shift left instruction:

ASLD	ASRD
------	------

“9-bit” rotate instructions.

ROL	opr	ROR	opr
ROLA		RORA	
ROLB		RORB	

Example 2.23 Write a program to count number of 0s in the 16-bit number stored at \$800-\$801 and save the result in \$805.

Solution:

- * The 16-bit number is shifted to the right 16 times.
- * If the bit shifted out is a 0 then increment the 0s count by 1.

```

                org $800
db              $23,$55          ; test data
                org $805
zero_cnt rmb 1
lp_cnt  rmb 1

                org $4000
clr zero_cnt      ; initialize the 0s count to 0
ldaa #16
staa lp_cnt
ldd $800        ; place the number in D
loop  lsr      ; shift the lsb of D to the C flag
      bcs chkend ; is the C flag a 0?
      inc zero_cnt ; increment 1s count if the lsb is a 1
chkend dec lp_cnt ; check to see if D is already 0
      bne loop
forever bra forever ; = bra *
end

```

Boolean Logic Instructions

- Changing a few bits are often done in I/O applications.
- Boolean logic operation can be used to change a few I/O port pins easily.

Table 2.8 Summary of Boolean logic instructions

Mnemonic	Function	Operation
ANDA <opr> ANDB <opr>	AND A with memory AND B with memory	$A \leftarrow (A) \bullet (M)$ $B \leftarrow (B) \bullet (M)$
EORA <opr> EORB <opr>	Exclusive OR A with memory Exclusive OR B with memory	$A \leftarrow (A) \oplus (M)$ $B \leftarrow (B) \oplus (M)$
ORAA <opr> ORAB <opr>	OR A with memory OR B with memory	$A \leftarrow (A) + (M)$ $B \leftarrow (B) + (M)$
COM <opr> COMA COMB	One's complement memory One's complement A One's complement B	$M \leftarrow \$FF - (M)$ $A \leftarrow \$FF - (A)$ $B \leftarrow \$FF - (B)$
NEG <opr> NEGA NEGB	Two's complement memory Two's complement A Two's complement B	$M \leftarrow \$00 - (M)$ $A \leftarrow \$00 - (A)$ $B \leftarrow \$00 - (B)$

Program Execution Time

- 68HC12 uses E clock as timing reference.
- Frequency of E clock is half of that of the crystal oscillator.
- Many applications require generation of time delays.

Creation of a time delay involves two steps:

1. Select a sequence of instructions that takes a certain amount of time to execute.
2. Repeat selected instruction sequence for an appropriate number of times.

Example: instruction sequence on following slide: 40 E cycles. Repeating this sequence certain number of times, other delays can be created.

Assume the 68HC12 runs under a crystal oscillator with a frequency of 16 MHz (E frequency: 8 MHz; clock period is 125 ns). The instruction sequence on next page will take 5 μ s.

Example 2.25 Write a program loop to create a delay of 100 ms.

Solution:

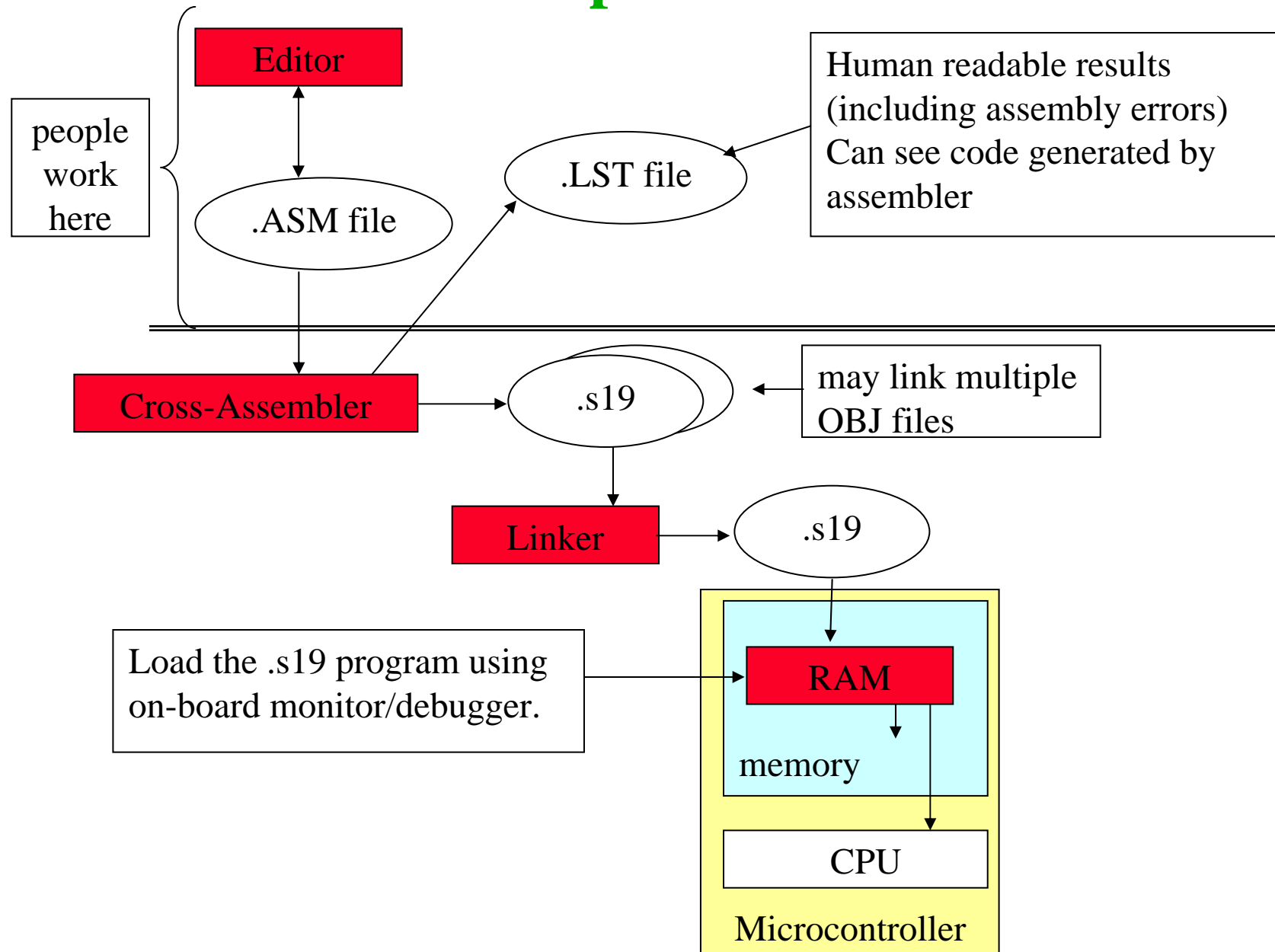
The following instruction sequence creates a delay of 100 ms.

```

        ldx      #20000
loop    psha                ; 2 E cycles
        pula                ; 3 E cycles
        psha
        pula
        psha
        pula
        psha
        pula
        psha
        pula
        psha
        pula
        psha
        pula
        nop                ; 1 E cycle
        nop                ; 1 E cycle
        dbne     x,loop     ; 3 E cycles

```

Development Process



2 Pass Assembly Process

Each **pass**: processes statements in **.ASM** file sequentially from start to finish

0th Pass: Macro Processor. Search/replace Macros and Constants (EQU)

1st Pass: for each statement:

1. check syntax
2. allocate any memory needed for image
 - memory declarations
 - instruction: opcode + operands
3. label definition: assign value to label; record (label, value) in Symbol Table
4. Generate instructions' opcodes/operands
5. Update .LST file

Any kind of errors: stop, else

2nd Pass:

- Associate addresses to not resolved symbols (forward jumps or variables; Symbol table)
- may require calculating offsets – may result in errors; e.g. trying to branch too far (target out of range) -
- write results to **.LST** file
- no errors? – write results to **.s19** file

Assembling: an example

After 1st Pass:

- syntax OK
- bytes have been allocated to statements
- Symbol Table constructed:

	<u>Symbol</u>	<u>Value</u>			
	start		0000H		
	<u>\$</u>	<u>binary image</u>			
1:		=00001000	org	\$1000	
2:	1000	79 ????	clr	zero_cnt	; initialize the zero count to 0
3:	1003	86 10	ldaa	#16	
4:	1005	7A ????	staa	lp_cnt	; initialize loop count to 16
5:	1008	FC 0800	ldd	\$800	; place the 16-bit number in D
6:	100B	49 again	lsrd		
7:	100C	25 ??	bcs	chk_end	; have we tested all 16 bits yet?
8:	100E	72 ????	inc	zero_cnt	
9:	1011	73 ???? chk_end	dec	lp_cnt	
10:	1014	26 F5	bne	again	
11:	1016	20 FE forever	bra	forever	
12:					
13:		=00000800	org	\$800	
14:	0800	23 55	db	\$23,\$55	
15:		=00000805	org	\$805	
16:	0805	+0001 zero_cnt	rmb	1	
17:	0806	+0001 lp_cnt	rmb	1	
18:			end		

Label definition: in 1st pass – put value in Symbol table

Note: Do not confuse \$ with PC!

\$ = artifact of assembler

PC = artifact of processor

Assembling: an example

After 2nd Pass:

- Missing symbols updated

	<u>\$</u>	<u>binary image</u>			
1:		=00001000	org	\$1000	
2:	1000	79 0805	clr	zero_cnt	; initialize the zero count to 0
3:	1003	86 10	ldaa	#16	
4:	1005	7A 0806	staa	lp_cnt	; initialize loop count to 16
5:	1008	FC 0800	ldd	\$800	; place the 16-bit number in D
6:	100B	49 again	lsrd		
7:	100C	25 03	bcs	chk_end	; have we tested all 16 bits yet?
8:	100E	72 0805	inc	zero_cnt	
9:	1011	73 0806 chk_end	dec	lp_cnt	
10:	1014	26 F5	bne	again	
11:	1016	20 FE forever	bra	forever	
12:					
13:		=00000800	org	\$800	
14:	0800	23 55	db	\$23,\$55	
15:		=00000805	org	\$805	
16:	0805	+0001 zero_cnt	rmb	1	
17:	0806	+0001 lp_cnt	rmb	1	
18:			end		