

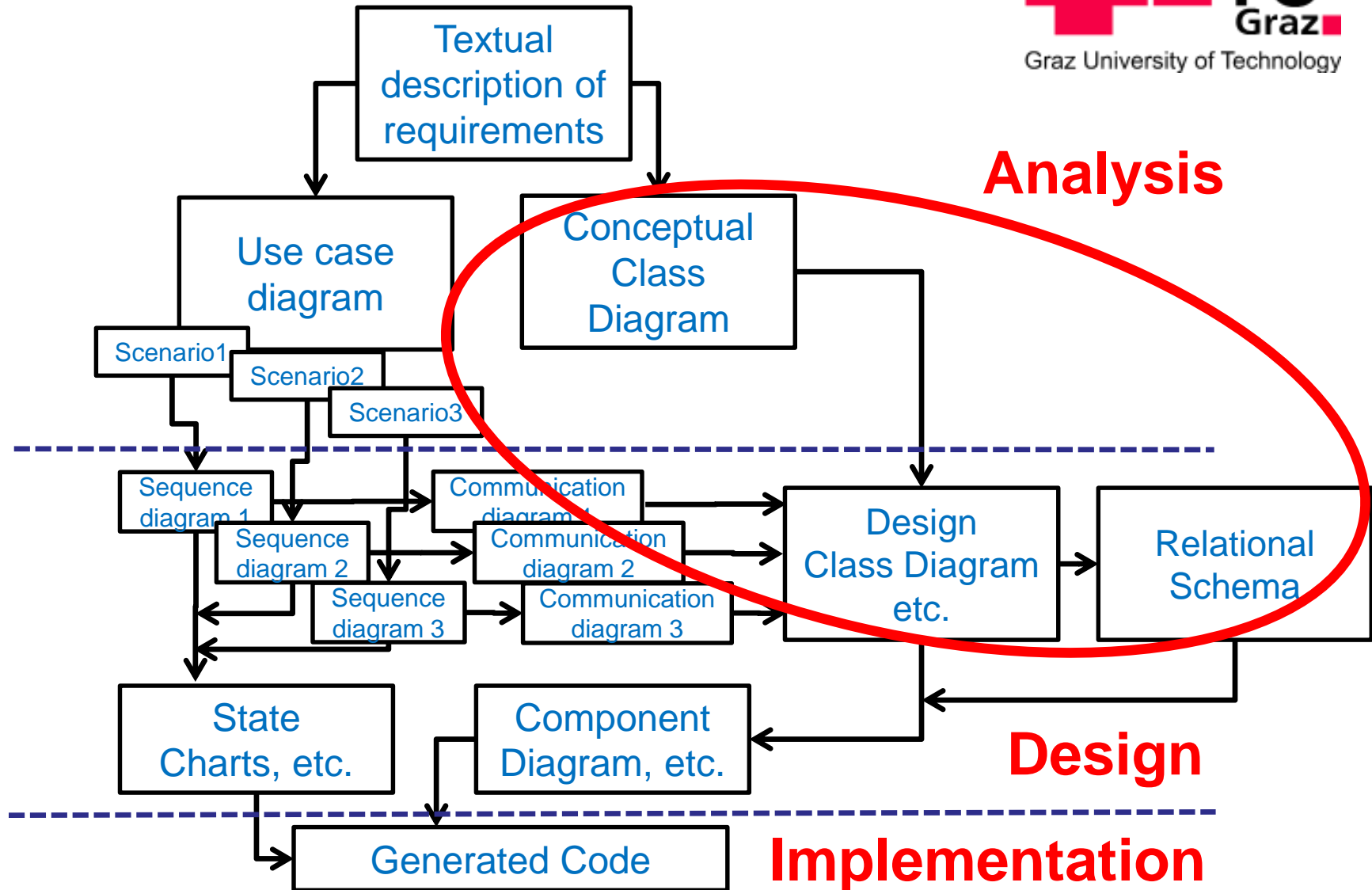
Object-Oriented Analysis & Design (OAD)

Object-Relational Mapping (ORM)

<https://youtu.be/HorIWBNoHEU>

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„Big Picture“



Database Systems

- Database (DB)
 - collection of data representing different aspects of the application domain
- Database management system (DBMS)
 - interface between user and DB
 - guarantees adequate and efficient data access, resistant to SW and HW failures, data available over a longer period
- Database system (DBS)
 - consists of a DBMS and one or more DBs

DBMS Languages

- Data description (definition) language (DDL)
 - definition of database schema
- Data manipulation language (DML)
 - creating (INSERT), adapting (UPDATE), and deleting database contents (DELETE)
- Query language (QL)
 - querying DB contents (SELECT) without changing the content itself

Relational Database

- Table (relation)
 - subset of tuples of the cross product of attribute domains together with column headings
 - e.g., persons →
- Attribute
 - ID, firstname, ...
- Domain
 - set of values of the same type
- Tuple
 - element of the cross product of attribute domains
 - e.g., (1, 'Albert', 'Einstein', 14.03.1879)

ID	firstname	lastname	date of birth
1	Albert	Einstein	14.03.1879
2	Nils	Bohr	07.10.1885
3	Werner	Heisenberg	05.12.1901
4	Max	Planck	23.04.1858
5	Christian	Doppler	23.11.1803
6	Ludwig	Boltzmann	20.02.1944
7	Erwin	Schrödinger	12.08.1887

Relational Database

- Superkey of a table (relation) R
 - set of attributes of R which allow the unique identification of each tuple in R
 - for example: {orderId,orderDate} in table “order”
- Key of a table
 - superkey with “minimal” set of attributes
 - often artificial key, for example: {orderId} or simply {ID}
- Candidate key
 - one key of a relation R , where R can have different keys, for example, {custID}, {SSN}

Relational Database

- Key attribute
 - attribute that belongs to a (candidate) key
 - for example: bookID in {authorID, bookID}
- Primary key
 - key used for identifying tuples in the table
 - for example: candidate keys {authorID} and {firstname, lastname, date of birth} → {authorID} selected as primary key
- Foreign key
 - set of attributes that occurs as primary key in another table (e.g., order.custID → customer.custID)

- Referential integrity
 - each value (set of values) that occurs in table R as a foreign key instance referencing table R' **must also exist as a primary key instance in R'**
 - for example: `order.custID = 5` → exists:
`customer.custID = 5`
- Entity integrity
 - each tuple of R must have a complete primary key instance, i.e., there must be **a value for each attribute of the primary key**

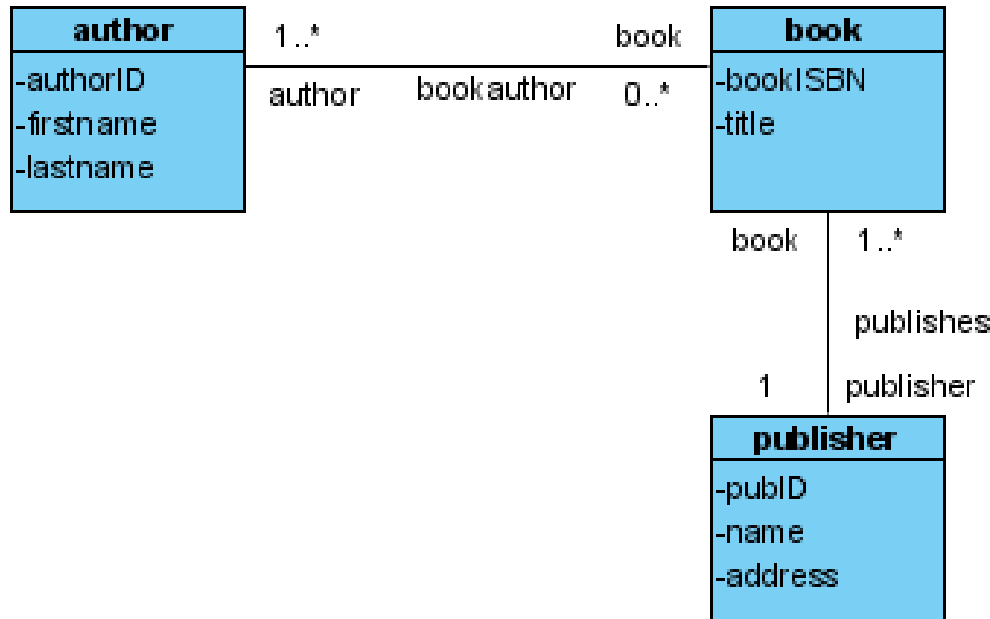
Example Tables (1)

- **customer**(custID, firstname, lastname, date of birth, address)
- **order**(orderID, date, *custID*)
- **orderpos**(posID, orderID, *productID*)
- **product**(productID, name, price)
- Used notation: **tablename**, primary key attribute, *foreign key attribute*

Example Tables (2)

- **author**(authorID, firstname, lastname)
- **book**(bookISBN, title, *pubID*)
- **bookauthor**(bookISBN, authorID)
- **publisher**(pubID, name, address)

Database Mapping (ORM)



rules of thumb
(R1 .. R7)

- **author**(authorID, firstname, lastname)
- **book**(bookISBN, title, *pubID*)
- **bookauthor**(bookISBN, authorID)
- **publisher**(pubID, name, address)

R1 (Classes)

- Rules R1..R7: translation of UML class diagrams to relational schemas
- For each class C (not association classes AC!) of the class diagram, a table C is created
 - table name = class name
 - all attributes of class C with cardinality max. 1 (~"single-valued") become attributes of table C
- Note: classes that are related to other classes by generalization might be treated differently (R7)
- Rules applied until no more R_i applicable

R1 (Example)



R2 (Primary Key)

- Identify the candidate keys in each table formed and specify one of them as primary (underline)
- Recommendation: if all candidate keys consist of more than one attribute, introduce an additional “artificial” attribute ID and specify it as primary key

R2 (Example)



→ • **author**(authorID, firstname, lastname)

R3 (Multivalued Attribute)

- For each attribute A of a class C (e.g. “customer”) with max. cardinality greater than 1 (e.g. “telno”) **create a separate table**:
 - table name = C_A (e.g. “**customer_telno**”)
 - primary key of C to be indicated as foreign key in C_A (e.g. *custID*)
 - attribute A (e.g. “**telno**”)
 - foreign key and A forms a minimal primary key of C_A (e.g. “*custID*, *telno*”)

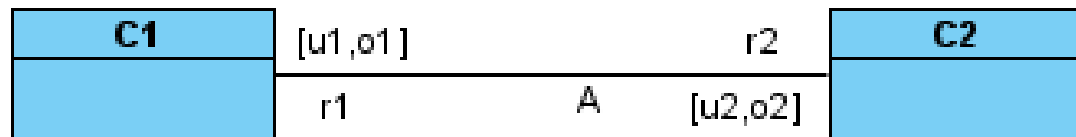
R3 (Example)

customer
-name [1]
-custID [1]
-address [1]
-telno [1..*]

- 
- **customer**(custID, name, address)
 - **customer_telno**(custID, telno)

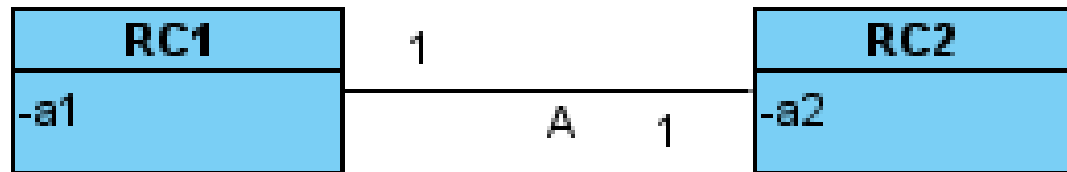
Association Connectivity

- Given an association A between the classes C1 and C2, with roles r1, r2 and the multiplicities [u1,o1] and [u2,o2]
 - Then A is called
 - **1:1 association** if $o1 = 1$ and $o2 = 1$
 - **1:n association** if $o1 = 1$ and $o2 \neq 1$
 - **n:m association** if $o1 \neq 1$ and $o2 \neq 1$
- } **connectivities**



R4 (1:1 association)

- An 1:1 association A between the classes RC1 and RC2 is mapped by extending the resulting tables RC1 and RC2 by the attributes a1 and a2 (foreign keys)



- **RC1(a1)**
 - **RC2(a2)**
-
- **RC1(a1, a2)**
 - **RC2(a2, a1)**

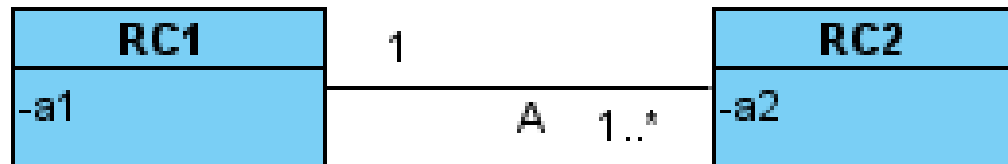
R4 (Example)



- **department**(deptID, name, *empID*)
- **employee**(empID, name, *deptID*)

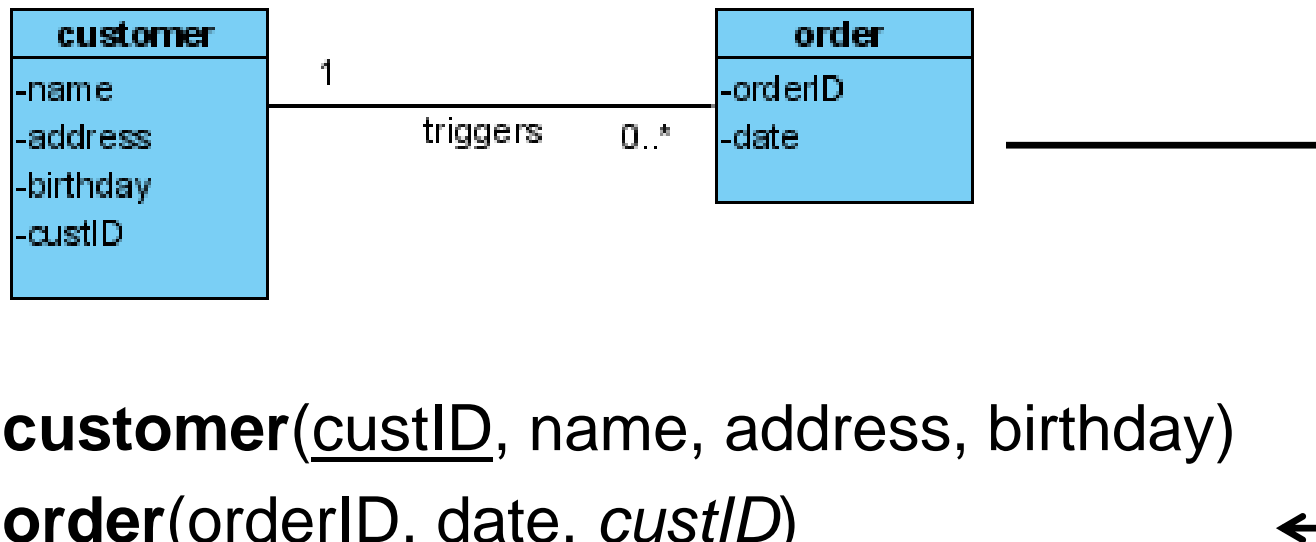
R5 (1:n association)

- An 1:n association A is mapped by extending the resulting table RC2 by the attribute a1 as foreign key



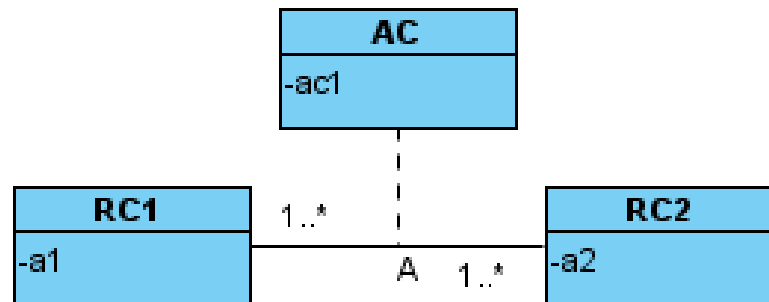
- | | | |
|-------------------------|---|-----------------------------|
| • rc1(<u>a1</u>) | → | • rc1(<u>a1</u>) |
| • rc2(<u>a2</u>) | | • rc2(<u>a2</u>, a1) |

R5 (Example)



R6 (m:n association)

- An **m:n association A** with the attributes $ac1, ac2, \dots, acn$ **between the classes RC1 and RC2** is mapped by creating the table A that includes aci and the primary key attributes of RC1 and RC2 as foreign keys



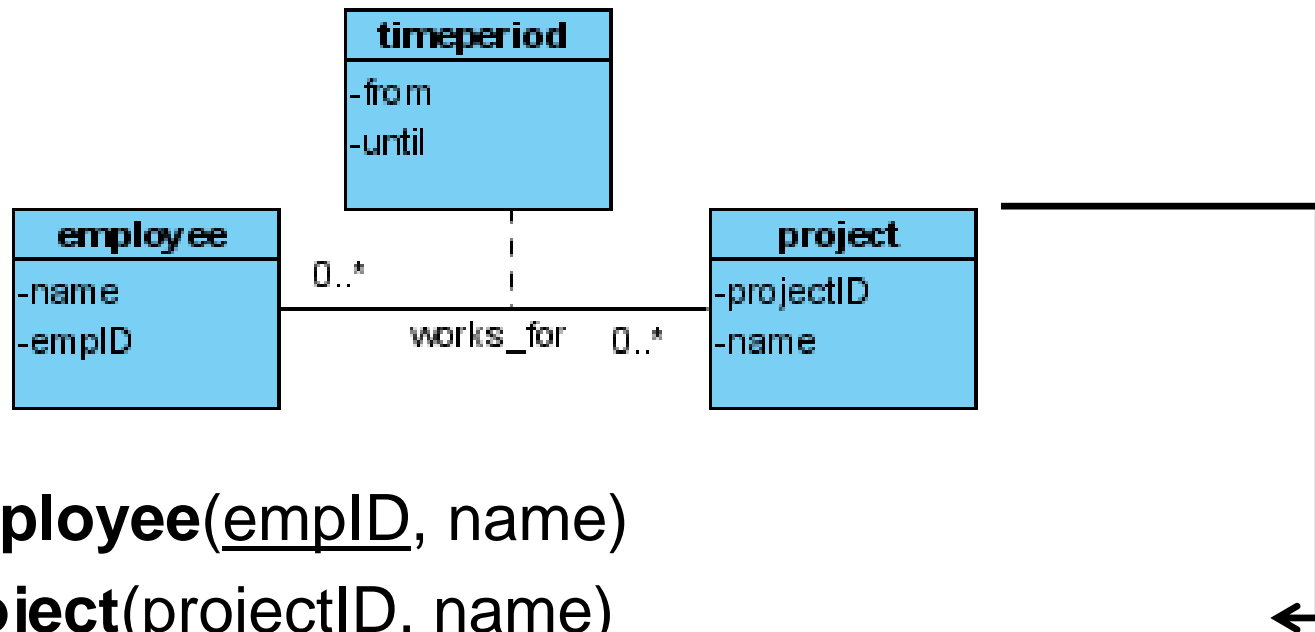
- **RC1(a1)**
 - **RC2(a2)**
-
- **RC1(a1)**
 - **RC2(a2)**
 - **A(a1, a2, ac1)**

R6 (Example 1)



- **author**(authorID, firstname, lastname)
- **book**(bookISBN, title, publisher)
- **bookauthor**(authorID, bookISBN)

R6 (Example 2)

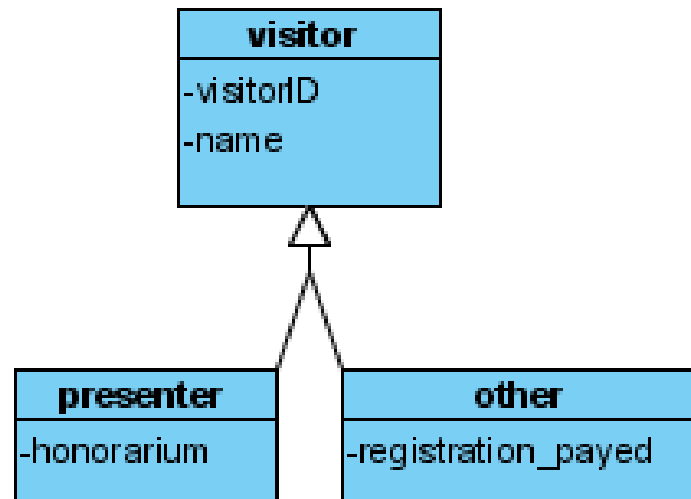


- **employee**(emplID, name)
- **project**(projectID, name)
- **works_for**(emplID, projectID, **from**, **until**)?
- **works_for**(emplID, projectID, **from**, **until**)?
- **works_for**(ID, emplID, projectID, from, until)

Rule 7 (Generalization)

- A generalization relationship between RCs and RCg can be mapped as follows:
- **Alternative 1 (ALT1):**
 - One table per class: extend RCs with the primary key of RCg (as foreign key)
- **Alternative 2 (ALT2):**
 - One table per concrete class: extend RCs with all attributes from RCg and delete RCg
- **Alternative 3 (ALT3):**
 - One table per generalization hierarchy: extend RCg by the attributes of RCs and delete RCs

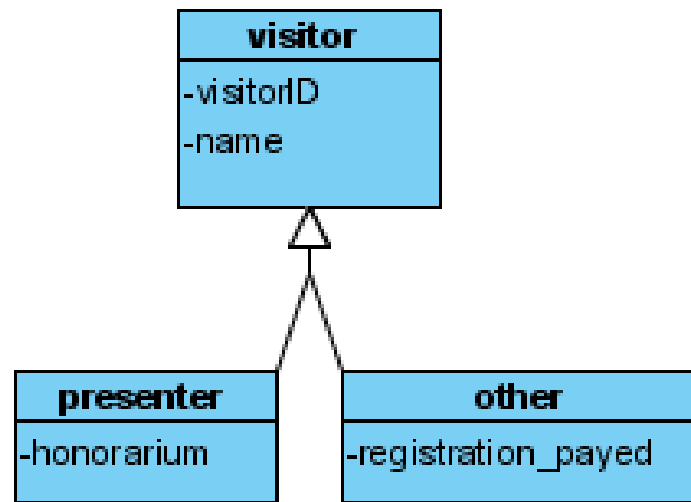
R7 (Example – ALT1)



Alternative 1:

- **visitor**(visitorID, name)
- **presenter**(visitorID, honorarium)
- **other**(visitorID, registration_payed)
- JOINS (-) but generalization hierarchy (+)

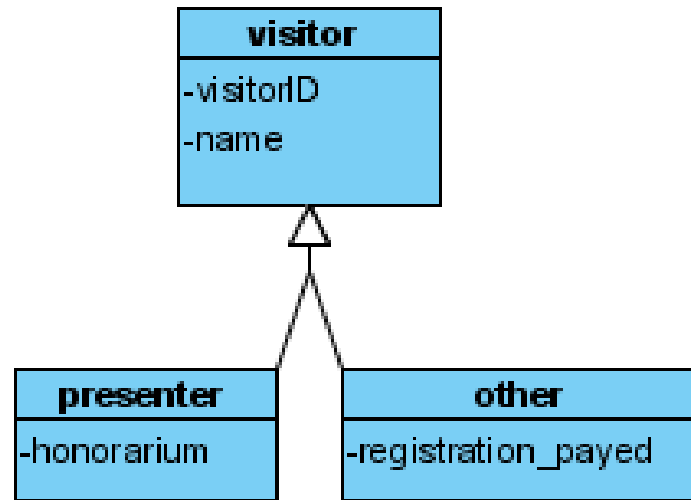
R7 (Example – ALT2)



Alternative 2:

- **presenter**(visitorID, name, honorarium)
- **other**(visitorID, name, registration_payed)
- No JOINS (+) but schema redundancy (-)

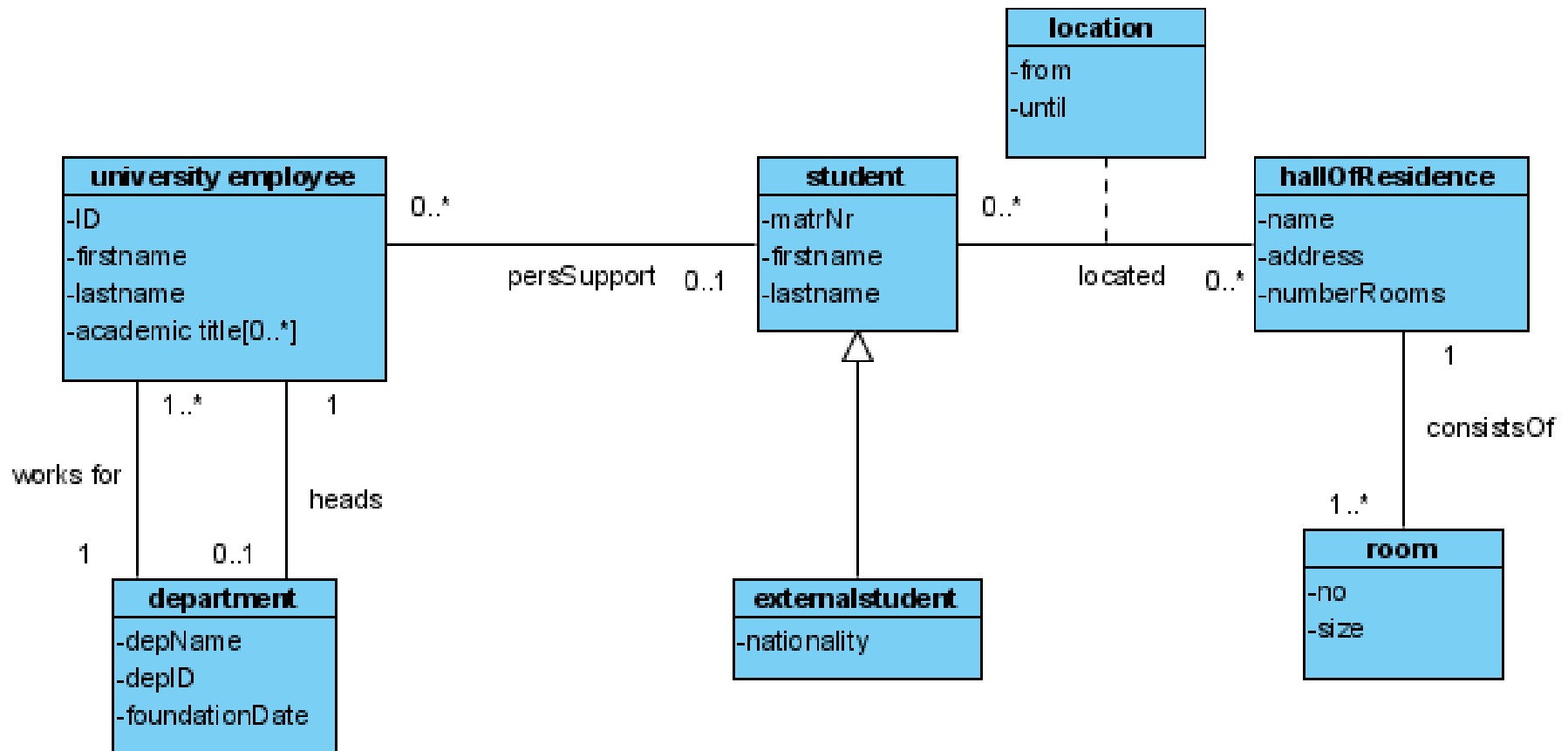
R7 (Example – ALT3)



Alternative 3:

- **visitor**(visitorID, name, honorarium, registration_payed)
- No JOINS (+) but more memory needed (-)

Integrated Example (Class Diagram)



Integrated Example (Database Structure)

- **universityemployee**(ID, firstname, lastname, *dep_w*, *dep_h*, *matrNr*)
- **universityemployee_academictitle**(academictitle, ID)
- **department**(depID, depName, foundationDate, *ID*)
- **student**(matrNr, firstname, lastname)
- **externalstudent**(matrNr, nationality)
- **hallOfResidence**(name, address, numberRooms)
- **room**(no, size, name)
- **located**(LID, *matrNr*, *name*, from, until)

Thanks!

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