



BlockSec

Security Audit Report for ExoSwap Bridge Contracts

Date: July 3, 2023

Version: 1.0

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Report Manifest

Item	Description
Client	ExoSwap
Target	ExoSwap Bridge Contracts

Version History

Version	Date	Description
1.0	July 3, 2023	First Release

About BlockSec BlockSec focuses on the security of the blockchain ecosystem and collaborates with leading DeFi projects to secure their products. BlockSec is founded by top-notch security researchers and experienced experts from both academia and industry. They have published multiple blockchain security papers in prestigious conferences, reported several zero-day attacks of DeFi applications, and successfully protected digital assets that are worth more than 5 million dollars by blocking multiple attacks. They can be reached at [Email](#), [Twitter](#) and [Medium](#).

Chapter 1 Introduction

1.1 About Target Contracts

Information	Description
Type	Smart Contract
Language	Solidity
Approach	Semi-automatic and manual verification

The target of this audit is the code repo of ExoSwap Bridge Contracts ¹ of the ExoSwap project. The ExoSwap project is a cross-chain bridge that bridges assets over multiple EVM-based chains.

The auditing process is iterative. Specifically, we would audit the commits that fix the discovered issues. If there are new issues, we will continue this process. The commit SHA values during the audit are shown in the following table. Our audit report is responsible for the code in the initial version ([Version 1](#)), as well as new code (in the following versions) to fix issues in the audit report.

Project	Version	Commit Hash
ExoSwap Contracts	Version 1	22400563e5cb685ddd68fea3485673e38fcfd50d
	Version 2	f51806bc48a1da04de6d709196c6e2ef7d8f01d0

1.2 Disclaimer

This audit report does not constitute investment advice or a personal recommendation. It does not consider, and should not be interpreted as considering or having any bearing on, the potential economics of a token, token sale or any other product, service or other asset. Any entity should not rely on this report in any way, including for the purpose of making any decisions to buy or sell any token, product, service or other asset.

This audit report is not an endorsement of any particular project or team, and the report does not guarantee the security of any particular project. This audit does not give any warranties on discovering all security issues of the smart contracts, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit cannot be considered comprehensive, we always recommend proceeding with independent audits and a public bug bounty program to ensure the security of smart contracts.

The scope of this audit is limited to the code mentioned in Section 1.1. Unless explicitly specified, the security of the language itself (e.g., the solidity language), the underlying compiling toolchain and the computing infrastructure are out of the scope.

1.3 Procedure of Auditing

We perform the audit according to the following procedure.

¹<https://github.com/exoswapio/monorepo>

- **Vulnerability Detection** We first scan smart contracts with automatic code analyzers, and then manually verify (reject or confirm) the issues reported by them.
 - **Semantic Analysis** We study the business logic of smart contracts and conduct further investigation on the possible vulnerabilities using an automatic fuzzing tool (developed by our research team). We also manually analyze possible attack scenarios with independent auditors to cross-check the result.
 - **Recommendation** We provide some useful advice to developers from the perspective of good programming practice, including gas optimization, code style, and etc.
- We show the main concrete checkpoints in the following.

1.3.1 Software Security

- * Reentrancy
- * DoS
- * Access control
- * Data handling and data flow
- * Exception handling
- * Untrusted external call and control flow
- * Initialization consistency
- * Events operation
- * Error-prone randomness
- * Improper use of the proxy system

1.3.2 DeFi Security

- * Semantic consistency
- * Functionality consistency
- * Permission management
- * Business logic
- * Token operation
- * Emergency mechanism
- * Oracle security
- * Whitelist and blacklist
- * Economic impact
- * Batch transfer

1.3.3 NFT Security

- * Duplicated item
- * Verification of the token receiver
- * Off-chain metadata security

1.3.4 Additional Recommendation

- * Gas optimization

* Code quality and style



Note The previous checkpoints are the main ones. We may use more checkpoints during the auditing process according to the functionality of the project.

1.4 Security Model

To evaluate the risk, we follow the standards or suggestions that are widely adopted by both industry and academy, including OWASP Risk Rating Methodology ² and Common Weakness Enumeration ³. The overall *severity* of the risk is determined by *likelihood* and *impact*. Specifically, likelihood is used to estimate how likely a particular vulnerability can be uncovered and exploited by an attacker, while impact is used to measure the consequences of a successful exploit.

In this report, both likelihood and impact are categorized into two ratings, i.e., *high* and *low* respectively, and their combinations are shown in Table 1.1.

Table 1.1: Vulnerability Severity Classification

Impact	High	High	Medium
	Low	Medium	Low
		High	Low
		Likelihood	

Accordingly, the severity measured in this report are classified into three categories: **High**, **Medium**, **Low**. For the sake of completeness, **Undetermined** is also used to cover circumstances when the risk cannot be well determined.

Furthermore, the status of a discovered item will fall into one of the following four categories:

- **Undetermined** No response yet.
- **Acknowledged** The item has been received by the client, but not confirmed yet.
- **Confirmed** The item has been recognized by the client, but not fixed yet.
- **Fixed** The item has been confirmed and fixed by the client.

²https://owasp.org/www-community/OWASP_Risk_Rating_Methodology

³<https://cwe.mitre.org/>

Chapter 2 Findings

In total, we find **four** potential issues. We have **one** recommendation.

- High Risk: 3
- Low Risk: 1
- Recommendation: 1

ID	Severity	Description	Category	Status
1	High	Invalid multi-signature verification	DeFi Security	Fixed
2	Low	Potential access control takeover in FeeManager	DeFi Security	Fixed
3	High	Potential price manipulation	DeFi Security	Fixed
4	High	Centralization Risk	DeFi Security	Fixed
5	-	Require minimum number of cosigners	Recommendation	Acknowledged

The details are provided in the following sections.

2.1 DeFi Security

2.1.1 Invalid multi-signature verification

Severity High

Status Fixed in [Version 2](#)

Introduced by [Version 1](#)

Description In the “exit” process, multiple signatures of the exit data by two-thirds of all the cosigners are required. This logic is implemented in the `verify` function of the [BridgeCosignerManager](#) contract.

However, this logic is implemented incorrectly. As shown in the following code segment, there is a `cached` variable that stores the cumulative cosigners. Inside the `for` loop, the `_inCache` is invoked to check if the `signer` has appeared. However, the `_inCache` is implemented incorrectly. For example, if two signatures are required (i.e. the `_required` variable equals to 2), and there are three signatures in the `signatures` parameter with the first signature signed by a non-signer (but this signature is still valid). Then the `cached[0] == address(0)` is always true, and the remaining two signatures can duplicate, so only one valid signature is needed.

```
122     function verify(  
123         bytes32 commitment,  
124         uint256 chainId,  
125         bytes[] calldata signatures  
126     ) external view override returns (bool) {  
127         uint8 _required = getCosignCount(chainId);  
128         if (_required > signatures.length) {  
129             return false;  
130         }  
131  
132         address[] memory cached = new address[](signatures.length);  
133         uint8 signersMatch;
```

```
134
135     for (uint8 i = 0; i < signatures.length; i++) {
136         address signer = recover(commitment, signatures[i]);
137         Cosigner memory cosigner = _cosigners[signer];
138
139         if (
140             cosigner.active &&
141             cosigner.chainId == chainId &&
142             !_inCache(cached, signer)
143         ) {
144             signersMatch++;
145             cached[i] = signer;
146             if (signersMatch == _required) return true;
147         }
148     }
149
150     return false;
151 }
152
153 function _inCache(
154     address[] memory cached,
155     address signer
156 ) internal pure returns (bool hasCache) {
157     for (uint8 j = 0; j < cached.length; j++) {
158         if (cached[j] == signer) {
159             hasCache = true;
160             break;
161         }
162         // prevent iteration if cache not updated in slot
163         if (cached[j] == address(0)) {
164             break;
165         }
166     }
167 }
```

Listing 2.1: BridgeCosignerManager.sol

Impact The multi-signature verification logic is broken. Only one valid signature is used to pass the verification.

Suggestion Fix the multi-signature verification logic.

2.1.2 Potential access control takeover in `FeeManager`

Severity Low

Status Fixed in [Version 2](#)

Introduced by [Version 1](#)

Description The access control in the `FeeManager` contract is in a “first come first served” manner. As shown in the following code segment, the `reserveFee` function assign the first caller of a specific `AppId` to the owner of the `AppId`. Therefore, the registration of the owner of a specific `AppId` can be front-run and the fee configuration for this `AppId` can be modified.


```
59 function reserveFee(  
60     bytes32 _appId,  
61     uint256 _chainId,  
62     uint256 _baseFee,  
63     uint256 _feePerByte  
64 ) public override {  
65     require(_appId != DEFAULT_APP_ID, Errors.R_RESERVED_ENTITY);  
66     bytes32 key = getAppOwnerKey(_appId);  
67     require(!_appOwners[key], Errors.R_RESERVED_OWNER);  
68     _appOwners[key] = true;  
69     _reserveFee(_appId, _chainId, _baseFee, _feePerByte);  
70 }
```

Listing 2.2: FeeManager.sol

Because the `AppId` for the `RelayerManager` is set in the constructor, the registration of the `RelayerManager` to the `FeeManager` can be front-run easily.

Impact The `AppId` registration in the `FeeManager` contract can be front-run.

Suggestion Revise the `AppId` registration logic.

2.1.3 Potential price manipulation

Severity High

Status Fixed in [Version 2](#)

Introduced by [Version 1](#)

Description In the `BridgeRouter` contract, when the relayer is used, some of the user deposits are charged to pay the fees. If the bridged token is not native token, then an oracle is used to calculate the amount of tokens to be charged. As shown on Line 420 in the following code segment, the `getAmountOut` function of the `priceOracle` is used to calculate the amount of user deposits to be charged.

This function is subject to price manipulation attacks.

```
413 function relayCharge(  
414     RToken.Token memory token,  
415     bytes32 commitment,  
416     uint256 dataLength  
417 ) private returns (uint256 fee, address relayer) {  
418     fee = feeManager.getFees(relayerManager.appId(), _chainId, dataLength);  
419     require(fee > 0, Errors.B_ZERO_AMOUNT);  
420     if (token.addr != address(0) && token.addr != address(weth9)) {  
421         fee = priceOracle.getAmountOut(address(weth9), token.addr, fee);  
422     }  
423     relayer = relayerManager.pickRelayer(commitment);  
424     require(relayer == _msgSender(), Errors.BR_WRONG_EXECUTOR);  
425     token.exit(address(this), relayer, fee);  
426     if (relayer.code.length > 0) {  
427         bool success = IRelayerProcessor(relayer).postAcquire(  
428             token.addr,  
429             fee  
430         );
```

```
431         require(success, Errors.R_ACQUIRE_FAILED);
432     }
433 }
```

Listing 2.3: BridgeRouter.sol

Specifically, in the current implementation of the oracle, the Uniswap V2 spot price calculated by the reserve ratio is used. This ratio is well-known to be subject to price manipulation attacks. The potential loss is up to the `amount - amountMin`. If the `amountMin` is not properly set, users can lose all their deposits.

```
53 function getAmountOut(
54     address tokenIn,
55     address tokenOut,
56     uint256 amountIn
57 ) external view override returns (uint256) {
58     IUniswapV2Factory factory = IUniswapV2Factory(uniRouter.factory());
59     address input = (tokenIn == address(0)) ? uniRouter.WETH() : tokenIn;
60     address output = (tokenOut == address(0)) ? uniRouter.WETH() : tokenOut;
61     (address token0, ) = sortTokens(input, output);
62     IUniswapV2Pair pair = IUniswapV2Pair(
63         token0 == input
64         ? factory.getPair(input, output)
65         : factory.getPair(output, input)
66     );
67     (uint256 reserve0, uint256 reserve1, ) = pair.getReserves();
68     (uint256 reserveInput, uint256 reserveOutput) = token0 == input
69         ? (reserve0, reserve1)
70         : (reserve1, reserve0);
71     return quote(amountIn, reserveInput, reserveOutput);
72 }
```

Listing 2.4: UniswapV2DynamicPriceOracle.sol

Impact The gas fee calculation is subject to price manipulation when the relay is used and fees are swapped from user deposit tokens.

Suggestion Fix the gas fee calculation logic.

2.1.4 Centralization Risk

Severity High

Status Fixed in [Version 2](#)

Introduced by [Version 1](#)

Description In the [BridgeRouter](#) contract, there is an `emergencyWithdraw` function which can withdraw all the balance of any token from the router to the owner. In the ExoSwap project, all users transfer their tokens to the router contract. So this function can result in severe centralization risk.

```
177 function emergencyWithdraw(address token) external onlyOwner {
178     if (token == address(0)) {
179         (bool success, ) = payable(_msgSender()).call{
180             value: address(this).balance
181         }("");
182     }
```

```
182         require(success, Errors.B_SEND_REVERT);
183     } else {
184         IERC20(token).transfer(
185             _msgSender(),
186             IERC20(token).balanceOf(address(this))
187         );
188     }
189 }
```

Listing 2.5: BridgeRouter.sol

Impact The `emergencyWithdraw` function can withdraw all the balance of any token from the router, which stores all user deposits.

Suggestion Remove this function.

2.2 Additional Recommendation

2.2.1 Require minimum number of cosigners

Status Acknowledged

Introduced by Version 1

Description In the `BridgeCosignerManager` contract, there is a minimum number (`MIN_COSIGNER_REQUIRED = 2`) of cosigners required in the `getCosignCount`.

```
105 function getCosignCount(
106     uint256 chainId
107 ) public view override returns (uint8) {
108     uint8 voteCount = (uint8(_cosaddrs[chainId].length) * 2) / 3; // 67%
109     return
110         MIN_COSIGNER_REQUIRED >= voteCount
111         ? MIN_COSIGNER_REQUIRED
112         : voteCount;
113 }
```

Listing 2.6: BridgeCosignerManager.sol

However, in the `removeCosigner` function, it is not checked if the number of cosigners is larger than the minimum required number of cosigners.

```
62 function removeCosigner(address cosaddr) public override onlyOwner {
63     require(cosaddr != address(0), Errors.B_ZERO_ADDRESS);
64     Cosigner memory cosigner = _cosigners[cosaddr];
65     require(cosigner.active, Errors.B_ENTITY_NOT_EXIST);
66
67     address[] storage addrs = _cosaddrs[cosigner.chainId];
68
69     if (addrs.length > 1) {
70         // move last to rm slot
71         addrs[cosigner.index] = _cosaddrs[cosigner.chainId][
72             addrs.length - 1
73         ];
74     }
```

```
74     addr.pop();
75
76     // change indexing
77     address cosaddrLast = addr[cosigner.index];
78     _cosigners[cosaddrLast].index = cosigner.index;
79 } else {
80     // just remove it as 1 left
81     addr.pop();
82 }
83
84 delete _cosigners[cosaddr];
85
86 emit CosignerRemoved(cosigner.addr, cosigner.chainId);
87 }
```

Listing 2.7: BridgeCosignerManager.sol

Impact The number of cosigners can be less than `MIN_COSIGNER_REQUIRED` (2) after cosigner removal, which can lead to failure of cosigner verification.

Suggestion Implement checks for the minimum number of cosigners.