

Props involving quantifiers (Come to me idea: ①)

$\forall n A(n)$

"For all natural numbers  $n$  some property  $A(n)$  holds"

- where the quantification occurs over all natural numbers  
↳ often proved by method called induction.

(Number Theory) ↳ studies the properties of nat. num: 1, 2, 3

what are methods? [methods of induction]

Let's see how to prove an existence statement ↳

To prove  $\exists x A(x)$

"There exists an  $x$  of  $A(x)$ "  $\neg (\forall x \neg A(x))$  (A $\exists x$ )

→ the obvious way is to find an object  $a$  of  $A(a)$

e.g. To show there is an irrational number,  
prove that  $\sqrt{2}$  is irrational [which we did]

But this does not always work.

Sometimes we use indirect proofs. ②

E.g.

Theorem: There are irrationals  $r, s$  such that  $r^2$  is rational

ASIDE:

Cannot exist! rational number can be written as a fraction (or ratio) using integers

$$\therefore \frac{3}{2} \text{ (two integer)} = \text{quotient}(1.5)$$

$$so \text{ too } 7, \quad 7_1 = 7$$

$$so \text{ too } 0.317 = \frac{317}{1000} = 0.317.$$

But some numbers CANNOT be written as ratio or fraction (using integers)

Proof: involves considering two cases:

Case 1: If  $\sqrt{2}^{\sqrt{2}}$  is rational, we take  $r=s=\sqrt{2}$

Cse 2: If  $\sqrt{2}^{\sqrt{2}}$  is rational, take

$$r = \sqrt{2}^{\sqrt{2}} \text{ and } s = \sqrt{2}$$

$$\text{then } r^2 = (\sqrt{2}^{\sqrt{2}})^{\sqrt{2}} = (\sqrt{2})^{2 \cdot \sqrt{2}} = (\sqrt{2})^2 = 2.$$

and 2 is rational ( $\frac{2}{1} = 2$ )

from the above we looked at Cse 1 and 2

- we don't need to know if they rational.

→ methods: Proof by Cases

(fixed) Now let's look at method involving universal quantifier

prove:  $\forall x A(x)$

One way is to take an arbitrary  $x$  and show  
that it satisfies  $A(x)$

Eg. To prove  $\forall n \exists m [m > n^2]$   
for all  $n$ ,  $\exists m$ ,  $m > n^2$

$(m, n \in \mathbb{N})$

→ how do we now handle the two  
quantifiers in proof.

→ where the first quantifier  $\forall$  is a universal quantifier.

Let  $n$  be an arbitrary natural number.

Set  $m = n^2 + 1$  [our trial argument]

$\therefore$  then  $m > n^2$

This proves the statement:

$$\exists m (m > n^2)$$

"There is an  $m$ ,  $m > n^2$ ?

And it follows that  $\forall n \exists m (m > n^2)$  (is true)

So how did we do it: (logical reasoning): we want to prove statement involving 2 quantifiers in  $\forall n \exists m$

- we eliminate one quantifier ( $\forall n$ )

by replacing it by arbitrary natural number

- why arbitrary?  $\Rightarrow$

- then we repeat same reasoning

or provide an argument ( $m = n^2 + 1$ )

-  $n^2 + 1$ , and let ~~call~~  $m$  be this number.

- so we explicitly find  $m$  that satisfies  $m > n^2$

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In practice, there may be pages of arguments

- Since the argument satisfies an arbitrary <sup>(works)</sup> natural number  
it follows it proves  $\forall n$  [ <sup>natural</sup> estimated quantific]

$\forall$  For all has been handled by arbitrary (argument)  
natural number.

Note: This works because the "n" we picked  
is arbitrary

Another method to use is method of Contradiction

To prove  $\forall x A(x)$ , assume  $\neg \forall x A(x)$

This is equivalent to:

$\exists x \neg A(x)$ .

$\therefore$  Let c be object such that  $\neg A(c)$   
 $\hookrightarrow$  not arbitrary object

Now we reason with c (and the fact that  $\neg A(c)$ )  
to derive a Contradiction.

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Qn3 Is this a valid proof?

To prove:  $(\forall x > 0)(\exists y > 0)[y < x]$ , where

variable ranges over natural numbers

(This says given any positive number, you can always find a smaller one)

To prove it: pick a positive ~~number~~ rational number  $p$  arbitrarily, say  $= p = 0.001$

take  $q = 0.0001$ .

$$\therefore 0 < q < p$$

" $q$  is less than  $p$  and greater than 0"

Since our choice of  $p$  was arbitrary,  
this proves the result!

But is this right?  
But picking a positive <sup>arbitrarily</sup> ~~arbitrary~~ number for  $p$

is not the same as letting  $p$  be arbitrary

$\Rightarrow$  the choice of  $p$  may be arbitrary, but once you made it 0.001, you have a specific  $p$ !

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$\Rightarrow$  we even said what it is!  $p=0.001 \leftarrow$  (specific)

$\Rightarrow$  that is different from saying let  $p$  be arbitrary.

$\Rightarrow$  subtle, but important point

Proof is not valid

- in category arguing with an arbitrary  $p$ ,  
we don't know what it is, and we not specific

Let's look at Proof by Induction

To prove statement of the form  $\forall n A(n)$ ,  
where  $n$  (quantifiers) ranges over  
the natural numbers

E.g. Prove that  $1+2+\dots+n = \frac{1}{2}n(n+1)$

First Step: check the first few cases  
(get some of that is going on)

$$n=1 \quad 1 = \frac{1}{2}(1)(1+1) = \frac{1}{2}(2) = 1 \quad \checkmark \text{ true for } n=1$$

$$\begin{aligned} n=2 & \quad 1+2 = \frac{1}{2}(2)(2+1) = 3 \\ & \quad 3 = 1(3) = 3 \quad (\text{true}) \end{aligned}$$

$$\begin{aligned} 2 \times 2 & \\ 2 \times 1 & \\ - \frac{2}{2} & = 1 \end{aligned}$$

$n=3$

$$1+2+3 = 1(3(3+1))$$

$$6 = \frac{3}{2}(4) = \frac{12}{2} = 6 \text{ (True)}$$

⑧

So far this is not a proof, it's just  
checking the first two cases. [Beweis Øf this]

Consider formula:

$$P(n) = n^2 - nt + 41, \text{ all values of } P(n)$$

for  $n=1, 2, \dots$  etc are prime numbers

⇒ Not quite, until you reach  $n=41$

$$P(41) = 1681 = 41^2$$

↳ Leonard Euler, 1772 ⇒

→ Method of induction depends principle of, known as

principle of mathematical induction

Here is what it says

To prove  $\forall n (A_n)$ , establish the following

2 statements:

①  $A(1)$  initial step / base

② Proof following statement  $(\forall n)[A(n) \Rightarrow A(n+1)]$

② Called induction step

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Intuitively, this gives  $V_n A(n)$  as follows:

By initial step ;  $A(1)$ .

If we now apply the induction step

$$A(n) \Rightarrow A(n+1)$$

$$A(1) \Rightarrow A(2)$$

So from  $A(1)$  we can conclude  $A(2)$

Next By  $A(2)$  (first step) and then induction step

$$A(2) \Rightarrow A(3), \text{ we can}$$

Conclude  $A(3)$  --- etc.

You need axiom (or principle) to make this to work, called the principle of mathematical induction.

The PMI is what tells you that Step 1 and 2 above

yield  $V_n A(n)$

Now lets apply the method.

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Let prove theorem:

for any natural  $n$ ,  $1+2+3+\dots+n = \frac{1}{2}n(n+1)$

Calling it theorem, I am going to prove by  
mathematical induction.

Proof: By mathematical induction [tell reader how you  
try to prove & method used]

\* So I have to prove show that this is true for  $n=1$ , then  
I have to show that it is true for  $n$ , it follows for  $n+1$  [step ①]

→ notion equation, [identity;  
two sides are always equal]

\* Step ① for  $n=1$ , the identity reduces to:

$$1 = \frac{1}{2}(1)(1+1) = 1, (\text{true})$$

(since both sides equal 1)

\* Step ② [new induction step]: Assume the  
identity holds for  $n$

i.e.  $1+2+\dots+n = \frac{1}{2}n(n+1)$  [original] ⑩

[induction step] → [want to deduce:  $1+2+\dots+n+1 = \frac{1}{2}(n+1)(n+1)$ ]

lets add  $(n+1)$  to both sides ⑩:

$$1+2+\dots+n+(n+1) = \frac{1}{2}n(n+1) + (n+1)$$

remove  $\frac{1}{2}$  as common factor:

⑦

$$\begin{aligned}
 H_{2t} + n + (n+1) &= \frac{1}{2}n(n+1) + (n+1) = \frac{1}{2}[n(n+1) + 2(n+1)] \\
 &= \frac{1}{2}[n^2 + n + 2n + 2] = \frac{1}{2}[n^2 + 3n + 2] \\
 &= \frac{1}{2}[(n+1)(n+2)] = \text{Same as above} \\
 \therefore \text{Same as} &= \frac{1}{2}(n+1)(n+1+1)
 \end{aligned}$$

which is the identity with  $n+1$  in place of  $n$ .  
Hence, by principle of mathematical induction,  
the identity holds for all  $n$



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Another example of induction proof:

Theorem: If  $x > 0$  (a positive real number), then

for any  $n$  (natural number),  $(1+x)^{n+1} > 1+(n+1)x$

Proof: By mathematical induction

Let  $A(n)$  be the statement  $(1+x)^{n+1} > 1+(n+1)x$ .

To prove:  $\forall n A(n)$ .

Step 1: Prove  $A(1)$

$A(1)$  is the statement  $(1+x)^2 > 1+2x$  True!

True by binomial theorem,

$$(1+x)^2 = 1+2x+x^2 > 1+2x \text{ (since } x \geq 0\text{)}$$

Step 2 (Induction Step):

Proof:  $\forall n [A(n) \Rightarrow A(n+1)]$

(we pick an arbitrary  $n$  and prove

$$A(n) \Rightarrow A(n+1)$$

$\therefore$  we assume antecedent  
and deduce a consequence.

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∴ We assume  $A(n)$  and deduce  $A(n+1)$

$$(1+x)^{n+1} > 1 + (A_{n+1})x$$

"Assuming this"

$$(1+x)^{n+2} > 1 + (A_{n+2})x$$

"and deducing this"

(Concentrate on this now)

$$\therefore (1+x)^{n+2} = (1+x)(1+x)^{n+1}$$

↓ pull out (refactor)

$$\Rightarrow (1+x)[1 + (n+1)x]$$

$$= 1 + (n+1)x + x + (n+1)x^2$$

$$= 1 + (n+2)x + (n+1)x^2$$

$$1 + (n+2)x$$

This proves  $A(n+1)$

## Induction Summary:

- ① You want to prove that some statement  $A(n)$  is valid for all natural numbers  $n$ .
  - ② First prove  $A(1)$ . Usually matter of simple observation.
  - ③ Give an algebraic argument to establish the induction  $A(n) \Rightarrow A(n+1)$ 
    - reduce  $A(n+1)$  to a form where you can use  $A(n)$
  - ④ Conclusion: By principle of Mathematical Induction  $\boxed{A(n)}$
- another variant [Common] of Induction.
- We sometimes need to prove a statement of form  $(\forall n \geq n_0) A(n)$
- where  $n_0$  is some fixed number e.g.  $J$ . In this case, the first step is to verify  $A(n_0)$
  - $A(1)$  may not be true

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(whereas for  $n_0$ , the first occurrence  $A(i)$  is normally not true, we have to find some other natural number as starting point)

$\Rightarrow$  general induction starts at 1, But.  
 This version starts at some point beyond 1.  
 (other than that, the argument is the same)

Now

The induction step is to prove:

$$(B_n \rightarrow n_0) [A_n \Rightarrow A(n+1)]$$

This induction is part of Fermat's Induction technique  
 Called "The Fundamental theorem of Arithmetic"

But let's first look at Ques.

"Mod"

Is 1 a prime number?

$\Rightarrow$  it is not a prime number, b/c of definition:

i) must be positive

2) greater than 1

3) and ~~exactly~~ one 1 and n.

Another definition:

0 is not a natural number,  
it's an integer, but ↑

⇒ natural numbers are Counting numbers

Fundamental Theorem of arithmetic to be proved:

Theorem. Every natural number greater than 1  
is either prime or a product of primes.  
[use variant of induction [n ≥ 2] or (n ≥ 1)]

Proof: By induction

The induction statement, A(n) is:

$\forall m [2 \leq m \leq n \Rightarrow m \text{ is either a prime or a product of primes}]$

For  $n=2$ , A(2) says:

"2 is either prime or a product of primes" [True]

⑦

Assume:  $A(n)$ ; & deduce  $A(n+1)$

Let  $m$  be natural number,  $2 \leq m \leq n+1$

If  $m \leq n$ , then by  $A(n)$ ,  $m$  is either power product of primes

If  $m = n+1$ , and if  $n+1$  is prime, then  $m$  is prime

If  $m = n+1$  and  $n+1$  is not prime, then  
 there are natural numbers  $p, q$  such that we can find two smaller numbers whose product is prime  
 that  $1 \leq p, q \leq n+1$  and  $n+1 = pq$ .

Since  $2 \leq p, q \leq n$  (by  $A(n)$ ),  $p, q$  are either power product of primes.

Hence  $n+1$  is a product of primes.

The theorem follows by induction.