

4. 기호 숙식

문제 2, 4, 6

\* 문제 2  $T(n) = T(n-1) + n$

$$T(n) = T(n-1) + n \quad \dots \quad (7)$$

$$T(n-1) = T(n-2) + (n-1) \quad \dots \quad (L)$$

$$T(n-2) = T(n-3) + (n-2) \quad \dots \quad (E)$$

$\vdots$

$$T(2) = T(1) + 2$$

(7) 에 (L) 대입,

$$T(n) = (T(n-2) + (n-1)) + n$$

$$= T(n-2) + n + (n-1)$$

(E) 대입,

$$T(n) = \{T(n-3) + (n-2)\} + n + (n-1)$$

$$= T(n-3) + n + (n-1) + (n-2)$$

$\vdots$

$$T(n) = T(1) + n + (n-1) + (n-2) + \dots + 2$$

$$= (n+1) \times \frac{n}{2}$$

$$= \frac{n^2 + n}{2}$$

$\therefore \underline{T(n) = O(n^2)}$  ( $\because n^2$ 보다 작은 차수와 계수들 무시가능)

문제 4  $T(n) = T(\frac{n}{2}) + 1$

$$\begin{aligned} T(n) &= T(\frac{n}{2}) + 1 \\ T(\frac{n}{2}) &= T(\frac{n}{4}) + 1 \\ T(\frac{n}{4}) &= T(\frac{n}{8}) + 1 \quad \dots (L) \end{aligned} \Rightarrow T(n) = (T(\frac{n}{4}) + 1) + 1 = T(\frac{n}{4}) + 2 \quad \dots (7)$$

(7)에 (L) 대입

$$\begin{aligned} T(n) &= (T(\frac{n}{8}) + 1) + 2 \\ &= T(\frac{n}{8}) + 3 \end{aligned}$$

$$T(n) = T(\frac{n}{2^k}) + k \quad \dots (k \text{ 번째})$$

$$\boxed{\frac{n}{2^k} = 1 \Rightarrow n = 2^k \Rightarrow k = \log n}$$

$$T(n) = T(1) + \log n \quad \dots (\log n \text{ 번째})$$

$\therefore T(n) = O(\log n)$   $(\because T \text{ 속의 상수가 들어가는 상수 기호})$

문제 6  $T(n) = 2T\left(\frac{n}{2}\right) + n$

$$T(n) = 2T\left(\frac{n}{2}\right) + n \quad \dots (1)$$

$$T\left(\frac{n}{2}\right) = 2T\left(\frac{n}{4}\right) + \frac{n}{2} \quad \dots (2)$$

$$T\left(\frac{n}{4}\right) = 2T\left(\frac{n}{8}\right) + \frac{n}{4} \quad \dots (3)$$

(1) or (2) 대입,

$$\begin{aligned} T(n) &= 2\left(2T\left(\frac{n}{4}\right) + \frac{n}{2}\right) + n \\ &= 4T\left(\frac{n}{4}\right) + 2n \end{aligned}$$

(3) 대입,

$$\begin{aligned} T(n) &= 4\left(2T\left(\frac{n}{8}\right) + \frac{n}{4}\right) + 2n \\ &= 8T\left(\frac{n}{8}\right) + 3n \end{aligned}$$

$\vdots$

$$T(n) = 2^k T\left(\frac{n}{2^k}\right) + k \cdot n \quad \dots (k \text{ 번째})$$

$$\boxed{\frac{n}{2^k} = 1 \rightarrow n = 2^k \rightarrow k = \log n}$$

$$T(n) = n \cdot T(1) + (\log n) \times n \quad \dots (\log n \text{ 번째})$$

$$= n(\cancel{T(1)} + \log n) \quad (\because T(1) \Rightarrow \log n \text{ 보다 작기 때문})$$

$$\therefore T(n) = O(n \cdot \log n)$$