

Chapter 4 Convex optimization problems

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Convex optimization problems

- ▶ relevant concepts (for general optimization problems & for convex problems)
- ▶ properties of convex problems (local implies global & optimality condition)
- ▶ operations preserving convexity (construct new from old)
- ▶ many examples of convex problems (LP, QP, QCQP, SOCP, etc.)
- ▶ extensions (quasiconvex optimization & geometric programming)
- ▶ combination with generalized inequalities (in constraints & in objective functions)

Optimization problems

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Optimization problem in standard form

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p\end{array}$$

$x \in \mathbb{R}^n$	optimization variable
$f_0: \mathbb{R}^n \rightarrow \mathbb{R}$	objective function (cost function)
$f_i: \mathbb{R}^n \rightarrow \mathbb{R}$	inequality constraint functions
$h_i: \mathbb{R}^n \rightarrow \mathbb{R}$	equality constraint functions

- ▶ **implicit constraints**

$$x \in \mathcal{D} = \left(\bigcap_{i=0}^m \text{dom } f_i \right) \cap \left(\bigcap_{i=1}^p \text{dom } h_i \right)$$

- ▶ \mathcal{D} is called the **domain** of the problem
- ▶ **explicit constraints**

$$f_i(x) \leq 0 \text{ for } 1 \leq i \leq m \quad \text{and} \quad h_i(x) = 0 \text{ for } 1 \leq i \leq p$$

- ▶ problem is **unconstrained** if it has no explicit constraints ($m = p = 0$)

Example

$$\text{minimize} \quad f_0(x) = - \sum_{i=1}^k \log(b_i - a_i^T x)$$

is an unconstrained problem with implicit constraints

$$a_i^T x < b_i$$

for each $1 \leq i \leq k$.

Feasibility

- ▶ x is **feasible** if $x \in \mathcal{D}$ and x satisfies all constraints
- ▶ the set of all feasible points is called the **feasible set** of the problem
- ▶ the problem is **infeasible** if the feasible set is empty
- ▶ the **feasibility problem** is to determine whether the feasible set is nonempty

$$\begin{array}{ll}\text{find} & x \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p\end{array}$$

- ▶ it can be rephrased as an optimization problem

$$\begin{array}{ll}\text{minimize} & 0 \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p\end{array}$$

- ▶ The **optimal value** is

$$p^* = \inf \left\{ f_0(x) \mid \begin{array}{l} f_i(x) \leq 0 \text{ for } 1 \leq i \leq m \\ h_i(x) = 0 \text{ for } 1 \leq i \leq p \end{array} \right\} \in \mathbb{R} \cup \{\pm\infty\}$$

- ▶ Extreme situations

$$p^* = \infty$$

if problem is infeasible

$$p^* = -\infty$$

if problem is unbounded below

- ▶ Optimal value may not be achieved.

- ▶ x is **optimal** if it is feasible and $f_0(x) = p^*$
- ▶ x is **locally optimal** if there exists $R > 0$ such that x is optimal for

$$\begin{array}{ll}\text{minimize} & f_0(z) \\ \text{subject to} & f_i(z) \leq 0, \quad i = 1, \dots, m \\ & h_i(z) = 0, \quad i = 1, \dots, p \\ & \|z - x\|_2 \leq R\end{array}$$

Examples (when $n = 1, m = p = 0$)

$f_0(x) = x \log x$	$\mathbf{dom} f_0 = \mathbb{R}_{++}$	$p^* = -1/e$	$x = 1/e$ is optimal
$f_0(x) = -\log x$	$\mathbf{dom} f_0 = \mathbb{R}_{++}$	$p^* = -\infty$	no optimal point
$f_0(x) = 1/x$	$\mathbf{dom} f_0 = \mathbb{R}_{++}$	$p^* = 0$	no optimal point
$f_0(x) = x^3 - 3x$	$\mathbf{dom} f_0 = \mathbb{R}$	$p^* = -\infty$	$x = 1$ is locally optimal

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Convex optimization problem in standard form

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & a_i^T x = b_i, \quad i = 1, \dots, p\end{array}$$

- ▶ f_0, f_1, \dots, f_m are convex
- ▶ equality constraints are affine, often written as $Ax = b$
- ▶ important property: feasible set of a convex problem is convex
- ▶ problem is **quasiconvex** if f_0 is quasiconvex (and f_1, \dots, f_m convex)

Example

$$\begin{array}{ll}\text{minimize} & f_0(x) = x_1^2 + x_2^2 \\ \text{subject to} & f_1(x) = x_1/(1 + x_2^2) \leq 0 \\ & h_1(x) = (x_1 + x_2)^2 = 0\end{array}$$

- ▶ f_0 is convex
- ▶ not a convex problem: f_1 is not convex, h_1 is not affine
- ▶ equivalent (but not identical) to the convex problem

$$\begin{array}{ll}\text{minimize} & x_1^2 + x_2^2 \\ \text{subject to} & x_1 \leq 0 \\ & x_1 + x_2 = 0\end{array}$$

Proposition

Any locally optimal point of a convex optimization problem is globally optimal.

Proof

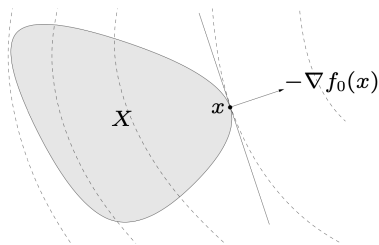
- ▶ suppose x is locally optimal, but there exists feasible y with $f_0(y) < f_0(x)$
- ▶ there exists $R > 0$ such that $f_0(z) \geq f_0(x)$ for all feasible z with $\|z - x\|_2 < R$
- ▶ consider $z = \theta y + (1 - \theta)x$ with $\theta = R/(2\|y - x\|_2)$, then $\|z - x\|_2 = R/2$
- ▶ $\|y - x\|_2 > R$ implies $0 < \theta < 1/2$, hence z is feasible by convexity of domain
- ▶ by convexity of objective $f_0(z) \leq \theta f_0(y) + (1 - \theta)f_0(x) < f_0(x)$, contradiction

Optimality criterion for differentiable objective

Optimality criterion

Suppose the problem is convex and f_0 is differentiable, then

$$x \text{ is optimal} \iff \begin{array}{l} x \text{ is feasible and} \\ \nabla f_0(x)^T(y - x) \geq 0 \text{ for all feasible } y \end{array}$$



Geometric interpretation

Either $\nabla f_0(x) = 0$ or $\nabla f_0(x)$ defines a supporting hyperplane to the feasible set at x .

Proof

(\Leftarrow) For any feasible y , since $y \in \mathbf{dom} f_0$, by the convexity of f_0

$$f_0(y) \geq f_0(x) + \nabla f_0(x)^T(y - x).$$

The assumption $\nabla f_0(x)^T(y - x) \geq 0$ implies $f_0(y) \geq f_0(x)$. Hence x is optimal.

(\Rightarrow) Assume on the contrary that $\nabla f_0(x)^T(y - x) < 0$ for some feasible y , then $z(t) = ty + (1 - t)x$ is feasible for $t \in [0, 1]$ since the feasible set is convex. Then

$$\left. \frac{d}{dt} f_0(z(t)) \right|_{t=0} = \nabla f_0(x)^T(y - x) < 0,$$

hence $f_0(z(t)) < f_0(x)$ for $0 < t \ll 1$, which contradicts the optimality of x . □

unconstrained problem

minimize $f_0(x)$

x is optimal $\iff x \in \mathbf{dom} f_0, \nabla f_0(x) = 0$

equality constrained problem

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & Ax = b\end{array}$$

$$x \text{ is optimal} \quad \Longleftrightarrow \quad \begin{array}{l} x \in \mathbf{dom} f_0, \quad Ax = b, \\ \nabla f_0(x) + A^T \nu = 0 \text{ for some vector } \nu \end{array}$$

minimization over nonnegative orthant

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & x \succeq 0\end{array}$$

$$x \text{ is optimal} \iff x \in \mathbf{dom} f_0, \quad x \succeq 0, \quad \begin{cases} \nabla f_0(x)_i \geq 0, & \text{if } x_i = 0 \\ \nabla f_0(x)_i = 0, & \text{if } x_i > 0 \end{cases}$$

Sample proof (for unconstrained problems)

- By optimality condition

$$x \text{ is optimal} \iff x \in \mathbf{dom} f_0, \nabla f_0(x)^T(y - x) \geq 0 \text{ for each } y \in \mathbf{dom} f_0$$

- $\nabla f_0(x) = 0$ is clearly sufficient for the above statement.
- Since f_0 is differentiable, $\mathbf{dom} f_0$ is open, hence

$$y = x - \varepsilon \nabla f_0(x) \in \mathbf{dom} f_0$$

for $0 < \varepsilon \ll 1$. For such y we have

$$\nabla f_0(x)^T(y - x) = -\varepsilon \|\nabla f_0(x)\|_2^2 \leq 0.$$

- Combining above gives

$$\nabla f_0(x) = 0$$

which proves necessity.

Equivalent convex problems

Two problems are (informally) **equivalent** if the solution of one is readily obtained from the solution of the other, and vice-versa.

Some common transformations that preserve convexity

- ▶ eliminating equality constraints
- ▶ introducing equality constraints
- ▶ introducing slack variables for linear inequalities
- ▶ epigraph form
- ▶ minimizing over some variables

eliminating equality constraints

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

is equivalent to

$$\begin{array}{ll}\text{minimize} & f_0(Fz + x_0) \quad (\text{over } z) \\ \text{subject to} & f_i(Fz + x_0) \leq 0, \quad i = 1, \dots, m\end{array}$$

where F and x_0 are such that

$$Ax = b \quad \Longleftrightarrow \quad x = Fz + x_0 \text{ for some } z$$

introducing equality constraints

$$\begin{array}{ll}\text{minimize} & f_0(A_0x + b_0) \\ \text{subject to} & f_i(A_ix + b_i) \leq 0, \quad i = 1, \dots, m\end{array}$$

is equivalent to

$$\begin{array}{lll}\text{minimize} & f_0(y_0) & (\text{over } x, y_i) \\ \text{subject to} & f_i(y_i) \leq 0, & i = 1, \dots, m \\ & y_i = A_ix + b_i, & i = 0, 1, \dots, m\end{array}$$

introducing slack variables for linear inequalities

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & a_i^T x \leq b_i, \quad i = 1, \dots, m\end{array}$$

is equivalent to

$$\begin{array}{lll}\text{minimize} & f_0(x) & (\text{over } x, s) \\ \text{subject to} & a_i^T x + s_i = b_i, & i = 1, \dots, m \\ & s_i \geq 0, & i = 1, \dots, m\end{array}$$

epigraph form

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

is equivalent to

$$\begin{array}{ll}\text{minimize} & t \quad (\text{over } x, t) \\ \text{subject to} & f_0(x) - t \leq 0 \\ & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

partial minimization

$$\begin{array}{ll} \text{minimize} & f_0(x_1, x_2) \\ \text{subject to} & f_i(x_1) \leq 0, \quad i = 1, \dots, m \end{array}$$

is equivalent to

$$\begin{array}{ll} \text{minimize} & \tilde{f}_0(x_1) \\ \text{subject to} & f_i(x_1) \leq 0, \quad i = 1, \dots, m \end{array}$$

where

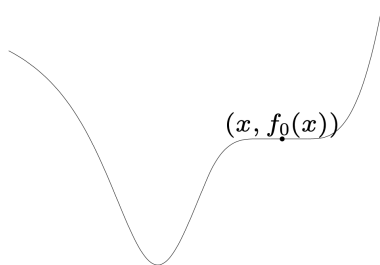
$$\tilde{f}_0(x_1) = \mathbf{\inf}_{x_2} f_0(x_1, x_2)$$

Quasiconvex optimization

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

with $f_0: \mathbb{R}^n \rightarrow \mathbb{R}$ quasiconvex, f_1, \dots, f_m convex.

Remark Locally optimal points may not be globally optimal



Convex representation of sublevel sets of f_0

For quasiconvex f_0 there exists a family of functions ϕ_t such that

- ▶ $\phi_t(x)$ is convex in x for each fixed t
- ▶ t -sublevel set of f_0 is 0-sublevel set of ϕ_t , i.e. $f_0(x) \leq t \iff \phi_t(x) \leq 0$
- ▶ $\phi_t(x)$ is nonincreasing in t for each fixed x , namely $\phi_s(x) \leq \phi_t(x)$ if $s \geq t$

In practice there are usually natural meaningful choices for ϕ_t .

Example

$$f_0(x) = \frac{p(x)}{q(x)}$$

with p convex, q concave, and $p(x) \geq 0$, $q(x) > 0$ on $\mathbf{dom} f_0$.

We can choose

$$\phi_t(x) = p(x) - tq(x)$$

- ▶ $\phi_t(x)$ convex in x for $t \geq 0$
- ▶ $f_0(x) \leq t \iff \phi_t(x) \leq 0$

Quasiconvex optimization via convex feasibility problems

$$\phi_t(x) \leq 0, \quad f_i(x) \leq 0, \quad i = 1, \dots, m, \quad Ax = b$$

- ▶ convex feasibility problem in x for each fixed t
- ▶ let p^* be the optimal value for the original quasiconvex problem, then

$$\begin{array}{ll} \text{above problem feasible} & \implies p^* \leq t \\ \text{above problem infeasible} & \implies p^* \geq t \end{array}$$

Bisection method

given $l \leq p^*, u \geq p^*, \text{ tolerance } \epsilon > 0$

repeat

1. $t := (l + u)/2$
2. solve the above convex feasibility problem
3. **if** feasible, $u := t$; **else** $l := t$

until $u - l \leq \epsilon$

requires exactly $\lceil \log_2((u - l)/\epsilon) \rceil$ iterations

Optimization problems

Convex optimization

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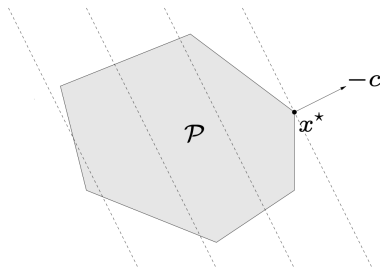
Generalized inequality constraints

Vector optimization

Linear program (LP)

$$\begin{array}{ll}\text{minimize} & c^T x + d \\ \text{subject to} & Gx \preceq h \\ & Ax = b\end{array}$$

- ▶ convex problem with affine objective and constraint functions
- ▶ feasible set is a polyhedron



Examples

Diet problem choose quantities x_1, \dots, x_n of n kinds of food

- ▶ one unit of food j costs c_j , contains amount a_{ij} of nutrient i
- ▶ healthy diet requires nutrient i in quantity at least b_i

to find cheapest healthy diet

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \succeq b \\ & x \succeq 0 \end{array}$$

Piecewise-linear minimization

$$\text{minimize} \quad \mathbf{max} \{a_i^T x + b_i \mid i = 1, \dots, m\}$$

equivalent to the LP

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & a_i^T x + b_i \leq t, \quad i = 1, \dots, m \end{array}$$

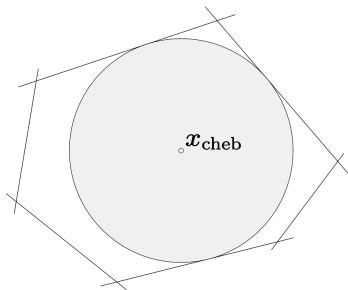
Chebyshev center of a polyhedron

Chebyshev center of

$$\mathcal{P} = \{x \mid a_i^T x \leq b_i, \ i = 1, \dots, m\}$$

is center of largest inscribed ball

$$\mathcal{B} = \{x_c + u \mid \|u\|_2 \leq r\}$$



$a_i^T x \leq b_i$ for all $x \in \mathcal{B}$ if and only if

$$\sup\{a_i^T(x_c + u) \mid \|u\|_2 \leq r\} = a_i^T x_c + r\|a_i\|_2 \leq b_i$$

hence x_c and r can be determined by solving the LP

$$\begin{array}{ll} \text{maximize} & r \\ \text{subject to} & a_i^T x_c + r\|a_i\|_2 \leq b_i, \quad i = 1, \dots, m \end{array}$$

Linear-fractional program

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & Gx \preceq h \\ & Ax = b\end{array}$$

where

$$f_0(x) = \frac{c^T x + d}{e^T x + f}, \quad \text{dom } f_0(x) = \{x \mid e^T x + f > 0\}$$

is a quasiconvex optimization problem; can be solved by bisection method.

If the feasible set is nonempty, then the linear-fractional problem is equivalent to the LP

$$\begin{array}{ll}\text{minimize} & c^T y + dz \\ \text{subject to} & Gy \preceq hz \\ & Ay = bz \\ & e^T y + fz = 1 \\ & z \geq 0\end{array}$$

Generalized linear-fractional program

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & Gx \preceq h \\ & Ax = b\end{array}$$

where

$$\begin{aligned} f_0(x) &= \mathbf{max} \left\{ \frac{c_i^T x + d_i}{e_i^T x + f_i} \mid i = 1, \dots, r \right\} \\ \mathbf{dom} \, f_0(x) &= \{x \mid e_i^T x + f_i > 0, \, i = 1, \dots, r\} \end{aligned}$$

is a quasiconvex optimization problem; can be solved by bisection method.

Example: von Neumann model of a growing economy

$$\begin{array}{ll} \text{maximize} & \min \{x_i^+/x_i \mid i = 1, \dots, n\} \quad (\text{over } x, x^+) \\ \text{subject to} & x^+ \succeq 0 \\ & Bx^+ \preceq Ax \end{array}$$

with domain $\{(x, x^+) \mid x \succ 0\}$

- ▶ $x, x^+ \in \mathbb{R}^n$: activity levels of n sectors, in current and next period
- ▶ $(Ax)_i, (Bx^+)_i$: produced resp. consumed amounts of good i
- ▶ x_i^+/x_i : growth rate of sector i

allocate activity to maximize growth rate of slowest growing sector

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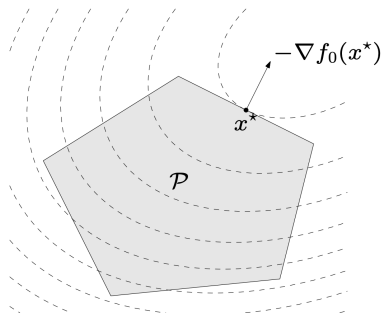
Generalized inequality constraints

Vector optimization

Quadratic program (QP)

$$\begin{array}{ll}\text{minimize} & (1/2)x^T Px + q^T x + r \\ \text{subject to} & Gx \preceq h \\ & Ax = b\end{array}$$

- ▶ $P \in \mathbb{S}_+^n$ thus objective is convex quadratic
- ▶ minimize a convex quadratic function over a polyhedron



Least-squares

$$\text{minimize} \quad \|Ax - b\|_2^2$$

- ▶ analytical solution $x^* = A^\dagger b$ (where A^\dagger is pseudo-inverse)
- ▶ can add linear constraints such as $l \preceq x \preceq u$

Linear program with random cost

Consider the linear program

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Gx \preceq h \\ & Ax = b\end{array}$$

- ▶ Assume c is random vector with mean \bar{c} and covariance Σ
- ▶ Then $c^T x$ is random variable with mean $\bar{c}^T x$ and variance $x^T \Sigma x$

$$\begin{aligned}\mathbf{E}(c^T x) &= \mathbf{E}(c)^T x = \bar{c}^T x \\ \mathbf{var}(c^T x) &= \mathbf{E}(c^T x - \bar{c}^T x)^2 = x^T \mathbf{E}((c - \bar{c})(c - \bar{c})^T) x = x^T \Sigma x\end{aligned}$$

We modify the above LP to the following QP

$$\begin{array}{ll}\text{minimize} & \bar{c}^T x + \gamma x^T \Sigma x \\ \text{subject to} & Gx \preceq h \\ & Ax = b\end{array}$$

- ▶ To keep both the expected cost and the cost variance (risk) under control, choose a linear combination of both as the new objective, called **risk-sensitive cost**.
- ▶ $\gamma > 0$ is the **risk-aversion parameter**, which controls the trade-off between expected cost and variance.
- ▶ Coefficient vector $(1, \gamma)$ lies in the interior of the dual cone of the nonnegative quadrant.

Quadratically constrained quadratic program (QCQP)

$$\begin{array}{ll}\text{minimize} & (1/2)x^T P_0 x + q_0^T x + r_0 \\ \text{subject to} & (1/2)x^T P_i x + q_i^T x + r_i \leq 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

- ▶ $P_i \in \mathbb{S}_+^n$ thus objective and constraints are convex quadratic
- ▶ feasible region is intersection of m ellipsoids and an affine set if $P_1, \dots, P_m \in \mathbb{S}_{++}^n$

Second-order cone program (SOCP)

$$\begin{array}{ll}\text{minimize} & f^T x \\ \text{subject to} & \|A_i x + b_i\|_2 \leq c_i^T x + d_i, \quad i = 1, \dots, m \\ & Fx = G\end{array}$$

with $A_i \in \mathbb{R}^{n_i \times n}$ and $F \in \mathbb{R}^{p \times n}$

- ▶ inequalities are called second-order cone constraints since

$$(A_i x + b_i, c_i^T x + d_i) \in \text{second-order cone in } \mathbb{R}^{n_i+1}$$

- ▶ if $n_i = 0$, reduces to LP
- ▶ if $c_i = 0$, reduces to QCQP (with linear objective)

Robust linear program

Parameters in optimization problems are often uncertain. Consider the LP

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & a_i^T x \leq b_i, \quad i = 1, \dots, m\end{array}$$

- ▶ There can be uncertainty in c, a_i, b_i (in a_i for example)
- ▶ There are two common approaches to handle uncertainty
 - ▶ deterministic model
 - ▶ stochastic model

- deterministic model: constraints must hold for all $a_i \in \mathcal{E}_i$

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & a_i^T x \leq b_i \text{ for all } a_i \in \mathcal{E}_i, \quad i = 1, \dots, m\end{array}$$

- stochastic model: a_i is random variable; constraints must hold with probability η

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & \mathbf{prob}(a_i^T x \leq b_i) \geq \eta, \quad i = 1, \dots, m\end{array}$$

deterministic approach via SOCP

- ▶ choose ellipsoid as \mathcal{E}_i with $\bar{a}_i \in \mathbb{R}^n$ and $P_i \in \mathbb{R}^{n \times n}$

$$\mathcal{E}_i = \{\bar{a}_i + P_i u \mid \|u\|_2 \leq 1\}$$

- ▶ robust LP

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & a_i^T x \leq b_i \text{ for all } a_i \in \mathcal{E}_i, \quad i = 1, \dots, m \end{array}$$

- ▶ equivalent SOCP

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & \bar{a}_i^T x + \|P_i^T x\|_2 \leq b_i, \quad i = 1, \dots, m \end{array}$$

which follows from

$$\sup_{\|u\|_2 \leq 1} (\bar{a}_i + P_i u)^T x = \bar{a}_i^T x + \|P_i^T x\|_2$$

stochastic approach via SOCP

- ▶ assume $a_i \sim \mathcal{N}(\bar{a}_i, \Sigma_i)$ is Gaussian, then $a_i^T x \sim \mathcal{N}(\bar{a}_i^T x, x^T \Sigma_i x)$ is also Gaussian

$$\text{prob}(a_i^T x \leq b_i) = \Phi \left(\frac{b_i - \bar{a}_i^T x}{\|\Sigma_i^{1/2} x\|_2} \right)$$

with $\Phi(x) = (1/\sqrt{2\pi}) \int_{-\infty}^x e^{-t^2/2} dt$ cumulative distribution function of $\mathcal{N}(0, 1)$

- ▶ robust LP

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & \text{prob}(a_i^T x \leq b_i) \geq \eta, \quad i = 1, \dots, m \end{array}$$

- ▶ equivalent SOCP when $\eta > 1/2$

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & \bar{a}_i^T x + \Phi^{-1}(\eta) \|\Sigma_i^{1/2} x\|_2 \leq b_i, \quad i = 1, \dots, m \end{array}$$

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Monomials and posynomials

► monomial function

$$f(x) = cx_1^{a_1} \cdots x_n^{a_n}, \quad \mathbf{dom} f = \mathbb{R}_{++}^n$$

with $c > 0$ and $a_i \in \mathbb{R}$

► posynomial function

$$f(x) = \sum_{k=1}^K c_k x_1^{a_{1k}} \cdots x_n^{a_{nk}}, \quad \mathbf{dom} f = \mathbb{R}_{++}^n$$

sum of monomials

change variables to $y_i = \log x_i$ and take logarithm

► monomial $f(x) = cx_1^{a_1} \cdots x_n^{a_n}$ transforms to

$$\log f(e^{y_1}, \dots, e^{y_n}) = a^T y + b, \quad (b = \log c)$$

► posynomial $f(x) = \sum_{k=1}^K c_k x_1^{a_{1k}} \cdots x_n^{a_{nk}}$ transforms to

$$\log f(e^{y_1}, \dots, e^{y_n}) = \log \left(\sum_{k=1}^K e^{a_k^T y + b_k} \right), \quad (b_k = \log c_k)$$

Geometric program (GP)

geometric program in standard form

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 1, \quad i = 1, \dots, m \\ & h_i(x) = 1, \quad i = 1, \dots, p \end{array}$$

with f_i posynomial, h_i monomial

geometric program in convex form

change variables to $y_i = \log x_i$ and take logarithm of objective and constraints

$$\begin{array}{ll}\text{minimize} & \log \left(\sum_{k=1}^K e^{a_{0k}^T y + b_{0k}} \right) \\ \text{subject to} & \log \left(\sum_{k=1}^K e^{a_{ik}^T y + b_{ik}} \right) \leq 0, \quad i = 1, \dots, m \\ & Gy + d = 0\end{array}$$

Example

Frobenius norm diagonal scaling

- ▶ Assume $M \in \mathbb{R}^{n \times n}$ defines a linear transformation. After scaling the coordinates by $D = \mathbf{diag}(d_1, \dots, d_n) \in \mathbb{R}^{n \times n}$, the resulting matrix becomes $DM D^{-1}$.
- ▶ How to choose D such that $DM D^{-1}$ is small under the Frobenius norm?

$$\|DM D^{-1}\|_F^2 = \sum_{i,j=1}^n (DM D^{-1})_{ij}^2 = \sum_{i,j=1}^n M_{ij}^2 d_i^2 / d_j^2.$$

- ▶ It is an unconstrained geometric program

$$\text{minimize} \quad \sum_{i,j=1}^n M_{ij}^2 d_i^2 / d_j^2$$

with variable $d = (d_1, \dots, d_n)$.

Optimization problems

Convex optimization

Linear optimization

Quadratic optimization

Geometric programming

Generalized inequality constraints

Vector optimization

Convex problem with generalized inequality constraints

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \preceq_{K_i} 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

- ▶ $f_0: \mathbb{R}^n \rightarrow \mathbb{R}$ is convex
- ▶ $f_i: \mathbb{R}^n \rightarrow \mathbb{R}^{k_i}$ is K_i -convex, where K_i is a proper cone
- ▶ same properties as standard convex problem
(convex feasible set, local optimum is global, etc)

Conic form problem (cone program)

special case of above with affine objective and constraints

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Fx + g \preceq_K 0 \\ & Ax = b\end{array}$$

extends linear programming ($K = \mathbb{R}_+^m$) to nonpolyhedral cones

Semidefinite program (SDP)

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & x_1 F_1 + \cdots + x_n F_n + G \preceq 0 \\ & Ax = b\end{array}$$

with $F_i, G \in \mathbb{S}^k$

- ▶ inequality constraint is called **linear matrix inequality** (LMI)
- ▶ includes problems with multiple LMI constraints:

$$x_1 F'_1 + \cdots + x_n F'_n + G' \preceq 0 \quad \text{and} \quad x_1 F''_1 + \cdots + x_n F''_n + G'' \preceq 0$$

is equivalent to single LMI

$$x_1 \begin{bmatrix} F'_1 & 0 \\ 0 & F''_1 \end{bmatrix} + \cdots + x_n \begin{bmatrix} F'_n & 0 \\ 0 & F''_n \end{bmatrix} + \begin{bmatrix} G' & 0 \\ 0 & G'' \end{bmatrix} \preceq 0$$

LP as equivalent SDP

LP

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Ax \preceq b\end{array}$$

equivalent SDP

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & \mathbf{diag}(Ax - b) \preceq 0\end{array}$$

note different interpretation of generalized inequality

SOCp as equivalent SDP

SOCp

$$\begin{array}{ll}\text{minimize} & f^T x \\ \text{subject to} & \|A_i x + b_i\|_2 \leq c_i^T x + d_i, \quad i = 1, \dots, m\end{array}$$

equivalent SDP

$$\begin{array}{ll}\text{minimize} & f^T x \\ \text{subject to} & \begin{bmatrix} (c_i^T x + d_i)I & A_i x + b_i \\ (A_i x + b_i)^T & c_i^T x + d_i \end{bmatrix} \succeq 0, \quad i = 1, \dots, m\end{array}$$

Eigenvalue minimization

$$\text{minimize} \quad \lambda_{\max}(A(x))$$

where $A(x) = A_0 + x_1 A_1 + \cdots + x_n A_n$ with given $A_i \in \mathbb{S}^k$

equivalent SDP with variables $(x, t) \in \mathbb{R}^{n+1}$

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & A(x) \preceq tI \end{array}$$

follows from

$$\lambda_{\max}(A) \leq t \quad \Longleftrightarrow \quad A \preceq tI$$

Matrix norm minimization

$$\text{minimize} \quad \|A(x)\|_2 = (\lambda_{\max}(A(x)^T A(x)))^{1/2}$$

where $A(x) = A_0 + x_1 A_1 + \cdots + x_n A_n$ with given $A_i \in \mathbb{R}^{p \times q}$

equivalent SDP with variables $x \in \mathbb{R}^n$ and $t \in \mathbb{R}$

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & \begin{bmatrix} tI & A(x) \\ A(x)^T & tI \end{bmatrix} \succeq 0 \end{array}$$

follows from

$$\begin{aligned} \|A\|_2 \leq t & \iff A^T A \preceq t^2 I, \quad t \geq 0 \\ & \iff \begin{bmatrix} tI & A \\ A^T & tI \end{bmatrix} \succeq 0 \end{aligned}$$

Optimization problems

Convex optimization

Linear optimization

Quadratic optimization

Geometric programming

Generalized inequality constraints

Vector optimization

Vector optimization

general vector optimization problem

$$\begin{array}{ll}\text{minimize (with respect to } K) & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p\end{array}$$

vector objective $f_0: \mathbb{R}^n \rightarrow \mathbb{R}^q$ minimized with respect to proper cone $K \subseteq \mathbb{R}^q$

convex vector optimization problem

$$\begin{array}{ll}\text{minimize (with respect to } K) & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

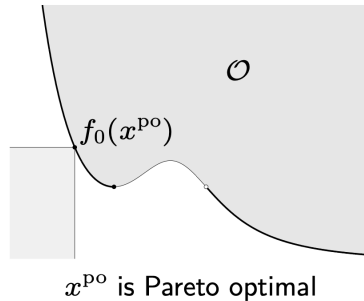
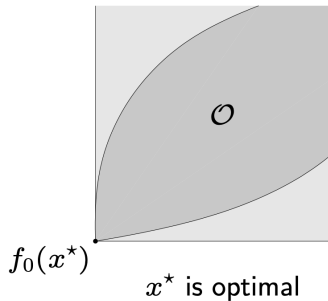
where f_0 is K -convex and f_1, \dots, f_m are convex

Optimal and Pareto optimal points

set of achievable objective values

$$\mathcal{O} = \{f_0(x) \mid x \text{ feasible}\}$$

- ▶ feasible x^* is optimal if $f_0(x^*)$ is the minimum value of \mathcal{O} (optimal value)
- ▶ feasible x^{po} is Pareto optimal if $f_0(x^{\text{po}})$ is a minimal value of \mathcal{O} (Pareto optimal value)



Multicriterion optimization

vector optimization problem with $K = \mathbb{R}_+^q$

$$f_0(x) = (F_1(x), \dots, F_q(x))$$

- ▶ q different objectives F_i , we want all of them to be small
- ▶ feasible x^* is optimal if

$$y \text{ feasible} \implies f_0(x^*) \preceq f_0(y)$$

if an optimal point exists, the objectives are noncompeting

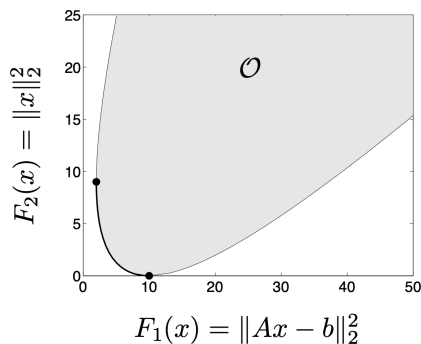
- ▶ feasible x^{po} is Pareto optimal if

$$y \text{ feasible, } f_0(y) \preceq f_0(x^{\text{po}}) \implies f_0(x^{\text{po}}) = f_0(y)$$

if multiple Pareto optimal values exist, there is a trade-off between the objectives

Regularized least-squares

minimize (with respect to \mathbb{R}_+^2) $(\|Ax - b\|_2^2, \|x\|_2^2)$

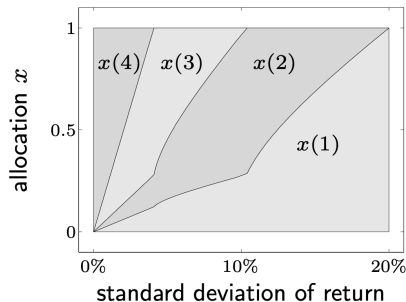
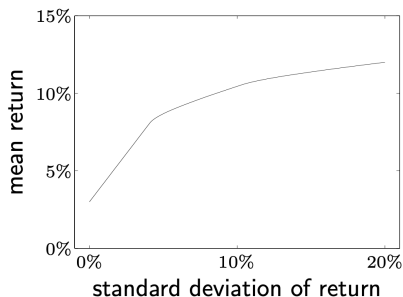


the optimal trade-off curve, shown darker, is formed by Pareto optimal points

Risk-return trade-off in portfolio optimization

$$\begin{array}{ll} \text{minimize (with respect to } \mathbb{R}_+^2) & (-\bar{p}^T x, x^T \Sigma x) \\ \text{subject to} & \mathbf{1}^T x = 1 \\ & x \succeq 0 \end{array}$$

- ▶ $x \in \mathbb{R}^n$ investment portfolio; x_i fraction invested in asset i
- ▶ $p \in \mathbb{R}^n$ (relative) asset price, random variable with mean \bar{p} and covariance Σ
- ▶ $r = p^T x$ (relative) return, random variable with mean $\bar{p}^T x$ and variance $x^T \Sigma x$



Scalarization

To find Pareto optimal points, choose $\lambda \succ_{K^*} 0$ and solve scalar problem

$$\begin{array}{ll} \text{minimize} & \lambda^T f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{array}$$

- ▶ if x is optimal for scalar problem, then it is Pareto optimal for vector optimization problem
- ▶ for convex vector optimization problem, can find (almost) all Pareto optimal points by varying $\lambda \succ_{K^*} 0$

Scalarization for multicriterion problems

In this more concrete situation

$$K = K^* = \mathbb{R}_+^q.$$

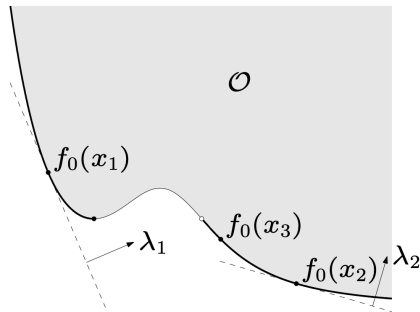
To find Pareto optimal points, write

$$\lambda = \begin{bmatrix} a_1 \\ \vdots \\ a_q \end{bmatrix} \in \mathbb{R}_{++}^q \quad \text{and} \quad f_0(x) = \begin{bmatrix} F_1(x) \\ \vdots \\ F_q(x) \end{bmatrix},$$

then minimize the positive weighted sum

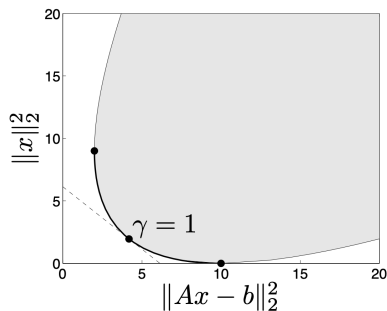
$$\lambda^T f_0(x) = a_1 F_1(x) + \cdots + a_q F_q(x)$$

Geometric interpretation



- ▶ \mathcal{O} is the set of achievable objective values
- ▶ Pareto optimal values $f_0(x_1)$ and $f_0(x_2)$ can both be obtained by scalarization: $f_0(x_1)$ minimizes $\lambda_1^T u$ and $f_0(x_2)$ minimizes $\lambda_2^T u$ over all $u \in \mathcal{O}$
- ▶ $f_0(x_3)$ is Pareto optimal, but cannot be found by scalarization

Regularized least-square problem



Take $\lambda = (1, \gamma)$ with $\gamma > 0$

$$\text{minimize} \quad \|Ax - b\|_2^2 + \gamma \|x\|_2^2$$

least-square problem for fixed $\gamma > 0$

Risk-return trade-off problem

Take $\lambda = (1, \gamma)$ with $\gamma > 0$

$$\begin{array}{ll}\text{minimize} & -\bar{p}^T x + \gamma x^T \Sigma x \\ \text{subject to} & \mathbf{1}^T x = 1 \\ & x \succeq 0\end{array}$$

quadratic program for each fixed $\gamma > 0$