

Persistent Homology

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Persistent Homology

Cech Complex

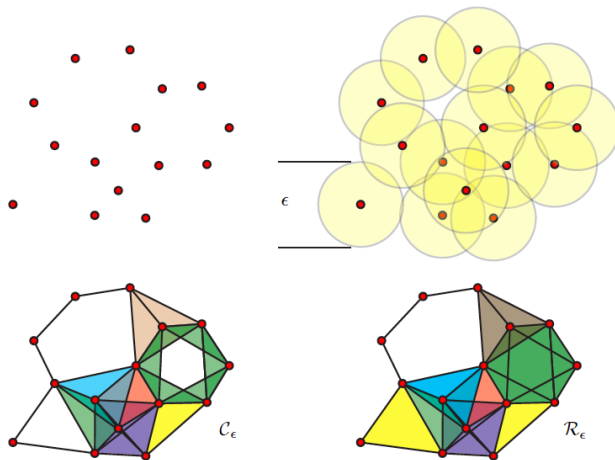


Figure: Čech complex.

Definition (Cech Complex)

Given a set of points $\{x_\alpha\}$ in Euclidean space \mathbb{R}^n , the Cech complex (also known as the nerve), \mathcal{C}_ε , is the abstract simplicial complex where a set of $k + 1$ vertices spans a k -simplex whenever the $k + 1$ corresponding closed $\varepsilon/2$ -ball neighborhoods have nonempty intersection.

Definition (Vietoris-Rips Complex)

Given a set of points $\{x_\alpha\}$ in Euclidean space \mathbb{R}^n , the Vietoris-Rips complex, \mathcal{R}_ε , is the abstract simplicial complex where a set S of $k + 1$ vertices spans a k -simplex whenever the distance between any pair of points in S is at most ε .

Cech Complex

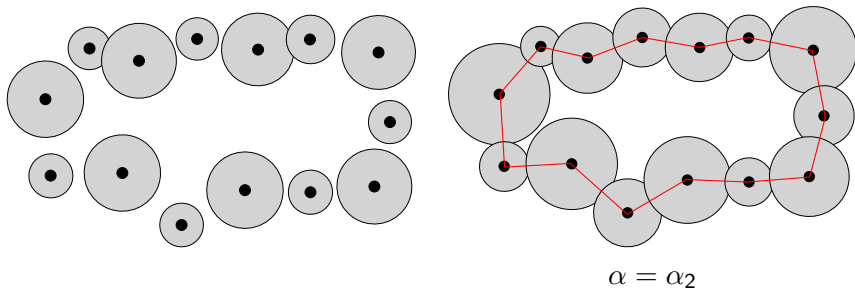
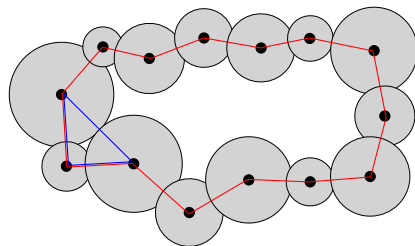
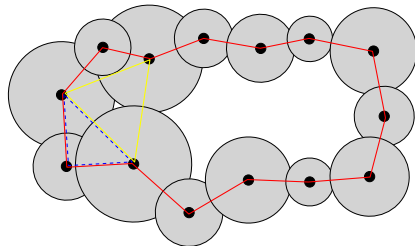


Figure: Cech complex.

Cech Complex



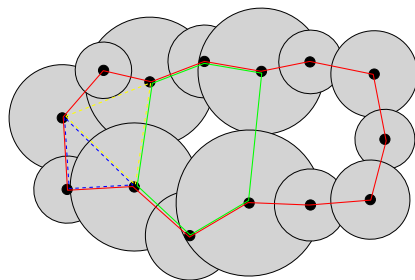
$$\alpha = \alpha_3$$



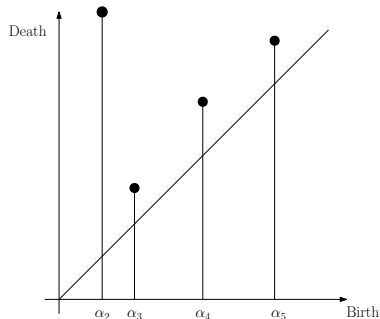
$$\alpha = \alpha_4$$

Figure: Cech complex.

Cech Complex



$$\alpha = \alpha_5$$



persistent diagram

Figure: Cech complex.

Filtration

Definition (filtration)

A filtration of a simplicial complex \mathbb{K} is a nested sequence of complexes,

$$\emptyset = \mathbb{K}_{-1} \subset \mathbb{K}_0 \subset \mathbb{K}_1 \subset \cdots \subset \mathbb{K}_n = \mathbb{K}.$$

Example

Suppose \mathbb{K} is a simplicial complex, we sort all the simplices in a sequence

$$\sigma_1^0, \sigma_2^0, \dots, \sigma_{n_0}^0, \sigma_1^1, \sigma_2^1, \dots, \sigma_{n_1}^1, \sigma_1^2, \sigma_2^2, \dots, \sigma_{n_2}^2.$$

where σ_i^k is the i -th k -simplex in \mathbb{K} . Then we relabel all the simplices as

$$\sigma^0, \sigma^1, \sigma^2, \dots,$$

We define \mathbb{K}_i as the union of $\sigma^0, \sigma^1, \dots, \sigma^i$.

Homology

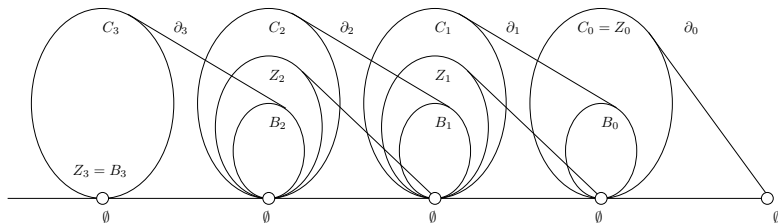


Figure: Chain, cycle, boundary groups and their images under the boundary operators.

$$H_k(\mathbb{K}, \mathbb{Z}_2) = \frac{\text{Ker} \partial_k}{\text{Im} \partial_{k+1}} = \frac{Z_k}{B_k}.$$

Persistent Homology

The inclusion map $f : \mathbb{K}_{i-1} \hookrightarrow \mathbb{K}_i$ defined by $f(x) = x$ induces a homomorphism $f_* : H_p(\mathbb{K}_{i-1}) \rightarrow H_p(\mathbb{K}_i)$. The nested sequence of complexes corresponds to a sequence of homology groups connected by the induced maps,

$$0 = H_p(\mathbb{K}_{-1}) \rightarrow H_p(\mathbb{K}_0) \rightarrow \cdots \rightarrow H_p(\mathbb{K}_n) = H_p(\mathbb{K})$$

Persistent homology studies how the homology groups change over the filtration.

Definition (positive simplex)

Given a filtration of \mathbb{K} , suppose $\mathbb{K}_i - \mathbb{K}_{i-1} = \sigma_i$, where σ_i is a $(k+1)$ -simplex. We call σ_i is **positive** if it belongs to a $(k+1)$ -cycle in \mathbb{K}_i and **negative** otherwise.

A positive simplex is also called a **generator**, a negative simplex a **killer**.

Definition (Betti Number)

Given a complex K , the i -th Betti number β_i is the rank of $H_i(K)$,

$$\beta_i = \text{Rank} H_i(K, \mathbb{Z}_2)$$

Suppose the number of positive k -simplexes is pos_k , and the number of negative k -simplexes is neg_k , then

$$\beta_k = \text{pos}_k - \text{neg}_{k+1}$$

Persistent Homology

Definition (Persistent Homology)

Define Z_k^l, B_k^l be the K -th cycle group and k -th boundary group respectively, of the l -complex K^l in a filtration. The p -persistent k -th homology group K^l is

$$H_k^{l,p} := \frac{Z_k^l}{B_k^{l+p} \cap Z_k^l}.$$

The p -persistent k -th Betti number $\beta_k^{l,p}$ of K^l is the rank of $H_k^{l,p}$.

Lemma

Consider the homomorphism $\eta_k^{l,p} : H_k^l \rightarrow H_k^{l+p}$, then

$$\text{img } \eta_k^{l,p} \cong H_k^{l,p}$$

Lemma

For each *positive* k -simplex σ^i , there exists a non-exact k -cycle c^i , c^i contains σ^i but no other positive k -simplices.

Proof.

Start with an arbitrary a k -cycle that contains σ^i and remove other positive k -simplices by adding their corresponding k -cycles. This method succeeds because each added cycle contains only one positive k -simplex by inductive assumption. \square

We use σ^i to represent c^i , and in turn the homologous class $[c^i] = c^i + B_k$.

$$\sigma^i \rightarrow c^i \rightarrow [c^i] = c^i + B_k. \quad \sigma^i \sim c^i$$

We add $[c^i]$ to the basis of $H_k(\mathbb{K}^i)$.

Generator

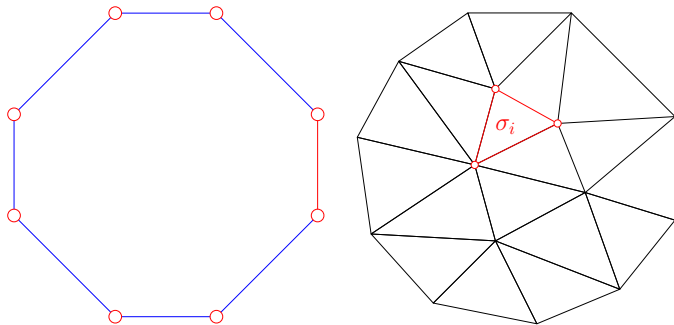


Figure: Generator, positive simplex.

For each **negative** $(k+1)$ -simplex σ^j , its boundary $d = \partial_{k+1}\sigma^j$ is a k -cycle, and can be represented as the linear combination of the basis of $H_k(\mathbb{K}_{j-1})$,

$$[d] = \sum_g [c^g], \{c^g\} \text{ basis } H_k(\mathbb{K}_{j-1}),$$

each $[c^g]$ is represented by a positive k -simplex σ^g , $g < j$, that is not yet paired. The collection of positive non-paired k -simplices is denoted as $\Gamma = \Gamma(d)$,

$$\Gamma(d) := \left\{ \sigma^g : [d] = \sum_g [c^g], \quad \sigma^g \sim c^g \right\}$$

Suppose the youngest positive simplex in $\Gamma(\partial_{k+1}\sigma^j)$ is σ^i , then we form the pair (σ^i, σ^j) , and remove $[c^i]$ from $H_k(\mathbb{K}_j)$.

$[c^i]$ is created by σ^i and killed by σ^j , the persistence life of the k -cycle $[c^i]$ is $j - i - 1$.

Example Filtration

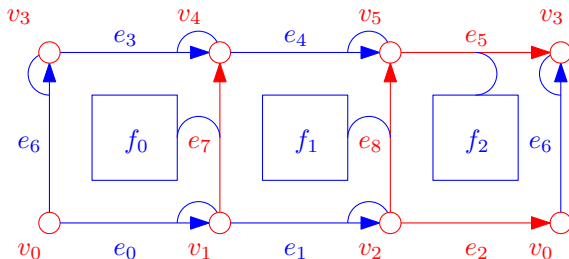


Figure: Generators and killers.

Filtration

$v_0, v_1, v_2, v_3, v_4, v_5, e_0, e_1, e_2, e_3, e_4, e_5, f_0, f_1, f_2$

Relabel them as

$\sigma^0, \sigma^1, \sigma^2, \sigma^3, \sigma^4, \sigma^5, \sigma^6, \sigma^7, \sigma^8, \sigma^9, \sigma^{10}, \sigma^{11}, \sigma^{12}, \sigma^{13}, \sigma^{14}, \sigma^{15}$

Example Generators

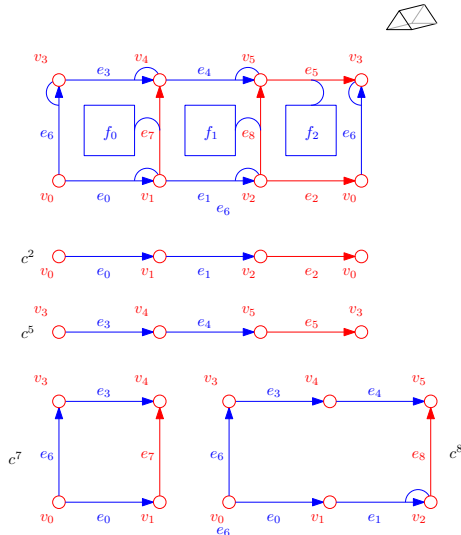


Figure: c_k contains a unique generator e_k .

Example Killers

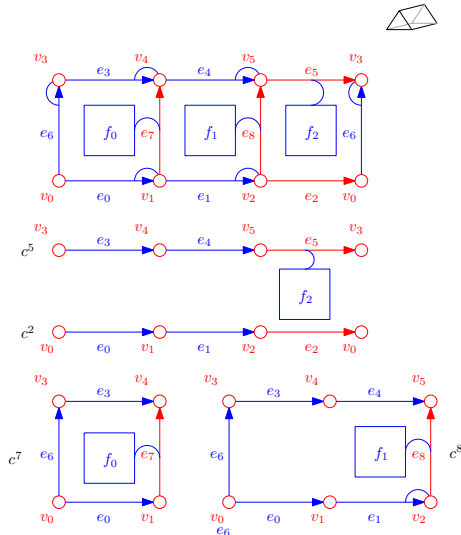


Figure: Killers.

Example Pairing

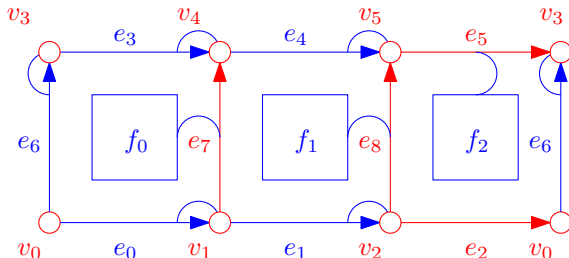


Figure: Generators and killers.

$$\begin{aligned}
 \partial_2 f_2 &= \textcolor{red}{e}_2 + \textcolor{red}{e}_5 + e_6 + e_8 = (\textcolor{red}{e}_5 + 2e_4 + 2e_3) + (\textcolor{red}{e}_2 + 2e_1 + 2e_0) + e_6 + e_8 \\
 &= (\textcolor{red}{e}_5 + e_4 + e_3) + (\textcolor{red}{e}_2 + e_1 + e_0) + \partial_2(f_0 + f_1) \\
 &= c_5 + c_2 + \partial_2(f_0 + f_1)
 \end{aligned}$$

Key Lemma

Definition (Collision Free Cycle)

A collision free cycle is one where the youngest positive simplex has not been paired (killed).

Lemma (Collision)

Given a filtration, $\mathbb{K}_j - \mathbb{K}_{j-1} = \sigma^j$, σ^j is the youngest positive simplex in $\Gamma(\partial_{k+1}\sigma^j)$. Let e be a collision free k -cycle in \mathbb{K}_{j-1} homologous to $\partial_{k+1}\sigma^j$. Suppose the youngest positive simplex in e is σ^g , then

$$\sigma^j = \sigma^g.$$

$$\max \Gamma(\partial_{k+1}\sigma^j) = \max(e) \quad \forall e \text{ collision free, } [e] = [\partial\sigma^j].$$

Key Lemma

Proof.

Let f be the sum of the basis cycles, homologous to $d = \partial_{k+1}\sigma^j$. By definition, f 's youngest positive simplex is σ^i , namely the youngest simplex in $\Gamma(\partial_{k+1}\sigma^j)$,

$$\sigma^i = \max \Gamma(\partial_{k+1}\sigma^j).$$

This implies that there are no cycles homologous to d in \mathbb{K}_{i-1} or earlier complexes. Let σ^g be the youngest positive simplex in e . $[e] = [d]$, therefore $g \geq i$.

If $g > i$, then $e = f + c$, where c bounds in \mathbb{K}^{j-1} . $\sigma^g \notin f$, implies $\sigma^g \in c$, and as σ^g is the youngest in e , it is also the youngest in c . \square

Key Lemma

continued.

Since e is collision free, the cycle created by σ^g , denoted as c^g , is still a non-boundary cycle in \mathbb{K}_{j-1} . Hence c^g can't be c , and can't be homologous to c when c becomes a boundary. Namely, when c is killed, σ^g is not paired yet.

It follows that the negative $(k+1)$ -simplex that kills c must pair a positive k -simplex in c , which is younger than σ^g , a contradiction.

This lemma shows, when σ^j is added to \mathbb{K}_{j-1} , we need to find any collision free cycle e homologous to $\partial_{k+1}\sigma^j$, and pair σ^j with the youngest positive simplex of e .

Pair Algorithm I

Pair(σ)

- 1 $c = \partial_p \sigma$
- 2 τ is the youngest positive $(p-1)$ -simplex in c .
- 3 **while** τ is paired and c is not empty **do**
- 4 find (τ, d) , d is the p -simplex paired with τ ;
- 5 $c \leftarrow \partial_p d + c$
- 6 Update τ to be the youngest positive $(p-1)$ -simplex in c
- 7 **end while**
- 8 **if** c is not empty **then**
- 9 σ is negative p -simplex and paired with τ
- 10 **else**
- 11 σ is a positive p -simplex
- 12 **endif**

Pair Algorithm II

- Suppose the simplices are sorted as $\sigma^0, \sigma_1, \dots, \sigma^{m-1}$, the filtration is constructed as $\mathbb{K}_i - \mathbb{K}_{i-1} = \sigma^i$.
- A linear array $T[0 \dots m-1]$, a pair (σ^i, σ^j) is stored in $T[i]$; together with a list of positive simplices Λ_i defining the cycle created by σ^i and destroyed by σ^j .
- Λ_i are not necessarily the same as the ones in $\Gamma(\partial\sigma^j)$, but the sum of cycles represented by the simplices in Λ_i is homologous to $\partial\sigma^j$.
- Λ_i are generated by the Pair Algorithm II.

Pair Algorithm II

T is initialized as empty;

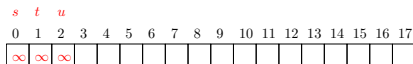
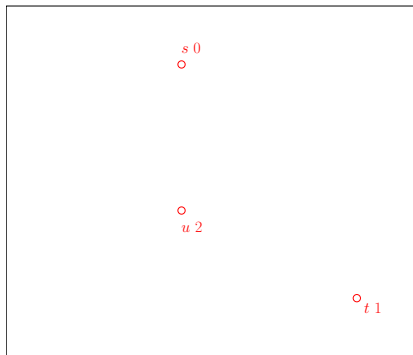
integer Pair(σ^j)

- 1 $\Lambda \leftarrow \{\sigma \in \partial_{k+1}\sigma^j \mid \sigma \text{ positive}\};$
- 2 **while** Λ is not empty **do**
- 3 $i \leftarrow \max(\Lambda);$
- 4 **if** $T[i]$ is unoccupied **then**
- 5 store j and Λ in $T[i];$
- 6 return i ; //killer
- 7 **end if**
- 8 $\Lambda \leftarrow \Lambda + \Lambda^i;$
- 9 **end while**
- 10 return ∞ ; //creator

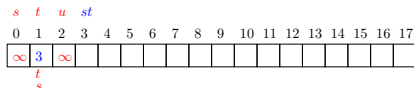
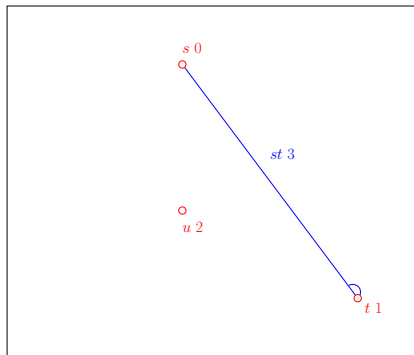
Cycle Search Halts

Observe the array T , we obtain if $\sigma^k \in \Lambda_i$, then $k \leq i$. Therefore in the while loop for cycle search, $\max(\Lambda)$ monotonously decreases, which implies that the search proceeds strictly from right to left in T . It necessarily ends at an unoccupied slot $T[g]$ of the hash table, namely Λ is a collision free cycle homologous to $\partial\sigma^j$, or Λ becomes empty, σ^j is a creator.

Example

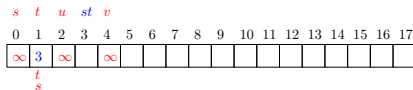
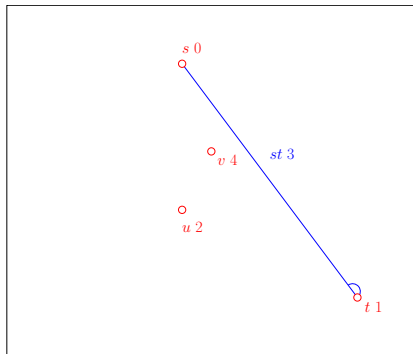


2. s_0, t_1, u_2

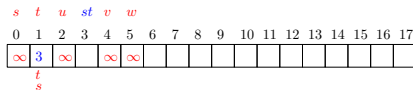
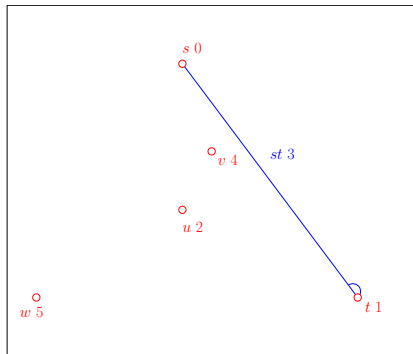


3. $\partial st_3 = s_0 + t_1, (t_1, st_3), \Lambda_1 = (t_1, s_0)$

Example

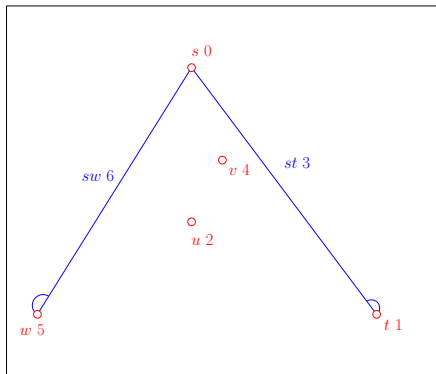


4. v_4



5. w_5

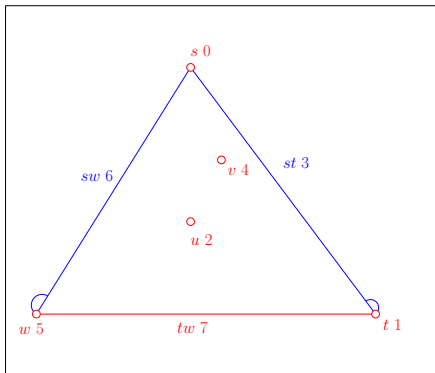
Example



$$6. \partial SW_6 = s_0 + w_5, (w_5, SW_6), \Lambda_5 = (w_5, s_0)$$

[illegible]

Example

[illegible]

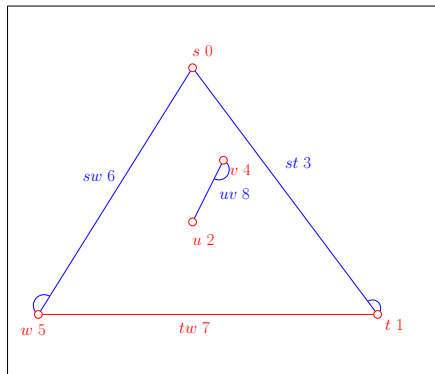
Alg 1.

$$\begin{aligned}\partial tw_7 &= w_5 + t_1 = w_5 + t_1 + \partial sw_6 \\ &= w_5 + t_1 + (s_0 + w_5) \\ &= t_1 + s_0 + \partial st_3 \\ &= t_1 + s_0 + (t_1 + s_0) \\ &= 0.\end{aligned}$$

Alg 2.

$$\begin{aligned}\partial tw_7 &= w_5 + t_1 = w_5 + t_1 + \Lambda_6 \\ &= w_5 + t_1 + (s_0 + w_5) \\ &= t_1 + s_0 + \Lambda_3 \\ &= t_1 + s_0 + (t_1 + s_0) \\ &= 0.\end{aligned}$$

Example

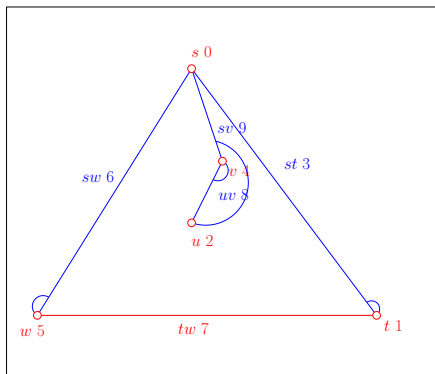


$$8. \partial uv_8 = u_2 + \boxed{v_4}, (v_4, uv_8)$$

$$\Lambda_4 = (v_4, u_2)$$

<i>s</i>	<i>t</i>	<i>u</i>	<i>st</i>	<i>v</i>	<i>w</i>	<i>sw</i>	<i>tw</i>	<i>uv</i>											
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17		
∞	3	∞		8	6		∞												
t_1	s_0			v_4	w_5														
				u_2	s_0														

Example



s	t	u	st	v	w	sw	tw	uv	sv										
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17		
∞	3	9		8	6		∞												
t_1	u_2			v_4	w_5														
s_0	s_0			u_2	s_0														

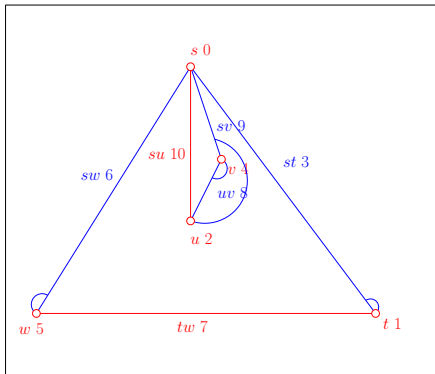
9. (u_2, sv_9) Alg. 1

$$\begin{aligned}
 \partial sv_9 &= s_0 + v_4 \\
 &= s_0 + v_4 + \partial uv_8 \\
 &= s_0 + v_4 + (u_2 + v_4) \\
 &= s_0 + \boxed{u_2} \\
 \Lambda_2 &= (s_0 + u_2)
 \end{aligned}$$

Alg. 2

$$\begin{aligned}
 \partial sv_9 &= s_0 + v_4 \\
 &= s_0 + v_4 + \Lambda_4 \\
 &= s_0 + v_4 + (u_2 + v_4) \\
 &= s_0 + \boxed{u_2} \\
 \Lambda_2 &= (s_0 + u_2)
 \end{aligned}$$

Example

[illegible]

10. SU_{10}

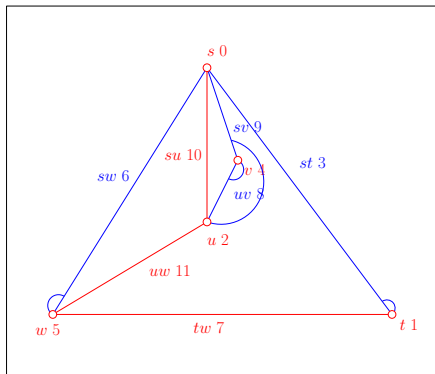
Alg. 1

$$\begin{aligned}\partial su_{10} &= s_0 + u_2 = s_0 + u_2 + \partial sv_9 \\ &= s_0 + u_2 + (s_0 + v_4) \\ &= u_2 + v_4 \\ &= u_2 + v_4 + \partial uv_8 \\ &= 0.\end{aligned}$$

Alg. 2

$$\begin{aligned}\partial su_{10} &= s_0 + u_2 = s_0 + u_2 + \Lambda_2 \\ &= s_0 + u_2 + (s_0 + u_2) \\ &= 0.\end{aligned}$$

Example

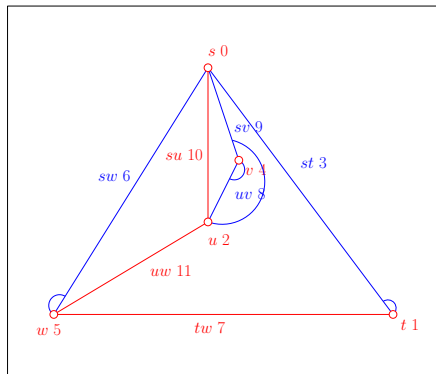


<i>s</i>	<i>t</i>	<i>u</i>	<i>st</i>	<i>v</i>	<i>w</i>	<i>sw</i>	<i>tw</i>	<i>uv</i>	<i>sv</i>	<i>su</i>	<i>uw</i>								
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17		
∞	3	9		8	6		∞			∞	∞								
t_1	u_2			v_4	w_5														
s_0	s_0			u_2	s_0														

11. uw_{11} Alg. 1

$$\begin{aligned}
 \partial uw_{11} &= u_2 + w_5 \\
 &= u_2 + w_5 + \partial sw_6 \\
 &= u_2 + w_5 + (s_0 + w_5) \\
 &= s_0 + u_2 + \partial sv_9 \\
 &= s_0 + u_2 + (s_0 + v_4) \\
 &= u_2 + v_4 + \partial uv_8 \\
 &= u_2 + v_4 + (u_2 + v_4) \\
 &= 0.
 \end{aligned}$$

Example

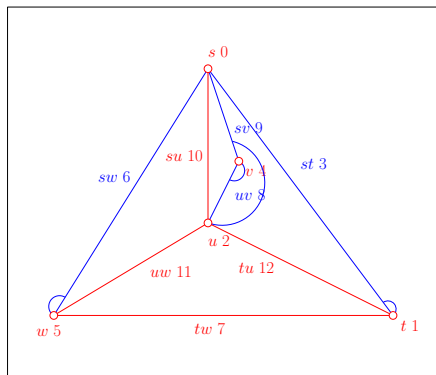


s	t	u	st	v	w	sw	tw	uv	sv	su	uw						
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
∞	3	9		8	6		∞			∞	∞						
t_1	u_2			v_4	w_5												
s_0	s_0			u_2	s_0												

11. uw_{11} Alg. 2

$$\begin{aligned}
 \partial uw_{11} &= u_2 + w_5 \\
 &= u_2 + w_5 + \Lambda_5 \\
 &= u_2 + w_5 + (s_0 + w_5) \\
 &= s_0 + u_2 + \Lambda_2 \\
 &= s_0 + u_2 + (s_0 + u_2) \\
 &= 0.
 \end{aligned}$$

Example



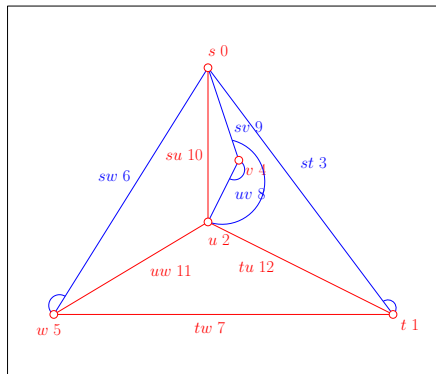
<i>s</i>	<i>t</i>	<i>u</i>	<i>st</i>	<i>v</i>	<i>w</i>	<i>sw</i>	<i>tw</i>	<i>uv</i>	<i>sv</i>	<i>su</i>	<i>uw</i>	<i>tu</i>							
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17		
∞	3	9		8	6		∞			∞	∞	∞							
<i>t</i> ₁	<i>u</i> ₂			<i>v</i> ₄	<i>w</i> ₅														
<i>s</i> ₀	<i>s</i> ₀			<i>u</i> ₂	<i>s</i> ₀														

12. tu_{12}

Alg. 1

$$\begin{aligned}
 \partial tu_{12} &= t_1 + u_2 = t_1 + u_2 + \partial sv_9 \\
 &= t_1 + u_2 + (s_0 + v_4) \\
 &= t_1 + u_2 + s_0 + v_4 + \partial uv_8 \\
 &= t_1 + u_2 + s_0 + v_4 + (u_2 + v_4) \\
 &= t_1 + s_0 \\
 &= s_0 + t_1 + \partial st_3 \\
 &= s_0 + t_1 + (t_1 + s_0) \\
 &= 0.
 \end{aligned}$$

Example



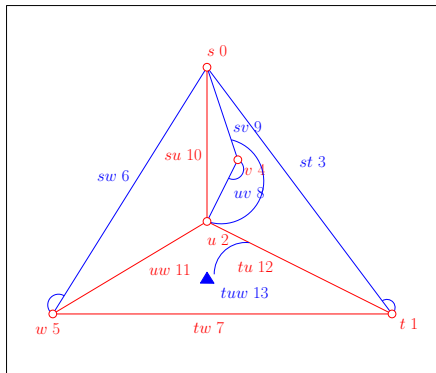
12. tu_{12}

Alg. 2

$$\begin{aligned}
 \partial tu_{12} &= t_1 + u_2 = t_1 + u_2 + \Lambda_2 \\
 &= t_1 + u_2 + (s_0 + u_2) \\
 &= t_1 + s_0 + \Lambda_1 \\
 &= s_0 + t_1 + (t_1 + s_0) \\
 &= 0.
 \end{aligned}$$

s	t	u	st	v	w	sw	tw	uv	sv	su	uw	tu						
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	
∞	3	9		8	6		∞			∞	∞	∞						
t_1	u_2			v_4	w_5													
s_0	s_0			u_2	s_0													

Example



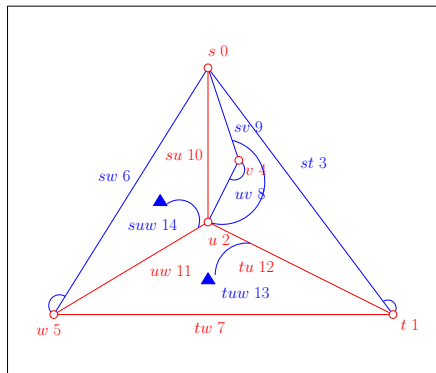
13. $tuw_{13}, (tu_{12}, tuw_{13})$

$$\partial tuw_{13} = tu_{12} + uw_{11} + wt_7$$

$$\Lambda_{12} = (\boxed{tu_{12}}, uw_{11}, wt_7) \\ (tu_{12}, tw_{13})$$

s	t	u	st	v	w	sw	tw	uv	sv	su	uw	tu	tw						
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17		
∞	3	9		8	6		∞			∞	∞	13							
	t_1	u_2		v_4	w_5							tu_{12}							
	s_0	s_0		u_2	s_0							uw_{11}							
												tw_7							

Example



14. suw_{14}

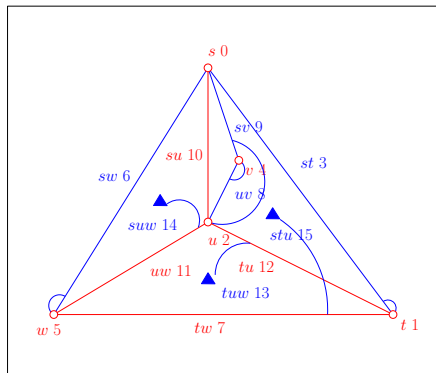
$$\partial suw_{14} = uw_{11} + su_{10} + sw_6$$

$$\Lambda_{12} = (uw_{11}, su_{10})$$

$$(uw_{11}, suw_{14})$$

s	t	u	st	v	w	sw	tw	uv	sv	su	uw	tu	$tuwsuw$				
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
∞	3	9		8	6		∞			∞	14	13					
t_1	u_2			v_4	w_5					uw	tu						
s_0	s_0			u_2	s_0					su	uw	tw					

Example

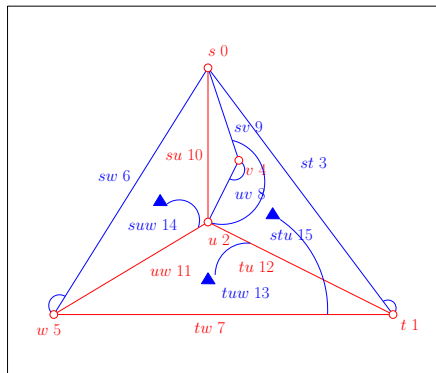


s	t	u	st	v	w	sw	tw	uv	sv	su	uw	tu	twsw	stuw	stus
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
∞	3	9		8	6		15			∞	14	13			
t_1	u_2			v_4	w_5		tw			uv	tu				
s_0	s_0			u_2	s_0					sw	tw				

15. $stu, (tw_7, stu_{15}), \Lambda_7 = (tw_7)$
Alg. 1

$$\begin{aligned}
 & \partial stu_{15} \\
 &= su_{10} + tu_{12} + st_3 \\
 &= su_{10} + st_3 + tu_{12} + \partial tuw_{13} \\
 &= su_{10} + st_3 + tu_{12} + \\
 & \quad (tu_{12} + uw_{11} + tw_7) \\
 &= su_{10} + st + uw_{11} + tw_7 \\
 &= su_{10} + st_3 + uw_{11} + tw_7 + \partial suw_{14} \\
 &= su_{10} + st_3 + uw_{11} + tw_7 + \\
 & \quad (sw_6 + su_{10} + uw_{11}) \\
 &= st_3 + \boxed{tw_7} + sw_6
 \end{aligned}$$

Example

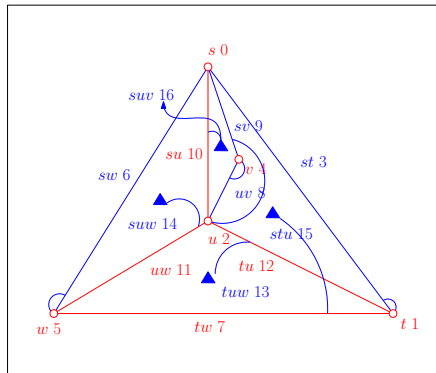


s	t	u	st	v	w	sw	tw	uv	sv	su	uw	tu	tw	sw	st	uw	st
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
∞	3	9		8	6		15			∞	14	13					
t_1	u_2			v_4	w_5		tw			uv	tu						
s_0	s_0			u_2	s_0					su	uw	tw					

15. $stu, (tw_7, stu_{15}), \Lambda_7 = (tw_7)$
Alg. 2

$$\begin{aligned}
 & \partial stu_{15} \\
 &= su_{10} + tu_{12} + st_3 \\
 &= su_{10} + st_3 + tu_{12} + \Lambda_{12} \\
 &= su_{10} + st_3 + tu_{12} + \\
 & \quad (tu_{12} + uw_{11} + tw_7) \\
 &= su_{10} + st + uw_{11} + tw_7 \\
 &= su_{10} + st_3 + uw_{11} + tw_7 + \Lambda_{11} \\
 &= su_{10} + st_3 + uw_{11} + tw_7 + \\
 & \quad (su_{10} + uw_{11}) \\
 &= st_3 + \boxed{tw_7}
 \end{aligned}$$

Example



<i>s</i>	<i>t</i>	<i>u</i>	<i>st</i>	<i>v</i>	<i>w</i>	<i>sw</i>	<i>tw</i>	<i>uv</i>	<i>sv</i>	<i>su</i>	<i>uw</i>	<i>tu</i>	<i>tuw</i>	<i>suw</i>	<i>stuv</i>	<i>sw</i>	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
∞	3	9		8	6		15			16	14	13					
<i>t</i> ₁	<i>u</i> ₂			<i>v</i> ₄	<i>w</i> ₅		<i>tw</i>			<i>su</i>	<i>uw</i>	<i>tu</i>					
<i>s</i> ₀	<i>s</i> ₀			<i>u</i> ₂	<i>s</i> ₀					<i>su</i>	<i>uw</i>	<i>tu</i>					

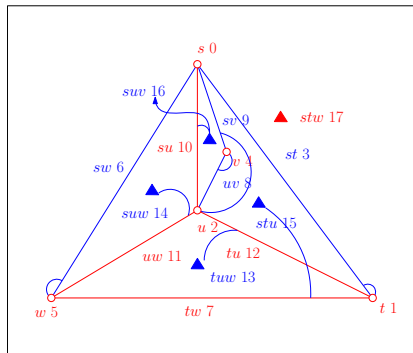
16. $su v_{16}$

$$\partial su v_{16} = \boxed{su_{10}} + uv_8 + sv_9$$

$$\Lambda_{10} = (su_{10})$$

$$(su_{10}, su v_{16})$$

Example



s	t	u	st	v	w	sw	tw	uv	sv	su	uw	tu	tw	suw	stu	svw	stw
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
∞	3	9		8	6		15			16	14	13					∞
t_1	u_2	v_4	w_5	u_2	s_0		tw			su	uw	tu					
s_0	s_0									su	uw	tu					
										su	uw	tw					

17. stw_{17}

Example

Alg. 1

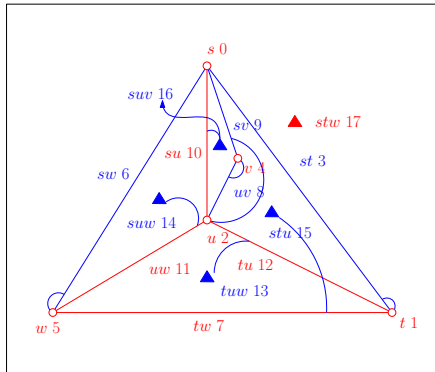
$$\begin{aligned} 17. \partial stw_{17} &= tw_7 + sw_6 + st_3 \\ &= sw_6 + st_3 + tw_7 + \partial stu_{15} \\ &= sw_6 + st_3 + tw_7 + (st_3 + tu_{12} + us_{10}) \\ &= sw_6 + st_3 + tw_7 + (st + tu_{12} + us_{10}) + \partial tuw \\ &= sw_6 + st_3 + tw_7 + (st_3 + tu_{12} + us_{10}) + (tu_{12} + uw_{11} + tw_7) \\ &= sw_6 + us_{10} + uw_{11} \\ &= sw_6 + us_{10} + uw_{11} + \partial suw_{14} \\ &= sw_6 + us_{10} + uw_{11} + (su_{10} + uw_{11} + sw_6) \\ &= 0. \end{aligned}$$

Example

Alg. 2

$$\begin{aligned} 17. \partial stw_{17} &= tw_7 + sw_6 + st_3 \\ &= sw_6 + st_3 + tw_7 + \Lambda_7 \\ &= sw_6 + st_3 + tw_7 + (tw_7) \\ &= sw_6 + st_3 \\ &= 0. \end{aligned}$$

Example



s	t	u	st	v	w	sw	tw	uv	sv	su	uw	tu	$twsw$	$stus$	sv	stw	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
∞	3	9		8	6		15			16	14	13					∞
	t_1	u_2		v_4	w_5		tw			su	uv	tu					
	s_0	s_0		u_2	s_0					su	uv	tw					

Creator	Killer	Λ
t_1	st_3	t_1, s_0
u_2	sv_9	u_2, s_0
v_4	uv_8	v_4, u_2
w_5	sw_6	w_5, s_0
tw_7	stu_{15}	tw_7
su_{10}	suv_{16}	t_1, su_{10}
uw_{11}	suw_{14}	uw_{11}, su_{10}
tu_{12}	tuw_{13}	tu_{12}, uw_{11}, tw_7

The unpaired creators are s_0 and stw_{17} .

Incidence Matrix

Assuming an ordering of the $(p - 1)$ simplices and of the p -simplices, the boundary of a p -chain can be obtained by multiplication of the corresponding vector with the incidence matrix,

$$\partial(c_p) = D_p c_p.$$

The incidence matrix is defined as

$$D_p[i, j] = \begin{cases} 1 & \sigma_i^{p-1} \in \sigma_j^p \\ 0 & \sigma_i^{p-1} \notin \sigma_j^p \end{cases}$$

Incidence Matrix and Betti Number

A classic algorithm computes the Betti numbers of K by reducing its incidence matrices to Smith normal form. It uses row and column operations to zero out all entries except along an initial portion of the diagonal.

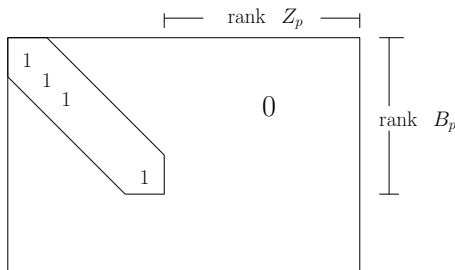


Figure: Smith norm of incidence matrix in \mathbb{Z}_2 .

The Betti number

$$\beta_p = \text{rank} Z_p - \text{rank} B_p.$$

Pairing Algorithm

Definition (Monotonous Filtering)

A filtering is monotonous, if in the ordering of K , any simplex σ is preceded by its faces.

An algorithm computes the persistence diagrams by pairing the simplices, which uses column operator to reduce D to another 0 – 1 matrix R . Let $\text{low}_R(j)$ be the row index of the last 1 in column j of R , and (undefined if the column is zero).

Definition (Reduced Matrix and Pairing)

We call R reduced and low_R a pairing function, if

$$\text{low}_R(j) \neq \text{low}_R(j'),$$

whenever $j \neq j'$ specify two non-zero columns.

Pairing Algorithm

Algorithm: Incidence matrix reduction

- ① $R \leftarrow D$
- ② **for** $j = 1$ **to** n **do**
- ③ **while** $\exists j' < j$ **with** $\text{low}_R(j') = \text{low}_R(j)$ **do**
- ④ add column j' to column j
- ⑤ **endwhile**
- ⑥ **endfor.**

The pairing is given by

$$(\sigma_i, \sigma_j) \iff i = \text{low}_R(j).$$

σ_i is positive, it generates a homology class; σ_j is negative, it kills a homology class.

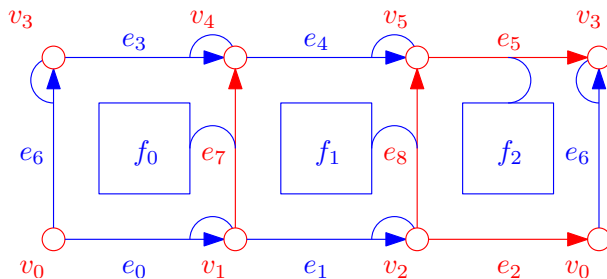


Figure: Generators and killers.

Pairing by matrix induction

Boundary operator ∂_1 , incidence matrix D_1 ,

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	1	0	0	0	1	0	0
v_1	1	1	0	0	0	0	0	1	0
v_2	0	1	1	0	0	0	0	0	1
v_3	0	0	0	1	0	1	1	0	0
v_4	0	0	0	1	1	0	0	1	0
v_5	0	0	0	0	1	1	0	0	1

Pairing by matrix induction

$1 + 2, 4 + 5, 3 + 7, 4 + 8$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	1	0	0	0	1	0	0
v_1	1	1	1	0	0	0	0	1	0
v_2	0	1	0	0	0	0	0	0	1
v_3	0	0	0	1	0	1	1	1	0
v_4	0	0	0	1	1	1	0	0	1
v_5	0	0	0	0	1	0	0	0	0

Pairing by matrix induction

$$1 + 2, 3 + 5, 3 + 8$$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	0	0	0	0	1	0	0
v_1	1	1	0	0	0	0	0	1	0
v_2	0	1	0	0	0	0	0	0	1
v_3	0	0	0	1	0	0	1	1	1
v_4	0	0	0	1	1	0	0	0	0
v_5	0	0	0	0	1	0	0	0	0

Pairing by matrix induction

$6 + 7, 6 + 8$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	0	0	0	0	1	1	1
v_1	1	1	0	0	0	0	0	1	0
v_2	0	1	0	0	0	0	0	0	1
v_3	0	0	0	1	0	0	1	0	0
v_4	0	0	0	1	1	0	0	0	0
v_5	0	0	0	0	1	0	0	0	0

Pairing by matrix induction

$$0 + 7, 1 + 8, 0 + 8$$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	0	0	0	0	1	0	0
v_1	1	1	0	0	0	0	0	0	0
v_2	0	1	0	0	0	0	0	0	0
v_3	0	0	0	1	0	0	1	0	0
v_4	0	0	0	1	1	0	0	0	0
v_5	0	0	0	0	1	0	0	0	0

Generators e_2, e_5, e_7, e_8 , corresponding to 0 columns. Killers corresponds to non-zero columns. Pairing

$$(e_0, v_1), (e_1, v_2), (e_3, v_4), (e_4, v_5), (e_6, v_3)$$

Pairing by matrix induction

	f_0	f_1	f_2
e_0	1	0	0
e_1	0	1	0
e_2	0	0	1
e_3	1	0	0
e_4	0	1	0
e_5	0	0	1
e_6	1	0	1
e_7	1	1	0
e_8	0	1	1

$$= \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix} \xrightarrow{2+3} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 1 \\ 0 & 1 & 0 \end{bmatrix} \xrightarrow{1+3} \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

The pairing is

$$(f_0, e_7), (f_1, e_8), (f_2, e_5)$$

Topological Annulus

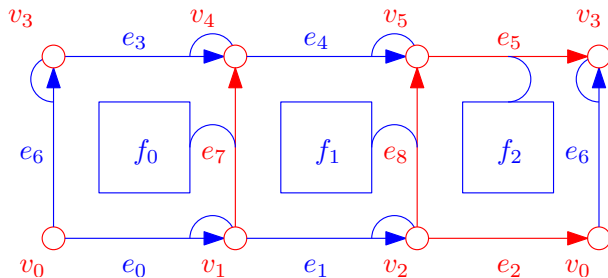
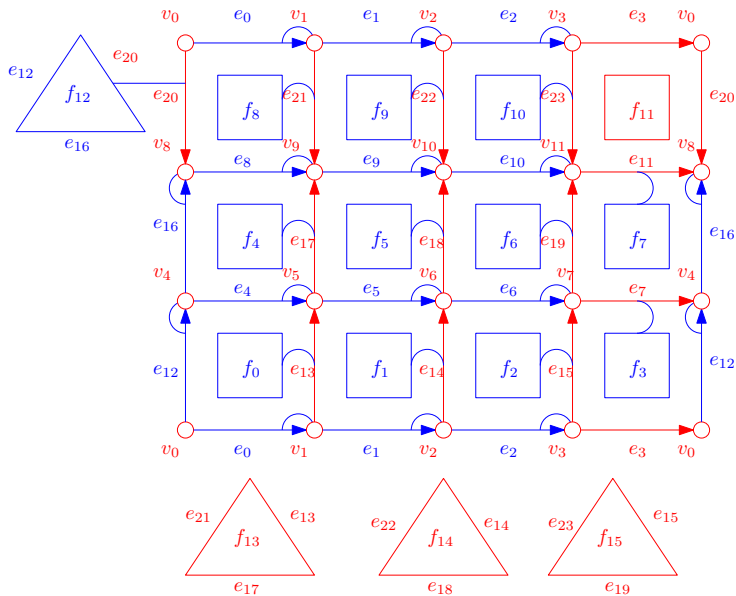
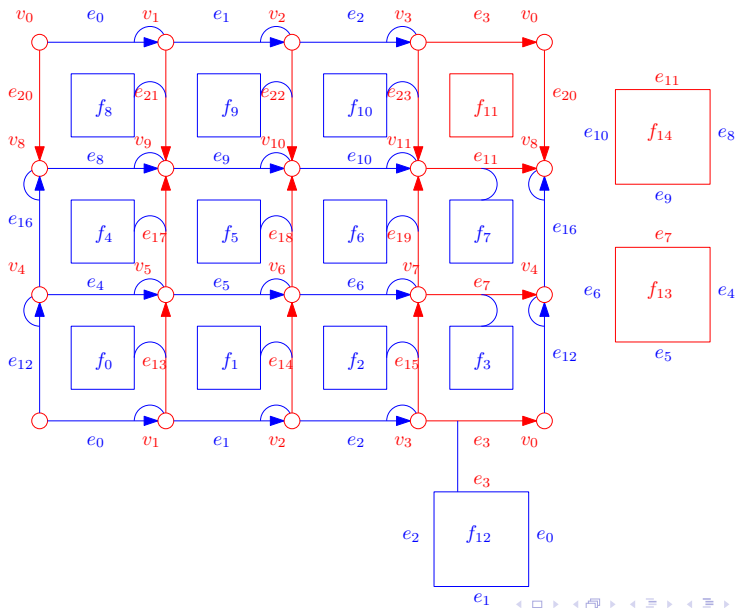


Figure: Topological Annulus. The unpaired creator is v_0 and e_2 .

Topological Torus



Topological Torus



Topological Torus

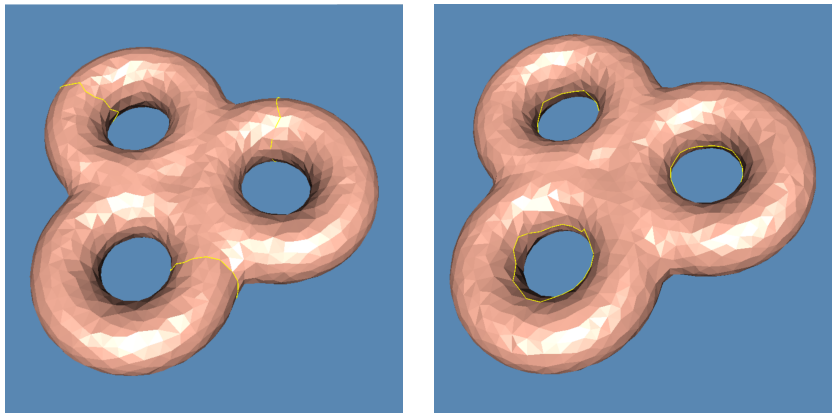


Figure: Handle and tunnel loops.

Handle Loop and Tunnel Loop

- 1 The simplices on the surface M are added into the filtration in any arbitrary order. Since $H_1(M)$ is of rank $2g$, the algorithm Pair generates $2g$ number of unpaired positive edges.
- 2 The simplices up to dimension 2 in I are added into the filtration. Since $H_1(I)$ of rank g , half of $2g$ positive edges generated in step 1 get paired with the negative triangles in I . Each pair corresponds to a killed loop, these g loops are handle loops.
- 3 Or the simplices up to dimension 2 in O are added into the filtration. Since $H_1(O)$ of rank g , half of $2g$ positive edges generated in step 2 get paired with the negative triangles in O . Each pair corresponds to a killed loop, these g loops are tunnel loops.

Handle Loops, Tunnel Loops

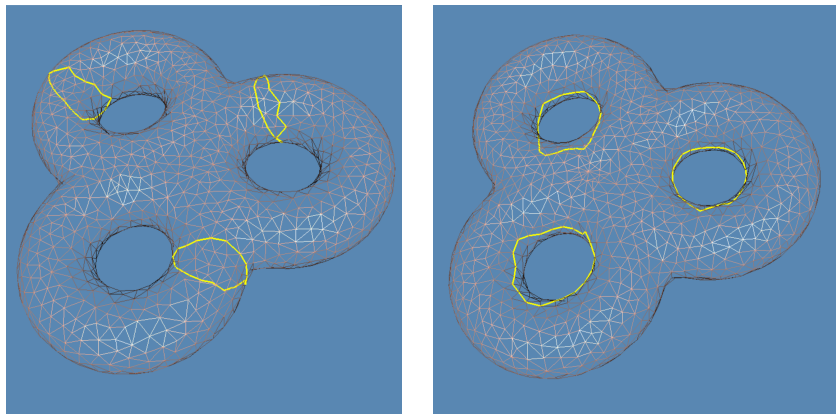


Figure: Handle and tunnel loops.

Topological Torus

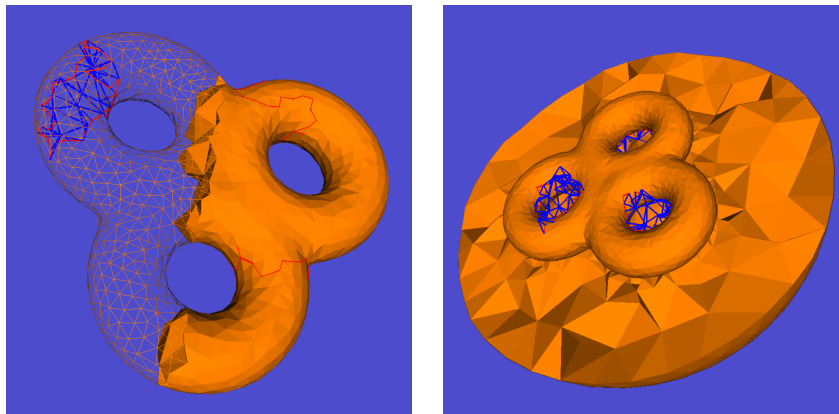


Figure: Interior and exterior volumes.

Topological Quadrilaterals

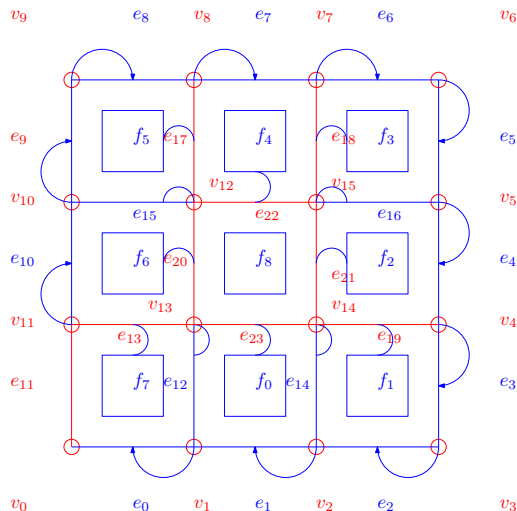
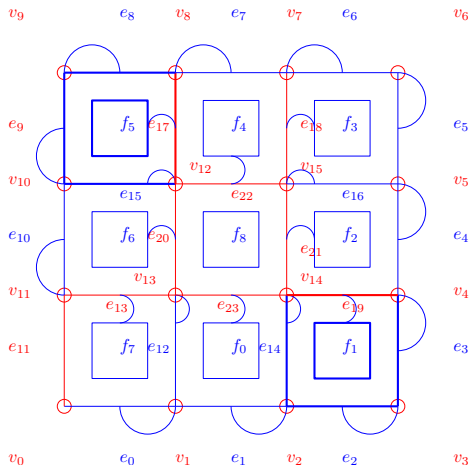


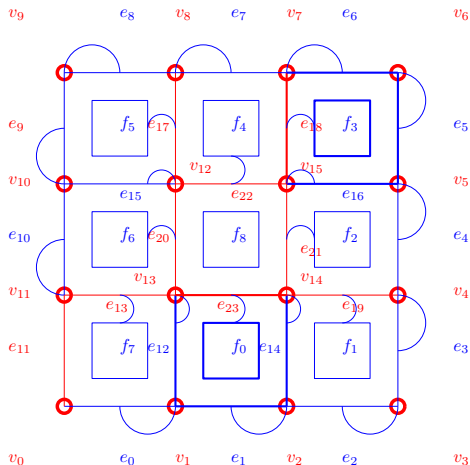
Figure: Quadrilateral example.

Topological Quadrilaterals



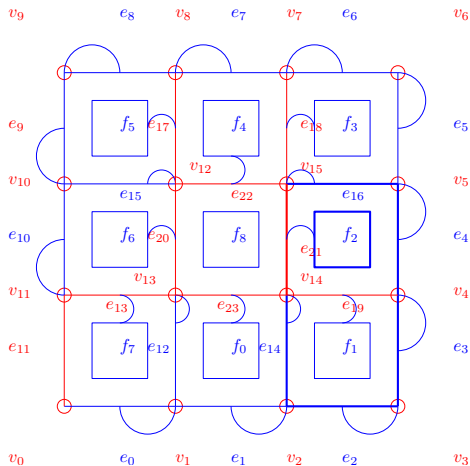
$$\partial f_5 = c_{17} = e_{15} + e_9 + e_8 + \textcolor{red}{e_{17}}, \quad \partial f_1 = c_{19} = e_3 + e_2 + e_{14} + \textcolor{red}{e_{19}}.$$

Topological Quadrilaterals



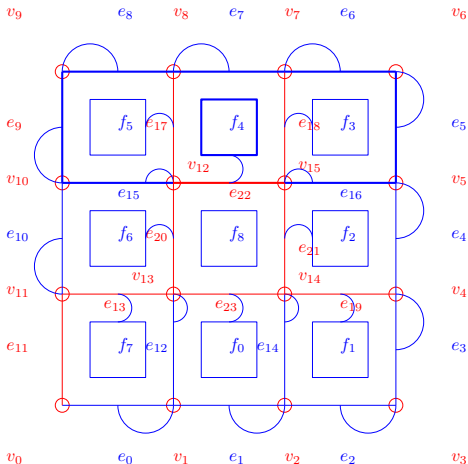
$$\partial f_3 = c_{18} = e_{16} + e_5 + e_6 + e_{18}, \quad \partial f_0 = c_{23} = e_{12} + e_{14} + e_1 + e_{23}.$$

Topological Quadrilaterals



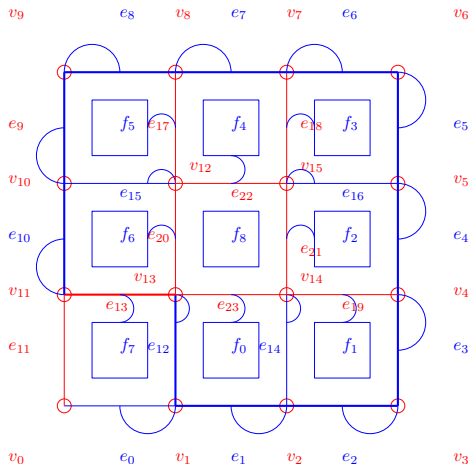
$$c_{21} = e_{16} + e_4 + e_3 + e_2 + e_{14} + e_{21}, \quad \partial f_2 = c_{19} + c_{21}$$

Topological Quadrilaterals



$$c_{22} = e_{15} + e_4 + e_{16} + e_9 + e_5 + e_8 + e_7 + e_6 + e_{22}, \partial f_4 = c_{22} + c_{18} + c_{17}$$

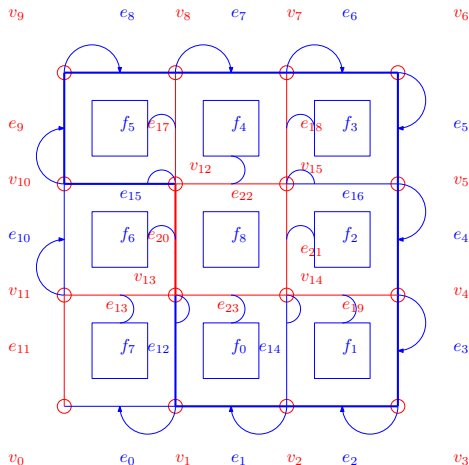
Topological Quadrilaterals



$$c_{13} = e_{12} + e_{10} + e_9 + e_8 + e_7 + e_6 + e_5 + e_4 + e_3 + e_2 + e_1 + e_{13}, \partial f_7 = c_{11} + c_{13}$$

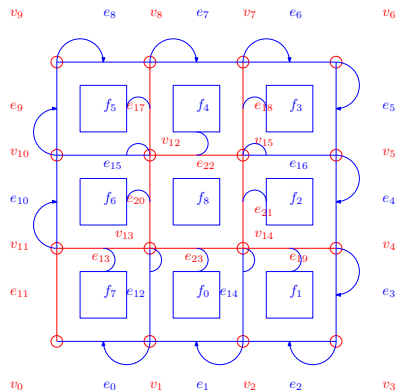
$$c_{11} = e_0 + e_1 + e_2 + e_3 + e_4 + e_5 + e_6 + e_7 + e_8 + e_9 + e_{10} + e_{11}$$

Topological Quadrilaterals



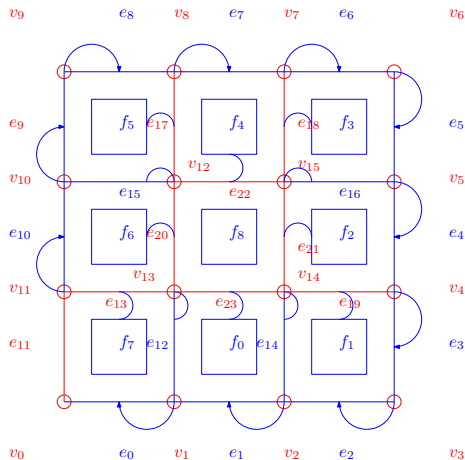
$$c_{20} = e_{15} + e_9 + e_8 + e_7 + e_6 + e_5 + e_4 + e_3 + e_2 + e_2 + e_{12} + e_{20}, \partial f_6 = c_{20} + c_{13}.$$

Topological Quadrilaterals



$$\begin{aligned}
 \partial f_0 &= c_{23}, (\textcolor{red}{e}_{23}, \textcolor{blue}{f}_0), \\
 \partial f_1 &= c_{19}, (\textcolor{red}{e}_{19}, \textcolor{blue}{f}_1), \\
 \partial f_2 &= c_{21} + \partial f_1, (\textcolor{red}{e}_{21}, \textcolor{blue}{f}_2), \\
 \partial f_3 &= c_{18}, (\textcolor{red}{e}_{18}, \textcolor{blue}{f}_3), \\
 \partial f_5 &= c_{17}, (\textcolor{red}{e}_{17}, \textcolor{blue}{f}_5), \\
 \partial f_4 &= \\
 &c_{17} + c_{18} + c_{22}, (\textcolor{red}{e}_{22}, \textcolor{blue}{f}_4), \\
 \partial f_6 &= c_{13} + c_{20}, (\textcolor{red}{e}_{20}, \textcolor{blue}{f}_6), \\
 \partial f_7 &= c_{13} + c_{11}, (\textcolor{red}{e}_{13}, \textcolor{blue}{f}_7),
 \end{aligned}$$

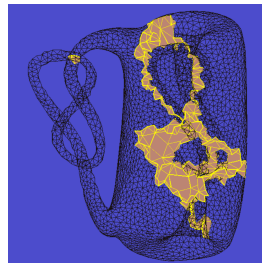
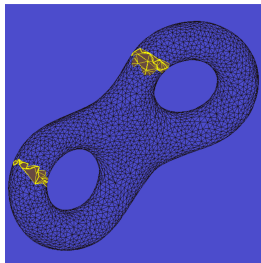
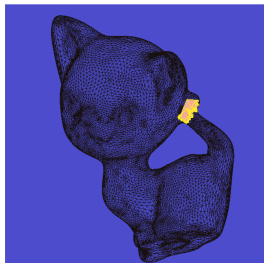
Topological Quadrilaterals



$$\partial f_8 \rightarrow \partial f_0, \partial f_4, \partial f_2, \partial f_6, \partial f_1, \partial f_3, \partial f_5, \partial f_7$$

$$\rightarrow e_0 + e_1 + e_2 + e_3 + e_4 + e_5 + e_6 + e_7 + e_8 + e_9 + e_{10} + e_{11}$$

Relative Homology



Compute $H_2(M, \partial M)$ to obtain meridians. From $H_2(M, \partial M)$ to compute $H_1(M)$.