



# UCS1302 DATA STRUCTURES

Shell sort



# Session Objectives

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- To learn about Shell sort algorithm

# Session Outcomes

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- At the end of this session, participants will be able to
  - Sort the numbers using shell sort

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# Shell sort

Dr. B. Bharathi  
SSNCE

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# Shellsort

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- Invented by Donald Shell in 1959.
- 1<sup>st</sup> algorithm to break the quadratic time barrier but few years later, a sub quadratic time bound was proven
- Shellsort works by comparing elements that are **distant** rather than adjacent elements in an array.

# Shellsort

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- Shellsort uses a sequence  $h_1, h_2, \dots, h_t$  called the **increment sequence**. Any increment sequence is fine as long as  $h_1 = 1$  and some other choices are better than others.
- Shellsort makes multiple passes through a list and sorts a number of equally sized sets using the insertion sort.
- Shellsort improves on the efficiency of insertion sort by **quickly** shifting values to their destination.

# Shellsort

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- Shellsort is also known as ***diminishing increment sort***.
- The distance between comparisons decreases as the sorting algorithm runs until the last phase in which adjacent elements are compared.
- After each phase and some increment  $h_k$ , for every  $i$ , we have  $a[i] \leq a[i + h_k]$  all elements spaced  $h_k$  apart are sorted.
- The file is said to be  $h_k$  – sorted.

# Empirical Analysis of Shellsort (Advantage)

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- Advantage of Shellsort is that its only efficient for medium size lists. For bigger lists, the algorithm is not the best choice. Fastest of all  $O(N^2)$  sorting algorithms.
- 5 times faster than the bubble sort and a little over twice as fast as the insertion sort, its closest competitor.



# Shellsort Best Case

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- Best Case: The best case in the shell sort is when the array is already sorted in the right order. The number of comparisons is less.

# Shellsort Worst Case

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- The running time of Shellsort depends on the choice of increment sequence.
- The problem with Shell's increments is that pairs of increments are not necessarily relatively prime and smaller increments can have little effect.

# Shellsort Examples

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Sort: 18 32 12 5 38 33 16 2

8 Numbers to be sorted, Shell's increment will be  $\text{floor}(n/2)$

\*  $\text{floor}(8/2) \rightarrow \text{floor}(4) = 4$

increment 4:    1            2            3            4

18   32   12   5   38   33   16   2

Step 1) Only look at 18 and 38 and sort in order ;  
18 and 38 stays at its current position because they are in order.

Step 2) Only look at 32 and 33 and sort in order ;  
32 and 33 stays at its current position because they are in order.

# Shellsort Examples

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Sort: 18 32 12 5 38 33 16 2

8 Numbers to be sorted, Shell's increment will be  $\text{floor}(n/2)$

\*  $\text{floor}(8/2) \rightarrow \text{floor}(4) = 4$

increment 4:    1            2            3            4

18   32   12   5   38   33   16   2

Step 3) Only look at 12 and 16 and sort in order ;  
12 and 16 stays at its current position because they are in order.

Step 4) Only look at 5 and 2 and sort in order ;  
2 and 5 need to be switched to be in order.

# Shellsort Examples (con't)

Sort: 18 32 12 5 38 33 16 2

Resulting numbers after increment 4 pass:

18 32 12 2 38 33 16 5

\*  $\text{floor}(4/2) \rightarrow \text{floor}(2) = 2$

increment 2: 1 2

18 32 12 2 38 33 16 5

Step 1) Look at 18, 12, 38, 16 and sort them in their appropriate location:

12 38 16 2 18 33 38 5

Step 2) Look at 32, 2, 33, 5 and sort them in their appropriate location:

12 2 16 5 18 32 38 33

## Shellsort Examples (con't)

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Sort: 18 32 12 5 38 33 16 2

\*  $\text{floor}(2/2) \rightarrow \text{floor}(1) = 1$

increment 1: 1

12	2	16	5	18	32	38	33
2	5	12	16	18	32	33	38

The last increment or phase of Shellsort is basically an Insertion Sort algorithm.