

# DATA STRUCTURES

All pair Shortest Path  
Floyd Warshall Algorithm



# Session Objectives

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- To learn about all pair shortest path algorithm using Floyd Warshall algorithm

# Session Outcomes

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- At the end of this session, participants will be able to
  - Understand all pair shortest path algorithm
  - Work on Floyd Warshall Algorithm

# Agenda

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- Floyd Warshall Algorithm
- Example

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# All pair Shortest Path (Floyd Warshall) Algorithm

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# Floyd Warshall Algorithm

The **All-Pairs Shortest Path Problem** asks to find the length of the shortest path between any pair of vertices in  $G$ .

Floyd Warshall algorithm is an **All-Pairs Shortest Path Problem**

# Floyd Warshall Algorithm

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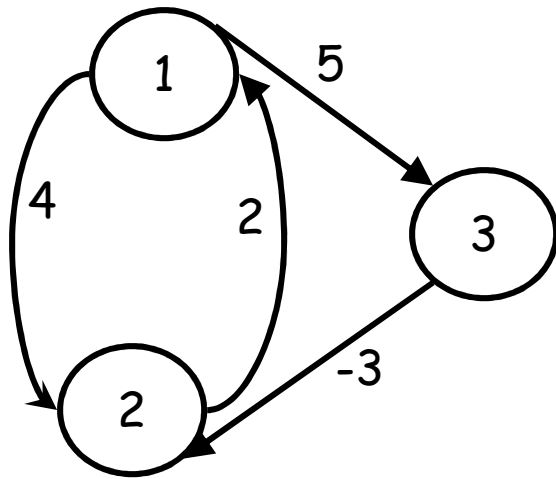
- D – Distance Matrix
- Pred – Predecessor Matrix with the intermediate information

# Floyd Warshall Algorithm

```
Floyd_Warshall(int n, int W[1..n, 1..n]) {
    array d[1..n, 1..n]
    for i = 1 to n do {
        for j = 1 to n do {
            d[i,j] = W[i,j]
            pred[i,j] = null}}
    for k = 1 to n do
        for i = 1 to n do
            for j = 1 to n do
                if (d[i,k] + d[k,j]) < d[i,j]) {
                    d[i,j] = d[i,k] + d[k,j]
                    pred[i,j] = k
                }
    return d
}
```



# Example

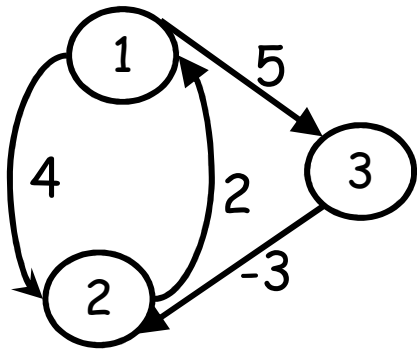


$$W = D^0 =$$

	1	2	3
1	0	4	5
2	2	0	$\infty$
3	$\infty$	-3	0

$$P =$$

	1	2	3
1	0	0	0
2	0	0	0
3	0	0	0



$$D^0 =$$

	1	2	3
1	0	4	5
2	2	0	$\infty$
3	$\infty$	-3	0

$k = 1$

Vertex 1 can be  
intermediate  
node

$$D^1 =$$

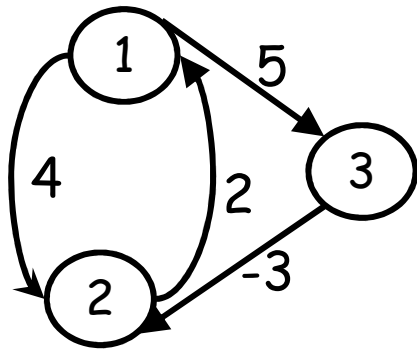
	1	2	3
1	0	4	5
2	2	0	7
3	$\infty$	-3	0

$$\begin{aligned} D^1[2,3] &= \min( D^0[2,3], D^0[2,1]+D^0[1,3] ) \\ &= \min( \infty, 7 ) \\ &= 7 \end{aligned}$$

$$P =$$

	1	2	3
1	0	0	0
2	0	0	1
3	0	0	0

$$\begin{aligned} D^1[3,2] &= \min( D^0[3,2], D^0[3,1]+D^0[1,2] ) \\ &= \min( -3, \infty ) \\ &= -3 \end{aligned}$$



$$D^1 =$$

	1	2	3
1	0	4	5
2	2	0	7
3	$\infty$	-3	0

$k = 2$

Vertices 1, 2  
can be  
intermediate

$$D^2 =$$

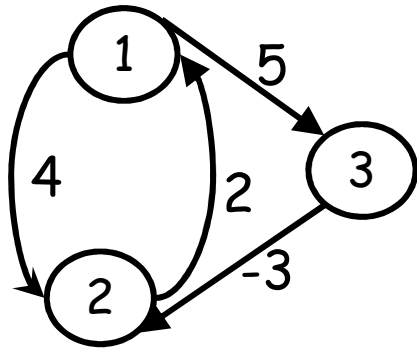
	1	2	3
1	0	4	5
2	2	0	7
3	-1	-3	0

$$\begin{aligned} D^2[1,3] &= \min( D^1[1,3], D^1[1,2]+D^1[2,3] ) \\ &= \min( 5, 4+7 ) \\ &= 5 \end{aligned}$$

$$P =$$

	1	2	3
1	0	0	0
2	0	0	1
3	2	0	0

$$\begin{aligned} D^2[3,1] &= \min( D^1[3,1], D^1[3,2]+D^1[2,1] ) \\ &= \min( \infty, -3+2 ) \\ &= -1 \end{aligned}$$



$$D^2 =$$

	1	2	3
1	0	4	5
2	2	0	7
3	-1	-3	0

$k = 3$   
 Vertices 1, 2,  
 3 can be  
 intermediate

$$D^3 =$$

	1	2	3
1	0	2	5
2	2	0	7
3	-1	-3	0

$$\begin{aligned} D^3[1,2] &= \min(D^2[1,2], D^2[1,3] + D^2[3,2]) \\ &= \min(4, 5 + (-3)) \\ &= 2 \end{aligned}$$

$$P =$$

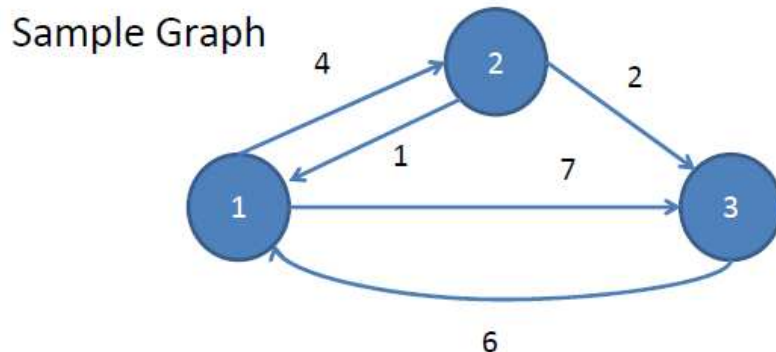
	1	2	3
1	0	3	0
2	0	0	1
3	2	0	0

$$\begin{aligned} D^3[2,1] &= \min(D^2[2,1], D^2[2,3] + D^2[3,1]) \\ &= \min(2, 7 + (-1)) \\ &= 2 \end{aligned}$$

# Distance & Predecessor Matrix Updates

D <sup>(0)</sup> - Initial			Pre. Matrix			D <sup>(1)</sup> – Intermediate Vertex			Pre. Matrix		
0	4	7	0	0	0	0	4	7	0	0	0
1	0	2	0	0	0	1	0	2	0	0	0
6	∞	0	0	0	0	6	10	0	0	1	0

Considering Intermediate **Vertex 1**



$$D(3,1) + D(1,2) < D(3,2)$$

$$10 < \infty$$

# Distance & Predecessor Matrix Updates

$D^{(2)}$  – Intermediate  
Vertex

0	4	6
1	0	2
6	10	0

Pre. Matrix

0	0	2
0	0	0
0	1	0

$D^{(3)}$  – Intermediate  
Vertex

0	4	6
1	0	2
6	10	0

Pre. Matrix

0	0	2
0	0	0
0	1	0

# How to find shortest path between any pair of vertices?

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## Algorithm:

```
Path(i,j)
{
    if pred[i,j] = null                // path is a single edge
        output(i,j)
    else
    {
        // path goes through
        pred
        Path(i, pred[i,j]); // print path from i to pred
        Path(pred[i,j], j); // print path from pred to j
    }
}
```

## Example – Shortest path from vertex 2 to vertex 3

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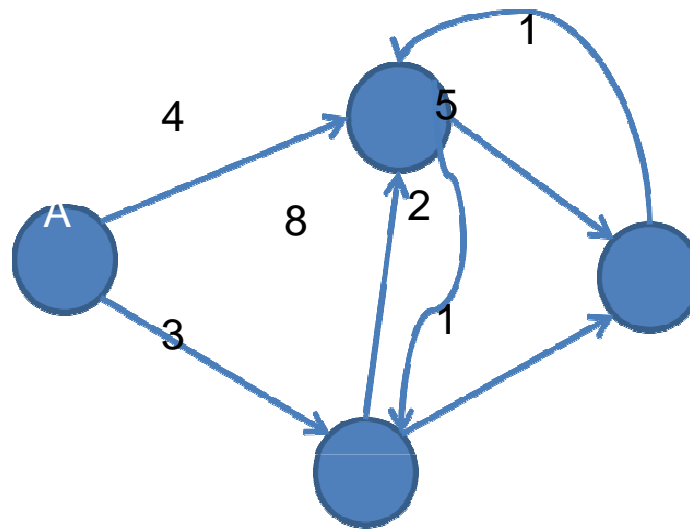
2..3	path(2,3)	pred(2,3) =4
2..4..3	path(2,4)	pred(2,4) =5
2..5..4..3	path(2,5)	pred(2,5) =nil output(2,5)
25..4..3	path(5,4)	pred(5,4) =nil output(5,4)
254..3	path(4,3)	pred(4,3)=nil output(4,3)

So Path is 2->5->4->3



# Excercise

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# Summary

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- All pair shortest path
- Floyd Warshall Algorithm
- Example