



UCS1302: DATA STRUCTURES

Graphs



Session Meta Data

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Revision History

Revision Date	Details	Version no.
22 September 2017	1. New SSN template applied	1.2

Session Objectives

- To learn about graph and its representations

Session Outcomes

- At the end of this session, participants will be able to
 - Understand graph terminologies
 - Represent the graphs using different methods

Agenda

- Graph introduction
- Terminologies
- Representation of graphs

Graphs

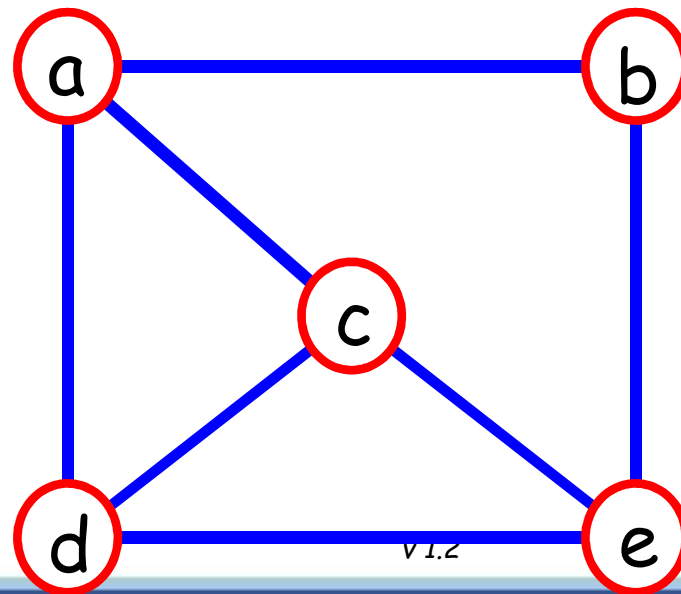
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SSNCE

August 20, 2019

What is a Graph?

- A graph $G = (V, E)$ is composed of:
 - V : set of vertices
 - E : set of edges or arcs connecting the vertices in V
- An edge $e = (u, v)$ is a pair of vertices belonging to E
- weight or cost

Example:

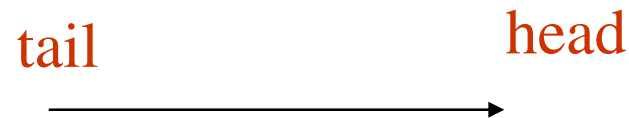


$V = \{a, b, c, d, e\}$
 $E = \{(a, b), (a, c), (a, d), (b, e), (c, d), (c, e), (d, e)\}$



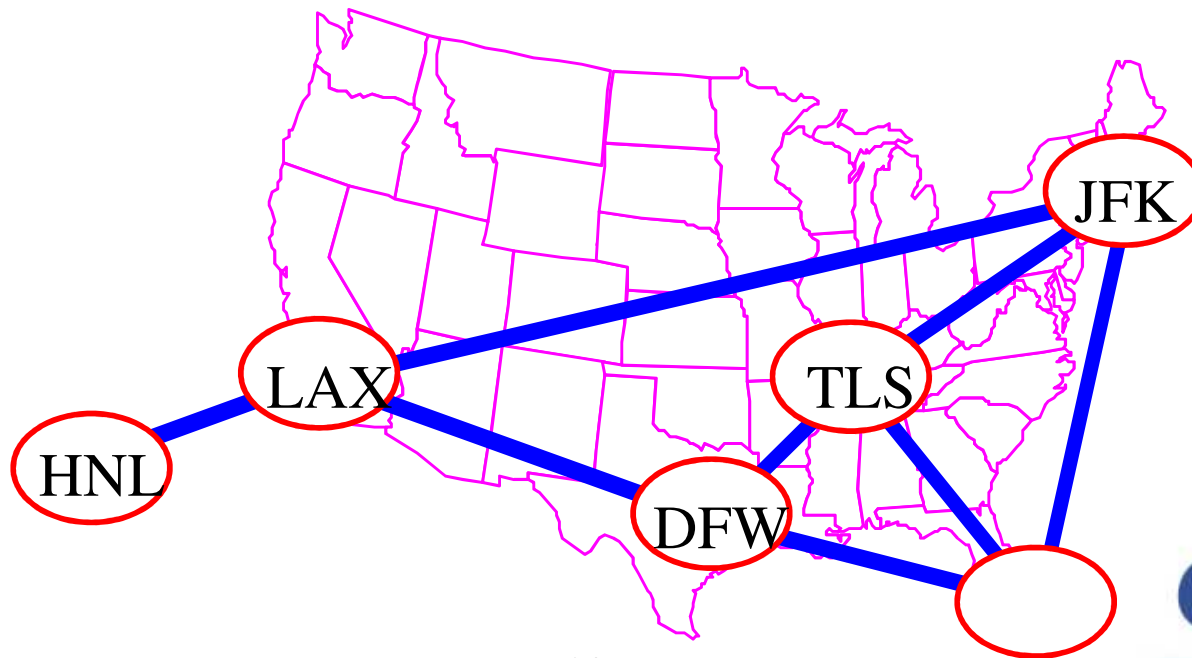
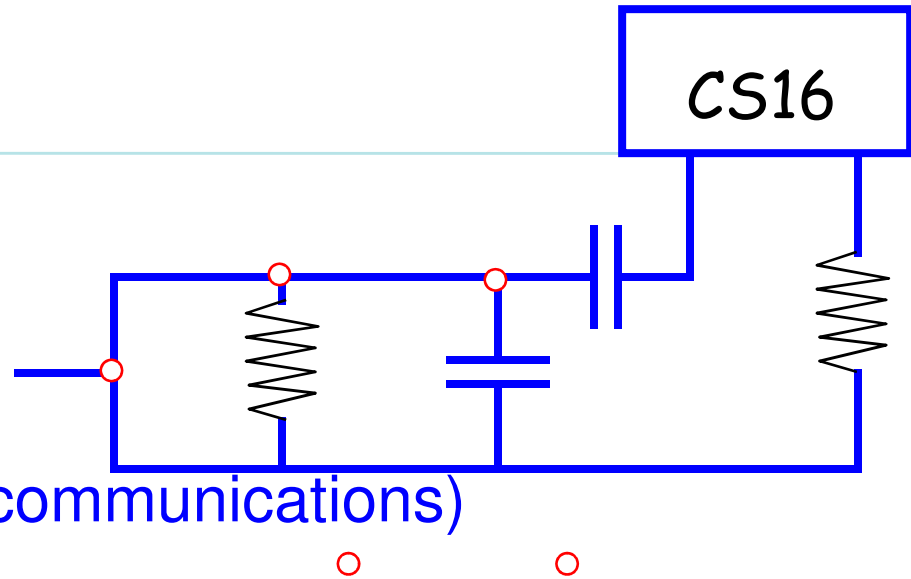
Directed vs. Undirected Graph

- An **undirected graph** is one in which the pair of vertices in a edge is unordered, $(v_0, v_1) = (v_1, v_0)$
- A **directed graph** is one in which each edge is a directed pair of vertices, $\langle v_0, v_1 \rangle \neq \langle v_1, v_0 \rangle$



Applications

- electronic circuits
- networks (roads, flights, communications)

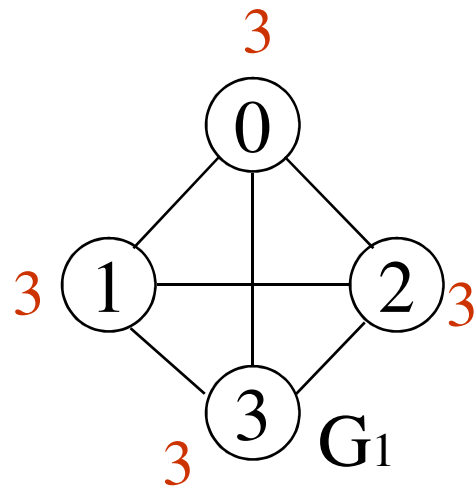


Terminology: Degree of a Vertex

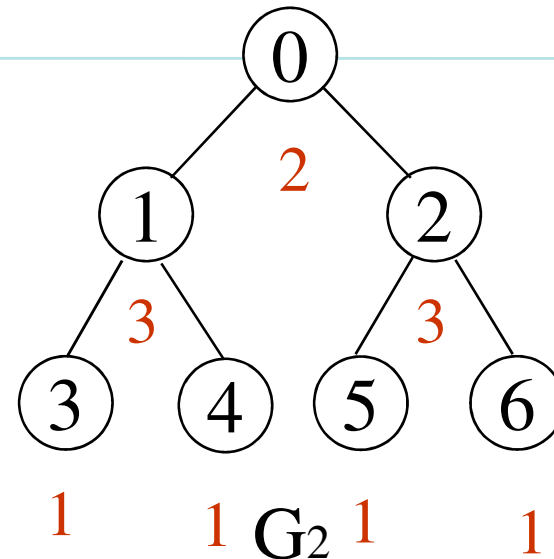
- The **degree** of a vertex is the number of edges incident to that vertex
- For directed graph,
 - the **in-degree** of a vertex v is the number of edges that have v as the head
 - the **out-degree** of a vertex v is the number of edges that have v as the tail
 - if d_i is the degree of a vertex i in a graph G with n vertices and e edges, the number of edges is

$$e = \left(\sum_{i=0}^{n-1} d_i \right) / 2$$

Examples



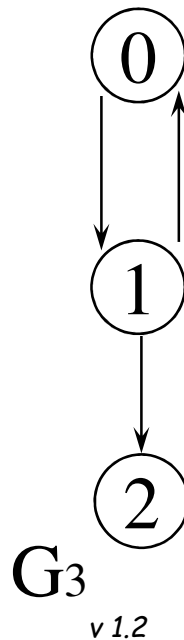
directed graph
in-degree
out-degree



in: 1, out: 1

in: 1, out: 2

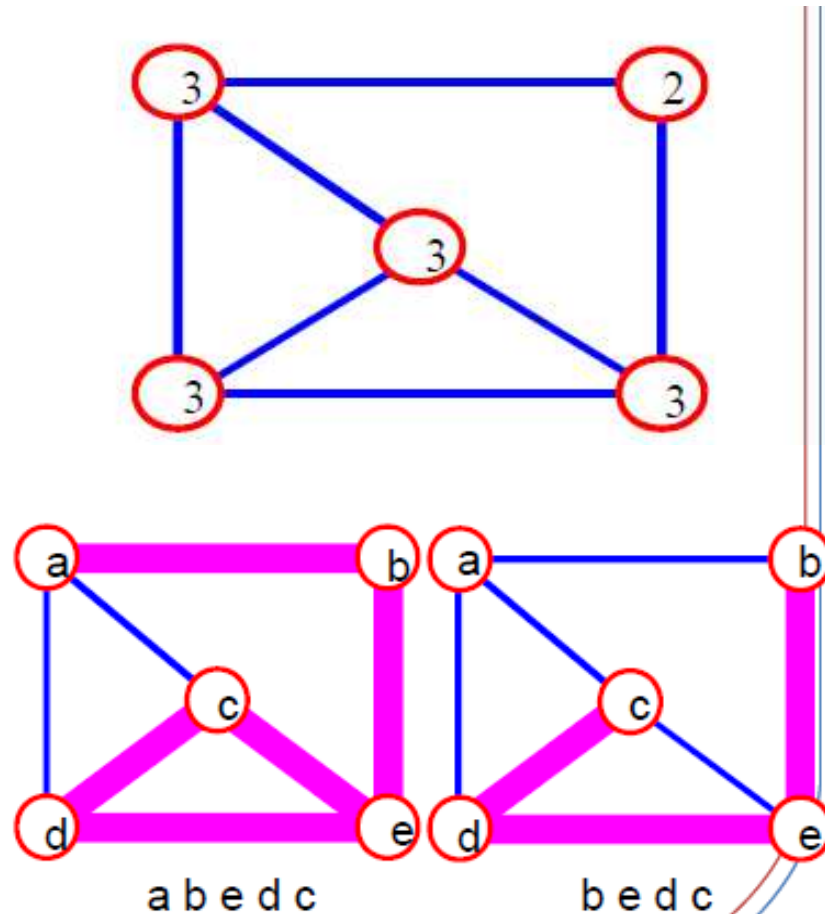
in: 1, out: 0



Path

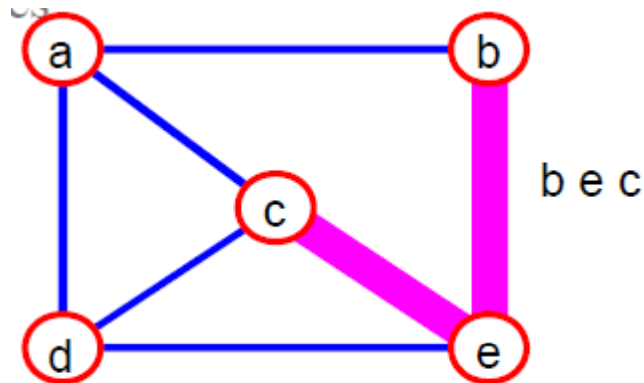
- **path**: sequence of vertices v_1, v_2, \dots, v_k such that consecutive vertices v_i and v_{i+1} are adjacent.

The **length** of the path is the number of edges along the path



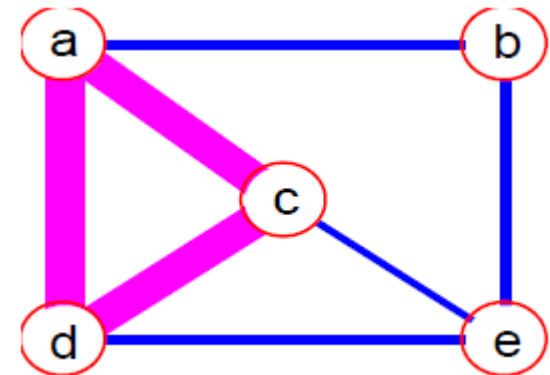
Terminology

Simple path: No repeated vertices



Cycle: simple path, except that the last vertex is the same as the first vertex

a c d a



Terminology

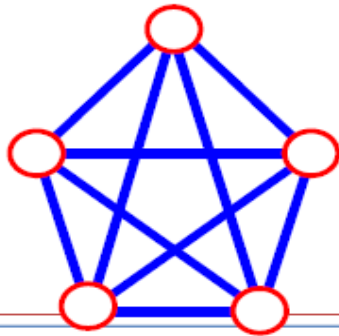
- A directed graph that has *no* cyclic paths is called a **DAG** (a Directed Acyclic Graph).
- An undirected graph that has an edge between every pair of vertices is called a **complete** graph.

Let n = no. of vertices, and m = no. of edges

- *How many total edges in a complete graph?*
 - Each of the n vertices is incident to $n-1$ edges, however, we would have counted each edge twice! Therefore, intuitively, $m = n(n-1)/2$.
- Therefore, if a graph is not complete, $m < n(n-1)/2$

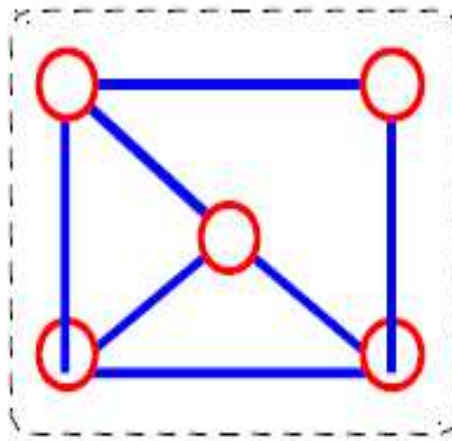
Note: A directed graph can also be a complete graph; in that case, there must be an edge from every vertex to every other vertex.

Terminology

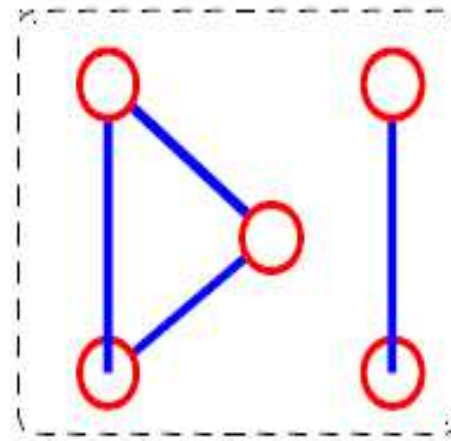


$$\mathbf{n} = 5$$
$$\mathbf{m} = (5 * 4)/2 = 10$$

- **connected graph**: any two vertices are connected by some path



connected



not connected

- **subgraph**: subset of vertices and edges forming a graph

Terminology

- An undirected graph **is connected** if a path exists from every vertex to every other vertex
- A directed graph is **strongly connected** if a path exists from every vertex to every other vertex
- A directed graph is **weakly connected** if a path exists from every vertex to every other vertex, disregarding the direction of the edge

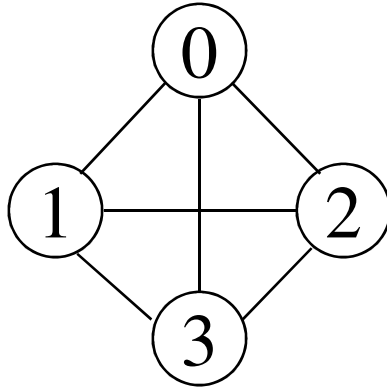
Graph representations

- **Adjacency matrix** - graph can be represented using a matrix of size total number of vertices by total number of vertices.
- **Adjacency lists** - every vertex of graph contains list of its adjacent vertices.

Adjacency matrix

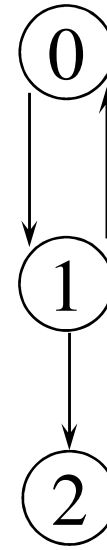
- Let $G=(V,E)$ be a graph with n vertices.
- The adjacency matrix of G is a two-dimensional n by n array, say `adj_mat`
- If the edge (v_i, v_j) is in $E(G)$, `adj_mat[i][j]=1`
- If there is no such edge in $E(G)$, `adj_mat[i][j]=0`
- The adjacency matrix for an undirected graph is symmetric; the adjacency matrix for a digraph need not be symmetric

Examples for Adjacency Matrix



$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

G_1



$$\begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$

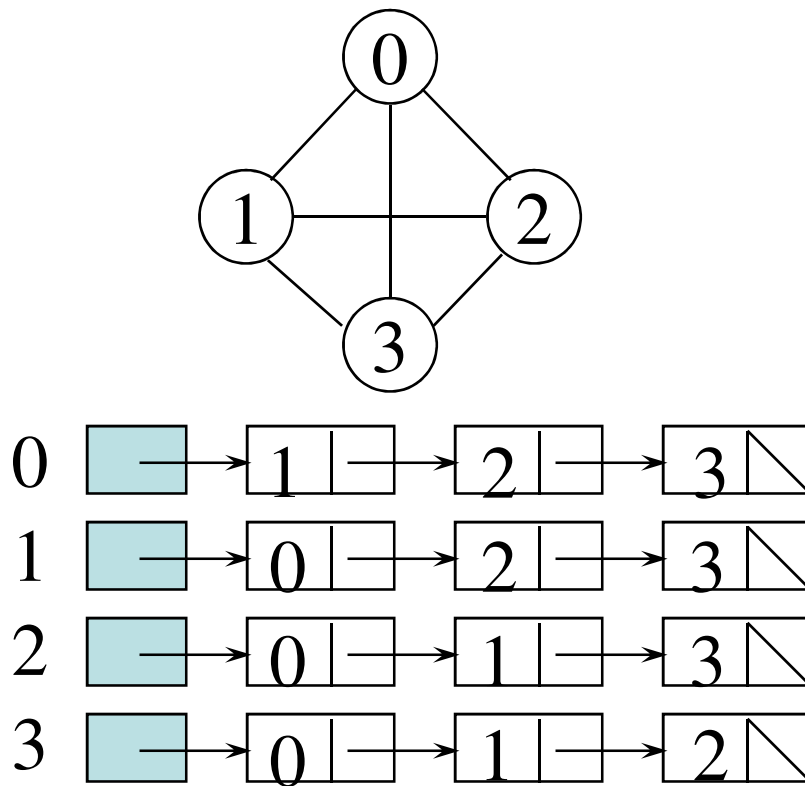
G_2

Adjacency Lists

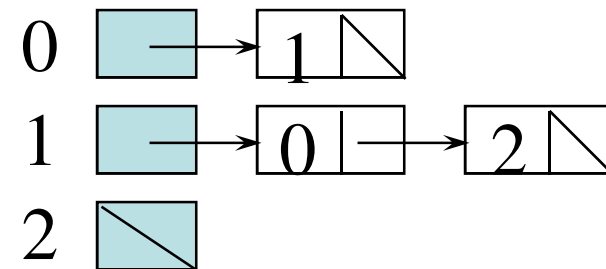
Each row in adjacency matrix is represented as an adjacency list

```
#define MAX_VERTICES 50
typedef struct node *node_pointer;
typedef struct node
{
    int vertex;
    struct node *link;
};
node_pointer graph[MAX_VERTICES];
```

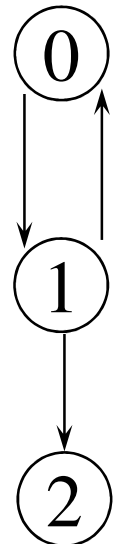
Examples for Adjacency Lists



G_1



G_2



Summary

- Introduction to graph
- Graph terminologies
- Representation of graph
 - Adjacency matrix
 - Adjacency list