

COMPUTER SCIENCE 5300

ADVANCED ALGORITHM DESIGN AND ANALYSIS

ASSIGNMENT # 1

Problem 1

Let $X(1..n)$ and $Y(1..n)$ contain two lists of n integers, each sorted in nondecreasing order. Give the best (worst-case complexity) algorithm that you can think of for finding

- (a) the largest integer of all $2n$ combined elements.
- (b) the second largest integer of all $2n$ combined elements.
- (c) the median (or the n th smallest integer) of all $2n$ combined elements.

For instance, $X = (4, 7, 8, 9, 12)$ and $Y = (1, 2, 5, 9, 10)$, then median = 7, the n th smallest, in the combined list $(1, 2, 4, 5, 7, 8, 9, 9, 10, 12)$. [Hint: use the concept similar to binary search]

Solution:

(a)

Algorithm 1 Calculate the maximum of two sorted arrays

$\text{max}(X[n-1], Y[n-1])$

(b)

Algorithm 2 Calculate the second largest element of two sorted arrays

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if  $X[n-1] == Y[n-1]$  then
    return  $\text{max}(X[n-2], Y[n-2])$ 
else if  $X[n-1] < Y[n-1]$  then
    return  $\text{max}(X[n-1], Y[n-2])$ 
else
    return  $\text{max}(X[n-1], Y[n-2])$ 
end if
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Problem 2

1-to-2 PARTITION:

Instance: A finite set of positive integers $Z = z_1, z_2, \dots, z_n$.

Question: Is there a subset Z' of Z such that $\text{Sum of all numbers in } Z' = 2 \times \text{Sum of all numbers in } Z - Z'$

- (a) Obtain the dynamic programming functional equation to solve the 1-to-2 PARTITION problem.
- (b) Give an algorithm to implement your functional equation.
- (c) Give an example of 5 numbers with a total of 21 as an input instance for 1-to-2 PARTITION problem, and show how your algorithm works on this input instance.
- (d) What is the complexity of your algorithm?

Solution

***** SOLUTION GOES HERE *****

Problem 3

Decide True or False for each of the followings. You MUST briefly justify your answer.

Satisfiability:

Instance: Set U of variables, collection C of clauses over U .

Question: Is there a satisfying truth assignment for C ?

- (a) If $P \neq NP$, then no problem in NP can be solved in polynomial time deterministically.
- (b) If a decision problem A is NP -complete, proving that A is reducible to B , in polynomial time, is sufficient to show that B is NP -complete.
- (c) It is known that SAT (Satisfiability) is NP -complete, and 3SAT (all clauses have size 3) is NP -complete. 1SAT (all clauses have size 1) is also NP -complete.

Solution

***** SOLUTION GOES HERE *****

Problem 4

Solution

***** SOLUTION GOES HERE *****

Problem 5

Solution

***** SOLUTION GOES HERE *****

Problem 6

Solution

***** SOLUTION GOES HERE *****