

## COMPUTER SCIENCE 5300

## ADVANCED ALGORITHM DESIGN AND ANALYSIS

## ASSIGNMENT # 1

**Problem 1**

Let  $X(1..n)$  and  $Y(1..n)$  contain two lists of  $n$  integers, each sorted in nondecreasing order. Give the best (worst-case complexity) algorithm that you can think of for finding

- (a) the largest integer of all  $2n$  combined elements.
- (b) the second largest integer of all  $2n$  combined elements.
- (c) the median (or the  $n$ th smallest integer) of all  $2n$  combined elements.

For instance,  $X = (4, 7, 8, 9, 12)$  and  $Y = (1, 2, 5, 9, 10)$ , then median = 7, the  $n$ th smallest, in the combined list  $(1, 2, 4, 5, 7, 8, 9, 9, 10, 12)$ . [Hint: use the concept similar to binary search]

**Solution:**

(a)

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**Algorithm 1** Calculates the largest element of two sorted arrays

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1: procedure LARGESTELEMENT (int X[n], int Y[n])
2:   return  $\max(X[n-1], Y[n-1])$ 
3: end procedure

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(b)

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**Algorithm 2** Calculate second largest element of two sorted arrays

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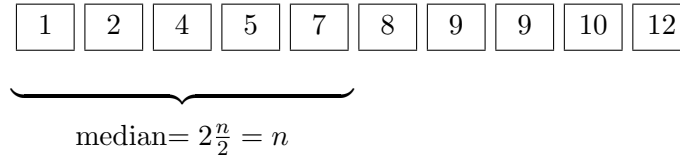
1: procedure SECONDLARGEST (int X[n], int Y[n])
2:   if  $X[n-1] == Y[n-1]$  then
3:     return  $\max(X[n-2], Y[n-2])$ 
4:   else if  $X[n-1] < Y[n-1]$  then
5:     return  $\max(X[n-1], Y[n-2])$ 
6:   else
7:     return  $\max(X[n-2], Y[n-1])$ 
8:   end if
9: end procedure

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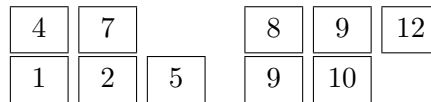
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(c)

By looking at the two sorted arrays, we see that the median value will correspond to the  $n$ th term in the list of length  $2n$ , and that this portion of the list contains part of  $X[]$  and  $Y[]$ . Also, note that the last element of any subset of the two sorted lists, starting at index zero, will always correspond to the  $k$ th element.



If we look at the two list individually, we find that the elements that lead up to the partition of the  $k$ th element are



Let the elements to the left and right of the partition of the first array be  $l_x$  and  $r_x$  respectively. And similarly, for the second array,  $l_y$  and  $r_y$ . Notice that for the solution to be valid, ie, for all the elements to the left of the partition to be less than all the elements to the right, requires

$$l_x \leq r_y$$

$$l_y \leq r_x$$

Only these conditions need to be checked because we are guaranteed to have  $l_x \leq r_x$  and  $l_y \leq r_y$ . Once these conditions are satisfied, the  $k$ th element will be  $\max(l_x, l_y)$  because we know that in the combined array, the largest element to the left of the partition corresponds to the  $k$ th element.

The question is now, how do we optimize the brute force method of finding the partition.

## Problem 2

### 1-to-2 PARTITION:

Instance: A finite set of positive integers  $Z = z_1, z_2, \dots, z_n$ .

Question: Is there a subset  $Z'$  of  $Z$  such that  $\text{Sum of all numbers in } Z' = 2 \times \text{Sum of all numbers in } Z - Z'$

- (a) Obtain the dynamic programming functional equation to solve the 1-to-2 PARTITION problem.
- (b) Give an algorithm to implement your functional equation.
- (c) Give an example of 5 numbers with a total of 21 as an input instance for 1-to-2 PARTITION problem, and show how your algorithm works on this input instance.
- (d) What is the complexity of your algorithm?

## Solution

\*\*\*\*\* SOLUTION GOES HERE \*\*\*\*\*

## Problem 3

Decide True or False for each of the followings. You MUST briefly justify your answer.

### Satisfiability:

Instance: Set  $U$  of variables, collection  $C$  of clauses over  $U$ .

Question: Is there a satisfying truth assignment for  $C$ ?

- (a) If  $P \neq NP$ , then no problem in  $NP$  can be solved in polynomial time deterministically.
- (b) If a decision problem  $A$  is  $NP$ -complete, proving that  $A$  is reducible to  $B$ , in polynomial time, is sufficient to show that  $B$  is  $NP$ -complete.
- (c) It is known that SAT (Satisfiability) is  $NP$ -complete, and 3SAT (all clauses have size 3) is  $NP$ -complete. 1SAT (all clauses have size 1) is also  $NP$ -complete.

## Solution

\*\*\*\*\* SOLUTION GOES HERE \*\*\*\*\*

## Problem 4

### Solution

\*\*\*\*\* SOLUTION GOES HERE \*\*\*\*\*

## Problem 5

### Solution

\*\*\*\*\* SOLUTION GOES HERE \*\*\*\*\*

## Problem 6

### Solution

\*\*\*\*\* SOLUTION GOES HERE \*\*\*\*\*