

Zippers and Derivatives

Shuwei Hu

Supervised by: Manuel Eberl

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Outline

Overview

Zipper Examples

Context as a Derivative

Combinatorial Species

Conclusion

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- ▶ Zipper = context type, which helps moving through and “modifying” a functional data structure [1]

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- ▶ Deriving context types vs. differentiating real-valued functions [2]
- ▶ Relation with combinatorial species [3]

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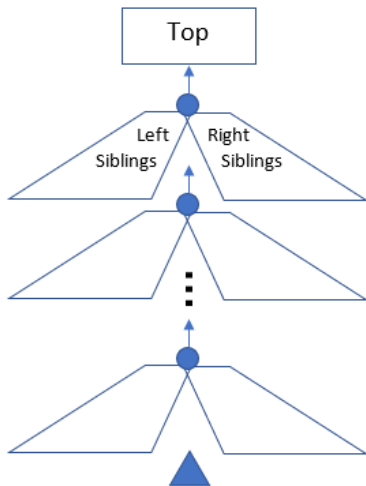
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Zipper!<sup>1</sup>



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<sup>1</sup>Image source: [www.pacifictrimming.com](http://www.pacifictrimming.com)

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- ▶ Math-ly notation, e.g.  $\mathbf{list}(a) = 1 + a \times \mathbf{list}(a)$

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- ▶ Context for an arbitrary algebraic data type?

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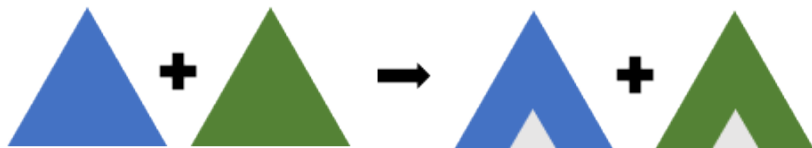


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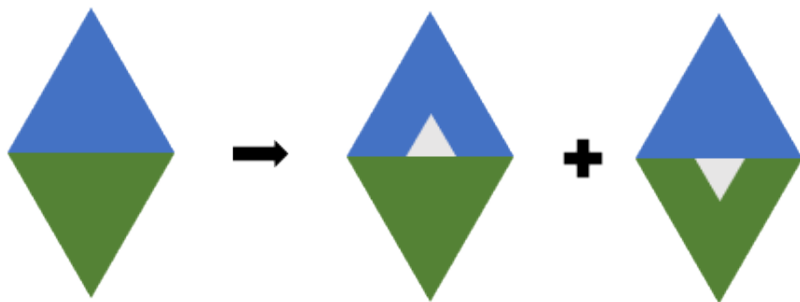
## Context of Sum Type

- ▶ Inside  $T_1 + T_2$ , type  $a$  occurs in either of them
  - ▶  $C[a](T_1 + T_2) = C[a](T_1) + C[a](T_2)$



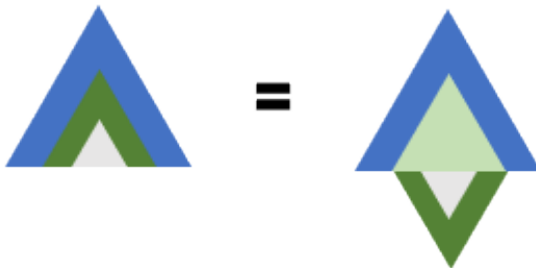
## Context of Product Type

- ▶ Inside  $T_1 \times T_2$ , type  $a$  occurs in one of them, while the other must be carried in the context
  - ▶  $C[a](T_1 \times T_2) = C[a](T_1) \times T_2 + T_1 \times C[a](T_2)$  content



## Context of Composed Type

- ▶ Inside composed type  $T(U(a))$ , type  $a$  occurs in one of  $U$ , which resides somewhere in  $T$ 
  - ▶  $C[a](T(U(a))) = C[b](T(b))|_{b=U(a)} \times C[a](U(a))$



# Context as Derivative

- ▶ Rules for context:

- ▶  $C[a](1) = 0$
- ▶  $C[a](a) = 1$
- ▶  $C[a](T_1 + T_2) = C[a](T_1) + C[a](T_2)$
- ▶  $C[a](T_1 \times T_2) = C[a](T_1) \times T_2 + T_1 \times C[a](T_2)$
- ▶  $C[a](T(U(a))) = C[b](T(b))|_{b=U(a)} \times C[a](U(a))$

- ▶ Rules for derivative:

- ▶  $\frac{\partial}{\partial x} c = 0$
  - ▶  $\frac{\partial}{\partial x} x = 1$
  - ▶  $\frac{\partial}{\partial x} (f + g) = \frac{\partial}{\partial x} f + \frac{\partial}{\partial x} g$
  - ▶  $\frac{\partial}{\partial x} (f \cdot g) = \frac{\partial}{\partial x} f \cdot g + f \cdot \frac{\partial}{\partial x} g$
  - ▶  $\frac{\partial}{\partial x} (f \circ g)|_{x_0} = \frac{\partial}{\partial u} f|_{u=g(x_0)} \cdot \frac{\partial}{\partial x} g|_{x=x_0}$
- ▶  $\frac{\partial}{\partial a}(T) \triangleq C[a](T)$

## Context of List, Revisit

$$\frac{\partial}{\partial a} \mathbf{list}(a) = \frac{\partial}{\partial a} (1 + a \times \mathbf{list}(a))$$

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$$\begin{aligned}\frac{\partial}{\partial \mathbf{a}} \mathbf{list}(a) &= \frac{\partial}{\partial \mathbf{a}} (1 + a \times \mathbf{list}(a)) \\ &= 0 + \frac{\partial}{\partial \mathbf{a}} (a \times \mathbf{list}(a))\end{aligned}$$

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## Binary Tree Context Revisit

$$\frac{\partial}{\partial a} \mathbf{btree}(a) = \frac{\partial}{\partial a} (1 + a \times \mathbf{btree}^2(a))$$

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$$\frac{\partial}{\partial \mathbf{a}} \mathbf{tree}(a) = \frac{\partial}{\partial \mathbf{a}} (1 + a \times \mathbf{list}(\mathbf{tree}(a)))$$

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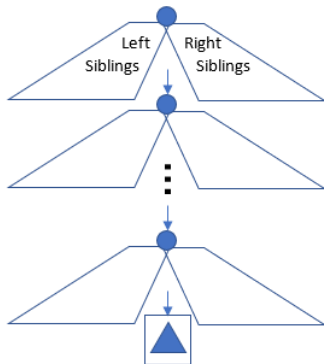
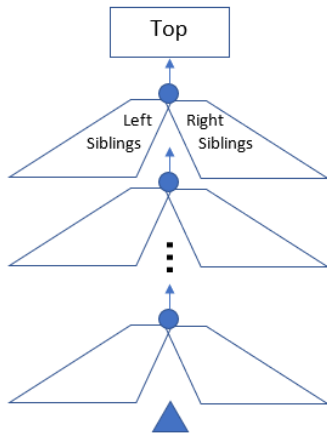
## Context of Tree, Revisit

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# Huet's Zipper, Revisit

- ▶  $\mathbf{ltree}(a) = a + \mathbf{list}(\mathbf{ltree}(a))$
- ▶ Differentiating against non-basic type is a bit tricky
- ▶  $\frac{\partial}{\partial \mathbf{ltree}}(\mathbf{ltree}) = 1?$
- ▶  $\frac{\partial}{\partial \mathbf{ltree}}(\mathbf{ltree}) = \frac{\partial}{\partial \mathbf{ltree}}(a + \mathbf{list}(\mathbf{ltree}))?$
- ▶ A hack: 
$$\begin{aligned}\frac{\partial}{\partial \mathbf{ltree}} &= ((1)) + \frac{\partial}{\partial \mathbf{ltree}}(a + \mathbf{list}(\mathbf{ltree})) \\ &= 1 + \frac{\partial}{\partial \mathbf{ltree}}(\mathbf{list}(\mathbf{ltree})) \\ &= 1 + \mathbf{list}^2(\mathbf{ltree}) \times \frac{\partial}{\partial \mathbf{ltree}}(\mathbf{ltree}) \\ &= \mathbf{ltree\_context}\end{aligned}$$

- ▶  $\frac{\partial}{\partial a} \mathbf{tree}(a) = \mathbf{list}(\mathbf{tree}(a)) \times \mathbf{list}(a \times \mathbf{list}^2(\mathbf{tree}(a)))$
- ▶ Reversed interpretation of the recursion path:





## Subtraction and Division?

►  $\mathbf{list}(a) = 1 + a \times \mathbf{list}(a)$

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- ▶  $\mathbf{list}(a) = 1 + a \times \mathbf{list}(a)$
- ▶  $\mathbf{list}(a) \leftrightarrow \frac{1}{1-a}?$

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- ▶  $\mathbf{list}(a) \leftrightarrow \frac{1}{1-a}?$
- ▶  $\frac{\partial}{\partial a} \mathbf{list}(a) \leftrightarrow \frac{\partial}{\partial a} \left( \frac{1}{1-a} \right) = \frac{1}{(1-a)^2} \leftrightarrow \mathbf{list}^2(a) \text{ ?!}$

# Outline

Overview

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**Combinatorial Species**

Conclusion

# Species

- ▶ General definition: endofunctor on the category of finite sets

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  - ▶ ...
- ▶ Species, as functors, preserve identity arrow and composition of arrows
  - ▶ should be obvious for *regular* species

## Regular Species

- Composition of  $0$ ,  $1$ ,  $X$ ,  $+$ ,  $\cdot$  and least fix-point



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- ▶  $F + G$ :  $S \rightarrow F(S) \sqcup G(S)$

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- ▶  $F \cdot G: S \rightarrow \bigcup_{S_1 \oplus S_2 = L} (F(S_1) \times G(S_2))$

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- ▶  $n \triangleq \underbrace{1 + (1 + (\dots + 1))}_n: \{\} \rightarrow \{\{\}_1, \{\}_2, \dots, \{\}_n\}$

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- ▶  $L \triangleq 1 + X \cdot L = 1 + X \cdot (1 + X \cdot (\dots)) \simeq 1 + X + X^2 + \dots$



# Exponential Generating Function

- ▶ Counting function:  $C_F : |L| \rightarrow |F(L)|$  for species  $F$
- ▶ Exponential Generating Function:  
$$E_F(x) = C_F(0) + C_F(1) * \frac{x}{1!} + C_F(2) * \frac{x^2}{2!} + \dots$$
- ▶  $E_0(x) = 0$
- ▶  $E_1(x) = 1$
- ▶ 
$$E_{F+G}(x) = (C_F(0) + C_G(0)) + (C_F(1) + C_G(1)) * \frac{x}{1!} + \dots$$
$$= E_F(x) + E_G(x)$$
- ▶ 
$$C_{F \cdot G}(i) * \frac{x^i}{i!} = \sum_{j=0}^i \frac{x^j}{j!} * \binom{i}{j} * C_F(j) * C_G(i-j)$$
$$= \sum_{j=0}^i \frac{x^j}{j!} * C_F(j) * \frac{x^{i-j}}{(i-j)!} * C_G(i-j)$$
- ▶  $E_{F \cdot G}(x) = E_F(x) * E_G(x)$
- ▶  $E_F$  and  $F$  share the same expression!

# Derivative of Regular Species

- ▶  $F' : S \rightarrow F(S \cup \{\square\})$ 
  - ▶ e.g.  $(X^2)' : \{1\} \rightarrow X^2(\{1, \square\}) = \{(1, \square), (\square, 1)\}$
- ▶ Derivative rules apply:
  - ▶  $(F + G)' = F' + G'$
  - ▶  $(F \cdot G)' = F' \cdot G + F \cdot G'$
  - ▶ ...
- ▶  $(X^n)' = n * X^{n-1}$
- ▶  $E_{(X^n)'}(x) = E_{n * X^{n-1}}(x) = n * x^{n-1} = (x^n)' = (E_{X^n}(x))'$
- ▶  $E_{F'} = (E_F)'$
- ▶ Derivative preserves the consistency between regular species and its counting function!

## List, Revisit

- ▶  $L = 1 + X \cdot L$

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- ▶  $L = \frac{1}{1-X} = ?$

## List, Revisit

- ▶  $L = 1 + X \cdot L$
- ▶  $E_L(x) = 1 + x \cdot E_L(x)$
- ▶  $E_L(x) = \frac{1}{1-x}$
- ▶  $L = \frac{1}{1-X} = (\text{THE } F. E_F(x) = \frac{1}{1-x})$



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- ▶  $L = \frac{1}{1-X} = (\mathbf{THE} \ F. \ E_F(x) = \frac{1}{1-x})$ 
  - ▶  $E_F = E_G \implies F \simeq G$

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# Conclusion

- ▶ “Functional pointer”: context type
- ▶ The structure of context of ordered tree resembles a zipper
- ▶ Differentiating an algebraic datatype
- ▶ Combinatorial species and its EGF
- ▶ Regular species vs. algebraic datatype?

Thank you for listening!

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