

1. After PC-1 the swapped bit is at position 8

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|--------------------------|-----------------------|-------|
| a. Round 1: position 7 | PC-2: position 20 | S4 |
| b. Round 2: position 6 | PC-2: position 10 | S2 |
| c. Round 3: position 4 | PC-2: position 16 | S3 |
| d. Round 4: position 2 | PC-2: position 24 | S4 |
| e. Round 5: position 28 | PC-2: position 8 | S2 |
| f. Round 6: position 26 | PC-2: position 17 | S3 |
| g. Round 7: position 24 | PC-2: position 4 | S1 |
| h. Round 8: position 22 | PC-2: position unused | Snone |
| i. Round 9: position 21 | PC-2: position 11 | S2 |
| j. Round 10: position 19 | PC-2: position 14 | S3 |
| k. Round 11: position 17 | PC-2: position 2 | S1 |
| l. Round 12: position 15 | PC-2: position 9 | S2 |
| m. Round 13: position 13 | PC-2: position 23 | S4 |
| n. Round 14: position 11 | PC-2: position 3 | S1 |
| o. Round 15: position 9 | PC-2: position unused | Snone |
| p. Round 16: position 8 | PC-2: position 18 | S3 |
- q. For decryption the order would be reversed, so decryption round1 would have S3 affected, round 2 none, round 3 S1 and so on.
2. Given that decryption uses the same key schedule except in reverse, and that L_0^d/R_0^d during decryption are equal to R_{16}/L_{16} : the final encryption creating $R_{16} = L_{15} * \text{key16}$ and $L_{16} = R_{15}$ would clearly be the same as the first step in decryption, where $\text{key1}^d = \text{key16}$, so since the key is used in a reversible operation, $R_1^d = L_0^d * \text{key1}^d = R_{16} * \text{key16} = L_{15}$, and $L_1^d = R_0^d = R_{15}$

3.

a. Inverse of 25mod1033 is 124mod1033

i	q	r	s	t
0		1033	1	0
1	41	25	0	1
2	3	8	1	-41
3	8	1	-3	124
4		0	24	-1033

b. Inverse of $25 \bmod 1034$ is $455 \bmod 1034$

i	q	r	s	t
0		1034	1	0
1	41	25	0	1
2	2	9	1	-41
3	1	7	-2	83
4	3	2	3	-124
5		1	-11	455

c. $25 \bmod 1035$ has no inverse

i	q	r	s	t
0		1035	1	0
1	41	25	0	1
2	2	10	1	-41
3	2	5	-2	83
4		0	5	-207

d. Inverse of $25 \bmod 1036$ is $373 \bmod 1036$

i	q	r	s	t
0		1036	1	0
1	41	25	0	1
2	2	11	1	-41
3	3	3	-2	83
4	1	2	7	-290
5		1	-9	373

4. Python code to use the EEA to find the multiplicative inverse of $123456 \bmod 7111111$:

```
s = [1, 0]
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```
t = [0, 1]
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```
i = 1
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```
r = [7111111, 123456]
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```
q = [None]
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while(True):
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```
i+=1
r.append(r[i-2]%r[i-1])
q.append((r[i-2] - r[i])/r[i-1])
s.append(s[i-2] - q[i-1]*s[i-1])
t.append(t[i-2] - q[i-1]*t[i-1])
if(r[i-1] is 0):
    if r[i-2] is 1:
        print("multiplicative inverse of %s mod %s is %s"%(r[i-1], r[0], t[i-2]))
    else:
        print("no multiplicative inverse exists for %s mod %s"%(r[i-1], r[0]))
    break
```