
Functional Programming

<https://proglang.informatik.uni-freiburg.de/teaching/functional-programming/2022/>

Exam Sheet

March 13-15, 2023

How to do this exam

- You can earn 50 points from the questions on this sheet.
- You can earn 150 points from the programming task.
- You need ≥ 100 points to pass the exam and you must earn ≥ 20 points from this sheet. Otherwise you are free to choose.
- Do not change the names and types of the functions as given on this sheet and in the template files supplied.

Question 1 (Terms, 20 points, file: Question1.hs)

In the lecture, we defined signatures and terms. Consider the following datatypes for representing signatures where function symbols are identified by `Ident` and terms over an empty set of variables.

```
import qualified Data.Map as M
type Ident      = Char
data Term       = Term { symbol :: Ident, subterms :: [Term] }
type Signature = M.Map Ident Int
```

(a) `parseTerm :: Parser Char Term`

Write a parser that reads a string as a “raw term” according to the grammar of terms (derived from symbol T , without checking the correct use of symbols). Terminal symbols in the grammar are either literal strings in single quotes (e.g. `' , '`) or regular expressions in double quotes (e.g. `"[a-zA-Z]"`, which denotes an ASCII letter).

$$T ::= "[a-zA-Z]" \mid '(' A ')' \mid A ::= \varepsilon \mid T B \quad B ::= \varepsilon \mid ', ' T B$$

(b) `type Pos = [Int]`
`arityCheck :: Signature -> Term -> Either Pos Term`

Write an arity checker that checks that a raw term is well-formed according to a given signature. The implementation of the arity checker should use the `Either` monad to report errors in terms of a list of numbers that describes the location of the first error: `[]` for an error at the root of the term, `[0]` for the root of the first subterm, `[0,1]` for the root of the second subterm of the first subterm, and so on.

```
>> sig = M.fromList [('x', 0)]
>> arityCheck sig (Term 'x' [])
Right (Term 'x' [])
>> arityCheck sig (Term 'x' [Term 'x' []])
Left []
```

Question 2 (Substitution and reduction, 30 points, file: Question2.hs)

Consider the following datatypes for representing expressions in lambda calculus.

```
import qualified Data.Set as S
type Ident = String
data Expr
  = Var Ident
  | Lam Ident Expr
  | App Expr Expr
```

(a) `free :: Expr -> S.Set Ident`

Define a function to calculate the set of free variables of a lambda expression (using the `Set` type from `Data.Set`).

(b) `fresh :: S.Set Ident -> Ident`

Define a function to return a string that is not a member of a given set of strings. The function does not have to be efficient, but it must terminate for all inputs.

(c) `subst :: Expr -> Ident -> Expr -> Expr`

Define capture-avoiding substitution for arbitrary lambda expressions. `subst t' x t` substitutes `t'` for `x` in `t`.

```
>> ex0 = Lam "y" (App (Var "x") (Var "y"))
>> ex0' = Var "y"
>> subst ex0' "x" ex0
Lam "x0" (App (Var "y") (Var "x0"))
```

(d) `tryBeta :: Expr -> Maybe Expr`

Define a function that locates the leftmost beta redex in a lambda expression and reduces it. To find the leftmost redex, traverse the expression in preorder. As no such redex may exist, you should write this function using the `Maybe` applicative.

```
>> ex1 = App (Lam "x" (Var "x")) (Var "y")
>> ex2 = App (Var "x") ex1
>> ex3 = Lam "x" ex2
>> tryBeta ex1
Just (Var "y")
>> tryBeta ex2
Just (App (Var "x") (Var "y"))
>> tryBeta ex3
Just (Lam "x" (App (Var "x") (Var "y")))
```