GameCube DSP User's Manual

Reverse-engineered and documented by Duddie $\frac{\rm duddie@walla.com}{\rm duddie}$

 $\begin{array}{c} \text{September 27, 2024} \\ \text{v0.1.5} \end{array}$

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Contents

Co	onter	nts	3
Di	sclai	imer	7
Ve	ersio	n History	8
\mathbf{G}	NU I	Documentation License 1.2	9
1	Ove	erview	13
	1.1	DSP Memory Map	13
		1.1.1 Instruction Memory	13
		1.1.2 Data Memory	14
	1.2	Initialization	15
2	Reg		16
	2.1	Register names	16
	2.2	Accumulators	17
			17
			17
			17
		2.2.4 16-bit and 40-bit modes	17
	2.3		18
	2.4		19
	2.5		20
	2.6	Product register	21
3	Exc		22
	3.1		22
	3.2	Exception vectors	23
4	Har	rdware interface	24
	4.1	Hardware registers	24
	4.2	DMA	25
	4.3	Accelerator	26
	4.4	Interrupts	29
	4.5		30
			30
		4.5.2 CPU Mailbox	31
5	Opc		32
	5.1	1 0	32
	5.2	±	33
	5.3		34
	5.4	1	35
	5.5	9	36
	5.6		37
			38
		5.6.2 ADD	39

5.6.3	ADDARN	40
5.6.4	ADDAX	41
5.6.5	ADDAXL	42
5.6.6	ADDI	43
5.6.7	ADDIS	
5.6.8	ADDP	
5.6.9	ADDPAXZ	
0.0.0	ADDR	
	ANDC	
	ANDCF	
	ANDER	
	ANDI	
	ANDR	
	ASL	
	ASR	
	ASRN	
	ASRNR	
	ASRNRX	
	ASR16	
	BLOOP	
	BLOOPI	
	CALL	
5.6.25	CALLcc	62
5.6.26	CALLR	63
5.6.27	CALLRcc	64
	CLR15	
5.6.29	CLR	66
5.6.30	CLRL	67
5.6.31	CLRP	68
5.6.32	CMP	69
5.6.33	CMPAXH	70
5.6.34	CMPI	71
5.6.35	CMPIS	72
5.6.36	DAR	73
5.6.37	DEC	74
5.6.38	DECM	75
5.6.39	HALT	
5.6.40		77
	IFcc	
	ILRR	
	ILRRD	
	ILRRI	
	ILRRN	
	INC	
	INCM	
	JMP	
	Jcc	
0.00	JMPR	
0.0.00	JRcc	
	LOOP	
	LOOPI	
	LR	
	LRI	
	LRIS	
	LRR	
	LRRD	
	LRRI	
5.6.60	LRRN	97

5.6.61																				-
5.6.62																				
5.6.63	LSL16		 									 				 			. 10	0
5.6.64	LSR		 									 							. 10	1
5.6.65	LSRN		 									 				 			10	2
5.6.66	LSRNR		 									 				 			. 10	13
5.6.67	LSRNRX .		 									 				 			. 10	14
	LSR16																			
	M0																			
	M2																			
5.6.71	MADD																			
5.6.72																			10	
J.J			 		 •	 •		•	 •	 •	 •	 	•	•	 •	 	•	•	11	_
	MOV		 		 -	 -		-	 -	 -	 -	 	-	-	 -		-	-		
			 		 -	 -		-	 -	 -	 -	 	-	-	 -	 	-	-		_
5.6.75	MOVAX																			
5.6.76	MOVNP																			
5.6.77	MOVP																			
	MOVPZ																			
5.6.79	MOVR		 		 •	 -	• •	•	 •	 •	 •	 	•	•	 •	 •	•	•		
5.6.80	$MRR \dots$		 		 •	 •		•	 •	 •	 •	 	•	•	 •	 	•	•		•
5.6.81	MSUB		 									 							. 11	8
5.6.82	MSUBC		 									 							. 11	9
5.6.83	MSUBX		 									 							. 12	0
5.6.84	MUL		 									 							. 12	1
5.6.85	MULAC		 									 							. 12	2
5.6.86	MULAXH .		 									 							. 12	3
5.6.87	MULC		 									 				 			. 12	4
5.6.88	MULCAC .		 									 							. 12	5
5.6.89	MULCMV		 									 				 			. 12	6
5.6.90	MULCMVZ		 									 				 			. 12	7
5.6.91	MULMV .		 									 				 			. 12	8
5.6.92	MULMVZ .		 									 							. 12	9
5.6.93	MULX		 									 							13	0
5.6.94	MULXAC .		 									 				 			. 13	1
5.6.95	MULXMV		 									 				 			. 13	2
5.6.96	MULXMVZ																			
5.6.97	NEG																			-
5.6.98	_	•														·			13	
5 6 99	NOP	•	 	•	 •	 •		٠	 ·		 •	 	•			 ·	-	-		_
	NX																		_	-
	ORC																			
	ORI																			-
	ORR																			-
	RET																			-
	RETcc																			
5.6.106																				
0.000	'RTIcc																			-
	SSBCLR																			-
	SBSET																			
	SET15																			
	SET16																			
	2 SET40																			
5.6.113																				
	SR																			
	SRR																			
	SSRRD																			
	SRRI																			
5.6.118	SRRN		 									 				 			15	5

	5.6.119	SRS				 	 	 	 	 			. 156
	5.6.120	SRSH				 	 	 	 	 			. 157
	5.6.121	SUB				 	 	 	 	 			. 158
	5.6.122	SUBARN .				 	 	 	 	 			. 159
	5.6.123	SUBAX				 	 	 	 	 			. 160
	5.6.124	SUBP				 	 	 	 	 			. 161
	5.6.125	SUBR				 	 	 	 	 			. 162
	5.6.126	TST				 	 	 	 	 			. 163
	5.6.127	TSTAXH .				 	 	 	 	 			. 164
	5.6.128	TSTPROD				 	 	 	 	 			. 165
	5.6.129	XORC				 	 	 	 	 			. 166
	5.6.130	XORI				 	 	 	 	 			. 167
	5.6.131	XORR				 	 	 	 	 			. 168
5.7	Extend	led opcodes				 	 	 	 	 	 		. 169
5.8	Alphal	petical list of e	extende	ed op	codes	 	 	 	 	 	 		. 170
	5.8.1	'DR				 	 	 	 	 			. 171
	5.8.2	'IR				 	 	 	 	 			. 172
	5.8.3	${\rm 'L} \ldots \ldots$. 173
	5.8.4	'LN				 	 	 	 	 			. 174
	5.8.5	${\rm 'LD} \; . \; . \; . \; .$. 175
	5.8.6	'LDM											
	5.8.7	'LDNM				 	 	 	 	 			. 177
	5.8.8	'LDN											
	5.8.9	'LDAX											
		'LDAXM .											
		'LDAXNM											
	5.8.12	'LDAXN				 	 	 	 	 			. 182
		'LS											
		'LSM											
		'LSNM											
		'LSN											
		'MV											
		'NOP											
		'NR											
		'S											
		'SL											
		'SLM											
		'SLNM											
		'SLN											
		'SN											
5.9	Instruc	tions sorted b	v once	ode									196

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The purpose of this documentation is purely academic and it aims at understanding described hardware. It is based on academic reverse engineering of hardware.

Version History

Version	Date	Author	Change
0.0.1	2005.05.08	Duddie	Initial release
0.0.2	2005.05.09	Duddie	Added \$prod and \$config registers, table of opcodes, dis-
			claimer.
0.0.3	2005.05.09	Duddie	Fixed BLOOP and BLOOPI and added description of the
			loop stack.
0.0.4	2005.05.12	Duddie	Added preliminary DSP memory map and opcode syntax.
0.0.5	2018.04.09	Lioncache	Converted document over to LaTeX.
0.0.6	2018.04.13	BhaaL	Updated register tables, fixed opcode operations
0.0.7	Mid 2020	Tilka	Fixed typos and register names, and improved readability.
0.1.0	2021.08.21	Pokechu22	Added missing instructions, improved documentation of
			hardware registers, documented additional behaviors, and
			improved formatting.
0.1.1	2022.05.14	xperia64	Added tested DSP bootloading transfer size
0.1.2	2022.05.21	Pokechu22	Fixed "ILLR" typo in Instruction Memory section
0.1.3	2022.05.27	Pokechu22	Renamed CMPAR instruction to CMPAXH
0.1.4	2022.06.02	Pokechu22	Fixed typos; added sections on 16-bit and 40-bit modes and
			on main and extended opcode writing to the same register.
0.1.5	2022.09.29	vpelletier	Fixed BLOOP and BLOOPI suboperation order

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Chapter 1

Overview

1.1 DSP Memory Map

The DSP has two address spaces, one for data and one for instructions. The DSP accesses memory in words, so all addresses refer to words. A DSP word is 16 bits in size.

1.1.1 Instruction Memory

Instruction Memory (IMEM) is divided into instruction RAM (IRAM) and instruction ROM (IROM).

Exception vectors are located at the top of the RAM and occupy the first 16 words, with 2 words available for each exception (enough for a JMP instruction for each exception).

There are no DSP instructions that write to IMEM; however, the ILRR family of instructions can read from it. This is sometimes used for jump tables or indexing into a list of pointers (which may point into either IMEM or DMEM).

0x0000	IRAM
0x0FFF	
0x8000	IROM
0x8FFF	

1.1.2 Data Memory

Data Memory (DMEM) is divided into data RAM (DRAM) and resampling coefficient data (COEF). Hardware registers (IFX) are also mapped into this space.

It is possible to both read and write to DMEM, but coefficient data cannot be written to.

0x0000	DRAM
0x1000 0x17FF	COEF
0xFFF00	IFX

1.2 Initialization

The DSP is initialized before it is used. This is done by copying a small program to physical address 0x01000000 (virtual 0x81000000) in GameCube/Wii main memory, and then writing to DSP_CONTROL_STATUS at 0xCC00500A with the 11th and 0th bits set (SDK titles write 0x08ad). The 11th bit being set appears to cause data from 0x01000000 to be DMAd to the start of IMEM; a basic hardware test (sending increasingly larger payloads until they fail) indicates 1024 bytes of data (512 DSP words) are transferred. (None of this has been extensively hardware tested, and is instead based on libogc's __dsp_bootstrap.)

The program that SDK titles send does the following:

- 1. Reads all 0x1000 words of IROM from 0x8000 through 0x8FFF (using ILRRI)
- 2. Writes zero to all 0x1000 words of DRAM from 0x0000 through 0x0FFF (using SRRI)
- 3. Reads all 0x0800 words of COEF data from 0x1000 through 0x17FF (using LRRI)
- 4. Waits for the top bit of DMBH to be clear (indicating the CPU is ready to receive mail)
- 5. Writes 0x0054 to DMBH and 0x4348 to DMBL, sending 0x00543448 ("TCH"?) to the CPU. The CPU does not check for this value, but it does wait for mail to be sent.

It is not clear why this is done, as the values read from IROM and COEF are not used; perhaps it works around a hardware bug where incorrect values are read from ROM initially.

Chapter 2

Registers

2.1 Register names

The DSP has 32 16-bit registers, although their individual purpose and their function differ from register to register.

\$0	\$r00	\$ar0	Addressing register 0
\$1	\$r01	\$ar1	Addressing register 1
\$2	\$r02	\$ar2	Addressing register 2
\$3	\$r03	\$ar3	Addressing register 3
\$4	\$r04	\$ix0	Indexing register 0
\$5	\$r05	\$ix1	Indexing register 1
\$6	\$r06	\$ix2	Indexing register 2
\$7	\$r07	\$ix3	Indexing register 3
\$8	\$r08	\$wr0	Wrapping register 0
\$9	\$r09	\$wr1	Wrapping register 1
\$10	\$r0A	\$wr2	Wrapping register 2
\$11	\$r0B	\$wr3	Wrapping register 3
\$12	\$r0C	\$st0	Call stack register
\$13	\$rOD	\$st1	Data stack register
\$14	\$r0E	\$st2	Loop address stack register
\$15	\$r0F	\$st3	Loop counter register
\$16	\$r10	\$ac0.h	40-bit Accumulator 0 (high)
\$17	\$r11	\$ac1.h	40-bit Accumulator 1 (high)
\$18	\$r12	\$config	Config register
\$19	\$r13	\$sr	Status register
\$20	\$r14	\$prod.1	Product register (low)
\$21	\$r15	\$prod.m1	Product register (mid 1)
\$22	\$r16	\$prod.h	Product register (high)
\$23	\$r17	\$prod.m2	Product register (mid 2)
\$24	\$r18	\$ax0.1	32-bit Accumulator 0 (low)
\$25	\$r19	\$ax1.1	32-bit Accumulator 1 (low)
\$26	\$r1A	\$ax0.h	32-bit Accumulator 0 (high)
\$27	\$r1B	\$ax1.h	32-bit Accumulator 1 (high)
\$28	\$r1C	\$ac0.1	40-bit Accumulator 0 (low)
\$29	\$r1D	\$ac1.1	40-bit Accumulator 1 (low)
\$30	\$r1E	\$ac0.m	40-bit Accumulator 0 (mid)
\$31	\$r1F	\$ac1.m	40-bit Accumulator 1 (mid)

2.2 Accumulators

The DSP has two long 40-bit accumulators (\$acX) and their short 24-bit forms (\$acsX) that reflect the upper part of 40-bit accumulator. There are additional two 32-bit accumulators (\$axX).

The high parts of the 40-bit accumulators (acX.h) are sign-extended 8-bit registers. Writes to the upper 8 bits are ignored, and the upper 8 bits read the same as the 7th bit. For instance, 0x007F reads back as 0x007F, but 0x0080 reads back as 0xFF80.

2.2.1 Accumulators \$acX

40-bit accumulator \$acX (\$acX.hml) consists of registers:

$$acX = acX.h << 32 | acX.m << 16 | acX.l$$

2.2.2 Short accumulators \$acsX

24-bit accumulator \$acsX (\$acX.hm) consists of the upper 24 bits of accumulator \$acX.

$$acs X = ac X \cdot h << 16 \mid ac X \cdot m$$

2.2.3 Additional accumulators \$axX

32-bit accumulators \$axX (\$axX.hl) consist of registers:

$$xX = xX.h << 16 | xX.l$$

2.2.4 16-bit and 40-bit modes

Depending on the value of \$sr.SXM (bit 14), loading to \$acX.m may also update \$acX.h and \$acX.l, and stores from \$acX.m may experience saturation based on \$acX.h. Regardless of the value of \$sr.SXM, arithmetic operations such as ADDI, INCM, MOVR, and LSRN will still affect the entire accumulator.

If \$sr.SXM is set to 0, then 16-bit mode (SET16) is in use. Loads to \$acX.m will only change \$acX.m, and storing \$acX.m will use the value directly contained in \$acX.m; the same applies to loads to and stores from \$acX.h or \$acX.l or any other register.

If \$sr.SXM is set to 1, then 40-bit mode (SET40) is in use. Loads to \$acX.m will set \$acX.1 to 0 and will sign-extend into \$acX.h (setting it to 0xFF if the sign bit is set (\$acX.m & 0x8000 != 0), and to 0 otherwise). This means that in 40-bit mode, loads to \$acX.m are effectively loads to the whole accumulator \$acX. Loads to \$acX.h and \$acX.1 do not have this special behavior; they only modify the specified register (as in 16-bit mode).

Additionally, if \$sr.SXM is set to 1, then moving or storing from \$acX.m may instead result in 0x7fff or 0x8000 being used. This happens if \$acX.hml is not the same as sign-extending \$acX.ml; 0x7fff is used if \$acX is positive and 0x8000 is used if \$acX is negative.

The conditions for this saturation are the same as the conditions for \$sr.AS (bit 4, above s32) to be set when flags are updated. (This does not mean that the saturation happens if and only if \$sr.AS is set, as the flags might have been set after an operation on a different register.)

The following instructions perform sign-extension when writing to \$acX.m: ILRR, ILRRD, ILRRI, and ILRRN; LR; LRI; LRIS; LRR, LRRD, LRRI, and LRRN; LRS; MRR; and 'L and 'LN.

The following instructions experience saturation when reading from \$acX.m: BLOOP; LOOP; MRR; SR; SRR, SRRD, SRRI, and SRRN; SRS; 'LS, 'LSM, 'LSM, and 'LSNM; 'MV; 'SL, 'SLM, 'SLN, and 'SLNM; and 'S and 'SN

2.3 Stacks

The GameCube DSP contains four stack registers:

- \$st0 Call stack register
- \$st1 Data stack register
- \$st2 Loop address stack register
- \$st3 Loop counter register

Stacks are implemented in hardware and have limited depth. The data stack is limited to four values and the call stack is limited to eight values. The loop stack is limited to four values. Upon underflow or overflow of any of the stack registers exception STOVF is raised.

The loop stack is used to control execution of repeated blocks of instructions. Whenever there is a value in \$st2 and the current PC is equal to the value in \$st2, then the value in \$st3 is decremented. If the value is not zero, then the PC is modified by the value from call stack \$st0. Otherwise values from the call stack \$st0 and both loop stacks, \$st2 and \$st3, are popped and execution continues at the next opcode.

2.4 Config register

Serves as a base offset for SRS, SRSH, and LRS. Zelda uCode writes it with 0x0004, but otherwise it is usually 0x00FF.

This is an 8-bit register. Writes to the upper 8 bits are ignored and those bits always read back as 0.

2.5 Status register

Status register \$sr reflects flags computed on accumulators after logical or arithmetic operations. Furthermore, it also contains control bits to configure the flow of certain operations.

Bit	Name	Comment
15	SU	Multiplication operands are signed $(1 = unsigned)$
14	SXM	Sign extension mode $(1 = 40\text{-bit}, \text{ see } 16\text{-bit} \text{ and } 40\text{-bit} \text{ modes})$
13	AM	Product multiply result by 2 (when AM = 0)
12		
11	EIE	External interrupt enable
10		
9	IE	Interrupt enable
8		Unknown, always reads back as 0
7	OS	Overflow (sticky)
6	LZ	Logic zero (used by ANDCF and ANDF)
5	TB	Top two bits are equal
4	AS	Above s32
3	S	Sign
2	Z	Arithmetic zero
1	0	Overflow
0	C	Carry

2.6 Product register

The product register is a register containing the intermediate product of a multiply or multiply and accumulation operation. Its result should never be used for calculation although the register can be read or written. It reflects the state of the internal multiply unit. The product is 40 bits with 1 bit of overflow.

$$prod = (prod.h \ll 32) + ((prod.m1 + prod.m2) \ll 16) + prod.l$$

It needs to be noted that \$prod.m1 + \$prod.m2 overflow bit (bit 16) will be added to \$prod.h.

Bit \$sr.AM affects the result of the multiply unit. If \$sr.AM is equal 0 then the result of every multiply operation will be multiplied by two.

prod.h is 8 bits. The upper 8 bits always read back as 0.

Chapter 3

Exceptions

3.1 Exception processing

Exception processing happens by setting the program counter to different exception vectors. At exception time, the exception program counter is stored at call stack \$st0 and status register \$sr is stored at data stack \$st1.

Operation:

```
PUSH_STACK($st0);
$st0 = $pc;
PUSH_STACK($st1);
$st1 = $sr;
$pc = exception_nr * 2;
```

3.2 Exception vectors

Exception vectors are located at address 0x0000 in Instruction RAM.

Level	Address	Name	Description
0	0x0000	RESET	
1	0x0002	STOVF	Stack under/overflow
2	0x0004		
3	0x0006		
4	0x0008		
5	0x000A	ACCOV	Accelerator address overflow
6	0x000C		
7	0x000E	INT	External interrupt (from CPU)

Chapter 4

Hardware interface

4.1 Hardware registers

Hardware registers (IFX) occupy the address space at \mathtt{OxFFxx} in the Data Memory space. Each register is 16 bits in width.

Address	Name	Description					
ADPCM (Coefficients						
OxFFAO	COEF_A1_0	A1 Coefficient # 0					
0xFFA1	COEF_A2_0	A2 Coefficient # 0					
	ı	:					
OxFFAE	COEF_A1_7	A1 Coefficient # 7					
OxFFAF	COEF_A2_7	A2 Coefficient # 7					
DMA Inte	rface						
0xFFC9	DSCR	DMA control					
0xFFCB	DSBL	Block length					
OxFFCD	DSPA	DSP memory address					
OxFFCE	DSMAH	Memory address H					
OxFFCF	DSMAL	Memory address L					
Accelerator	r	-					
0xFFD1	FORMAT	Accelerator sample format					
0xFFD2	ACUNK1	Unknown, usually 3					
0xFFD3	ACDATA1	Alternative ARAM interface					
0xFFD4	ACSAH	Accelerator start address H					
0xFFD5	ACSAL	Accelerator start address L					
0xFFD6	ACEAH	Accelerator end address H					
0xFFD7	ACEAL	Accelerator end address L					
0xFFD8	ACCAH	Accelerator current address H					
0xFFD9	ACCAL	Accelerator current address L					
OxFFDA	SCALE	ADPCM predictor and scale					
OxFFDB	YN1	ADPCM YN1					
OxFFDC	YN2	ADPCM YN2					
OxFFDD	ACDAT	Accelerator data					
OxFFDE	GAIN	Gain					
OxFFDF	ACUNK2	Unknown, usually 0x0C					
OxFFED	AMDM	ARAM DMA Request Mask					
Interrupts	1						
OxFFFB	DIRQ	IRQ request					
Mailboxes	1						
0xFFFC	DMBH	DSP Mailbox H					
OxFFFD	DMBL	DSP Mailbox L					
OxFFFE	CMBH	CPU Mailbox H					
0xFFFF	CMBL	CPU Mailbox L					
	1						

4.2 DMA

The GameCube DSP is connected to the memory bus through a DMA channel. DMA can be used to transfer data between DSP memory (both instruction and data) and main memory.

0xFFC9	DSCR	DMA Control
		tid

Bit	Name	R/W	Action
2	t	R	Transfer currently in progress if set
1	i	R/W	1 - DMA to/from IMEM 0 - DMA to/from DMEM
0	d	R/W	1 - DMA to CPU from DSP 0 - DMA from CPU to DSP

0xFFCB	DSBL	Block lengt	
dddd	dddd	dddd	dddd

Bit	Name	R/W	Action
15-0	d	W	Length in bytes to transfer. Writing to this register starts a DMA
			transfer.

OxFFCD	DSPA	DSP Address
dddd	dddd	dddd dddd

Bit	Name	R/W	Action
15-0	d	R/W	Bits 15–0 of the DSP memory address

0xFFCE	DSM	IAH	Memo	ry Address H
(dddd	dddd	dddd	dddd

Bit	Name	R/W	Action
15-0	d	R/W	Bits 31–16 of the main memory address

0xFFCF	DSM	ÍAL	Memo	ry Address L
	dddd	dddd	dddd	dddd

Bit	Name	R/W	Action
15-0	d	R/W	Bits 15–0 of the main memory address

4.3 Accelerator

The accelerator is used to transfer data from accelerator memory (ARAM) to DSP memory. The accelerator area can be marked with ACSA (start) and ACEA (end) addresses. Current address for the accelerator can be set or read from the ACCA register. Reading from accelerator memory is done by reading from the ACDAT register. This register contains data from ARAM pointed to by the ACCA register. After reading the data, ACCA is incremented by one. After ACCA grows bigger than the area pointed to by ACEA, it gets reset to a value from ACSA and the ACCOV interrupt is generated.

0xFFD1	FORMAT	Accelerator sample format
	dddd d	lddd dddd dddd

Bit	Name	R/W	Action
			0x00 - ADPCM audio
			0x05 - u8 reads (D3)
15-0	d	R/W	0x06 - u16 reads (D3)
			0x0A - 16-bit PCM audio, u16 writes (D3)
			0x19 - 8-bit PCM audio

0xFFD2	ACUNK	(1 U	nknown	1
dddd	dddd	dddd	dddd	

Bit	Name	R/W	Action
15-0	d	R/W	Usually 3

0xFFD3	ACDATA1	Alternative ARAM interface
	dddd	dddd dddd dddd

Bit	Name	R/W	Action
15-0	d	R/W	Reads from or writes to data pointed to by current accelerator
			address, and then increments the current address. It is unclear whether this respects the start and end addresses.

0xFFD4	ACSAH	Acce	elerato	r Start	Address H
	dddd	dddd	dddd	dddd	

Bit	Name	R/W	Action
15-0	d	R/W	Bits 31–16 of the accelerator start address

0xFFD5	ACSAL	Acce	lerato	r Start	Address L
	dddd	dddd	dddd	dddd	

${f Bit}$	Name	R/W	Action
15-0	d	R/W	Bits 15–0 of the accelerator start address

0xFFD6	ACEAH	Acce	elerato	r End	Address H
	dddd	dddd	dddd	dddd	

Bit	Name	R/W	Action
15-0	d	R/W	Bits 31–16 of the accelerator end address

0xFFD7	ACEAL	Acce	elerato	r End	Address L
	dddd	dddd	dddd	dddd	

Bit	Name	R/W	Action
15-0	d	R/W	Bits 15–0 of the accelerator end address

0xFFD8	ACCAH	Accelerator Current Address H
	dddo	dddd dddd dddd

Bit	Name	R/W	Action
15-0	d	R/W	Bits 31–16 of the accelerator current address

0xFFD9	ACSAH	Accelerator Current Address L
	dddd	dddd dddd dddd

Bit	Name	R/W	Action
15-0	d	R/W	Bits 15–0 of the accelerator current address

OxFFDA	SCALE	ADPCM predictor and scale
		ppp ssss

Bit	Name	R/W	Action
6–4	d	R/W	Used to decide which pair of coefficients to use (COEF_A1_p and
			$\mathtt{COEF_A2_p}$, at $\mathtt{OxFFAO} + 2p$ and $\mathtt{OxFFAO} + 2p + 1$
3-0	s	R/W	The scale to use, as 2^s

0xFFDB	YN1	ADPCM YN1
dddd	dddd	dddd dddd

Bit	Name	R/W	Action
15-0	d	R/W	Last value read by the accelerator, updated to the new value of
			ACDAT when ACDAT is read. Used when calculating ADPCM, but
			updated for all sample formats.

OxFFDC	YN1	ADP	CM YN2
dddd	dddd	dddd	dddd

Bit	Name	R/W	Action
15-0	d	R/W	Second-last value read by the accelerator, updated to the previous
			value of YN1 when ACDAT is read. Used when calculating ADPCM,
			but updated for all sample formats. Writing this value starts the
			accelerator.

OxFFDD	A	CDAT	Acce	lerator	data
dd	dd	dddd	dddd	dddd	

Bit	Name	R/W	Action
15-0	d	R	Reads new data from the accelerator. When there is no data left,
			returns 0.

0xFFI	E G	AIN	Gain
dddd	dddd	dddd	dddd

Bit	Name	R/W	Action	
15-0	d	R/W	Exact behavior unknown	

OxFFDF	ACUNK2 U		Unknown 2	
dddd	dddd	dddd	dddd	

Bit	Name	R/W	Action
15-0	d	R/W	Usually 0x0C

OxFFEF	AMDM	ARAM DMA Request Mask
		m

Bit	Name	R/W	Action	
		דו/ מ	0 - DMA with ARAM unmasked	
0	m	R/W	1 - masked	

4.4 Interrupts

The DSP can raise interrupts at the CPU. Interrupts are usually used to signal that a DSP mailbox has been filled with new data.

0xFFFB	DIRQ	IRQ Request
		I

Bit	Name	R/W	Action
0	I	W	1 - Raise interrupt at CPU

4.5 Mailboxes

4.5.1 DSP Mailbox

The DSP mailbox (DMB) is an interface to send 31 bits of information from the DSP to the CPU.

0xFFFC	DMBH	DSP	Mailbox H
Md	dd dddd.	dddd	dddd

Bit	Name	R/W	Action			
		D	1 - Mailbox has not been received by CPU			
15	15 M R W		0 - Mailbox empty			
			Does not matter. It will be set when DMBL is written to			
14-0	d	W	Bits 30–16 of mail sent from the DSP to the CPU			

OxFFFD		OMBL	DSP	Mailbox L
dd	dd	dddd	dddd	dddd

Bit	Name	R/W	Action
15-0	d	W	Bits 15–0 of mail sent from the DSP to the CPU. Writing to this
			register by the DSP causes the DMBH.M bit to be set, indicating that the mail is ready.

Operation:

Sending mail from the DSP to the CPU can be achieved by writing mail to register DMBH and then to register DMBL in that order. After writing to DMBL, bit DMBH.M will be set, signaling that the mail is ready to be received by the CPU. If the DSP needs to receive a response from the CPU, then it usually waits for the M bit to be cleared after sending a mail. If the DSP does processing when the CPU receives a mail, then it waits for the M bit to be cleared before issuing another mail to the CPU.

4.5.2 CPU Mailbox

The CPU Mailbox (CMB) is a register that allows sending 31 bits of information from the CPU to the DSP.

0xFFF	E C	MBH	CPU	Mailbox H
N	lddd	dddd	dddd	dddd

Bit	Name	R/W	Action
15	М	R	1 - Mailbox contains mail from the CPU0 - Mailbox empty
14-0	d	R	Bits 30–16 of the mail sent from the CPU

0xFF	FF C	CMBL	CPU	Mailbox L
	dddd	dddd	dddd	dddd

Bit	Name	R/W	Action
15-0	d	R	Bits 15–0 of mail sent from the CPU. Reading of this register by
			the DSP causes the CMBH.M bit to be cleared.

Operation:

From the CPU side, software usually checks the M bit of CMBH. It takes action only in the case that this bit is 0. Said action is to write CMBH first and then CMBL. After writing to CMBL, the mail is ready to be received by the DSP.

From the DSP side, the DSP loops by probing the M bit. When this bit is 1, the DSP reads CMBH first and then CMBL. After reading CMBH. M will be cleared.

Chapter 5

Opcodes

5.1 Opcode syntax

Basic opcode syntax:

```
OPC opcode parameters>
```

Above syntax is correct for all opcodes.

EXAMPLES:

JMP 0x0300 CALL loop HALT

Extended syntax:

```
OPC'EXOPC <opcode parameters> : <extended opcode parameters>
```

Above syntax is correct only for arithmetic opcodes, because those can be extended with additional load/store unit behavior.

EXAMPLES:

DECM'L \$acs0 : \$acl.m, @ar0
NX'MV : \$acx1.h, \$ac0.l

5.2 Operation — Used Functions

Functions used for describing opcode operation.

PUSH_STACK(\$stR)

Description:

Pushes value onto given stack referenced by stack register \$stR. Operation moves values down in internal stack.

Operation:

```
stack_stR[stack_ptr_stR++] = $stR;
```

POP_STACK(\$stR)

Description:

Pops value from stack referenced by stack register \$stR. Operation moves values up in internal stack.

Operation:

```
$stR = stack_stR[--stack_ptr_stR];
```

FLAGS(val)

Description:

Calculates flags depending on given value or result of operation and sets corresponding bits in status register \$sr.

EXECUTE_OPCODE(new_pc)

Description:

Executes opcode at the given new_pc address.

5.3 Bit meanings

Opcode decoding uses special naming for bits and their decimal representations to provide easier understanding of bit fields in the opcode.

Binary form	Decimal form	Meaning		
d, dd, ddd, dddd	D	Destination register		
s, ss, sss, ssss	S	Source register		
t, tt, ttt, tttt	T	Source register		
r, rr, rrr, rrrr	R	Register (either source or destination)		
Aaaaa(a)	A, addrA	Address in either instruction or data memory		
xxxx xxxx	Х	Extended opcode		
mmm(m)	M, addrM	Address in memory		
iii(i)	I, Imm	Immediate value		
сссс	СС	Condition (see conditional opcodes)		

5.4 Conditional opcodes

Conditional opcodes are executed only when the condition described by their encoded conditional field has been met. The groups of conditional instructions are: CALLCC, Jcc, IFcc, RETCC, RTICC, JRCC, and CALLRCC.

Bits	СС	Name	Evaluated expression
0b0000	GE	Greater than or equal	\$sr.0 == \$sr.S
0b0001	L	Less than	\$sr.O != \$sr.S
0b0010	G	Greater than	(\$sr.O == \$sr.S) && (\$sr.Z == 0)
0b0011	LE	Less than or equal	(\$sr.0 != \$sr.S) (\$sr.Z != 0)
0b0100	NZ	Not zero	\$sr.Z == 0
0b0101	Z	Zero	\$sr.Z != 0
0b0110	NC	Not carry	\$sr.C == 0
0b0111	С	Carry	\$sr.C != 0
0b1000	x8	Below s32	\$sr.AS == 0
0b1001	x9	Above s32	\$sr.AS != 0
0b1010	хA		((\$sr.AS != 0) (\$sr.TB != 0)) && (\$sr.Z == 0)
0b1011	хB		((\$sr.AS == 0) && (\$sr.TB == 0)) (\$sr.Z != 0)
0b1100	LNZ	Not logic zero	\$sr.LZ == 0
0b1101	LZ	Logic zero	\$sr.LZ != 0
0b1110	0	Overflow	\$sr.0 != 0
0b1111		<always></always>	

Note:

There are two pairs of conditions that work similarly: Z/NZ and LZ/LNZ. Z/NZ pair operates on arithmetic zero flag (arithmetic 0) while LZ/LNZ pair operates on logic zero flag (logic 0). The logic zero flag is only set by ANDCF and ANDF.

5.5 Flags

Most opcodes update flags in the status register (\$sr) based on their result. (Extended opcodes do not update flags.)

Overflow (0) occurs when the result has wrapped around. The expression C=A+B has overflown if A>0 and B>0 but $C\leq 0$ or if A<0 and B<0 but $C\geq 0$. Any instruction that sets the 0 flag will also set the 0S flag; when the 0 flag is set, 0S is also set, but 0S is not cleared when 0 is cleared.

Carry (C) occurs when an arithmetic carry occurs and should be added to the next most significant word. The expression C = A + B generates a carry if A > C. The DSP uses different logic for subtraction: the expression C = A - B generates a carry if $A \ge C$ (so if B = 0, a carry is generated for all A). This is because the DSP uses a carry flag, not a borrow flag.

Each instruction has a table showing what flags it updates, such as this:

Flags:

OS	LZ	TB	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

A "-" indicates that the flag retains its previous value, a "0" indicates that the flag is set to 0, and a "X" indicates that the value of the flag changes depending on what the instruction did.

5.6 Alphabetical list of opcodes

5.6.1 ABS

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1010	d0	01			хх	хx	:		хх	ХX	:

Format:

ABS \$acD

Description:

Sets \$acD to the absolute value of \$acD.

Operation:

C	\overline{S}	LZ	ТВ	S32	S	AZ	О	С
	-	-	X	X	X	X	0	0

5.6.2 ADD

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0100	11	0d	l		хх	ХX	:		хх	ХX	

Format:

Description:

Adds accumulator ac(1-D) to accumulator register acD.

Operation:

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	X	-	X	X	X	X	X	X

5.6.3 ADDARN

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	00)		00	01			SS	dd	l

Format:

${\bf Description:}$

Adds indexing register $\$ to an addressing register $\$

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.4 ADDAX

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0100	10	sd	l		хх	хx	:		хх	ХX	:

Format:

Description:

Adds secondary accumulator $\$ to accumulator register $\$

Operation:

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	X	_	X	X	X	X	X	X

5.6.5 ADDAXL

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
0111	00:	sd			хх	хx	:		хх	ХX	:

Format:

Description:

Adds secondary accumulator \$axS.1 to accumulator register \$acD. \$axS.1 is treated as an unsigned value.

Operation:

```
$acD += $axS.1
FLAGS($acD)
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.6 ADDI

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	001d	0000	0000
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Adds a 16-bit sign-extended immediate to mid accumulator $\c acD.\,hm.$

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.7 ADDIS

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	01	0d	l		ii	ii			ii	ii	

Format:

Description:

Adds an 8-bit sign-extended immediate to mid accumulator accomplex.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.8 ADDP

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
0100	111d			хх	хx	:		хх	хx	

Format:

ADDP \$acD

Description:

Adds the product register to the accumulator register.

Operation:

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	X	_	X	X	X	X	X	X

5.6.9 ADDPAXZ

$15\ 14\ 13\ 12$	11 10	9 8	7	6	5	4	3	2	1	0
1111	10s	sd		хх	ХX	:		ХX	ХX	:

Format:

Description:

Adds secondary accumulator axS to product register and stores result in accumulator register. Low 16-bits of acD (acD.1) are set to 0.

Operation:

```
$acD.hm = $prod.hm + $ax.h
$acD.l = 0
FLAGS($acD)
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.10 ADDR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0100	0s	sċ	l		хх	ХX	:		хх	ХX	:

Format:

ADDR
$$\$acD$$
, $\$(0x18+S)$

Description:

Adds register (0x18+S) to the accumulator acD register.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
X	_	X	X	X	X	X	X

5.6.11 ANDC

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	11	0d	l		0x	ХX	:		хх	ХX	:

Format:

Description:

Logic AND middle part of accumulator \$acD.m with middle part of accumulator \$ax(1-D).m.

Operation:

```
$acD.m &= $ac(1-D).m
FLAGS($acD)
$pc++
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.12 ANDCF

$15\ 14\ 13\ 12$	11 10 9 8	7 6 5 4	3 2 1 0
0000	001d	1100	0000
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Sets the logic zero (LZ) flag in status register \$sr if the result of the logical AND operation involving the mid part of accumulator \$acD.m and the immediate value I is equal to immediate value I. If the logical AND operation does not result in a value equal to I, then the LZ flag is cleared.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	X	-	-	-	-	-	-

5.6.13 ANDF

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	001d	1010	0000
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Sets the logic zero (LZ) flag in status register \$sr if the result of the logic AND operation involving the mid part of accumulator \$acD.m and the immediate value I is equal to zero. If the result is not equal to zero, then the LZ flag is cleared.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	X	-	-	-	-	-	-

5.6.14 ANDI

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	001d	0100	0000
iiii	iiii	iiii	iiii

Format:

Description:

Performs a logical AND with the mid part of accumulator acd.m and the immediate value I.

Operation:

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	-	-	X	X	X	X	0	0

5.6.15 ANDR

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
0011	01	sd	l		0x	ХX	:		ХX	ХX	:

Format:

Description:

Performs a logical AND with the middle part of accumulator \$acD.m and the high part of secondary accumulator, \$axS.h.

Operation:

```
$acD.m &= $axS.h
FLAGS($acD)
$pc++
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

)S	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

$\mathbf{5.6.16} \quad \mathbf{ASL}$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	01	0r	•		10	ii			ii	ii	

Format:

Description:

Arithmetically left shifts the accumulator acR by the amount specified by immediate I.

Operation:

	OS	LZ	ТВ	S32	S	AZ	О	С
Γ	-	-	X	X	X	X	0	0

$\mathbf{5.6.17}\quad\mathbf{ASR}$

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
0001	0101	<u>.</u>		11	ii			ii	ii	

Format:

Description:

Arithmetically right shifts accumulator \$acR specified by the value calculated by negating sign-extended bits 0-6.

Operation:

OS	LZ	ТВ	S32	\mathbf{S}	AZ	О	С
-	-	X	X	X	X	0	0

5.6.18 ASRN

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
0000	001	0		11	01			10	11	

Format:

ASRN

Description:

Arithmetically shifts accumulator \$ac0 either left or right based on \$ac1.m: if bit 6 is set, a right by the amount calculated by negating sign-extended bits 0–5 occurs, while if bit 6 is clear, a left shift occurs by bits 0–5.

Operation:

```
IF ($ac1.m & 64)
    IF ($ac1.m & 63) != 0
        $ac0 >>= (64 - ($ac1.m & 63))
    ENDIF
ELSE
        $ac0 <<= $ac1.m
ENDIF
FLAGS($ac0)
$pc++</pre>
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.19 ASRNR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	11	1d	l		1x	ХX	:		хх	ХX	:

Format:

ASRNR \$acD

Description:

Arithmetically shifts accumulator \$acD either left or right based on \$ac(1-D).m: if bit 6 is set, a right by the amount calculated by negating sign-extended bits 0-5 occurs, while if bit 6 is clear, a left shift occurs by bits 0-5.

Operation:

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.20 **ASRNRX**

15 14 13 12	$11\ 10$	9	8	7	6	5	4	3	2	1	0
0011	10	sd	l		1x	хx			хх	ХX	:

Format:

```
ASRNRX $acD, $axS.h
```

Description:

Arithmetically shifts accumulator \$acD either left or right based on \$axS.h: if bit 6 is set, a right by the amount calculated by negating sign-extended bits 0–5 occurs, while if bit 6 is clear, a left shift occurs by bits 0–5.

Operation:

```
IF ($axS.h & 64)
    IF ($axS.h & 63) != 0
        $acD >>= (64 - ($axS.h & 63))
    ENDIF
ELSE
        $acD <<= $axS.h
ENDIF
FLAGS($acD)
$pc++</pre>
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.21 ASR16

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1001	r0	01			хх	хx	:		хх	ХX	:

Format:

ASR16 \$acR

Description:

Arithmetically right shifts accumulator \$acR by 16.

Operation:

\$acR >>= 16
FLAGS(\$acR)
\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
_	-	X	X	X	X	0	0

5.6.22 BLOOP

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0		
0000	0000	011r	rrrr		
aaaa	aaaa	aaaa	aaaa		

Format:

```
BLOOP $R, addrA
```

Description:

Repeatedly execute a block of code starting at the following opcode until the counter specified by the value from register \$R reaches zero. Block ends at specified address addrA inclusive. i.e. opcode at addrA is the last opcode included in loop. Counter is pushed on loop stack \$st3, end of block address is pushed on loop stack \$st2 and the repeat address is pushed on call stack \$st0. Up to 4 nested loops are allowed.

When using \$ac0.m or \$ac1.m as the initial counter value, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.23 BLOOPI

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0001	0001	iiii	iiii
aaaa	aaaa	aaaa	aaaa

Format:

```
BLOOPI #I, addrA
```

Description:

Repeatedly execute a block of code starting at the following opcode until the counter specified by the immediate value I reaches zero. Block ends at specified address addrA inclusive. i.e. opcode at addrA is the last opcode included in loop. Counter is pushed on loop stack \$st3, end of block address is pushed on loop stack \$st2 and the repeat address is pushed on call stack \$st0. Up to 4 nested loops are allowed.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.24 CALL

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0		
0000	0010	1011	1111		
aaaa	aaaa	aaaa	aaaa		

Format:

CALL addressA

${\bf Description:}$

Call function. Push program counter of the instruction following "call" to call stack \$st0. Set program counter to address represented by the value that follows this CALL instruction.

Operation:

```
// Must skip value that follows "call"
PUSH_STACK($st0)
$st0 = $pc + 2
$pc = addressA
```

	OS	LZ	ТВ	S32	S	AZ	О	С
Γ	-	-	-	-	-	-	-	-

5.6.25 CALLcc

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	0010	1011	сссс
aaaa	aaaa	aaaa	aaaa

Format:

```
{\tt CALLcc\ addressA}
```

${\bf Description:}$

Call function if condition cc has been met. Push program counter of the instruction following "call" to call stack \$st0. Set program counter to address represented by the value that follows this CALL instruction.

Operation:

```
// Must skip value that follows "call"
IF (cc)
    PUSH_STACK($st0)
    $st0 = $pc + 2
    $pc = addressA
ELSE
    $pc += 2
ENDIF
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.26 CALLR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	01	11	-		rr	r1			11	11	

Format:

CALLR \$R

Description:

Call function. Push program counter of the instruction following "call" to call stack \$\$t0. Set program counter to register \$R.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.27 CALLRcc

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	01	11			rr	r1			сс	cc	;

Format:

CALLRcc \$R

Description:

Call function if condition cc has been met. Push program counter of the instruction following "call" to call stack \$st0. Set program counter to register \$R.

Operation:

```
IF (cc)
          PUSH_STACK($st0)
          $st0 = $pc + 1
          $pc = $R

ELSE
          $pc++
ENDIF
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	_	-	-	-	-	-	-

$\boldsymbol{5.6.28}\quad \mathbf{CLR15}$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	11	00)		хх	ХX			хх	ХX	

Format:

CLR15

Description:

Sets \$sr.SU (bit 15) to 0, causing multiplication to treat its operands as signed.

Operation:

See also:

SET15

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.29 CLR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	r0	01			хх	ХX	:		хх	ХX	

Format:

CLR \$acR

${\bf Description:}$

Clears accumulator \$acR.

Operation:

\$acR = 0
FLAGS(\$acR)
\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	1	0	0	1	0	0

5.6.30 CLRL

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1111	11	0r	•		хх	хx	:		хх	ХX	:

Format:

```
CLRL $acR.1
```

Description:

Rounds \$acR such that \$acR.1 is 0. This is a round-to-even operation.

Operation:

```
IF ($acR & 0x10000) != 0
        $acR = ($acR + 0x8000) & ~0xffff
ELSE
        $acR = ($acR + 0x7fff) & ~0xffff
ENDIF
FLAGS($acR)
$pc++
```

Note:

An alternative interpretation is that if \$acR.m is odd, then increment \$acsR if \$acR.l is greater than or equal to 0x8000; if \$acR.m is even, then increment \$acsR if \$acR.l is greater than or equal to 0x7ffff. Afterwards set \$acR.l to 0.

OS	LZ	ТВ	S32	S	AZ	О	C
-	-	X	X	X	X	0	0

5.6.31 CLRP

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	01	00)		хх	ХX	:		хх	ХX	

Format:

CLRP

Description:

Clears product register \$prod.

Operation:

```
$prod = 0 // See note below
$pc++
```

Note:

Actually product register gets cleared by setting registers with following values:

\$prod.1 = 0x0000
\$prod.m1 = 0xfff0
\$prod.h = 0x00ff
\$prod.m2 = 0x0010

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.32 CMP

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	00	10			хх	хx	:		хх	хx	

Format:

CMP

Description:

Compares accumulator ac0 with accumulator ac1.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.33 CMPAXH

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
110r	s00	1		хх	хx	:		хх	ХX	:

Format:

Description:

Compares accumulator \$acS with high part of secondary accumulator \$axR.h.

Operation:

$$sr = FLAGS(sacS - (saxR.h << 16))$$

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.34 CMPI

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	001d	1000	0000
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Compares accumulator with immediate. Comparison is performed by subtracting the immediate (16-bit sign-extended) from mid accumulator \$acD.hm and computing flags based on whole accumulator \$acD.

Operation:

О	\overline{S}	LZ	ТВ	S32	S	AZ	О	С
X		-	X	X	X	X	X	X

5.6.35 CMPIS

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
0000	01	1d	L		ii	ii			ii	ii	

Format:

Description:

Compares accumulator with short immediate. Comparison is performed by subtracting the short immediate (8-bit sign-extended) from mid accumulator \$acD.hm and computing flags based on whole accumulator \$acD.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.36 DAR

15 14 13 12	11 10	9 8	7	6	5	4	3	2	1	0
0000	000	00		00	00)		01	dd	l

Format:

DAR \$arD

${\bf Description:}$

Decrement address register \$arD.

Operation:

\$arD-\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.37 DEC

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0111	10	1d	l		хх	хx	:		хх	ХX	:

Format:

DEC \$acD

${\bf Description:}$

Decrements accumulator \$acD.

Operation:

\$acD-FLAGS(\$acD)
\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
X	_	X	X	X	X	X	X

5.6.38 **DECM**

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0111	10	0d	l		хх	хx	:		хх	ХX	:

Format:

DECM \$acsD

Description:

Decrements 24-bit mid-accumulator \$acsD.

Operation:

\$acsD-FLAGS(\$acD)
\$pc++

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	X	-	X	X	X	X	X	X

5.6.39 HALT

15 14 13 12	11 10	9 8	7	6	5	4	3	2	1	0
0000	000	00		00	10)		00	01	

Format:

HALT

Description:

Stops execution of DSP code. Sets bit ${\tt DSP_CR_HALT}$ in register ${\tt DREG_CR}.$

Operation:

OS	LZ	ТВ	S32	\mathbf{S}	AZ	О	С
-	-	-	-	-	-	-	-

5.6.40 IAR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	00)		00	00)		10	dd	l

Format:

IAR \$arD

${\bf Description:}$

Increment address register \$arD.

Operation:

\$arD++
\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.41 IFcc

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	10)		01	11			сс	cc	;

Format:

IFcc

Description:

Executes the following opcode if the condition described by cccc has been met.

${\bf Operation:}$

```
IF (cc)
    EXECUTE_OPCODE($pc + 1)
ELSE
    $pc += 2
ENDIF
```

О	\overline{S}	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-	-

5.6.42 ILRR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	1d	l		00	01			00	SS	,

Format:

Description:

Move value from instruction memory pointed by addressing register $\$ to mid accumulator register $\$ acD.m.

Optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.43 ILRRD

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	1d	l		00	01			01	SS	,

Format:

Description:

Move value from instruction memory pointed by addressing register \$arS to mid accumulator register \$acD.m. Decrement addressing register \$arS.

Optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$acD.m = MEM[$arS]
$arS--
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.44 ILRRI

15 14 13 12	11 10	9 8	7	6	5	4	3	2	1	0
0000	001	ld		00	01			10	SS	}

Format:

Description:

Move value from instruction memory pointed by addressing register \$arS to mid accumulator register \$acD.m. Increment addressing register \$arS.

Optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$acD.m = MEM[$arS]
$arS++
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.45 ILRRN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	1 d	l		00	01			11	SS	,

Format:

Description:

Move value from instruction memory pointed by addressing register \$arS to mid accumulator register \$acD.m. Add corresponding indexing register \$ixS to addressing register \$arS.

Optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.46 INC

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0111	01	1d	l		хх	хx	:		хх	ХX	:

Format:

INC \$acD

${\bf Description:}$

Increments accumulator \$acD.

Operation:

\$acD++
FLAGS(\$acD)
\$pc++

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	X	-	X	X	X	X	X	X

5.6.47 INCM

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0111	01	0d			хх	ХX			хх	ХX	:

Format:

INCM \$acsD

Description:

Increments 24-bit mid-accumulator \$acsD.

Operation:

\$acsD++
FLAGS(\$acD)
\$pc++

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	X	-	X	X	X	X	X	X

5.6.48 JMP

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	0010	1001	1111
aaaa	aaaa	aaaa	aaaa

Format:

 ${\tt JMP\ addressA}$

${\bf Description:}$

Jumps to address. Set program counter to the address represented by the value that follows this JMP instruction.

Operation:

\$pc = addressA

О	S	LZ	ТВ	S32	S	AZ	О	С
-		-	-	-	-	-	-	-

5.6.49 Jcc

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	0010	1001	сссс
aaaa	aaaa	aaaa	aaaa

Format:

 $\operatorname{\mathsf{Jcc}}$ addressA

${\bf Description:}$

Jumps to addressA if condition cc has been met. Set program counter to the address represented by the value that follows this Jcc instruction.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.50 JMPR

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
0001	0111		r	rr	0:			11	11	

Format:

JMPR \$R

Description:

Jump to address; set program counter to a value from register $\$ R.

Operation:

$$pc = R$$

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.51 JRcc

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	01	11			rr	r0)		сс	cc	;

Format:

Description:

Jump to address if condition cc has been met; set program counter to a value from register \$R.

${\bf Operation:}$

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.52 LOOP

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	00		(01	0r	•		rr	rr	•

Format:

LOOP \$R

Description:

Repeatedly execute the following opcode until the counter specified by the value from register \$R reaches zero. Each execution decrements the counter. Register \$R remains unchanged. If register \$R is set to zero at the beginning of loop then the looped instruction will not get executed.

When using \$ac0.m or \$ac1.m as the initial counter value, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
counter = $R
WHILE (counter--)
        EXECUTE_OPCODE($pc + 1)
END
$pc += 2
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.53 LOOPI

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
0001	000	00			ii	ii			ii	ii	

Format:

LOOPI #I

Description:

Repeatedly execute the following opcode until the counter specified by immediate value I reaches zero. Each execution decrements the counter. If immediate I is set to zero at the beginning of loop then the looped instruction will not get executed.

${\bf Operation:}$

```
counter = I
WHILE (counter--)
     EXECUTE_OPCODE($pc + 1)
END
$pc += 2
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.54 LR

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	0000	110d	dddd
mmmm	mmmm	mmmm	mmmm

Format:

${\bf Description:}$

Move value from data memory pointed by address M to register \$D.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.55 LRI

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	0000	100d	dddd
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Load immediate value I to register \$D.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

$$D = I$$

 $pc += 2$

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.56 LRIS

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	1d	dd	l		ii	ii			ii	ii	

Format:

LRIS
$$$(0x18+D)$$
, #I

Description:

Load immediate value I (8-bit sign-extended) to accumulator register (0x18+D) .

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

$$(0x18+D) = I$$

\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.57 LRR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	10	00)		0s	sd	l		dd	dd	L

Format:

Description:

Move value from data memory pointed by addressing register \$arS to register \$D.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.58 LRRD

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
0001	1000)		1s	sd	L		dd	dd	L

Format:

Description:

Move value from data memory pointed by addressing register \$arS to register \$D. Decrements register \$arS.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$D = MEM[$arS]
$arS--
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.59 LRRI

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	10	01			0s	sd	l		dd	dd	L

Format:

Description:

Move value from data memory pointed by addressing register \$arS to register \$D. Increments register \$arS.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$D = MEM[$arS]
$arS++
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.60 LRRN

15 14 13 12	11 10	9 8	7	6	5	4	3	2	1	0
0001	100)1		1s	sd	L		dd	dd	l

Format:

LRRN \$D, @\$arS

Description:

Move value from data memory pointed by addressing register \$arS to register \$D. Add indexing register \$ixS to register \$arS.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$D = MEM[$arS]
$arS += $ixS
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.61 LRS

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0010	0d	dċ	l	1	mm	mm	1	1	mm	mn	1

Format:

LRS
$$$(0x18+D)$$
, @M

Description:

Move value from data memory pointed by address ($cr \ll 8$) | M to register (0x18+D).

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
(0x18+D) = MEM[($cr << 8) | M]
$pc++
```

Note:

LRS can use \$axD, but cannot use \$acD.h, while SRS and SRSH only work on \$acS.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.62 LSL

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
0001	0101	-		00	ii			ii	ii	

Format:

${\bf Description:}$

Logically left shifts accumulator acR by the amount specified by value ${\tt I}.$

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.63 LSL16

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1111	00	0r	•		хх	хx	:		хх	ХX	:

Format:

LSL16 \$acR

Description:

Logically left shifts accumulator \$acR by 16.

Operation:

\$acR <<= 16
FLAGS(\$acR)
\$pc++</pre>

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.64 LSR

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
0001	0101	:		01	ii			ii	ii	

Format:

Description:

Logically right shifts accumulator \$acR by the amount calculated by negating sign-extended bits 0–6.

${\bf Operation:}$

OS	LZ	TB	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.65 LSRN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	10)		11	00)		10	10)

Format:

LSRN

Description:

Logically shifts accumulator \$ac0 either left or right based on \$ac1.m: if bit 6 is set, a right by the amount calculated by negating sign-extended bits 0–5 occurs, while if bit 6 is clear, a left shift occurs by bits 0–5.

${\bf Operation:}$

```
IF ($ac1.m & 64)
    IF ($ac1.m & 63) != 0
        $ac0 >>= (64 - ($ac1.m & 63))
    ENDIF
ELSE
        $ac0 <<= $ac1.m
ENDIF
FLAGS($ac0)
$pc++</pre>
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.66 LSRNR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	11	0d	l		1x	ХX			хх	ХX	:

Format:

LSRNR \$acD

Description:

Logically shifts accumulator \$acD either left or right based on \$ac(1-D).m: if bit 6 is set, a right by the amount calculated by negating sign-extended bits 0-5 occurs, while if bit 6 is clear, a left shift occurs by bits 0-5.

Operation:

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.67 LSRNRX

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
0011	01	sd			1x	ХX	:		хх	ХX	

Format:

```
LSRNRX $acD, $axS.h
```

Description:

Logically shifts accumulator \$acD either left or right based on \$axS.h: if bit 6 is set, a right by the amount calculated by negating sign-extended bits 0–5 occurs, while if bit 6 is clear, a left shift occurs by bits 0–5.

Operation:

```
IF ($axS.h & 64)
    IF ($axS.h & 63) != 0
        $acD >>= (64 - ($axS.h & 63))
    ENDIF
ELSE
        $acD <<= $axS.h
ENDIF
FLAGS($acD)
$pc++</pre>
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.68 LSR16

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1111	01	0r	•		хх	хx	:		хх	ХX	:

Format:

LSR16 \$acR

Description:

Logically right shifts accumulator \$acR by 16.

Operation:

\$acR >>= 16
FLAGS(\$acR)
\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.69 M0

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
1000	1013	1		хх	ХX	:		хх	ХX	

Format:

МО

Description:

Sets \$sr.AM (bit 13) to 1, **disabling** the functionality that doubles the result of every multiply operation.

${\bf Operation:}$

See also:

M2

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

$5.6.70 \quad \mathbf{M2}$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	10	10)		хх	хx	:		хх	ХX	

Format:

M2

Description:

Sets \$sr.AM (bit 13) to 0, enabling the functionality that doubles the result of every multiply operation.

${\bf Operation:}$

See also:

МО

О	$^{\circ}$ S	LZ	ТВ	S32	S	AZ	О	С
	-	-	-	-	-	-	-	-

5.6.71 MADD

$15\ 14\ 13\ 12$	11 10	9 8	3	7	6	5	4	3	2	1	0
1111	00:	1s		:	хх	ХX			хх	ХX	:

Format:

Description:

Multiply low part \$ax\$.1 of secondary accumulator \$ax\$ by high part \$ax\$.h of secondary accumulator \$ax\$ (treat them both as signed) and add result to product register.

Operation:

See also:

\$sr.AM bit affects multiply result.

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	-	-	-	-	-	-	-	-

5.6.72 MADDC

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
1110	10	st	;		хх	ХX	:		хх	ХX	:

Format:

Description:

Multiply middle part of accumulator acs.m by high part of secondary accumulator accumulator (treat them both as signed) and add result to product register.

Operation:

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.73 MADDX

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
1110	00s1	t		хх	хx	:		хх	ХX	:

Format:

MADDX
$$$(0x18+S*2), $(0x19+T*2)$$

Description:

Multiply one part of secondary accumulator ax0 (selected by S) by one part of secondary accumulator ax1 (selected by T) (treat them both as signed) and add result to product register.

Operation:

$$prod += (0x18+S*2) * (0x19+T*2)$$

 $pc++$

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.74 MOV

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0110	11	0d	l		хх	хx	:		хх	ХX	:

Format:

$$\texttt{MOV}$$
 \$acD, \$ac(1-D)

Description:

Moves accumulator \$ac(1-D) to accumulator \$acD.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	0	X	X	0	0

5.6.75 MOVAX

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0110	10	sd	l		хх	хx	:		хх	ХX	:

Format:

Description:

Moves secondary accumulator accumulator

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.76 MOVNP

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0111	11	1d	l		хх	хx			хх	ХX	

Format:

MOVNP \$acD

Description:

Moves negated multiply product from the \$prod register to the accumulator register \$acD.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.77 MOVP

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0110	11	1d	l		хх	ХX	:		хх	ХX	

Format:

MOVP \$acD

Description:

Moves multiply product from the \$prod register to the accumulator register \$acD.

Operation:

\$acD = \$prod
FLAGS(\$acD)
\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

$\boldsymbol{5.6.78}\quad \boldsymbol{MOVPZ}$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1111	11	1d	l		хх	хx	:		хх	ХX	:

Format:

MOVPZ \$acD

Description:

Moves multiply product from the prod register to the accumulator acD and sets acD.1 to acD.

Operation:

```
$acD.hm = $prod.hm
$acD.l = 0
FLAGS($acD)
$pc++
```

O	\overline{S}	LZ	ТВ	S32	S	AZ	О	С
-	-	-	X	X	X	X	0	X

5.6.79 MOVR

$15\ 14\ 13\ 12$	11 10 9	9 8	7	6	5	4	3	2	1	0
0110	0ss	d		хх	хx	:		хх	ХX	:

Format:

MOVR
$$\$acD$$
, $\$(0x18+S)$

Description:

Moves register \$(0x18+S) (sign-extended) to middle accumulator \$acD.hm. Sets \$acD.l to 0.

Operation:

```
acD.hm = $(0x18+S)

acD.1 = 0

FLAGS($acD)

pc++
```

Γ	OS	LZ	ТВ	S32	S	AZ	О	С
	-	-	X	X	X	X	0	0

5.6.80 MRR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	11	dd	l		dd	ds	;		SS	SS	,

Format:

Description:

Move value from register \$5 to register \$D.

When moving to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

When moving from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.81 MSUB

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1111	01	1s	5		хх	хx	:		хх	ХX	:

Format:

Description:

Multiply low part \$ax\$.1 of secondary accumulator \$ax\$ by high part \$ax\$.h of secondary accumulator \$ax\$ (treat them both as signed) and subtract result from product register.

Operation:

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.82 MSUBC

$15\ 14\ 13\ 12$	11 10	9 8	7	6	5	4	3	2	1	0
1110	11s	st		хх	хx			хх	ХX	:

Format:

Description:

Multiply middle part of accumulator acs.m by high part of secondary accumulator accumulator (treat them both as signed) and subtract result from product register.

Operation:

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.83 MSUBX

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1110	01	st	;		хх	ХX	:		хх	ХX	

Format:

MSUBX
$$$(0x18+S*2), $(0x19+T*2)$$

Description:

Multiply one part of secondary accumulator ax0 (selected by S) by one part of secondary accumulator ax1 (selected by T) (treat them both as signed) and subtract result from product register.

Operation:

$$prod -= (0x18+S*2) * (0x19+T*2)$$

 $pc++$

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.84 MUL

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1001	s 0	00)		хх	ХX	:		хх	ХX	:

Format:

Description:

Multiply low part \$ax\$.1 of secondary accumulator \$ax\$ by high part \$ax\$.h of secondary accumulator \$ax\$ (treat them both as signed).

Operation:

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.85 MULAC

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1001	s1	0r	:		хх	хx	:		хх	ХX	:

Format:

Description:

Add product register to accumulator register \$acR. Multiply low part \$axS.1 of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

Operation:

See also:

\$sr.AM bit affects multiply result.

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	-	-	X	X	X	X	0	X

$\mathbf{5.6.86} \quad \mathbf{MULAXH}$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	00	11			хх	хx	:		хх	ХX	:

Format:

MULAXH

Description:

Multiplies \$ax0.h by itself.

Operation:

See also:

sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.87 MULC

$15\ 14\ 13\ 12$	11 10	9 8	7	6	5	4	3	2	1	0
110s	t00	00		хх	хx	:		хх	ХX	:

Format:

Description:

Multiply mid part of accumulator register \$acS.m by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

Operation:

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.88 MULCAC

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
110s	t1	0r	•		хх	ХX	:		ХX	ХX	:

Format:

Description:

Multiply mid part of accumulator register \$acS.m by high part \$axS.h of secondary accumulator \$axS (treat them both as signed). Add product register before multiplication to accumulator \$acR.

Operation:

```
temp = $prod
$prod = $acS.m * $axS.h
$acR += temp
$pc++
```

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.89 MULCMV

15 14 13 12	11 10	9 8	7	6	5	4	3	2	1	0
110s	t1	1r		хх	хx	:		хх	ХX	:

Format:

Description:

Multiply mid part of accumulator register \$acS.m by high part \$axT.h of secondary accumulator \$axT (treat them both as signed). Move product register before multiplication to accumulator \$acR.

Operation:

```
temp = $prod
$prod = $acS.m * $axT.h
$acR = temp
$pc++
```

See also:

 $\$ affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.90 MULCMVZ

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
110s	t0	1r	•		хх	ХX	:		хх	ХX	:

Format:

```
MULCMVZ $acS.m, $axT.h, $acR
```

Description:

Multiply mid part of accumulator register \$acS.m by high part \$aTS.h of secondary accumulator \$axT (treat them both as signed). Move product register before multiplication to accumulator \$acR. Set low part of accumulator \$acR.1 to zero.

Operation:

```
temp = $prod
$prod = $acS.m * $axT.h
$acR.hm = temp.hm
$acR.l = 0
$pc++
```

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.91 MULMV

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
1001	s1	1r			хх	хx	:		хх	ХX	:

Format:

Description:

Move product register to accumulator register \$acR. Multiply low part \$axS.1 of secondary accumulator Register\$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

Operation:

See also:

\$sr.AM bit affects multiply result.

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	-	-	X	X	X	X	0	X

5.6.92 MULMVZ

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1001	s0	1r	•		хх	хx	:		хх	ХX	:

Format:

Description:

Move product register to accumulator register \$acR and clear low part of accumulator register \$acR.1. Multiply low part \$axS.1 of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

${\bf Operation:}$

```
$acR.hm = $prod.hm
$acR.l = 0
$prod = $axS.l * $axS.h
$pc++
```

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	\mathbf{S}	AZ	Ο	С
-	-	X	X	X	X	0	X

5.6.93 MULX

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
101s	t0	00)		хх	хx	:		хх	ХX	:

Format:

Description:

Multiply one part ax0 by one part ax1 (treat them both as signed). Part is selected by ax1 and ax1 bits. Zero selects low part, one selects high part.

Operation:

$$prod = (S == 0) ? ax0.1 : ax0.h * (T == 0) ? ax1.1 : ax1.h pc++$$

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.94 MULXAC

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
101s	t1	0r	•		хх	ХX	:		хх	ХX	:

Format:

Description:

Add product register to accumulator register \$acR. Multiply one part \$ax0 by one part \$ax1 (treat them both as signed). Part is selected by S and T bits. Zero selects low part, one selects high part.

Operation:

```
acR += prod

prod = (S == 0) ? ax0.1 : ax0.h * (T == 0) ? ax1.1 : ax1.h

<math>pc++
```

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.95 MULXMV

$15\ 14\ 13\ 12$	11 10	9 8	7	6	5	4	3	2	1	0
101s	t1:	1r		хх	ХX	:		хх	ХX	:

Format:

Description:

Move product register to accumulator register \$acR. Multiply one part \$ax0 by one part \$ax1 (treat them both as signed). Part is selected by S and T bits. Zero selects low part, one selects high part.

Operation:

```
acR = prod

prod = (S == 0) ? ax0.1 : ax0.h * (T == 0) ? ax1.1 : ax1.h

<math>pc++
```

See also:

\$sr.AM bit affects multiply result.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.96 MULXMVZ

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
101s	t0	1r			хх	ХX	:		хх	ХX	:

Format:

```
MULXMVZ $ax0.S, $ax1.T, $acR
```

Description:

Move product register to accumulator register \$acR and clear low part of accumulator register \$acR.1. Multiply one part \$ax0 by one part \$ax1 (treat them both as signed). Part is selected by S and T bits. Zero selects low part, one selects high part.

Operation:

```
$acR.hm = $prod.hm
$acR.l = 0
$prod = (S == 0) ? $ax0.l : ax0.h * (T == 0) ? $ax1.l : $ax1.h
$pc++
```

See also:

\$sr.AM bit affects multiply result.

	$\overline{\text{OS}}$	LZ	ТВ	S32	S	AZ	О	С
Г	-	-	X	X	X	X	0	X

5.6.97 NEG

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0111	11	0d	l		хх	хx	:		хх	ХX	:

Format:

NEG \$acD

Description:

Negates accumulator \$acD.

Operation:

Note:

The carry flag is set only if \$acD was zero. The overflow flag is set only if \$acD was 0x800000000 (the minimum value), as -INT_MIN is INT_MIN in two's complement. In both of these cases, the value of \$acD after the operation is the same as it was before.

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.98 NOT

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	00	1d			1x	хx	:		хх	ХX	

Format:

NOT \$acD.m

Description:

Invert all bits in the middle part of accumulator acc.m (i.e. XOR with c.m).

${\bf Operation:}$

```
$acD.m = ~acD.m
FLAGS($acD)
$pc++
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.99 NOP

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	00)		00	00)		00	00)

Format:

NOP

Description:

No operation.

Operation:

\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.100 NX

15 14 13 12	11 10 9	8	7	6	5	4	3	2	1	0
1000	-000			хх	хx	:		хх	ХX	

Format:

NX

${\bf Description:}$

No operation, but can be extended with extended opcode.

Operation:

\$pc++

OS	LZ	TB	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.101 ORC

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	11	1d	l		0x	ХX	:		хх	ХX	:

Format:

Description:

 $\label{logic or middle part of accumulator $acD.m$ with middle part of accumulator $ax(1-D).m.}$

Operation:

```
$acD.m |= $ac(1-D).m
FLAGS($acD)
$pc++
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

$\boldsymbol{5.6.102}\quad \mathbf{ORI}$

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	001d	0110	0000
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Logical OR of accumulator mid part acD.m with immediate value I.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	X

5.6.103 ORR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	10	sd			0x	хx	:		хх	ХX	

Format:

Description:

Logical OR middle part of accumulator \$acD.m with high part of secondary accumulator \$axS.h.

Operation:

```
$acD.m |= $axS.h
FLAGS($acD)
$pc++
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	-	-	X	X	X	X	0	X

$\mathbf{5.6.104} \quad \mathbf{RET}$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	10)		11	01			11	11	

Format:

RET

Description:

Return from subroutine. Pops stored PC from call stack \$st0 and sets \$pc to this location.

${\bf Operation:}$

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

$\bf 5.6.105 \quad RETcc$

15 14 13 12	11 10 9	8 7	6	5	4	3	2	1	0
0000	0010		11	01			сс	cc	;

Format:

RETcc

Description:

Return from subroutine if condition cc has been met. Pops stored PC from call stack st0 and sets cc to this location.

Operation:

```
IF (cc)
    POP_STACK($st0)
ELSE
    $pc++
ENDIF
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.106 RTI

15 14 13 12	11 10	9 8	3 7	6	5	4	3	2	1	0
0000	00	10		11	11			11	11	

Format:

RTI

Description:

Return from exception. Pops stored status register \$sr from data stack \$st1 and program counter PC from call stack \$st0 and sets \$pc to this location.

Operation:

```
$sr = $st1
POP_STACK($st1)
$pc = $st0
POP_STACK($st0)
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.107 RTIcc

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	10)		11	11			сс	cc	;

Format:

RTIcc

Description:

Return from exception if condition cc has been met. Pops stored status register \$sr from data stack \$st1 and program counter PC from call stack \$st0 and sets \$pc to this location.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

$\mathbf{5.6.108} \quad \mathbf{SBCLR}$

15 14 13 12	11 10	9 8	3 7	6	5	4	3	2	1	0
0001	00	10		00	00)		0i	ii	

Format:

SBCLR #I

Description:

Clear bit of status register \$sr. Bit number is calculated by adding 6 to immediate value I; thus, bits 6 through 13 (LZ through AM) can be cleared with this instruction.

Operation:

O	S	LZ	ТВ	S32	S	AZ	О	С
-		-	-	-	-	-	-	-

5.6.109 SBSET

$15\ 14\ 13\ 12$	11 10 9	8	7	6	5	4	3	2	1	0
0001	001	1		00	00)		0i	ii	

Format:

Description:

Set bit of status register \$sr. Bit number is calculated by adding 6 to immediate value I; thus, bits 6 through 13 (LZ through AM) can be set with this instruction.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.110 SET15

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	11	01			хх	хx	:		хх	ХX	

Format:

SET15

Description:

Sets \$sr.SU (bit 15) to 1, causing multiplication to treat its operands as unsigned.

Operation:

See also:

CLR15

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.111 SET16

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	11	10)		хх	хx	:		хх	ХX	:

Format:

SET16

Description:

Sets \$sr.SXM (bit 14) to 0, resulting in 16-bit sign extension.

Operation:

See also:

SET40

16-bit and 40-bit modes

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.112 SET40

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	11	11			хх	хx			хх	ХX	

Format:

SET40

Description:

Sets sr.SXM (bit 14) to 1, resulting in 40-bit sign extension.

Operation:

See also:

SET16

16-bit and 40-bit modes

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.113 SI

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0001	0110	mmmm	mmmm
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Store 16-bit immediate value I to a memory location pointed by address OxFFOO | M.

Operation:

Note:

Unlike SRS, SRSH, and LRS, SI does not use \$cr to decide the base address and instead always uses 0xFF00.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.114 SR

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	0000	111s	SSSS
mmmm	mmmm	mmmm	mmmm

Format:

${\bf Description:}$

Store value from register \$S to a memory pointed by address M.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.115 SRR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	10	10)		0d	ds	,		SS	SS	,

Format:

Description:

Store value from source register \$S to a memory location pointed by addressing register \$arD.

When storing from ac0.m or ac1.m, optionally apply saturation depending on the value of ac1.m, (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.116 SRRD

$15\ 14\ 13\ 12$	11 10 9	8	7	6	5	4	3	2	1	0
0001	1010)		1d	ds	}		ss	SS	,

Format:

Description:

Store value from source register \$S to a memory location pointed by addressing register \$arD. Decrement register \$arD.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
MEM[$arD] = $S
$arD--
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.117 SRRI

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	10	11			0d	ds	;		SS	SS	,

Format:

Description:

Store value from source register \$S to a memory location pointed by addressing register \$arD. Increment register \$arD.

When storing from ac0.m or ac1.m, optionally apply saturation depending on the value of ac1.m, (see 16-bit and 40-bit modes).

Operation:

```
MEM[$arD] = $S
$arD++
$pc++
```

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.118 SRRN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0001	10	11			1d	ds			SS	SS	,

Format:

Description:

Store value from source register \$\$\$ to a memory location pointed by addressing register \$arD. Add indexing register \$ixD to register \$arD.

When storing from ac0.m or ac1.m, optionally apply saturation depending on the value of ac1.m, (see 16-bit and 40-bit modes).

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

5.6.119 SRS

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0010	11	SS	5	1	mm	mn	1	1	mm	mn	1

Format:

SRS
$$OM, $(Ox1C+S)$$

Description:

Store value from register (0x1C+S) to a memory pointed by address $(cr \ll 8)$ | M.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

$$MEM[($cr << 8) | M] = $(0x1C+S)$$

 $$pc++$

Note:

Unlike LRS, SRS and SRSH only work on \$acS. The pattern 101s is unused and does not perform any write.

(OS	LZ	ТВ	S32	S	AZ	О	С
	-	-	-	-	-	-	-	-

5.6.120 SRSH

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0010	10	0ຣ	5	1	mm	mn	1	1	mm	mn	1

Format:

Description:

Store value from register acs.h to a memory pointed by address (cr « 8) | M.

Operation:

Note:

Unlike LRS, SRS and SRSH only work on \$acS. The pattern 101s is unused and does not perform any write.

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	-	-	-	-	-	-

$\mathbf{5.6.121} \quad \mathbf{SUB}$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0101	11	0d	l		хх	ХX			хх	ХX	:

Format:

Description:

Subtracts accumulator ac(1-D) from accumulator register acD.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.122 SUBARN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0000	00	00)		00	00)		11	dd	l

Format:

SUBARN \$arD

Description:

Subtracts indexing register $\$ from addressing register $\$

Operation:

OS	S	LZ	ТВ	S32	S	AZ	О	С
-		-	-	-	-	-	-	-

$\bf 5.6.123 \quad SUBAX$

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0101	10	sd	l		хх	хx	:		хх	ХX	:

Format:

Description:

Subtracts secondary accumulator $\$ from accumulator register $\$ CD.

${\bf Operation:}$

OS	3	LZ	ТВ	S32	S	AZ	О	С
X		-	X	X	X	X	X	X

5.6.124 SUBP

15 14 13	3 12 1	1 10	9	8	7	6	5	4	3	2	1	0
010	1	11	1d			хх	хx			хх	хx	

Format:

SUBP \$acD

Description:

Subtracts product register from accumulator register.

${\bf Operation:}$

\$acD -= \$prod
FLAGS(\$acD)
\$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.125 SUBR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0101	0s	sċ	l		хх	хx	:		хх	ХX	:

Format:

SUBR
$$$acD, $(0x18+S)$$

Description:

Subtracts register (0x18+S) from accumulator acD register.

${\bf Operation:}$

OS	LZ	ТВ	S32	S	AZ	О	С
X	-	X	X	X	X	X	X

5.6.126 TST

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1011	r0	01			хх	ХX	:		хх	ХX	:

Format:

TST \$acR

Description:

Test accumulator \$acR.

Operation:

FLAGS(\$acR) \$pc++

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.127 TSTAXH

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	01	1r	•		хх	хx	:		хх	ХX	:

Format:

TSTAXH \$axR.h

Description:

Test high part of secondary accumulator axR.h.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	0	X	X	0	0

5.6.128 TSTPROD

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
1000	01	01			хх	хx	:		хх	ХX	:

Format:

TSTPROD

${\bf Description:}$

Test the product register \$prod.

Operation:

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	0	X	X	0	X

5.6.129 XORC

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	00	0d	l		1x	ХX	:		хх	ХX	:

Format:

Description:

Logical XOR (exclusive OR) middle part of accumulator acD.m with middle part of accumulator ac(1-D).m.

Operation:

```
$acD.m ^= $ac(1-D).m
FLAGS($acD)
$pc++
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

OS	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.6.130 XORI

15 14 13 12	11 10 9 8	7 6 5 4	3 2 1 0
0000	001d	0010	0000
iiii	iiii	iiii	iiii

Format:

${\bf Description:}$

Logical XOR (exclusive OR) of accumulator mid part acd.m with immediate value I.

Operation:

	OS	LZ	ТВ	S32	S	AZ	О	С
ſ	-	-	X	X	X	X	0	0

5.6.131 XORR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
0011	00	sċ	l		0x	ХX	:		хх	ХX	:

Format:

Description:

Logical XOR (exclusive OR) middle part of accumulator \$acD.m with high part of secondary accumulator \$axS.h.

Operation:

```
$acD.m ^= $axS.h
FLAGS($acD)
$pc++
```

Note:

The main opcode is 9 bits and the extension opcode is 7 bits. The extension opcode is treated as if the 8th bit was 0 (i.e. it is 0xxxxxxx).

)S	LZ	ТВ	S32	S	AZ	О	С
-	-	X	X	X	X	0	0

5.7 Extended opcodes

Extended opcodes do not exist on their own. These opcodes can only be attached to opcodes that allow extending. Specifically, opcodes where the first nybble is 0, 1, or 2 cannot be extended. Opcodes where the first nybble is 4 or higher can be extended, using the 8 lower bits. Opcodes where the first nybble is 3 can also be extended, but the main opcode is 9 bits and the extension opcode is 7 bits. For these instructions, the extension opcode is treated as if the first bit were 0 (i.e. Oxxxxxxx). (NX has no behavior of its own, so it can be used to get an extended opcode's behavior on its own.)

Extended opcodes do not modify the program counter (\$pc register).

Extended opcodes are run *in parallel* with the main opcode; they see the same register state as the input. (For instance, MOVR'MV \$ac1, \$ax0.1: \$ax0.1, \$ac1.m (encoded as 0x6113) *swaps* the values of \$ac1.m and \$ax0.1 (and also extends the new value of \$ac1.m into \$ac1.1 and \$ac1.h).)

Since they are executed in parallel, the main and extension opcodes could theoretically write to the same registers. All opcodes that support extension only modify a main accumulator \$acD, as well as \$prod, \$sr, and/or \$pc, while the extension opcodes themselves generally only modify an additional accumulator \$axD and addressing registers \$arS. The exception is 'L and 'LN, which has the option of writing to \$acD. Thus, INC'L \$acO: \$acO.1, @\$arO (encoded as 0x7660) increments \$acO (and thus \$acO.1), but also sets \$acO.1 to the value in data memory at address \$arO and increments \$arO.

When the main and extension opcodes write to the same register, the register is set to the two values bitwise-or'd together. For the above example, \$ar0.1 would be set to (\$ar0.1 + 1) | MEM[\$ar0]. Note that no official uCode writes to the same register twice like this.

CHAPTER 5. OPCODES 170

5.8 Alphabetical list of extended opcodes

5.8.1 'DR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		00	00)		01	rr	

Format:

'DR \$arR

${\bf Description:}$

Decrement addressing register \$arR.

Operation:

\$arR--

5.8.2 'IR

15 14 13 12	11 10 9	8	7	6	Э	4	3	2	1	U
xxxx	xxxx			00	00)		10	rr	•

Format:

'IR \$arR

${\bf Description:}$

Increment addressing register \$arR.

Operation:

\$arR++

5.8.3 'L

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	хx			01	dd	l	·	d0	SS	,

Format:

```
'L $(0x18+D), @$arS
```

Description:

Load register (0x18+D) with value from memory pointed by register ars. Post increment register ars.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
(0x18+D) = MEM[$arS]
arS++
```

5.8.4 'LN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		01	dd	l		d1	SS	}

Format:

```
'LN $(0x18+D), 0$arS
```

Description:

Load register (0x18+D) with value from memory pointed by register ars. Add indexing register ix to register ars.

When loading to \$ac0.m or \$ac1.m, optionally perform sign extension depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$arS]
$arS += $ixS
```

5.8.5 'LD

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		11	dr			00	SS	}

Format:

```
'LD $ax0.D, $ax1.R, @$arS
```

Description:

Load register \$ax0.D (either \$ax0.1 or \$ax0.h, as \$(0x18+D*2)) with value from memory pointed by register \$arS. Load register \$ax1.R (either \$ax1.1 or \$ax1.h, as \$(0x19+R*2)) with value from memory pointed by register \$ar3. Increment both \$arS and \$ar3.

Operation:

```
$ax0.D = MEM[$arS]
$ax1.R = MEM[$ar3]
$arS++
$ar3++
```

Note:

S cannot be 3, as that instead encodes 'LDAX. Thus, \$arS is guaranteed to be distinct from \$ar3.

5.8.6 'LDM

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		11	dr			10	SS	}

Format:

```
'LDM $ax0.D, $ax1.R, @$arS
```

Description:

Load register \$ax0.D (either \$ax0.1 or \$ax0.h, as \$(0x18+D*2)) with value from memory pointed by register \$arS. Load register \$ax1.R (either \$ax1.1 or \$ax1.h, as \$(0x19+R*2)) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix3 to addressing register \$ar3 and increment \$arS.

Operation:

```
$ax0.D = MEM[$arS]
$ax1.R = MEM[$ar3]
$arS++
$ar3 += $ix3
```

Note:

S cannot be 3, as that instead encodes 'LDAXM. Thus, \$arS is guaranteed to be distinct from \$ar3.

5.8.7 'LDNM

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		11	dr			11	SS	}

Format:

```
'LDNM $ax0.D, $ax1.R, @$arS
```

Description:

Load register \$ax0.D (either \$ax0.1 or \$ax0.h, as \$(0x18+D*2)) with value from memory pointed by register \$arS. Load register \$ax1.R (either \$ax1.1 or \$ax1.h, as \$(0x19+R*2)) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ixS to addressing register \$arS and add corresponding indexing register \$ix3 to addressing register \$ar3.

Operation:

```
$ax0.D = MEM[$arS]
$ax1.R = MEM[$ar3]
$arS += $ixS
$ar3 += $ix3
```

Note:

S cannot be 3, as that instead encodes 'LDAXNM. Thus, \$arS is guaranteed to be distinct from \$ar3.

5.8.8 'LDN

$15\ 14\ 13\ 12$	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		11	dr			01	SS	;

Format:

```
'LDN $ax0.D, $ax1.R, @$arS
```

Description:

Load register \$ax0.D (either \$ax0.1 or \$ax0.h, as \$(0x18+D*2)) with value from memory pointed by register \$arS. Load register \$ax1.R (either \$ax1.1 or \$ax1.h, as \$(0x19+R*2)) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ixS to addressing register \$arS and increment \$ar3.

Operation:

```
$ax0.D = MEM[$arS]
$ax1.R = MEM[$ar3]
$arS += $ixS
$ar3++
```

Note:

S cannot be 3, as that instead encodes 'LDAXN. Thus, \$arS is guaranteed to be distinct from \$ar3.

5.8.9 'LDAX

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		11	sr			00	11	

Format:

```
'LDAX $axR, @$arS
```

Description:

Load register \$axR.h with value from memory pointed by register \$arS. Load register \$axR.l with value from memory pointed by register \$ar3. Increment both \$arS and \$ar3.

Operation:

```
$axR.h = MEM[$arS]
$axR.l = MEM[$ar3]
$arS++
$ar3++
```

Note:

S can be either 0 or 1, corresponding to ar0 or ar1. Thus, arS is guaranteed to be distinct from ar3. ar2 cannot be used with this instruction.

5.8.10 'LDAXM

15 14 13 12	11 10	9 8	7	6	5	4	3	2	1	0
xxxx	xx	xx		11	sr			00	11	

Format:

```
'LDAXM $axR, @$arS
```

Description:

Load register \$axR.h with value from memory pointed by register \$arS. Load register \$axR.1 with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix3 to addressing register \$ar3 and increment \$arS.

Operation:

```
$axR.h = MEM[$arS]
$axR.l = MEM[$ar3]
$arS++
$ar3 += $ix3
```

Note:

S can be either 0 or 1, corresponding to \$ar0 or \$ar1. Thus, \$arS is guaranteed to be distinct from \$ar3. \$ar2 cannot be used with this instruction.

5.8.11 'LDAXNM

15 14 13 12	11 10	9 8	7	6	5	4	3	2	1	0
xxxx	XXX	x		11	sr			00	11	

Format:

'LDAXNM \$axR, @\$arS

Description:

Load register \$axR.h with value from memory pointed by register \$arS. Load register \$axR.1 with value from memory pointed by register \$ar3. Add corresponding indexing register \$ixS to addressing register \$arS and add corresponding indexing register \$ix3 to addressing register \$ar3.

Operation:

```
$axR.h = MEM[$arS]
$axR.l = MEM[$ar3]
$arS += $ixS
$ar3 += $ix3
```

Note:

S can be either 0 or 1, corresponding to \$ar0 or \$ar1. Thus, \$arS is guaranteed to be distinct from \$ar3. \$ar2 cannot be used with this instruction.

5.8.12 'LDAXN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		11	sr			00	11	

Format:

```
'LDAXN $axR, @$arS
```

Description:

Load register \$axR.h with value from memory pointed by register \$arS. Load register \$axR.1 with value from memory pointed by register \$ar3. Add corresponding indexing register \$ixS to addressing register \$arS and increment \$ar3.

Operation:

```
$axR.h = MEM[$arS]
$axR.l = MEM[$ar3]
$arS += $ixS
$ar3++
```

Note:

S can be either 0 or 1, corresponding to \$ar0 or \$ar1. Thus, \$arS is guaranteed to be distinct from \$ar3. \$ar2 cannot be used with this instruction.

5.8.13 'LS

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		00	0s	}

Format:

```
'LS $(0x18+D), $acS.m
```

Description:

Load register \$(0x18+D) with value from memory pointed by register \$ar0. Store value from register \$acs.m to memory location pointed by register \$ar3. Increment both \$ar0 and \$ar3.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar0]
MEM[$ar3] = $acS.m
$ar0++
$ar3++
```

Note:

Differs from 'SL in that \$(0x18+D) is associated with \$ar0 instead of \$ar3 and \$acS.m is associated with \$ar3 instead of \$ar0. In both cases, \$(0x18+D) is loaded and \$acS.m is stored.

5.8.14 'LSM

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		10	0s	}

Format:

```
'LSM $(0x18+D), $acS.m
```

Description:

Load register \$(0x18+D) with value from memory pointed by register \$ar0. Store value from register \$acS.m to memory location pointed by register \$ar3. Add corresponding indexing register \$ix3 to addressing register \$ar3 and increment \$ar0.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar0]
MEM[$ar3] = $acS.m
$ar0++
$ar3 += $ix3
```

Note:

Differs from 'SLM in that (0x18+D) is associated with ar0 instead of ar3 and ar3 instead of ar3 instead of ar3. In both cases, (0x18+D) is loaded and ar3.

5.8.15 'LSNM

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		11	0s	}

Format:

```
'LSNM $(0x18+D), $acS.m
```

Description:

Load register \$(0x18+D) with value from memory pointed by register \$ar0. Store value from register \$acS.m to memory location pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and add corresponding indexing register \$ix3 to addressing register \$ar3.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar0]
MEM[$ar3] = $acS.m
$ar0 += $ix0
$ar3 += $ix3
```

Note:

Differs from 'SLNM in that (0x18+D) is associated with ar0 instead of ar3 and ar3 instead of ar3 instead or ar3 instead of ar3 instead or ar3

5.8.16 'LSN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		01	0s	}

Format:

```
'LSN $(0x18+D), $acS.m
```

Description:

Load register \$(0x18+D) with value from memory pointed by register \$ar0. Store value from register \$acS.m to memory location pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and increment \$ar3.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar0]
MEM[$ar3] = $acS.m
$ar0 += $ix0
$ar3++
```

Note:

Differs from 'SLN in that \$(0x18+D) is associated with \$ar0 instead of \$ar3 and \$acS.m is associated with \$ar3 instead of \$ar0. In both cases, \$(0x18+D) is loaded and \$acS.m is stored.

5.8.17 'MV

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	XX	2		00	01			dd	ss	,

Format:

'MV
$$(0x18+D)$$
, $(0x1c+S)$

Description:

Move value of register (0x1c+S) to the register (0x18+D).

When moving from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

$$(0x18+D) = (0x1c+S)$$

5.8.18 'NOP

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		00	00)		00	хx	:

Format:

'NOP

 ${\bf Description:}$

No operation.

Note:

Generally written by not including any extension operation, such as writing INC ac0 instead of writing INC ac0.

5.8.19 'NR

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		00	00)		11	rr	•

Format:

'NR \$arR

${\bf Description:}$

Add corresponding indexing register $\$ to addressing register $\$ register.

Operation:

\$arR += \$ixR

5.8.20 'S

15 14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
xx	ХX	:		XX	ХX	:		00	1s	,		s0	dd	L

Format:

```
'S @$arD, $(0x1c+S)
```

Description:

Store value of register \$(0x1c+S) in the memory pointed by register \$arD. Post increment register \$arD.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
MEM[\$arD] = \$(0x1c+S)
\$arD++
```

5.8.21 'SL

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		00	1s	}

Format:

```
'SL $acS.m, $(0x18+D)
```

Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Increment both \$ar0 and \$ar3.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar3]
MEM[$ar0] = $acS.m
$ar0++
$ar3++
```

Note:

Differs from 'LS in that \$(0x18+D) is associated with \$ar3 instead of \$ar0 and \$acS.m is associated with \$ar0 instead of \$ar3. In both cases, \$(0x18+D) is loaded and \$acS.m is stored.

5.8.22 'SLM

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		10	1s	,

Format:

```
'SLM $acS.m, $(0x18+D)
```

Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix3 to addressing register \$ar3 and increment \$ar0.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar3]
MEM[$ar0] = $acS.m
$ar0++
$ar3 += $ix3
```

Note:

Differs from 'LSM in that (0x18+D) is associated with ar3 instead of ar0 and ar3 in stead of ar3. In both cases, (0x18+D) is loaded and ar3.

5.8.23 'SLNM

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		11	1s	}

Format:

```
'SLNM $acS.m, $(0x18+D)
```

Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and add corresponding indexing register \$ix3 to addressing register \$ar3.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar3]
MEM[$ar0] = $acS.m
$ar0 += $ix0
$ar3 += $ix3
```

Note:

Differs from 'LSNM in that (0x18+D) is associated with ar3 instead of ar3 and ar3 instead of ar3. In both cases, (0x18+D) is loaded and ar3.

5.8.24 'SLN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		10	dd	L		01	1s	}

Format:

```
'SLN $acS.m, $(0x18+D)
```

Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and increment \$ar3.

When storing from \$ac0.m or \$ac1.m, optionally apply saturation depending on the value of \$sr.SXM (see 16-bit and 40-bit modes).

Operation:

```
$(0x18+D) = MEM[$ar3]
MEM[$ar0] = $acS.m
$ar0 += $ix0
$ar3++
```

Note:

Differs from 'LSN in that \$(0x18+D) is associated with \$ar3 instead of \$ar0 and \$acS.m is associated with \$ar0 instead of \$ar3. In both cases, \$(0x18+D) is loaded and \$acS.m is stored.

5.8.25 'SN

15 14 13 12	11 10	9	8	7	6	5	4	3	2	1	0
xxxx	xx	ХX	:		00	1s	}		s1	dd	l

Format:

```
'SN @ arD, $(0x1c+S)
```

Description:

Store value of register (0x1c+S) in the memory pointed by register arD. Add indexing register ixD to register arD.

When storing from ac0.m or ac1.m, optionally apply saturation depending on the value of ac1.m, (see 16-bit and 40-bit modes).

Operation:

```
MEM[$arD] = $(0x1c+S)
$arD += $ixD
```

5.9 Instructions sorted by opcode

Instruction	Opco	de							Page
NOP			0000	0000					136
DAR	0000	0000	0000	01dd					73
IAR	0000	0000	0000	10dd					77
SUBARN	0000	0000	0000	11dd					159
ADDARN	0000	0000	0001	ssdd					40
HALT	0000	0000	0010	0001					76
LOOP	0000	0000	010r	rrrr					89
BL00P	0000	0000	011r	rrrr	aaaa	aaaa	aaaa	aaaa	59
LRI	0000	0000	100d	dddd	iiii	iiii	iiii	iiii	92
Unknown	0000	0000	101x	xxxx					
LR	0000	0000	110d	dddd	mmmm	mmmm	mmmm	mmmm	91
SR	0000	0000	111s	ssss	mmmm	mmmm	mmmm	mmmm	151
IFcc	0000	0010	0111	сссс					78
Jcc	0000	0010	1001	сссс	aaaa	aaaa	aaaa	aaaa	86
CALLcc	0000	0010	1011	cccc	aaaa	aaaa	aaaa	aaaa	62
RETcc	0000	0010	1101	cccc					142
RTIcc	0000	0010	1111	сссс					144
ADDI	0000	001d	0000	0000	iiii	iiii	iiii	iiii	43
XORI	0000	001d	0010	0000	iiii	iiii	iiii	iiii	167
ANDI	0000	001d	0100	0000	iiii	iiii	iiii	iiii	51
ORI	0000	001d	0110	0000	iiii	iiii	iiii	iiii	139
CMPI	0000	001d	1000	0000	iiii	iiii	iiii	iiii	71
ANDF	0000	001d	1010	0000	iiii	iiii	iiii	iiii	50
ANDCF	0000	001d	1100	0000	iiii	iiii	iiii	iiii	49
LSRN	0000	0010	1100	1010					102
ASRN	0000	0010	1100	1011					55
ILRR	0000	001d	0001	00ss					79
ILRRD	0000	001d	0001	01ss					80
ILRRI	0000	001d	0001	10ss					81
ILRRN	0000	001d	0001	11ss					82
ADDIS	0000	010d	iiii	iiii					44
CMPIS	0000	011d	iiii	iiii					72
LRIS	0000	1ddd	iiii	iiii					93
LOOPI			iiii						90
BLOOPI	0001	0001	iiii	iiii	aaaa	aaaa	aaaa	aaaa	60
SBCLR	0001	0010	xxxx	xiii					145
SBSET	0001	0011	xxxx	xiii					146
LSL			00ii						99
LSR			01ii						101
ASL			10ii						53
ASR			11ii						54
SI					iiii	iiii	iiii	iiii	150
JRcc			rrr0						88
CALLRcc	0001	0111	rrr1	cccc					64

Instruction	Opco	de			Page
LRR	0001	1000	0ssd	dddd	94
LRRD	0001	1000	1ssd	dddd	95
LRRI	0001	1001	0ssd	dddd	96
LRRN	0001	1001	1ssd	dddd	97
SRR	0001	1010	0dds	ssss	152
SRRD	0001	1010	1dds	ssss	153
SRRI	0001	1011	0dds	ssss	154
SRRN	0001	1011	1dds	SSSS	155
MRR	0001	11dd	ddds	SSSS	117
LRS	0010	0ddd	mmmm	mmmm	98
SRSH	0010	100s	mmmm	mmmm	157
Unknown	0010	101x	mmmm	mmmm	
SRS	0010	11ss	mmmm	mmmm	156
XORR	0011	00sd	0xxx	xxxx	168
ANDR	0011	01sd	0xxx	xxxx	52
ORR	0011	10sd	0xxx	xxxx	140
ANDC	0011	110d	0xxx	xxxx	48
ORC	0011	111d	0xxx	xxxx	138
XORC	0011	000d	1xxx	xxxx	166
NOT	0011	001d	1xxx	xxxx	135
LSRNRX	0011	01sd	1xxx	xxxx	104
ASRNRX	0011	10sd	1xxx	xxxx	57
LSRNR	0011	110d	1xxx	xxxx	103
ASRNR	0011	111d	1xxx	xxxx	56
ADDR	0100	0ssd	xxxx	xxxx	47
ADDAX	0100	10sd	xxxx	xxxx	41
ADD	0100	110d	xxxx	xxxx	39
ADDP	0100	111d	xxxx	xxxx	45
SUBR	0101	0ssd	xxxx	xxxx	162
SUBAX	0101	10sd	xxxx	xxxx	160
SUB	0101	110d	xxxx	xxxx	158
SUBP	0101	111d	xxxx	xxxx	161
MOVR	0110	0ssd	xxxx	xxxx	116
MOVAX	0110	10sd	xxxx	xxxx	112
MOV	0110	110d	xxxx	xxxx	111
MOVP	0110	111d	xxxx	xxxx	114
ADDAXL	0111	00sd	xxxx	xxxx	42
INCM	0111	010d	xxxx	xxxx	84
INC	0111	011d	xxxx	xxxx	83
DECM	0111	100d	xxxx	xxxx	75
DEC	0111	101d	xxxx	xxxx	74
NEG	0111	110d	xxxx	xxxx	134
MOVNP	0111	111d	xxxx	xxxx	113

Instruction	Opcode	Page
NX	1000 x000 xxxx xxxx	137
CLR	1000 r001 xxxx xxxx	66
CMP	1000 0010 xxxx xxxx	69
MULAXH	1000 0011 xxxx xxxx	123
CLRP	1000 0100 xxxx xxxx	68
TSTPROD	1000 0101 xxxx xxxx	165
TSTAXH	1000 011r xxxx xxxx	164
	1000 1010	10
M2	1000 1010 xxxx xxxx	107
MO	1000 1011 xxxx xxxx	106
CLR15	1000 1100 xxxx xxxx	65
SET15	1000 1101 xxxx xxxx	147
SET16	1000 1110 xxxx xxxx	148
SET40	1000 1111 xxxx xxxx	149
MUL	1001 s000 xxxx xxxx	121
ASR16	1001 s000 xxxx xxxx 1001 r001 xxxx xxxx	58
MULMVZ	1001 1001 XXXX XXXX 1001 s01r xxxx xxxx	129
MULAC	1001 S011 XXXX XXXX 1001 S10r XXXX XXXX	129 122
MULMV	1001 s101 xxxx xxxx 1001 s11r xxxx xxxx	128
MOLMV	1001 5111 XXXX XXXX	120
MULX	101s t000 xxxx xxxx	130
ABS	1010 d001 xxxx xxxx	38
TST	1011 r001 xxxx xxxx	163
MULXMVZ	101s t01r xxxx xxxx	133
MULXAC	101s t10r xxxx xxxx	131
MULXMV	101s t11r xxxx xxxx	132
MULC	110s t000 xxxx xxxx	124
CMPAXH	110r s001 xxxx xxxx	70
MULCMVZ	110s t01r xxxx xxxx	127
MULCAC	110s t10r xxxx xxxx	125
MULCMV	110s t11r xxxx xxxx	126
MADDX	1110 00st xxxx xxxx	110
MSUBX	1110 01st xxxx xxxx	120
MADDC	1110 10st xxxx xxxx	109
MSUBC	1110 11st xxxx xxxx	119
LSL16	1111 000r xxxx xxxx	100
MADD	1111 0001 XXXX XXXX 1111 001s XXXX XXXX	100
LSR16	1111 0015 XXXX XXXX 1111 010r xxxx xxxx	105
MSUB	1111 010r xxxx xxxx 1111 011s xxxx xxxx	118
ADDPAXZ	1111 Olis xxxx xxxx 1111 10sd xxxx xxxx	46
CLRL		$\frac{46}{67}$
	1111 110r xxxx xxxx	
MOVPZ	1111 111d xxxx xxxx	115

Extension Opcodes

Instruction	Opco	de			Page
'NOP	xxxx	xxxx	0000	00xx	188
'DR	xxxx	xxxx	0000	01rr	171
'IR	xxxx	xxxx	0000	10rr	172
'NR	xxxx	xxxx	0000	11rr	189
' MV	xxxx	xxxx	0001	ddss	187
'S	xxxx	xxxx	001s	s0dd	190
'SN	xxxx	xxxx	001s	s1dd	195
'L	xxxx	xxxx	01dd	d0ss	173
'LN	xxxx	xxxx	01dd	d1ss	174
'LS	xxxx	xxxx	10dd	000s	183
'SL	xxxx	xxxx	10dd	001s	191
'LSN	xxxx	xxxx	10dd	010s	186
'SLN	xxxx	xxxx	10dd	011s	194
'LSM	xxxx	xxxx	10dd	100s	184
'SLM	xxxx	xxxx	10dd	101s	192
'LSNM	xxxx	xxxx	10dd	110s	185
'SLNM	xxxx	xxxx	10dd	111s	193
'LD	xxxx	xxxx	11dr	00ss	175
'LDAX	xxxx	xxxx	11sr	0011	179
'LDN	xxxx	xxxx	11dr	01ss	178
'LDAXN	xxxx	xxxx	11sr	0111	182
'LDM	xxxx	xxxx	11dr	10ss	176
'LDAXM	xxxx	xxxx	11sr	1011	180
'LDNM	xxxx	xxxx	11dr	11ss	177
'LDAXNM	xxxx	xxxx	11dr	1111	181