



Lab Report-12

(Dijkstra's Algorithm)

CSE-2212 (Design and Analysis of Algorithms Lab)

Submitted By:

Name: Eyasir Ahamed
Exam Roll: 413
Class Roll: 15
Registration No:
202004017

Submitted To:

Sharad Hasan
Ex. Lecturer
Dept. of CSE
Sheikh Hasina University,
Netrokona

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING
SHEIKH HASINA UNIVERSITY
NETROKONA, BANGLADESH

#12_Dijkstra's Algorithm

Problem Definition: Given a weighted directed graph and a source vertex, Dijkstra's algorithm is used to find the shortest distance from the source vertex to all other vertices in the graph.

Formal Statement of Algorithm (Dijkstra's Algorithm):

- Initialize a priority queue pq to store pairs (distance, vertex) sorted by distance in non-decreasing order.
- Create a 1-indexed array distTo to store the shortest distances from the source vertex to all other vertices. Initialize all distances to infinity, except for the source vertex distance set to 0.
- Push the pair (0, source) onto the priority queue.
- While the priority queue is not empty, do the following:
 - Extract the pair (dist, prev) from the priority queue, where prev is the vertex with the shortest known distance from the source.
 - For each adjacent vertex next of prev, calculate the distance nextDist from the source through prev and update distTo[next] if it's shorter than the current distance.

- Push the pair $(\text{distTo}[\text{next}], \text{next})$ onto the priority queue if $\text{distTo}[\text{next}]$ was updated.
- After the algorithm terminates, the array distTo contains the shortest distances from the source vertex to all other vertices in the graph.

Complexity Analysis:

- Time Complexity: $O((V + E) \log V)$, where V is the number of vertices and E is the number of edges in the graph. This complexity arises from the fact that each edge is processed once, and each edge addition to the priority queue takes logarithmic time.
- Space Complexity: $O(V + E)$, where V is the number of vertices and E is the number of edges. The space complexity is dominated by the adjacency list representation of the graph and the priority queue.

Actual Code and Output

```
1  #include<bits/stdc++.h>
2  using namespace std;
3
4  int main(){
5      int n=5,m=6,source=1;
6      vector<pair<int,int>> g[n+1]; // assuming 1 based indexing of graph
7      // Constructing the graph
8      g[1].push_back({2,2});
9      g[1].push_back({4,1});
10     g[2].push_back({1,2});
11     g[2].push_back({5,5});
12     g[2].push_back({3,4});
13     g[3].push_back({2,4});
14     g[3].push_back({4,3});
15     g[3].push_back({5,1});
16     g[4].push_back({1,1});
17     g[4].push_back({3,3});
18     g[5].push_back({2,5});
19     g[5].push_back({3,1});
20     // Dijkstra's algorithm begins from here
21     priority_queue<pair<int,int>,vector<pair<int,int>>,greater<pair<int,int>>> pq;
22     vector<int> distTo(n+1,INT_MAX); //1-indexed array for calculating shortest paths
23     distTo[source] = 0;
24     pq.push(make_pair(0,source)); // (dist,source)
25     while( !pq.empty() ){
26         int dist = pq.top().first;
27         int prev = pq.top().second;
28         pq.pop();
29         for( auto it = g[prev].begin() ; it != g[prev].end() ; it++){
30             int next = it->first;
31             int nextDist = it->second;
32             if( distTo[next] > distTo[prev] + nextDist){
33                 distTo[next] = distTo[prev] + nextDist;
34                 pq.push(make_pair(distTo[next], next));
35             }
36         }
37     }
38     cout << "The distances from source " << source << " are : \n";
39     for(int i = 1 ; i<=n ; i++)    cout << distTo[i] << " ";
40     cout << "\n";
41     return 0;
42 }
43
```

The distances from source 1 are :
0 2 4 1 5
[Finished in 1.2s]