



Lab Report-11

(Longest_Common_Subsequence)

CSE-2212 (Design and Analysis of Algorithms Lab)

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#11_Longest Common Subsequence

Problem Definition: Given two strings, the problem is to find the length of the Longest Common Subsequence (LCS) between them.

Formal Statement of Algorithm (Longest Common Subsequence using Dynamic Programming):

- Define a function lcs that takes two strings s1 and s2 as input parameters.
- Initialize two variables n and m to store the lengths of s1 and s2 respectively.
- Create a 2D DP table dp of dimensions $(n + 1) \times (m + 1)$ and initialize all cells to -1.
- Initialize the base cases:
 - Set $dp[i][0] = 0$ for all i from 0 to n.
 - Set $dp[0][i] = 0$ for all i from 0 to m.
- Iterate over the characters of s1 and s2 using two nested loops:
 - If the characters at indices ind1 and ind2 match ($s1[ind1 - 1] == s2[ind2 - 1]$), increment the length of LCS by 1 ($dp[ind1][ind2] = 1 + dp[ind1 - 1][ind2 - 1]$).
 - Otherwise, set the length of LCS at the current indices to the maximum of the LCS lengths obtained by excluding each

character ($dp[ind1][ind2] = \max(dp[ind1 - 1][ind2], dp[ind1][ind2 - 1])$).

- The final result is stored in $dp[n][m]$, which represents the length of the Longest Common Subsequence.

Complexity Analysis:

- Time Complexity: The algorithm fills in a 2D DP table of size $(n + 1) \times (m + 1)$, so the time complexity is $O(nm)$, where n and m are the lengths of the input strings $s1$ and $s2$ respectively.
- Space Complexity: The space complexity is also $O(nm)$ due to the DP table.

Actual Code and Output

```
1  #include <bits/stdc++.h>
2  using namespace std;
3
4  int lcs(string s1, string s2) {
5      int n = s1.size();
6      int m = s2.size();
7
8      vector<vector<int>> dp(n + 1, vector<int>(m + 1, -1));
9
10     for (int i = 0; i <= n; i++) {
11         dp[i][0] = 0;
12     }
13     for (int i = 0; i <= m; i++) {
14         dp[0][i] = 0;
15     }
16
17     for (int ind1 = 1; ind1 <= n; ind1++) {
18         for (int ind2 = 1; ind2 <= m; ind2++) {
19             if (s1[ind1 - 1] == s2[ind2 - 1])
20                 dp[ind1][ind2] = 1 + dp[ind1 - 1][ind2 - 1];
21             else
22                 dp[ind1][ind2] = max(dp[ind1 - 1][ind2], dp[ind1][ind2 - 1]);
23         }
24     }
25
26     return dp[n][m];
27 }
28
29 int main() {
30     string s1 = "acd";
31     string s2 = "ced";
32
33     cout << "The Length of Longest Common Subsequence is " << lcs(s1, s2) << endl;
34
35     return 0;
36 }
37
```

The Length of Longest Common Subsequence is 2
[Finished in 1.2s]