

10/31/18 415

Project 3 is up
- Due Tuesday 11/13

Compare efficiency Z-3 v. BST

main goal to implement Z-3 tree

probability 2 people $\frac{1}{365}$ $\frac{364}{365}$
3 people

$$\frac{365}{365} \frac{364}{365}$$

$$\frac{n!}{n} = \boxed{(n-1)!}$$

$$\frac{n!}{365^n}$$

$$\frac{n!}{365^n}$$

no two people

$$\frac{365!}{(365-n)! 365^n}$$

at least two people $1 - \left(\frac{365!}{(365-n)! 365^n} \right)$

Collision Types

Open Hashing

each cell is header of linked list

$$h(c_0 \dots c_n) = \left(\sum_{i=0}^n p w(c_i) \right) \% 13$$

if evenly distributed

length of list = load factor = $\alpha \approx n/m$

Average probes in S (successful) or U (unsuccessful)

$$S \approx 1 + \frac{\alpha}{2} \quad U = \alpha$$

$$\begin{array}{r} 30 \\ + 18 \\ \hline 48 \end{array}$$

1	2	3	4	5	6	7	8	9	10	11	12	13
A	B	C	D	E	F	G	H	I	J	K	L	M

14	15	16	17	18	19	20	21	22	23	24	25	26
N	O	P	Q	R	S	T	U	V	W	X	Y	Z

$$\begin{array}{r} 3 \\ 13 \overline{) 48} \\ \underline{39} \\ 9 \end{array}$$

$$13 \overline{) 24} \quad 1$$

$$\boxed{11} = h(\text{key})$$

load α is typically kept small (ideally close to 1)

Closed Hashing:

each slot used 1x

linear probing \rightarrow probe til next open slot

$$S = \frac{1}{2} \left(1 + \frac{1}{1-\alpha} \right) \quad \text{and} \quad U \approx \frac{1}{2} \left(1 + \frac{1}{(1-\alpha)^2} \right)$$

11/5 40'

Announcements

Project 3 Due next week

HW 5 TBD

Hashing

- A very efficient way of implementing a dictionary with:
 - insert
 - delete
 - find

How to implement

closed: don't

→ if

Hashing Application

- symbol tables (ie table of a compiler's generation of program's symbols)
- AI applications
store predictions and probabilities of
- Databases (extendible hashing)

Open
Closed $S(x) = 1 + \frac{1}{1-x}$

Dynamic Programming

main idea

set up recurrence relationships
and compute and store subproblems
into a table instead of recomputing them

→
Coin-row problem

DP solution

$F(n)$ maximum value of non adjacent coins
last coin is C_i

divide into two groups

$$\begin{array}{l|l} F(n) \subseteq C_i & F(n) \notin C_i \\ F(n) = C_i + F(n-2) & F(n) = C_{i-1} + F(n-3) \end{array}$$

Rec Rel

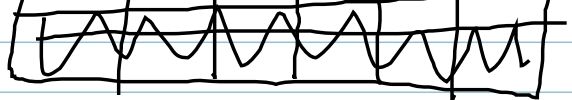
$$F(n) = \max\{C_n + F(n-2), F(n-1)\} \quad n > 1$$

index	0	1	2	3	4	5	6
coins		5	1	2	10	6	2
$F(n)$	0	5					

max of non adj.

Ex 3 Coin-collecting by robot

	1	2	3	4	5	6
1	0	0	0	0	1	1
2	0	1	1	2	2	2
3	0	1	1	3	3	4
4	0	1	2	3	3	5
5	1	1	2	3	4	5



$$F(i, j) = \max(F(i-1, j), F(i, j-1)) + C_{ij}$$

Rec Rel $F(i, j) = \max(F(i-1, j), F(i, j-1)) + C_{ij}$ for $1 \leq i \leq n, 1 \leq j < M$
 $C_{1,1} = 1$

11/19 415

Greedy

Prim's - network

Greedy strategy = Optimal solutions

- Huffman Codes = simple scheduling
- Normal coin denominations -
- MST
- single source shortest path

Approximations

- TSP
- Knapsack
- other combinatorial

Dijkstra's need a source

Prim's

Goal = not from source. Is shortest path (Minimal Spanning Tree)

✓	✓	X	X	✓	✓	✓	X	X	✓
1	2	3	4	5	6	7	8	9	10
E	C	C	C	F	F	B	T	C	A

11/

Announcements

Project 4 due Sun 12/2

3 more lectures!

HW6 will be up 11/28
- due 12/5

CS 330 Game Programming

Plaza question → ~~what~~ quite a bit vector math
- if haven't taken linear algebra still ok
- will give tools for vector math needed

All games:

- engine

→ input

→ collision physics

- AI

Game Play

Games. What are rules?

What are characters?

Labs will focus on gameplay using unreal game engine

Lecture focus on both

1 midterm → 3 Projects → weekly labs, one lab VR

No final → Final Project instead, will make game

1st Two projects individual
Final project 3 or 4

Greedy Technique Optimal Solutions

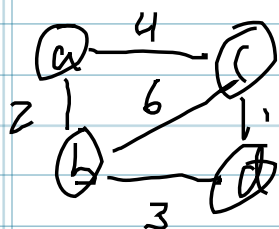
→ Project
creat bottom up ^{max} heap

show greedy speed at cost of optimality

P.I Extra Credit involves ^{approach} getting $O(n \cdot k \log n)$

PI Build path on go instead of backtracing

Prim's MST algorithm
use for solving TSP



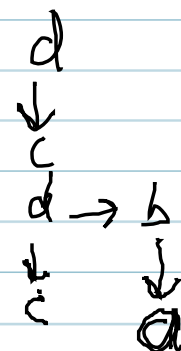
start w/ a and

start w/ a

Tree

$a(-, 0)$	$c(a, 4)$	$b(a, 2)$	$d(-, \infty)$	a ↓ b ↓ c ↓ d
$b(a, 2)$	$c(a, 4)$	$d(b, 3)$		

$d(-, 0)$	$c(d, 1), b(d, 3), a(-, \infty)$
$c(d, 1)$	$b(d, 3), a(c, 4)$
$b(d, 3)$	$a(b, 2)$



Same complexity as Dijkstra's

$|V| * (\text{Find min} + \text{update distances})$

How many

Unordered array
min or Heap
adjacency list
Weight matrix

$|E| \approx |V|^2$

Same as Dijkstra's

Complete graph $|E| \approx |V|^2$

1st option better for complete graph
2nd much better for sparse graph

$|V|^2$

Practice Drawing lines on graph from
slide show

- Shading inequalities B4 next class

