### Control Flow Edward Z. Yang

1. Structured Programming

3. Continuations

2. Procedural abstraction — the Stack

X = A

Y = B-A

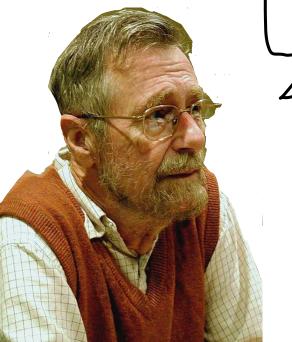
GO TO 11

#### A Case against the GO TO Statement.

by Edsger W.Dijkstra
Technological University
Eindhoven, The Netherlands

Since a number of years I am familiar with the observation that the quality of programmers is a decreasing function of the density of go to statements in the programs they produce. Later I discovered why the use of the go to statement has such disastrous effects and did I become convinced that the go to statement should be abolished from all "higher level" programming languages (i.e. everything except -perhaps- plain machine code). At that time I did not attach too much importance to this discovery; I now submit my considerations for publication because in very recent discussions in which the subject turned up, I have been urged to do so.

Go To Statement Considered Harmful



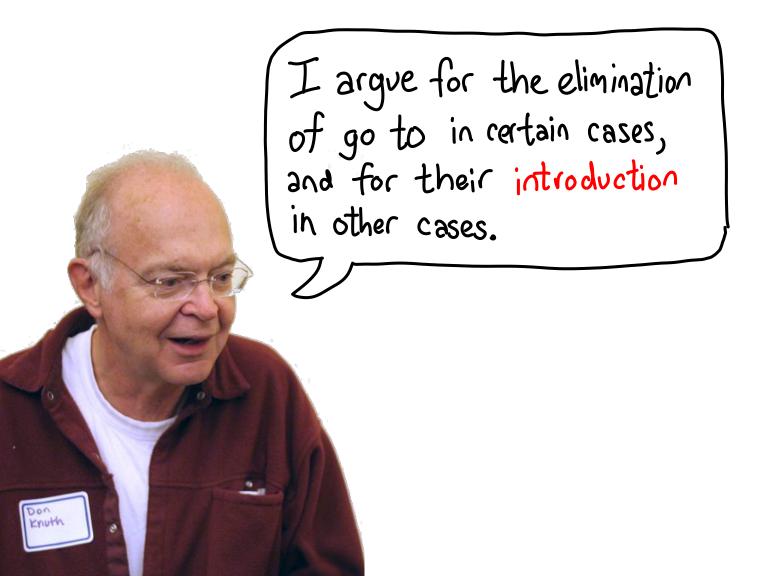
### Structured Programming with go to Statements DONALD E. KNUTH

Stanford University, Stanford, California 94805

A consideration of several different examples sheds new light on the problem of creating reliable, well-structured programs that behave efficiently. This study focuses largely on two issues: (a) improved syntax for iterations and error exits, making it possible to write a larger class of programs clearly and efficiently without go to statements; (b) a methodology of program design, beginning with readable and correct, but possibly inefficient programs that are systematically transformed if necessary into efficient and correct, but possibly less readable code. The discussion brings out opposing points of view about whether or not go to statements should be abolished; some merit is found on both sides of this question. Finally, an attempt is made to define the true nature of structured programming, and to recommend fruitful directions for further study.

Keywords and phrases: structured programming, go to statements, language design, event indicators, recursion, Boolean variables, iteration, optimization of programs, program transformations, program manipulation systems searching, Quicksort, efficiency

CR categories: 4.0, 4.10, 4.20, 5.20, 5.5, 6.1 (5.23, 5.24, 5.25, 5.27)



### Structured Programming

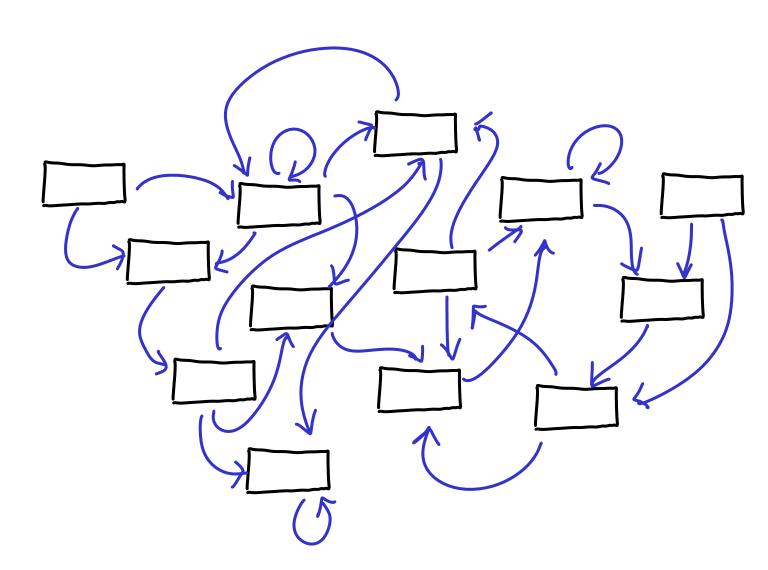
if...then ... else ...

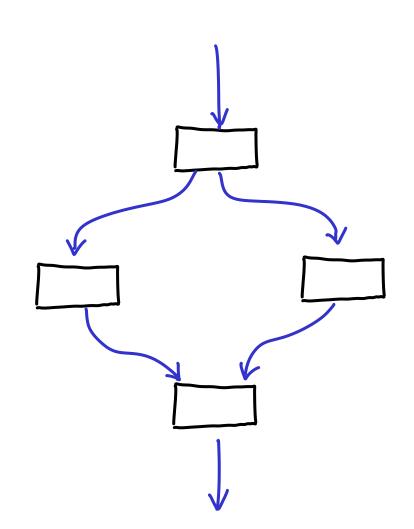
while...do...

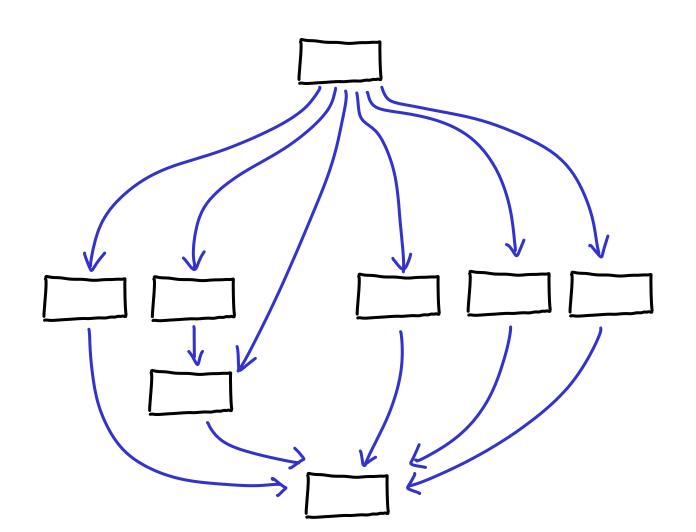
Structured Programming

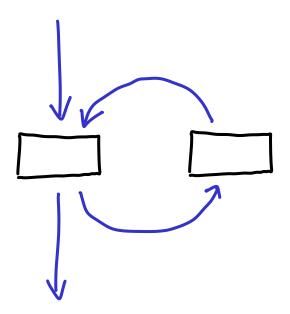
Case...

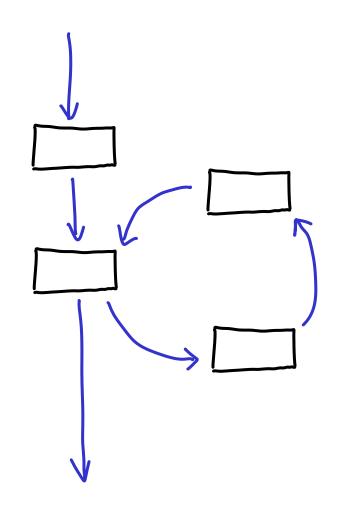
for... \{ ... \}











## Procedural Abstraction The Stack

```
function f(x) {
  return h(x) + 1;
function g(x) {
  return h(x) - 1;
function h(x) {
  return x * 2;
```

```
function f(x) {
  return h(x) + 1;
function g(x) {
  return h(x) - 1;
function h(x) {
  return x * 2;
```

```
function f(x) {
  return h(x) + 1;
function g(x) {
  return h(x) - 1;
function h(x) {
  return x * 2;
```

control link parameters local vallables return to main control link parameters return to f Stack!

return pointer 7 on the stack

Dynamic control flow

# Exceptions Also The Stack

raise/throw handler/catch

### Dynamic control flow

throw handler frame on the stack

### Dynamic control flow

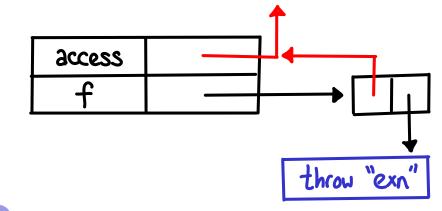
function f(y) { throw "exn";}
try {

} catch(e) { show Error();}

f(1);

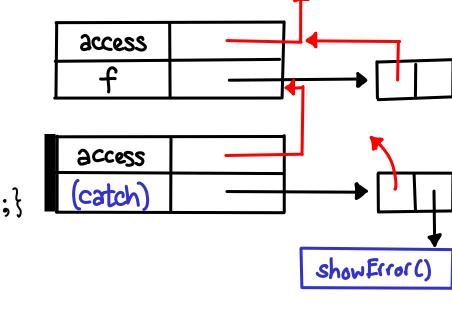
```
function f(y) { throw "exn";}

try {
  f(1);
} catch(e) { show Error();}
```



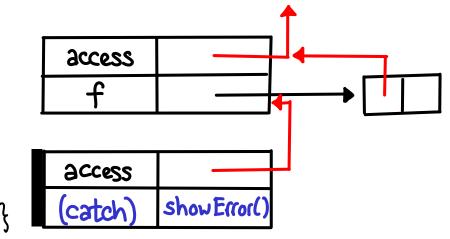
function f(y) { throw "exn";}

try {
 f(1);
} catch(e) { show Error();}



function f(y) { throw "exn";} f(1); 3 catch (e) { show Error(); }

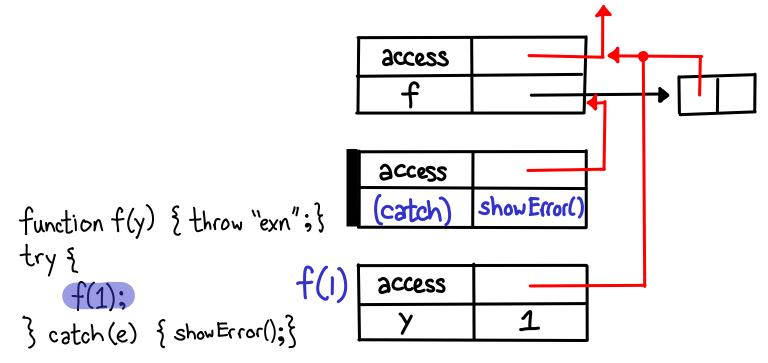
NB: try does not actually introduce scope

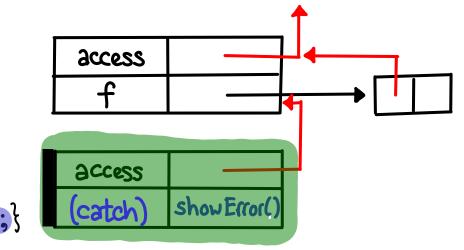


function f(y) { throw "exn";}

f(1);

3 catch (e) { show Error(); }





function f(y) { throw "exn";}

try {
 f(1);
} catch(e) { show Error();}

```
try {
    function f(y) { throw "exn";}
    function g(h) { try {h(1);}
                     catch(e) {(3)} }
     try 1
     g(f);
} catch(e) { 1 }
 } catch(e) { 2 }
```

} catch(e) { 2}

```
function f(y) { throw "exn";}

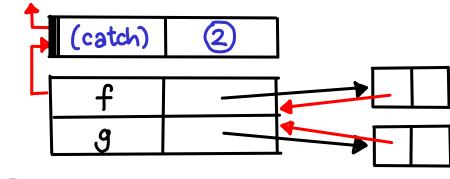
function g(h) { try {h(1);}

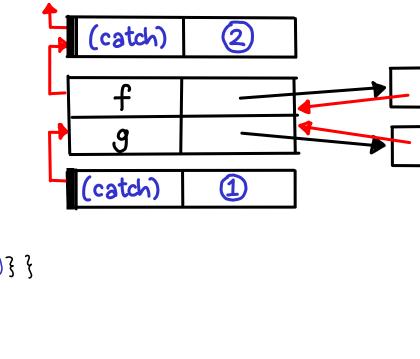
catch(e) { 3 } }

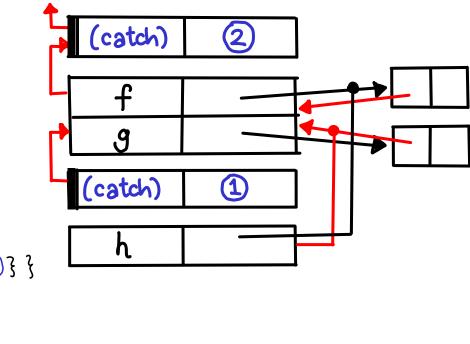
try {

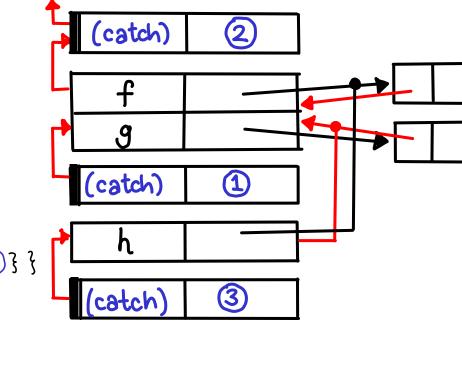
g(f);
} catch(e) { 1 }

} catch(e) { 2 }
```

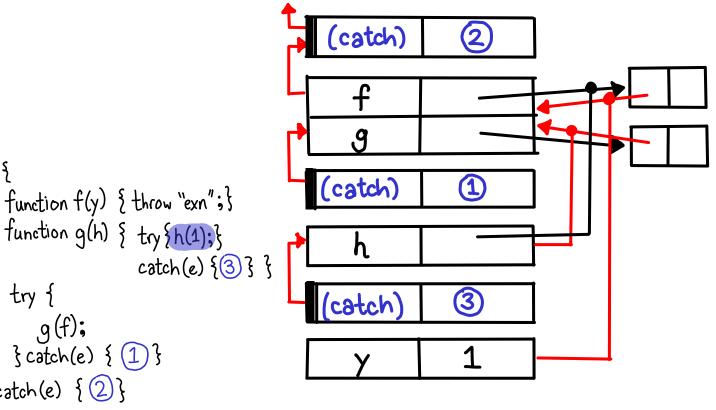








```
try {
    function f(y) { throw "exn";}
    function g(h) { try { h(1); }
                    catch(e) {3}}
     try {
     g(f);
} catch(e) { 1 }
 } catch(e) { 2}}
```

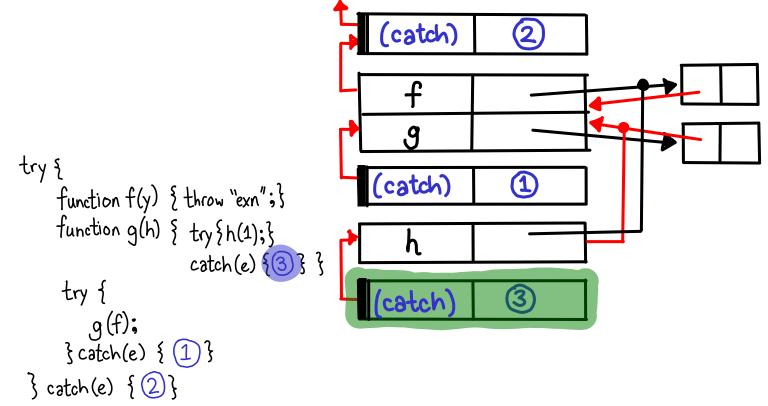


try {

try {

} catch(e) { 2}}

g(f); } catch(e) { 1 }



try { function f(y) { throw "exn";} function g(h) { try { h(1); } catch(e) {3}} try 1 g(f);
} catch(e) { 1 } } catch(e) { 2 }

DYNAMIC

function f(y) { throw "exn";} function g(h) { try {h(1);} catch(e) {(3)} } try 1 g(f);
} catch(e) { 1 } } catch(e) { 2 }

LEXICAL

```
{ let e=2;
     function f(y) { print (e); }
     function g(h) { { let e=3; h(1); }
     { let e = 1;
    g(f);
```

LEXICAL

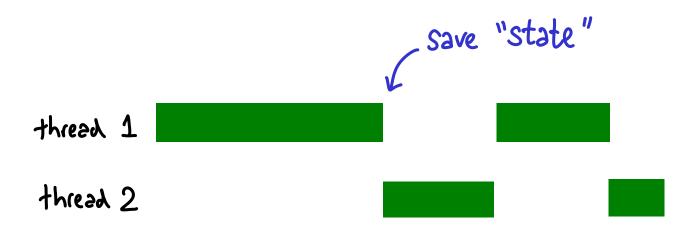
```
} let e=2;
  function f(e,y) { print(e); }
  function g(e,h) { {let e=3; h(e,1); }
   { let e = 1;
  g (e,f);
```

#### Continuations (it's more likely than you think!)

### Example: "Async" Programming

```
function getPhoto(tag, handlerCallback) {
    asyncGet(requestTag(tag), function(photoList) {
        asyncGet(requestOneFrom(photoList),
            function(photoSizes) {
                  handlerCallback(sizesToPhoto(photoSizes));
            });
    });
}
```

## Example: Cooperative Multithreading



# Example: GUI/Web Programming

```
waitForButton();
firstOperation();
sleep(1000);
secondOperation();
```

```
button.onclick = function(e) {
  firstOperation();
  setTimeout(function() {
    secondOperation();
  }, 1000);
}
```

# Example: Debugger

Var 
$$x = y + 3$$

Hit Breakpoint

 $x \mid 2$ 
 $y \mid 4$ 

What is a continuation?

Continuation-passing Style

Implementing control flow with continuations

### (2\*x+1/y)\*2

(2\*x+1/y)\*2

1. Multiply 2 and &

2. Divide 1 by y

3. Add (1) and (2)

4. Multiply 3 and 2

$$(2*x+1/y)*2$$

1. Multiply 2 and &

current -> 2. Divide 1 by y computation

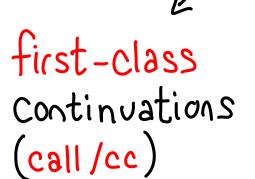
3. Add (1) and (2) } continuation
4. Multiply (3) and 2

```
(2*x+1/y)*2
            var before = 2 * x;
just a function cont(r) {

return (before + r) *2;

function }
            cont(1/y)
```

# Continuations as an implicit notion



continuation passing Style to explicitly encode continuations

## node.js example

continuation

```
current computation
var data = fs.readFileSync ("foo.txt")
console.log(data);
process Data (data);
```

# node.js example

```
fs.readFile ("foo.txt", callback)
function callback (err, data) {
    var data =
    console.log(data);
    process Data (data);
```

# Continuation Passing Style

NEVER USE RETURN

Don't call me: I'll call you!

CPS

function zero() {
 return Ø;
}

CPS

function zero() {

return Ø;
}

continuation )
function Zero(cc) {

CPS

function zero() {
 return Ø;
}

continuation )
function Zero(cc) {

cc(Ø);

Call the continuation with the return value

```
function fact(n) {
  if (n == Ø) {
    return 1;
  } else {
    return n*fact(n-1);
```

function fact(n) {
if (n==ø) { return 1;
} else {
return n*fact(n-i

function fact(n, cc) {
if (n==\phi) {
}

} else {

CPS

function fact(n) {
if (n==ø) {
return 1;
} else {
return n*fact(n-1)
}
3

function fact(n, cc) {

if (n == \phi) {

 cc(1);

} else {

25

function fact(n) { if (n == Ø) } return 1; 3 else 3 return n\*fact(n-1);

function fact(n, cc) {
if (n == \phi) { cc(1);} else cc(... fact...)

fun	ction	fa	nt(n	) }
_	(n==	=Ø)	{	<i>J</i> (
	retu	m	1;	
3	else			
	retu	<b>m</b>	n*fa	act (n-
2				

function fact(n, cc) {
 if (n == \psi) {
 cc(1);
 } else {
 fact(n-1, ...cc...);
 }

function fact(n) { if (n == Ø) { return 1; 3 else 3 return n\*fact(n-1);

function fact(n, cc) {
if (n == \phi) {

cc(1);} else { fact (n-1, function (r) { cc((\*n);

CPS

function twice (f, x) { return f(f(x));

function twice (f, x, cc) { f(x, function(r)) { f(r, cc); Mangle of

DS

function twice $(f, x)$ {
varr = f(x);
return $f(r)$ ;
}

function twice (f, x, cc) {
 f(x, function(r)) {
 f(r, cc);
 f(r, cc);

#### The rules

function 
$$(x)$$
  $\S$   $\Rightarrow$  function  $(x, cc)$   $\S$  return  $x$   $\Rightarrow$   $cc(x)$ 

var  $r = g(x)$ ;  $\Rightarrow$   $g(x, function(r))$   $\S$   $\S$   $\S$   $\S$ ;  $\S$ ;

# Do-notation CPSes for you!

 $do \{x \leftarrow e; s\} = e \gg = |x \rightarrow do \{s\}$ 

do 
$$\{e; s\}$$
 =  $e \gg do \{s\}$   
do  $\{e\}$  =  $e$   
let  $m \gg = f = \langle cc - \rangle m(\langle r - \rangle fr cc)$   
return  $\chi = \langle cc - \rangle cc \chi$ 

# Do-notation CPSes for you!

do 
$$\{e; s\}$$
 =  $\lambda cc. e(\lambda_{-}.do \{s\}cc)$   
do  $\{e\}$  =  $\lambda cc. e cc$   
let  $m \gg = f = \langle cc - \rangle m(\langle r - \rangle f r cc)$   
return  $\chi = \langle cc - \rangle cc \chi$ 

do  $\{x \leftarrow e; s\} \equiv \lambda cc. e(\lambda x. do \{s\} cc)$ 

#### Tail call optimization

$$f x = if p x$$
 then  $g x$ 
else  $g x + 2$ 

$$t not tail position$$

```
return to top
frame for f

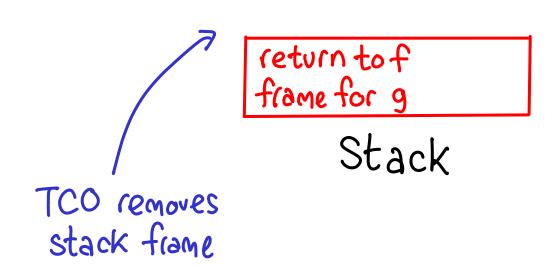
return to f
frame for g

Stack
```

```
return to top
frame for f

return to f
frame for g

Stack
```



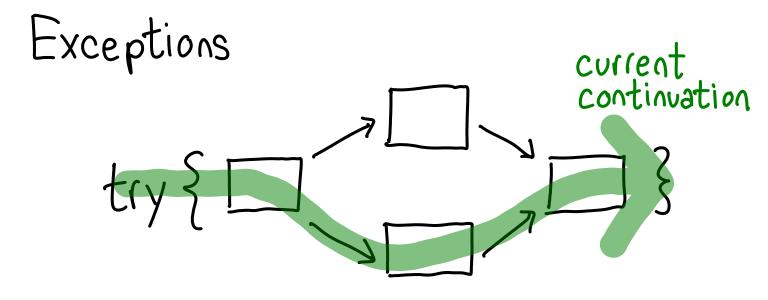
$$f \propto cc = if p \propto then g \propto (\lambda r. cc r)$$
  
else  $g \propto (\lambda r. cc (r+2))$ 

tail call optimization f x cc = if p x then g x ccelse g x (\(\alpha\rm \cap (r+2)\)

Can't eliminate new continuation/Stack frame

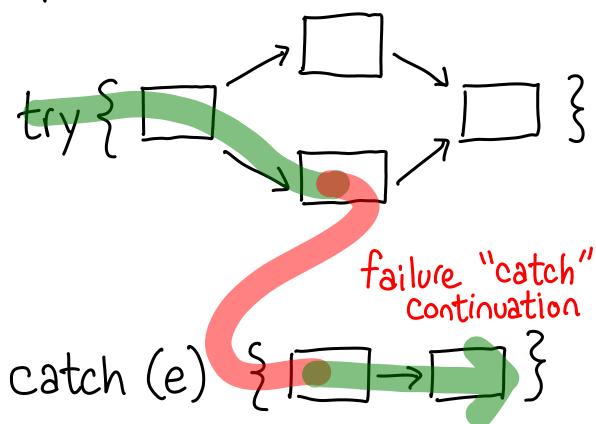
# Exceptions

catch (e) 
$$\{ \square \rightarrow \square \}$$



catch (e) 
$$\{ \Box \rightarrow \Box \}$$

## Exceptions



function f(y) { throw "exn"; } try & f(1); console.log ("No"); } catch(e) { show Error();} console.log("YES");

#### TryStatement

```
function f(y) { throw "exn"; }
     f(1); console.log ("NO");
} catch(e) { show Error();}
console.log("YES");
                       continuation
```

#### Apply Expression

```
function f(y) { throw "exn"; }
try &
    f(1); console.log ("NO");
} catch(e) { show Error();}
console.log("YES");
```

#### Apply Expression

composed w/ the previous current continuation!

```
function f(y) { throw "exn"; }
try {
     f(1); console.log ("No");
} catch(e) { show Error();}
console.log("YES");
                Current continuation
```

### Apply Expression

```
function f(y) { throw "exn";}
try {
    f(1); console.log ("NO");
} catch(e) { show Error();}
console.log ("YES");
```

tailure ("catch") continuation

#### Throw Statement

(abbr)
current
continuation

function f(y) { throw "exn";} try & f(1); console.log("NO"); } catch(e) { show Error();} console.log ("YES");

failure ("catch") continuation

#### CatchClause

```
function f(y) { throw "exn"; }
try {
    f(1); console.log("NO");
} catch(e) { show Error();}
console.log("YES");
                  current continuation
```

#### CatchClause

# failure continuation restored

```
function f(y) { throw "exn";}
try {
     f(1); console.log ("NO");
} catch(e) { show Error();}
console.log ("YES");
                  current continuation
```

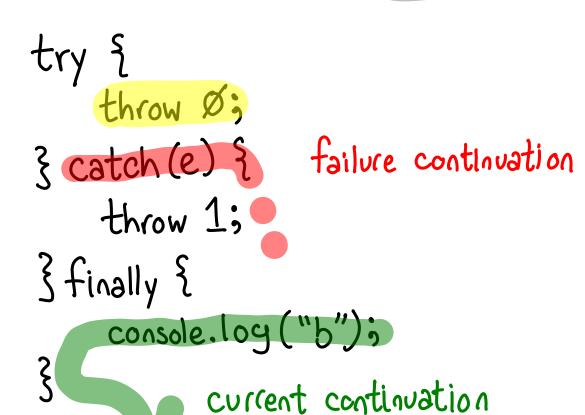
```
try ?
    throw Ø;
} catch (e) {
    throw 1;
3 finally {
    console.log("b");
```

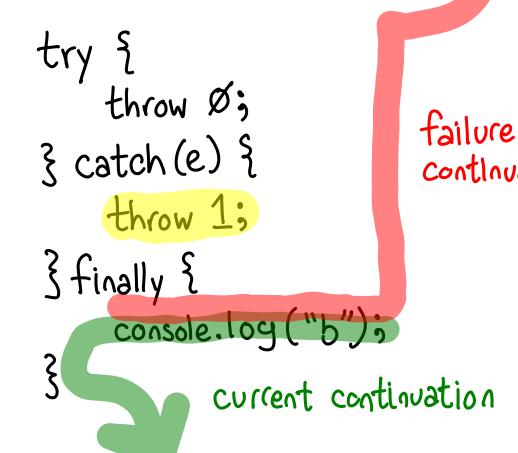
```
try ?
     throw Ø;
} catch (e) {
                       always runs, no matter how we exit scope
     throw 1;
3 finally ?
      console.log("b");
```



} catch(e) { throw 1; 3 finally { console.log("b");



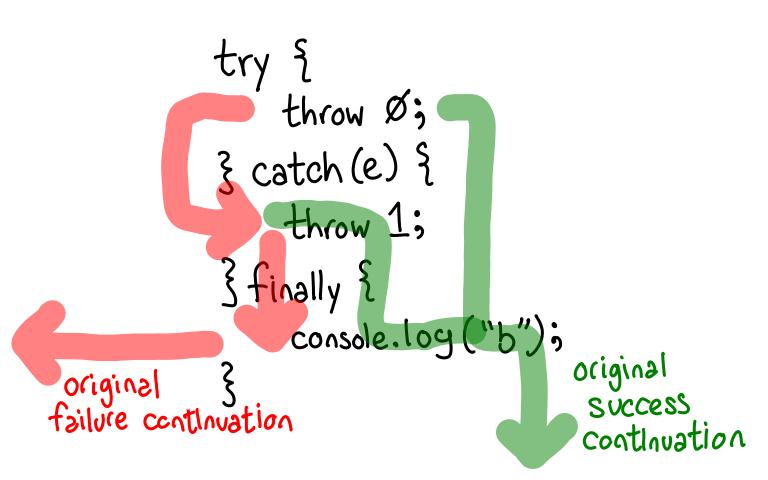




continuation

```
try ?
    throw Ø;
} catch(e) {
    throw 1;
3 finally {
    console.log("b");
```

current continuation



continuations are generalized goto!

try block catch block finally original failure continuation

original success continuation

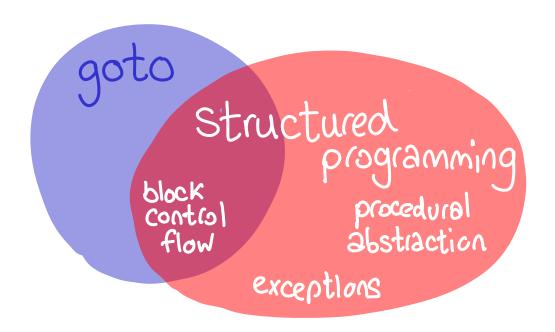
#### Other examples

- -Call with current continuation
- -Nondeterministic choice
- Generators
- -Coroutines

## Conclusion



### Conclusion



### Conclusion

