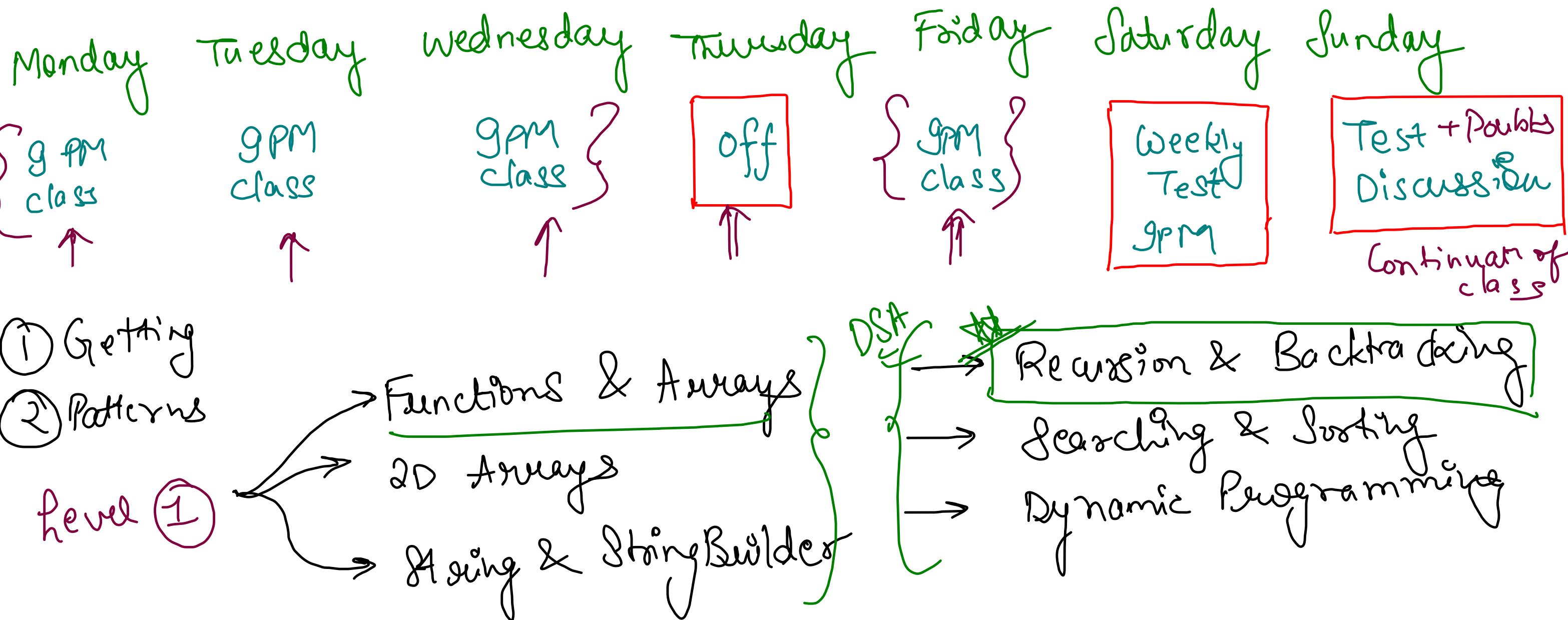


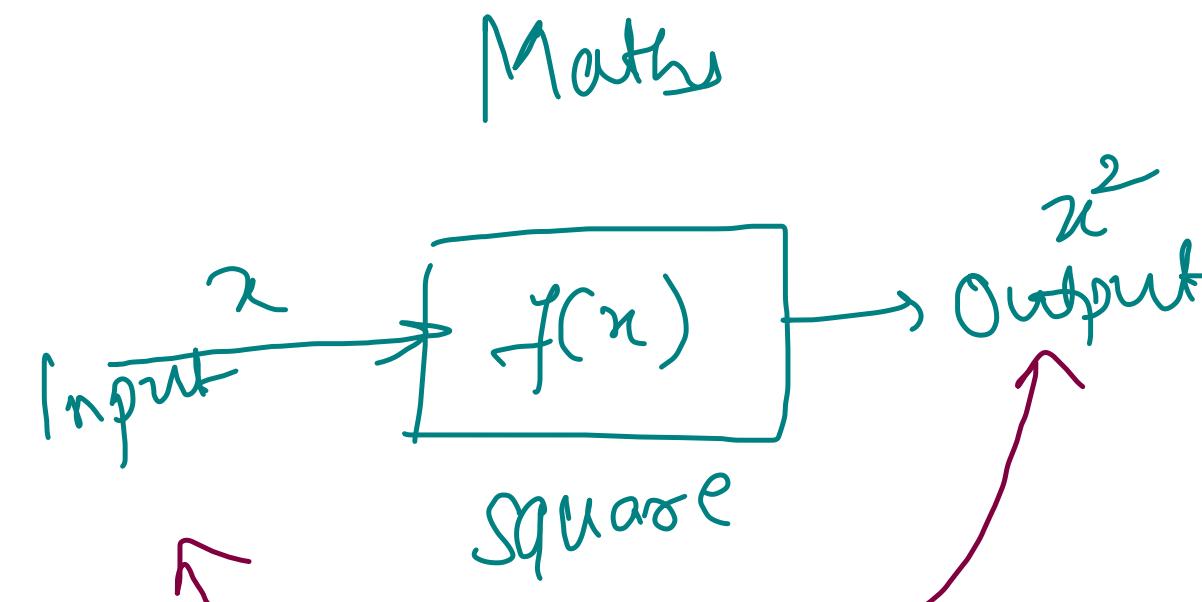
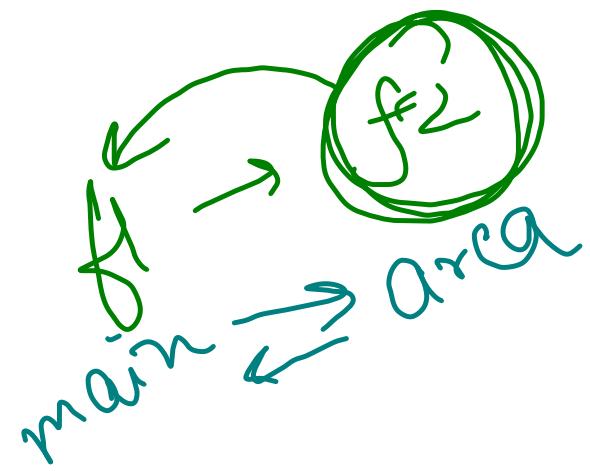
# Time Table { F-JP - 2 }



## My Introduction

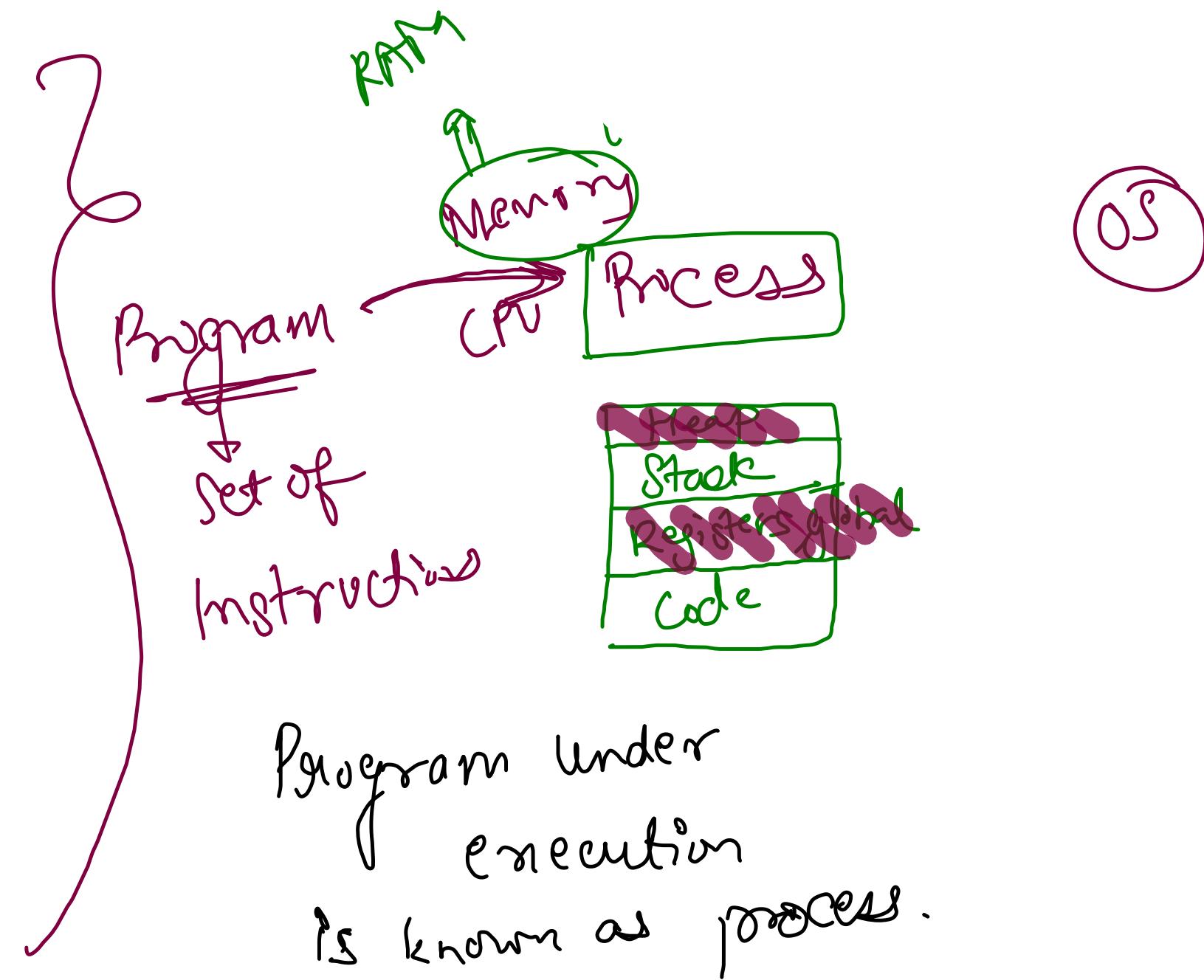
- Pointers
- Programming Articles (30 - 40+)
- Job-Switch program - 4 { 3-4 months }
- GeeksforGeeks
- LinkedIn
- Competitive Programming
- Product Based company
  - Google (3)
  - Salesforce (AMTS) (2)
  - Sprintel (1)

## # Functions



public static int volume( int l, int b, int h ) {  
 datatype function name

}



```

public static int factorial(int n){
    int res = 1;

    for(int i=1; i<=n; i++){
        res = res * i;
    }

    return res;
}

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);

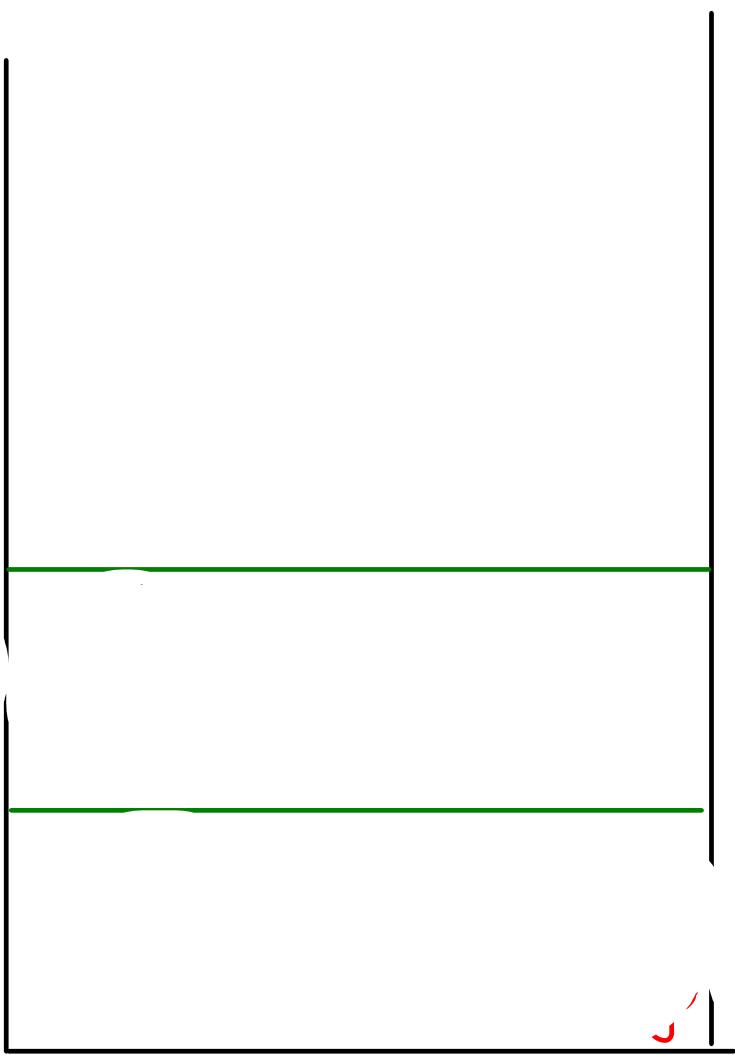
    int n = scn.nextInt();
    int r = scn.nextInt();

    int nfact = factorial(n);
    int rfact = factorial(r);
    int nmrfact = factorial(n - r);

    int ncr = (nfact / (rfact * nmrfact));
    System.out.println(ncr);
}

```

function call stack

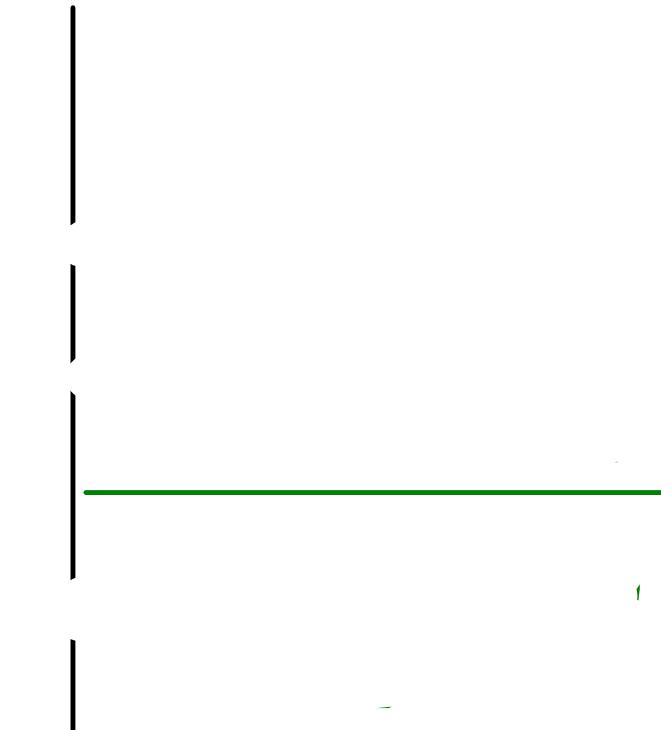


10

int  
void

$n =$ 
  
 $d =$

$\text{freq} =$



```

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    int d = scn.nextInt();
    int f = getDigitFrequency(n, d);
    System.out.println(f);
}

```

```

public static int getDigitFrequency(int n, int d){

    int freq = 0;

    // For Extracting Digits
    while(n > 0){
        int digit = n % 10;

        if(digit == d){
            freq++;
        }

        n = n / 10;
    }

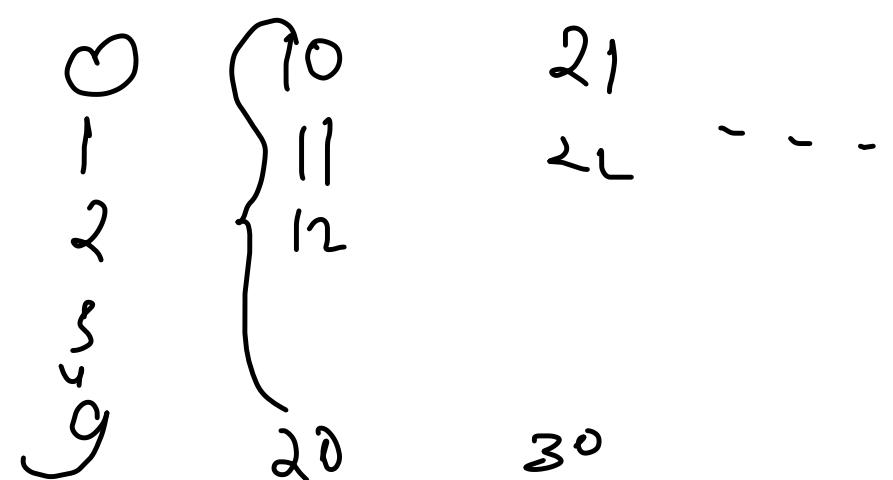
    return freq;
}

```

## Number system

~~#~~ Decimal No

0-9  
10 digits



$$10 = 1 * 10^1 + 0 * 10^0$$

$$100 = 99$$

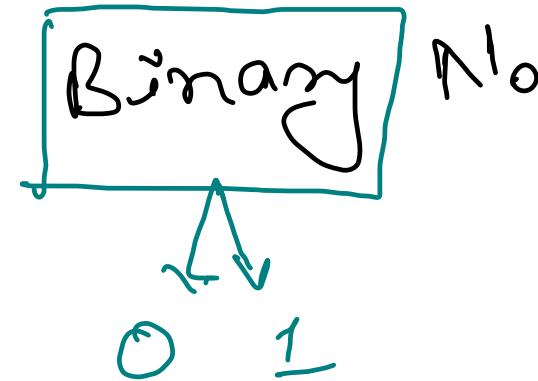
$$100 =$$

25 16.93

$$\begin{aligned}
 &= 6 * 10^2 + 9 * 10^1 + 3 * 10^0 \\
 &\quad + 1 * 10^3 + 5 * 10^4 \\
 &\quad + 2 * 1.5
 \end{aligned}$$

min no 0000  
 max no 9999  
 count of nos  $10^4$

#



0	00
1	01
	10
	11

000
001
010
011
100
101
110
111

$2^0$        $2^1$

$$11001011 = 1 * 2^0 + 1 * 2^1 + 0 * 2^2 + 1 * 2^3 + 0 * 2^4 + \dots$$

min 0000  
 $2 * 2 * 2 * 2$   
0/1 0/1 0/1 0/1 max 1111  
Count  $2^4$

## Octal Nb System

↳ 0, 1, 2, 3, 4, 5, 6, 7

$(7\ 5\ 3\ 4\ 6\ 1\ 2)_8$

$$= 2*8^0 + 1*8^1 + 7*8^2 \\ + 4*8^3 + 3*8^4 + 5*8^5 + 7*8^6$$

6 month +  
↳ electrode L2  
Octal

$8*8*8*8*8$   
 $\underline{d_7\ d_6\ d_5\ d_4}$  ↳ min  $\Rightarrow 0000$   
max  $\Rightarrow 7777$   
total  $\Rightarrow 8^4$

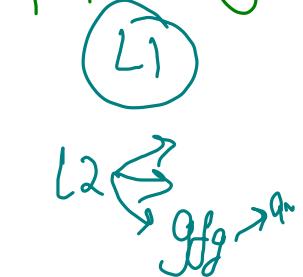
0	10	20	30	40	50	60
1	11	21	31	41	51	61
2	12	22	32	42	52	62
3	13	23	33	43	53	63
4	14	24	34	44	54	64
5	15	25	35	45	55	65
6	16	26	36	46	56	66
7	17	27	37	47	57	67

70	71	72	73	74	75	76	77	78	79
100	101	102	103	104	105	106	107	108	109
777	1000	7777	10000	77777	100000	777777	1000000	7777777	10000000
100000000	777777777	1000000000	7777777777	10000000000	77777777777	100000000000	777777777777	1000000000000	7777777777777
7777777777777	100000000000000	77777777777777	1000000000000000	777777777777777	10000000000000000	7777777777777777	100000000000000000	77777777777777777	1000000000000000000

$$(777)_8 = 1060 \\ 7*8^2 + 7*8^1 + 7*8^0 \\ \underline{1060} \\ 1*8^3$$

$$(gg)_10 = 1060 \\ 9*10^2 + 9*10^1 + 9*10^0 \\ \underline{1060}$$

preceding



service based

tier 3 ↲

tier 1 (3<sup>rd</sup>, 4<sup>th</sup>)

DSA

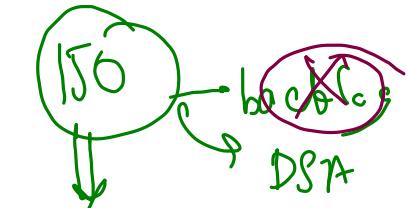


2<sup>nd</sup> Dev

DSA { 600 f lected }  
Dev { 50 + Proj }

2-3 hr

8 hr



80 - 90  
↓  
40 - 50

# Today's Target

- </> Decimal To Any Base ] → classwork
- </> Any Base To Decimal ]
- </> Any Base To Any Base ] → homework
- </> Any Base Addition ] → classwork
- </> Any Base Subtraction ]
- </> Any Base Multiplication ] → homework

10:20

$$\begin{array}{r} \text{dividend} \\ \downarrow \\ (\underline{\underline{634}})_{10} \end{array} \quad \begin{array}{l} \text{Decimal to Any Base} \\ = \\ (\underline{\underline{?}})_8 \end{array}$$

$$\begin{array}{r}
 \begin{array}{c} b \\ \xrightarrow{\quad} 8 \\ \xrightarrow{\quad} 8 \end{array} \left| \begin{array}{c} x \\ 634 \\ \hline 7g \end{array} \right. \begin{array}{l} \text{multiplier} \\ \hline \end{array} \end{array}$$

7g - 2 \* 10^1 = 1  
 9 - 7 \* 10^1 = 2  
 1 - 1 \* 10^2 = 2  
 0 - 1 \* 10^3 = 2

$\{0+2\} + 70\} + 100$   
 $= 1172$   
 $(1172)_g$

$$(634)_{10}$$

$$6 \cdot 10^2 + 3 \cdot 10^1 + 4 \cdot 10^0$$

- What ?
- how ?
- why ?

$$79 = 9 \cancel{4} 8 + 7$$

00 + 100 0

$$= 1172$$

```
public static int getValueInBase(int n, int b){  
    int res = 0;  
    int multiplier = 1;  
  
    while(n > 0){  
        int divisor = n / b;  
        int remainder = n % b;  
  
        n = divisor;  
        res = res + remainder * multiplier;  
        multiplier *= 10;  
    }  
  
    return res;  
}
```

(1172) 8

$$1 * 8^3 + 1 * 8^2 + 7 * 8^1 + 2 * 8^0$$

$$\begin{aligned} (634)_{10} &= 8^4 * (-) + 2 \\ &= 8^4 * 0 + 8^3 * 1 + 8^2 * 1 + 8^1 * 7 + 8^0 * 2 \\ &\quad \square \quad \square \quad \square \quad \square \quad \square \quad (8^2 * 0 + 8^1 * 1 + 8^0 * 1) \end{aligned}$$

$$f_9 = g^* \left( \frac{g_{10}}{g_{11}} \right)$$

$$= \left( g^3 + 0f g^2 + 1f g^1 + 1f g^0 \right)$$

$$g = g' + (g' \times 0 + 1) + 1$$

$$g = g^2 * 0 + g^1 * 1 + g^0 * 1$$

$$(173)_{10} = (?)_8$$

The diagram illustrates the conversion of the decimal number 173 into an octal number. It shows a series of horizontal divisions by 8, starting from the remainder 1 at the top.

- Step 1:  $1 \div 8$  remainder  $1$ . This is labeled  $173$  above the first line.
- Step 2:  $21 \div 8$  remainder  $5$ . This is labeled  $21 - 5 \cancel{*} 10^0$ .
- Step 3:  $2 \div 8$  remainder  $2$ . This is labeled  $2 - 2 \cancel{*} 10^1$ .
- Step 4:  $0 \div 8$  remainder  $0$ . This is labeled  $0 - 0 \cancel{*} 10^2$ .

The remainders 1, 5, 2, and 0 are connected by vertical lines to form the octal number  $(255)_8$ .

$$5 + 50 + 200 = (255)_8$$

# Any Base to Decimal

$$(1172)_8 = (?)_{10}$$

1172  
10  
117 - 2 \* 8<sup>0</sup>  
10  
11 - 7 \* 8<sup>1</sup>  
10  
1 - 1 \* 8<sup>2</sup>  
0 - 1 \* 8<sup>3</sup>

$$0 \cdot 2 * 8^0 + 7 * 8^1 + 1 * 8^2 + 1 * 8^3 = (634)_{10}$$

```
public static int getValueInDecimal(int n, int b){  
    int res = 0;  
    int multiplier = 1;  
  
    while(n > 0){  
  
        int divisor = n / 10;  
        int remainder = n % 10;  
  
        n = divisor;  
  
        res = res + remainder * multiplier;  
        multiplier *= b;  
    }  
  
    return res;  
}
```

$$(110010)_2 = (?)_{10}$$

10	110010
10	11001 - 0
10	1100 - 1
10	110 - 0
10	11 - 0
10	1 - 1
10	0 - 1

\*  $2^0$   
 \*  $2^1$   
 \*  $2^2$   
 \*  $2^3$   
 \*  $2^4$   
 \*  $2^5$

$$\begin{aligned}
 &= 0 + 2 + 0 + 0 + 16 + 32 \\
 &\quad = 18 + 32 \\
 &\quad = (50)_{10}
 \end{aligned}$$

# Any base to any base

$$(172)_8 = (?)_2$$

Any base to decimal

$$(?)_{10}$$

Decimal to Any base

# Any Base addition

Carry 1

$$\begin{array}{r}
 & 0 & 1 \\
 & \cancel{4} & \cancel{3} \\
 + & (\cancel{7} & \cancel{5})_{10} \\
 \hline
 & 1 & 9 & 0
 \end{array}$$

$$\begin{aligned}
 \text{temp} &= (0 + 6 + 4) \\
 &\xrightarrow[10]{10} 1 \quad \{ \text{quot} \} \\
 &\xrightarrow[10]{10} 0 \quad \{ \text{rem} \}
 \end{aligned}$$

$$1 + 3 + 5 = \frac{9}{10} \quad 0$$

$$0 + 4 + 7 = \frac{11}{10} \quad 1$$

Carry

$$\begin{array}{r}
 1 \\
 0 \\
 0
 \end{array}
 \begin{array}{r}
 1 \\
 3 \\
 5
 \end{array}
 \begin{array}{r}
 0 \\
 8 \\
 8
 \end{array}$$


---


$$\begin{array}{r}
 1 \\
 2 \\
 1 \\
 2
 \end{array}$$

$$236 \rightarrow 23$$

$$236/10 = 23$$

①<sup>st</sup>

$$c + d_1 + d_2$$

$$= 0 + 0 + 4$$

1 (quot)

2 (rem)

②<sup>nd</sup>

$$1 + 3 + 5$$

$$= \frac{9}{8}$$

③<sup>rd</sup>

$$1 + 2 + 7 = \frac{10}{8}$$

④\*

$$1 + 0 + 0 = \frac{1}{8}$$

$$\begin{aligned}
 1212 &= 1000 + 200 \\
 &\quad + 10 + 2
 \end{aligned}$$

$$\left\{ 2 \times 10^0 + 1 \times 10^1 + 2 \times 10^2 \right.$$

$$\left. + 1 \times 10^3 \right\}$$

base = 8

```
public static int getSum(int b, int n1, int n2){  
    int res = 0;  
    int multiplier = 1;  
    int carry = 0;  
  
    while(carry > 0 || n1 > 0 || n2 > 0){  
  
        int d1 = n1 % 10;  
        int d2 = n2 % 10;  
  
        int temp = carry + d1 + d2;  
        int quot = temp / b;  
        int rem = temp % b;  
  
        res = res + (rem * multiplier);  
        carry = quot;  
  
        multiplier *= 10;  
        n1 /= 10;  
        n2 /= 10;  
    }  
  
    return res;  
}
```

0 1 1 1 0  
0 0 ~~4~~ ~~6~~ ~~3~~  
0 0 ~~7~~ ~~2~~ ~~8~~  
-----  
res → 1 4 9 3  
-----

$$4 \times 100 + 93$$

$$\frac{1}{8} \leftrightarrow 0$$

$$\frac{12}{8} \leftrightarrow 4$$

[1000]  
mult

$$0+8+3=11$$



$$\begin{array}{r} 1 \\ 3 \\ \hline 9 \end{array}$$

## Pseudo Code

carry = 0

res = 0

multiplier = 1

while( $\frac{carry}{10} \parallel n_1 \parallel n_2 > 0$ ) {

$d_1 = n_1 \% 10$

$d_2 = n_2 \% 10$

temp = c + d1 + d2

quot = temp / base

rem = temp % base

carry = quot

res += temp \* multiplier

multiplier \*= 10

$n_1 /= 10$

$n_2 /= 10$

}

quotient will become  
next carry

88

rem will be  
Contributing in  
result.

# Any base subtraction

borrow

$$\begin{array}{r}
 & -1 & -1 & -1 \\
 & 1 & 2 & 1 & 2)_{10} \\
 - & 2 & 5 & 6)_{10} \\
 \hline
 0 & 9 & 5 & 6
 \end{array}$$

$$d_1 - d_2 + \text{borrow}$$

$$2 - 6 + 0 =$$

$$\circled{-4}$$

$$\begin{aligned}
 2 - 6 &= -4 + 10 \\
 &\Rightarrow \circled{6}
 \end{aligned}$$

951

$$\begin{aligned}
 2 - 2 + (-1) + 10 \\
 &= -1 + 10 = 9
 \end{aligned}$$

$$\begin{aligned}
 1 - 5 + (-1) \\
 &= -4 - 1 = \circled{-5} + 10 \\
 &= 8
 \end{aligned}$$

# Any base subtraction

borrow

$$\begin{array}{r}
 & -1 & -1 \\
 & 1 & 2 & 1 \\
 & & 2 & 5 & 6 \\
 - & & & & \\
 \hline
 & 0 & 9 & 5 & 6
 \end{array}$$

$-4 + 0$

$\boxed{Q - d_1 + \text{borrow}}$

$$2 - 6 + 0 = -4 + 10$$

$\approx 6$

$$\begin{aligned}
 2 - 6 &= -4 + 10 \\
 &\approx 6
 \end{aligned}$$

951

$$\begin{aligned}
 2 - 2 + (-1) + 10 \\
 &= -1 + 10 = 9
 \end{aligned}$$

$$\begin{aligned}
 1 - 5 + (-1) \\
 &= -4 - 1 = -5 + 10 \\
 &= 8
 \end{aligned}$$

## Pseudo code

```

borrow = 0      res = 0      multiplier = 1
while( n2 > 0) {
    d1 = n1 / 10, d2 = n2 / 10
    temp = d2 - d1 + borrow
    if(temp < 0) {
        res += (temp + base) * multiplier
        borrow = -1;
    } else {
        res += temp * multiplier;
        borrow = 0;
    }
    multiplier *= 10;
    n1 /= 10;   n2 /= 10;
}
  
```

$$n_2 > n_1$$

$$\begin{array}{r}
 0 \quad -1 \quad -1 \quad 0 \\
 (1) \quad 4 \quad 1 \quad 8 \\
 (2) \quad 5 \quad 6 \quad 8 \\
 \hline
 1 \quad | \quad 3 \quad 4
 \end{array}$$

$$\begin{aligned}
 & 1 - 0 + 0 \\
 & = 1
 \end{aligned}$$

$$\begin{aligned}
 & 4 - 2 + (-1) \\
 & = 2 - 1 = 1
 \end{aligned}$$

$$\begin{aligned}
 & 1 - 5 + (-1) \\
 & -4 - 1 = -5 + 8 \\
 & = 3
 \end{aligned}$$

# Any base Multiplication

$$1 \left( \begin{array}{cccc} 1 \\ 2 & 3 & 4 & 0 \\ \hline 1 & 5 & 6 \end{array} \right)_{10}$$

$$\times \left( \begin{array}{cc} 7 & 4 \end{array} \right)_{10}$$


---

$$\left\{ \begin{array}{cccc} 18 & 6 & 2 & 4 \\ \hline 15 & 0 & 9 & 2 \end{array} \right.$$

~~21 \* 6 \* 4~~

~~2156 \* 70~~

$$\left\{ \begin{array}{cccc} 18 & 6 & 2 & 4 \\ \hline 15 & 0 & 9 & 2 \end{array} \right. \times$$

$$\underline{15 \quad 9 \quad 5 \quad 4 \quad 4}$$

Any base  
Addition

$$7*2+1 = \frac{15}{10} \rightarrow 5$$

$$\frac{24}{10} \rightarrow 2$$

$$7*6 = \frac{42}{10} \rightarrow 4$$

$$\frac{39}{10} \rightarrow 3$$

$$5*4+2 = \frac{22}{10} \rightarrow 2$$

$$4*1+2 = \frac{6}{10} \rightarrow 0$$

$$\frac{8}{10} \rightarrow 0$$

$$\frac{10}{10} \rightarrow 1$$

Algorithm

while {  $n_1 > 0$  } {

- ① Extract digits of  $n_1$  :  $d_1$
- ② multiply  $d_1$  with  $n_2$
- ③  $res = \text{anybaseAddition}(res, \text{temp} * \text{multiplier})$   
multiplier  $\times = 10$

getProductWithDigit()

*pencil  
eraser*

{

$$\begin{array}{r}
 & 1 & 1 & 5 & 5 & 6) \\
 & (2 & 1 & 5 & 7 & 8 \\
 & \times & (7 & ) & 4) \\
 \hline
 & 1 & 0 & 6 & 7 & 0) \\
 & (1 & 7 & 4 & 0 & 2 \\
 \hline
 & 2 & 0 & 4 & 7 & 1 & 0
 \end{array}$$

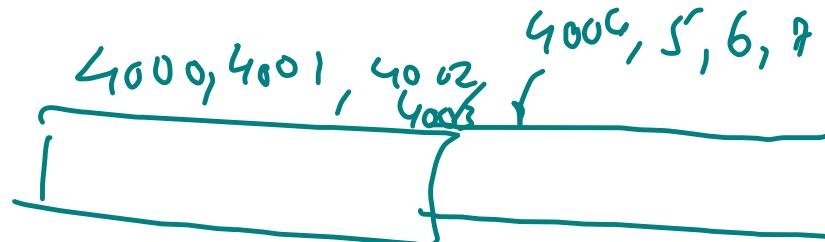
$\frac{2156 * 74}{2156 * 4}$

$$\begin{array}{r}
 = 2156 * 4 + \textcircled{+} \\
 \text{total}
 \end{array}$$

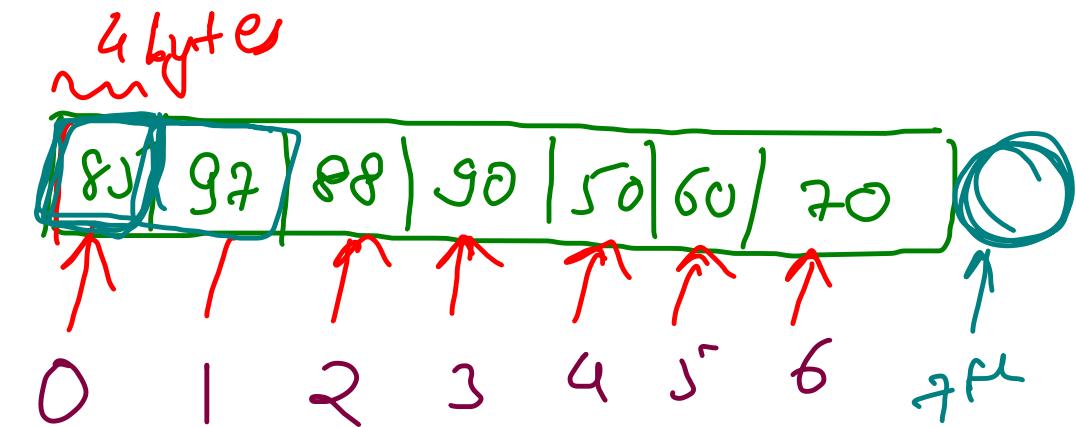
$$\begin{array}{r}
 \frac{24}{8} \rightarrow 3 \\
 20 + 3 = \frac{23}{8} \rightarrow 2 \\
 \frac{6}{8} \rightarrow 0 \\
 \frac{40}{8} \rightarrow 5 \\
 \frac{12}{8} \rightarrow 1 \\
 \frac{8}{8} \rightarrow 0 \\
 \frac{15}{8} \rightarrow 1
 \end{array}$$

## Data Structure

int m1 = 95;  
int m2 = 97;  
int m3 = 82;



## Arrays



contiguous  
memory  
allocation

7 size  
0 - 6

$$7 \times 4 = 28 \text{ bytes}$$

## Memory mapping of arrays

```
int[] marks = new int[n];
```

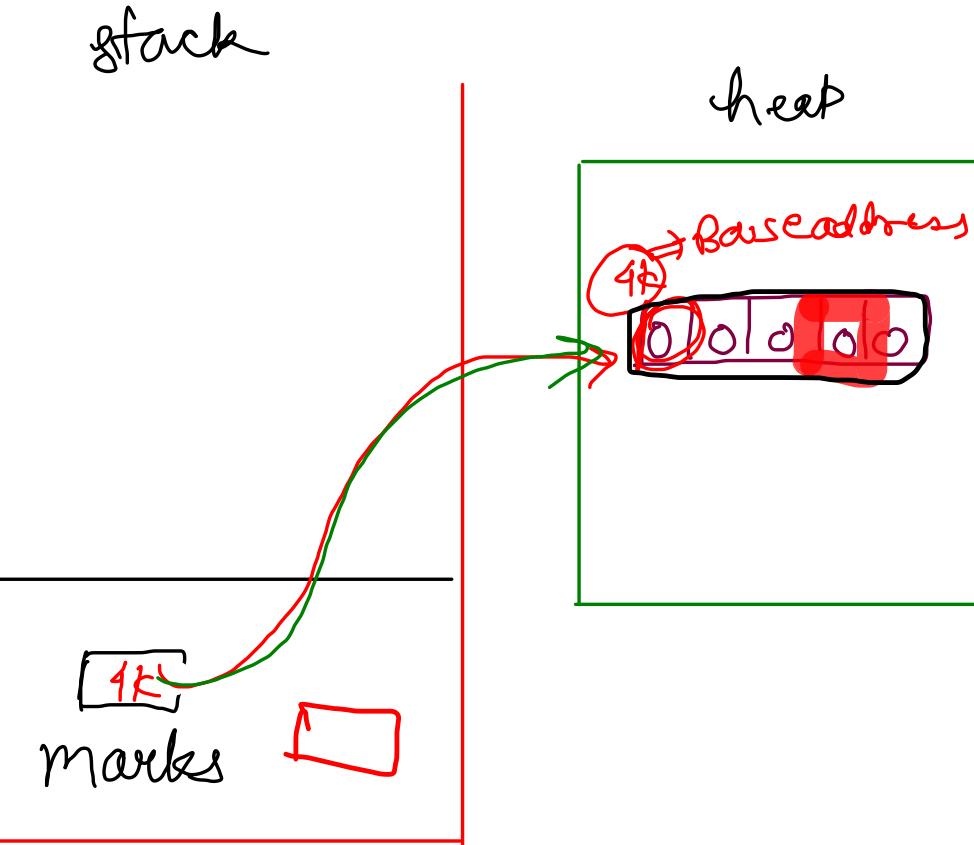
{Static array}

marks[1]

base address +  $i^{th}$  block  
of array  
 $\times$   
4 bytes

Random access {constant operation}

$\Rightarrow$  BA + index \* size of one block



stack  
primitive

referencing  
if

heap  
array  
objects

marks [0]

$$4k + 0 \times 4 = 4k$$

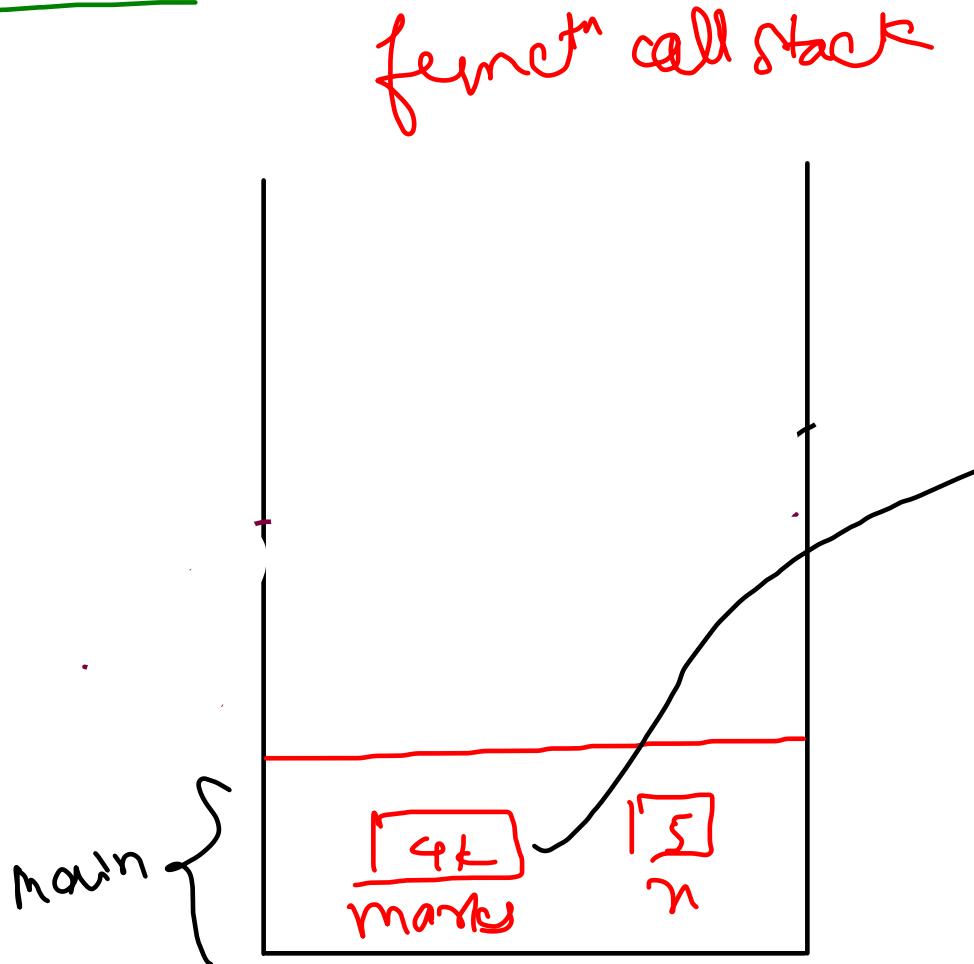
marks [3]

$$4k + 3 \times 4 = 4k12$$

## Passing array to function

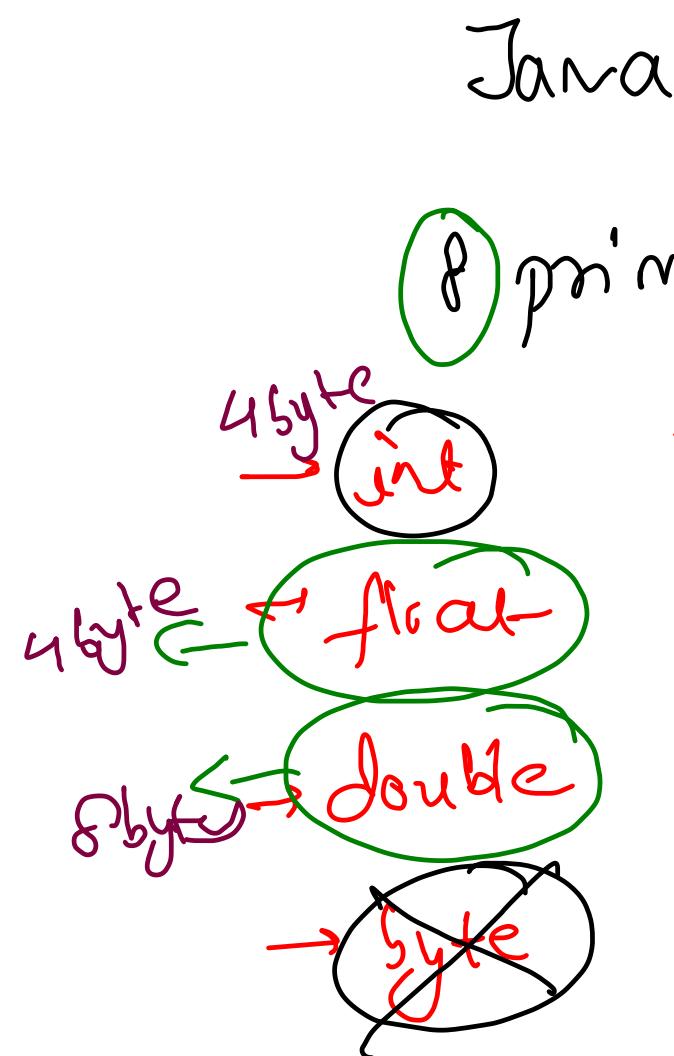
"Changes in the heap always persist"  
Shallow copy permanent

```
public static void display(int[] marks){  
    for(int i=0; i<marks.length; i++){  
        System.out.print(marks[i] + " ");  
    }  
  
}  
  
public static void main(String[] args) {  
    Scanner scn = new Scanner(System.in);  
  
    int n = scn.nextInt();  
  
    // Declaration & Initialization  
    int[] marks = new int[n];  
  
    // Random Access using Indexing  
    for(int i=0; i<n; i++){  
        marks[i] = scn.nextInt();  
    }  
  
    // Array Traversal  
    display(marks);  
}
```

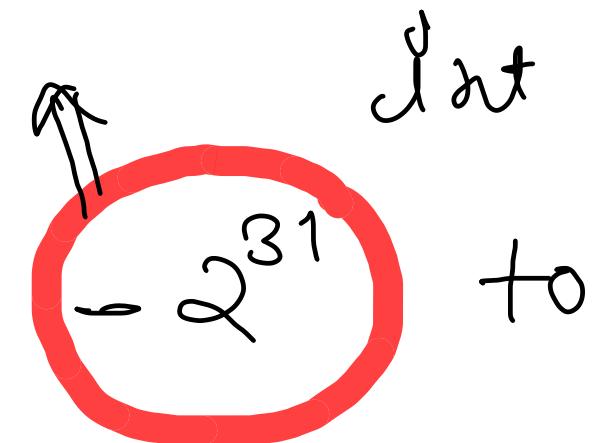


"Java is always

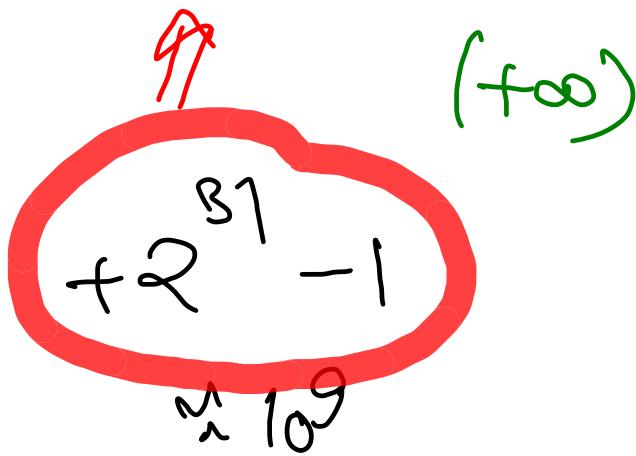
pass by  
value  
Shallow  
Copy



(-2<sup>31</sup>) Integer.MIN-VALUE



Integer.MAX-VALUE



long



$2^{63} - 1$

$\downarrow$								
63	78	32	46	26	92	16	104	2

$$\text{Span} = \underline{\max^m} - \underline{\min^m}$$

$$\max^m = -\infty \cancel{63} \cancel{78} \cancel{92} 104$$

$$\min^m = +\infty \cancel{63} \cancel{32} \cancel{26} 162$$

```

int min = Integer.MAX_VALUE;
int max = Integer.MIN_VALUE;

for(int i=0; i<n; i++){

    // is current element less than min of previous
    if(arr[i] < min){
        min = arr[i];
    }

    // is current ele greater than max
    if(arr[i] > max){
        max = arr[i];
    }
}

int span = max - min;
System.out.println(span);

```

$\max(x, -\infty) = x$   
 $\min(x, +\infty) = x$

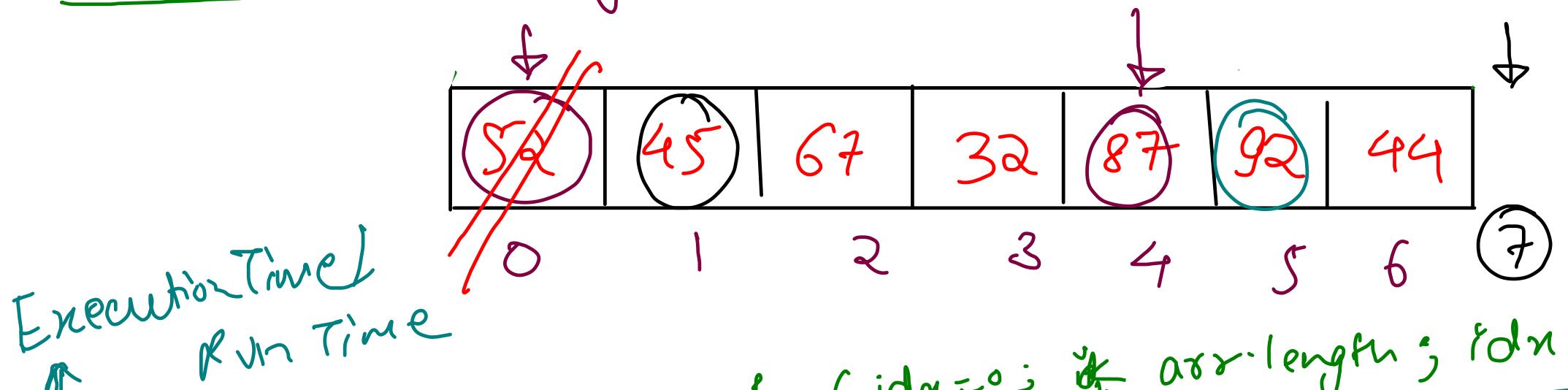
$x + 0 = x$

## Today's Questions

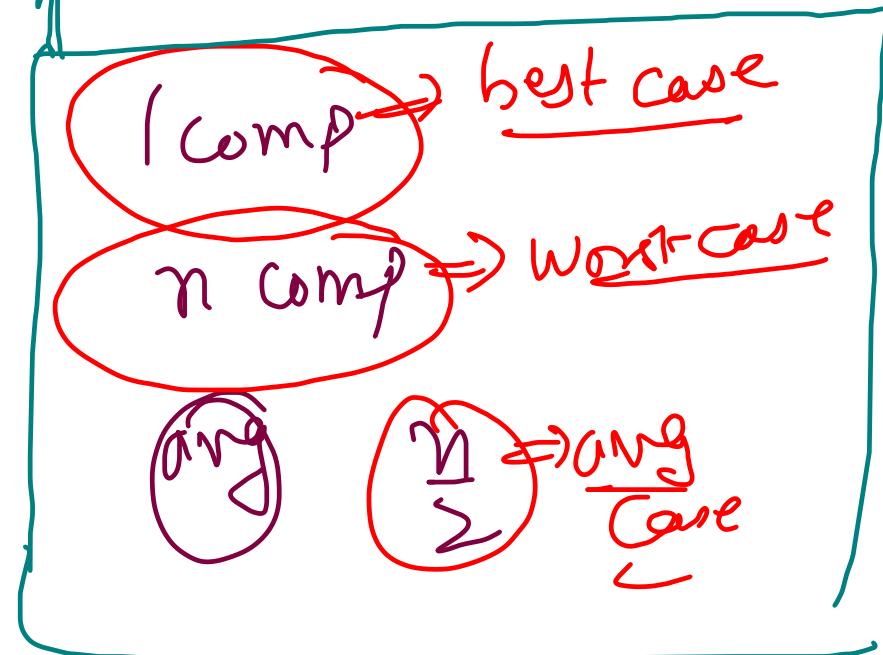
- </> Find Element In An Array
- </> Bar Chart
- </> Sum Of Two Arrays
- </> Difference Of Two Arrays
- </> Reverse An Array
- </> Rotate An Array

→ Linear Search Algorithm

## Linear Search Algorithm



Execution Time  
Run Time

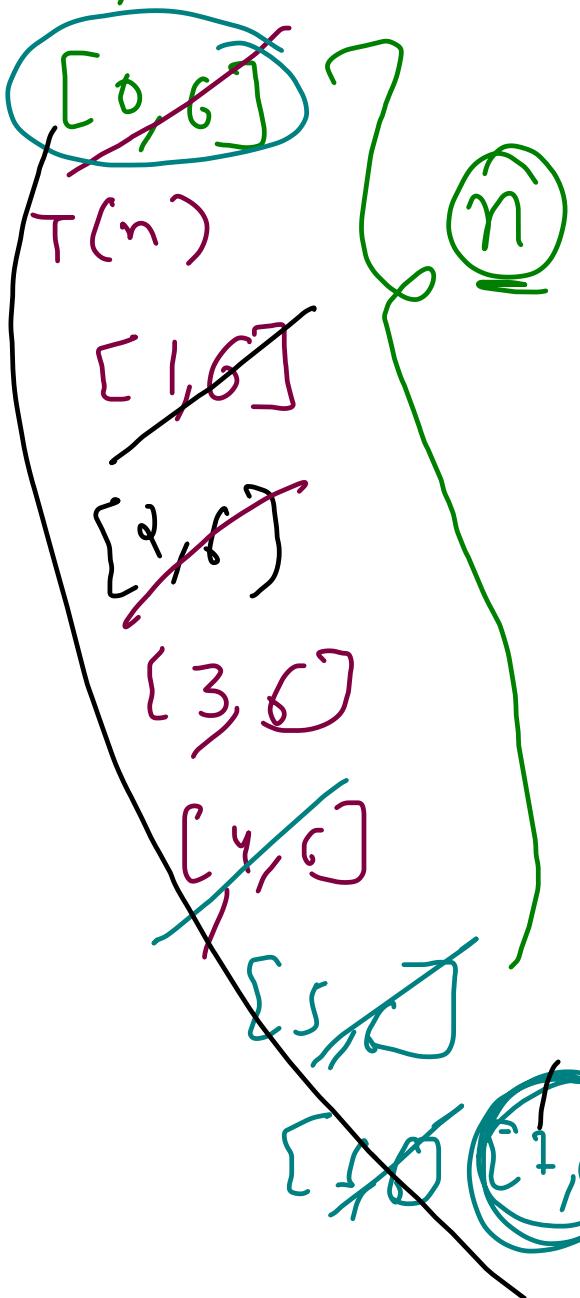


```
for( idx=0; idx < arr.length; idx++ ) {
    if( arr[idx] == ele ) {
        successful
        return idx;
    }
}
return -1;
            unsuccessful
            search

```

ele = ~~87~~ ~~45~~ ~~67~~ ~~32~~

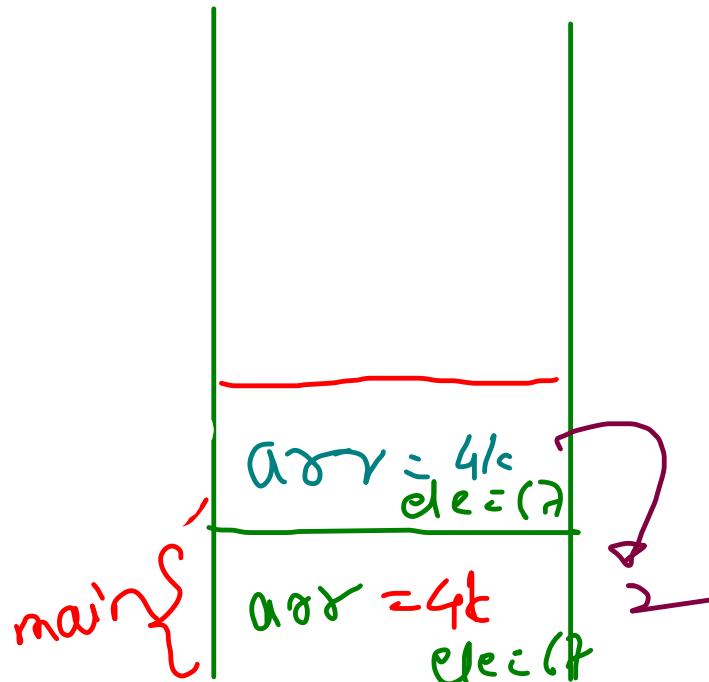
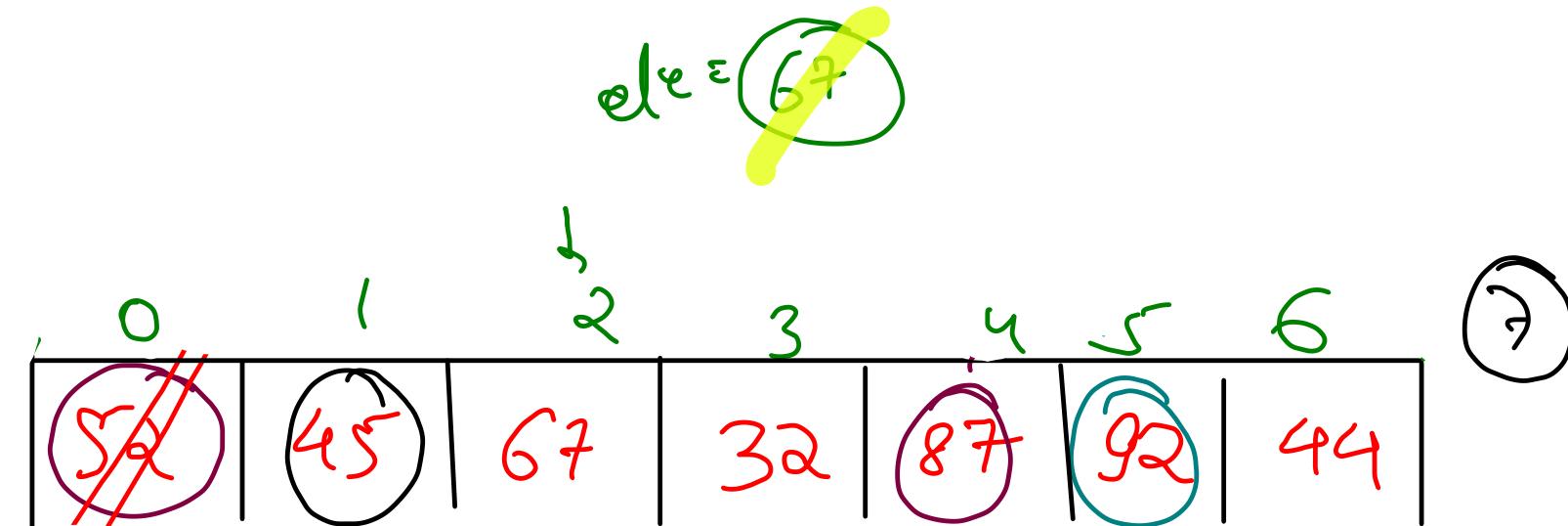
search space



```

public static int linearSearch(int[] arr, int ele) {
    for (int idx = 0; idx < arr.length; idx++) {
        if (arr[idx] == ele) {
            // successful search
            return idx;
        }
    }
    return -1; // unsuccessful search (element not found)
}

```



```

public static void main(String[] args) throws Exception {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();

    int[] arr = new int[n];

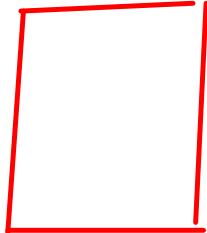
    for (int i = 0; i < n; i++) {
        arr[i] = scn.nextInt();
    }

    int ele = scn.nextInt();

    int ans = linearSearch(arr, ele);
    System.out.println(ans);
}

```

## Notes

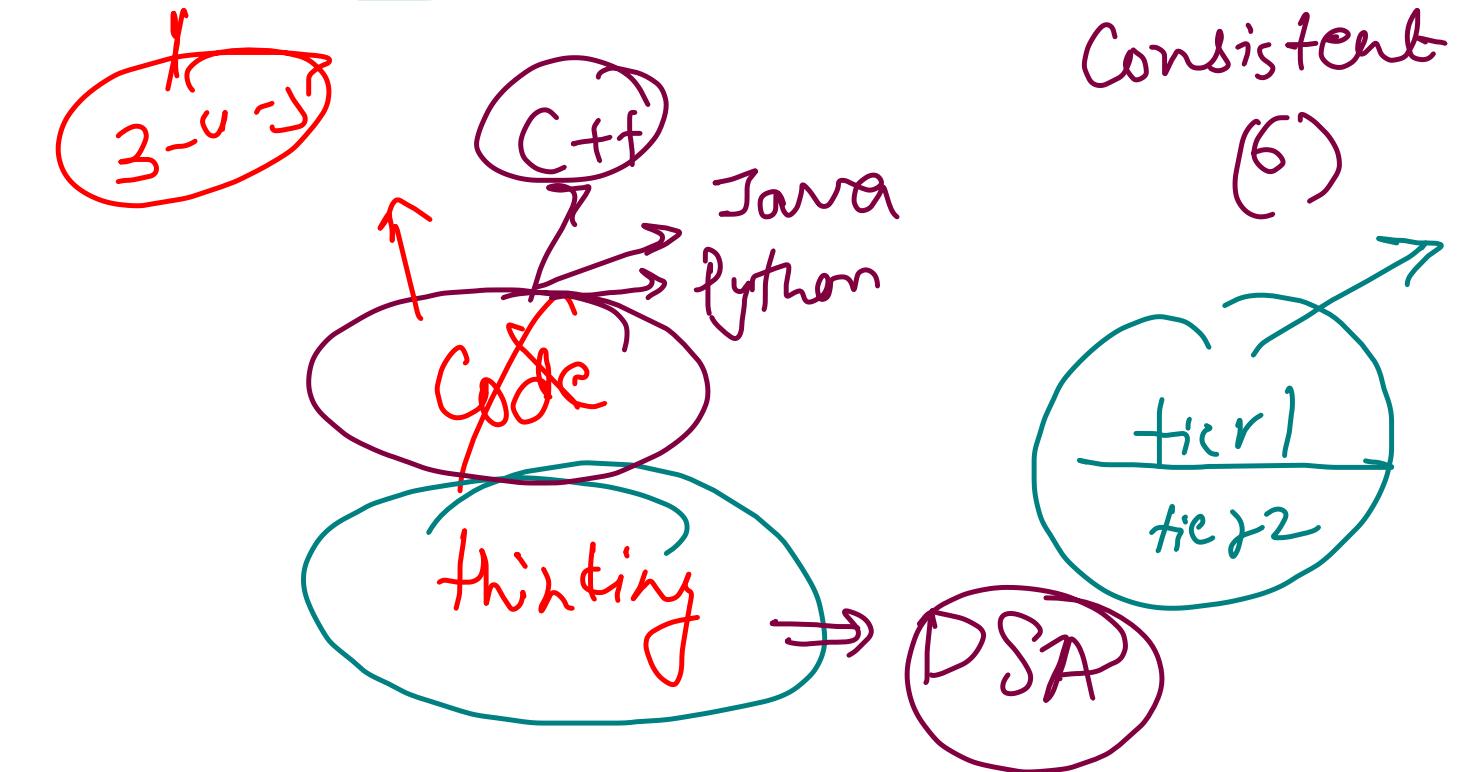
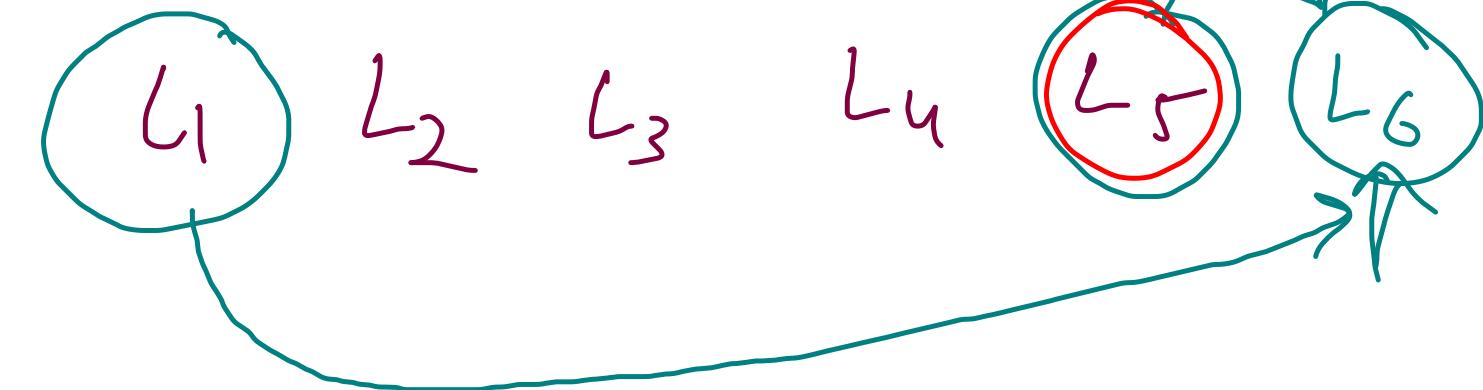


Question name  
1 liner logic

Important points



## Revision



Consistent  
(6)

DSA + Dev + Core

Hypercoding  
leetcode  
level 2  
MB  
Virtual context

SB-C.  
(3-6)

Aptitude  
Level 1

20-30%

Small Startups  
(6-12)

Level 1

Dev → 50-60%  
Core

big startups  
(12-30)

Level 1

+2  
Dev → 30-40%  
Core

70-80

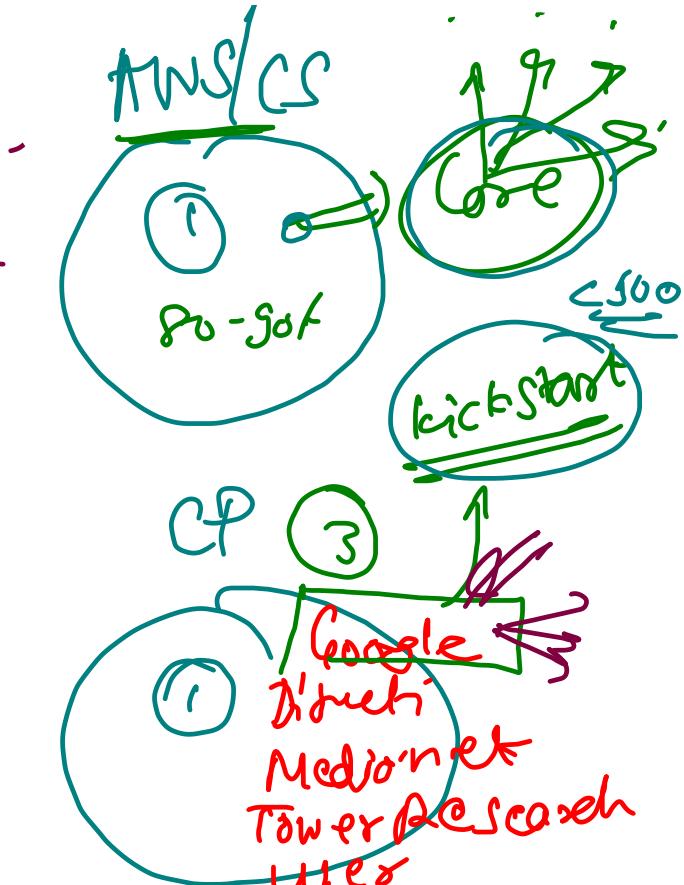
PBC  
SOF

1  
+2

new  
Core

tech  
giant

{FAANG}



10-30

Codechef  
Bit, stric etc  
DP/graphs

Tree, LL, STQ  
leetcode

fin-tech  
Aptitude  
Maths  
level 1/2  
Core

# Bar Chart

{ 2, 1, 0, 2, 5 }

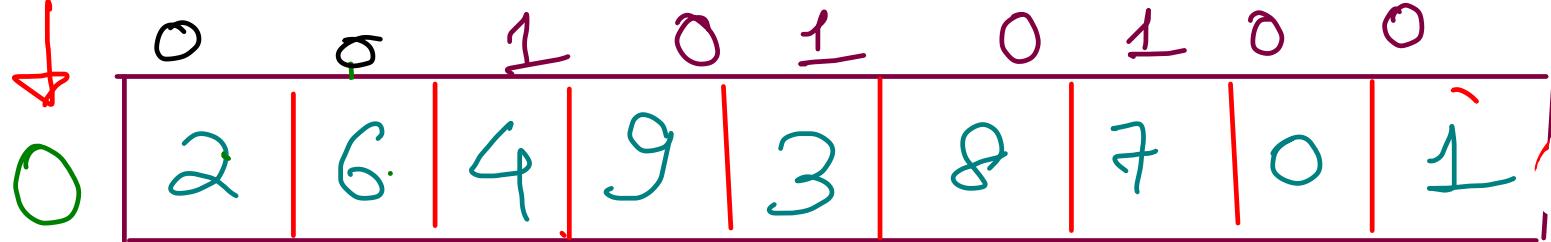
```
int totalCols = n;
int totalRows = maxEle(arr);
for(int i=totalRows - 1; i>=0; i--){
    for(int j=0; j<totalCols; j++){
        if(*arr[j] >= i)
            System.out.print("*\t");
        else {
            System.out.print("\t");
        }
    }
    System.out.println();
}
```

6	5	4	3	2	1
*	*	*	*	*	
*	*	*	*	*	
*	*	*	*	*	
*	*	*	*	*	
*	*	*	*	*	
*	*	*	*	*	

$i \geq arr[j]$  : Space  
else  $i < arr[j]$  : Star

# Sum of two arrays

$$i1 = n1 - 1$$



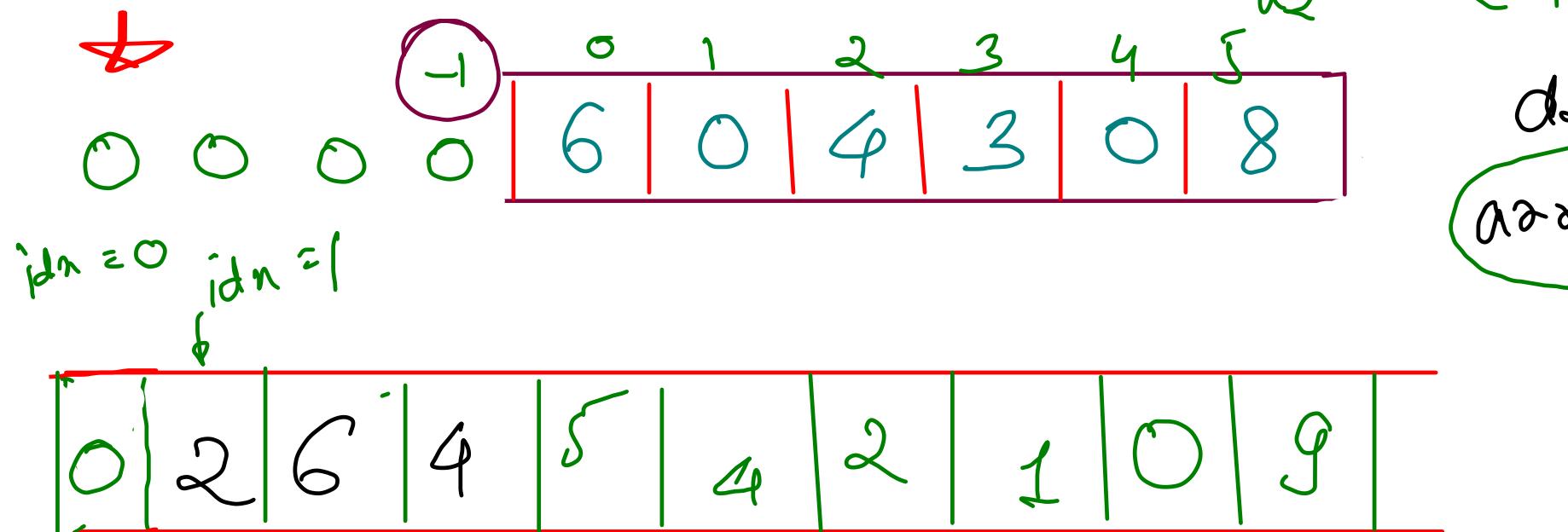
$$d_1 = \text{arr1}[i1]$$

long {8 bytes}

long → 2<sup>63</sup> - 1

→ 10<sup>18</sup>

$$\text{carry} = 0$$



$$d_2 = \text{arr2}[i2]$$

$$\text{arr2}[i2]$$

Index out of  
bound exception

> 10<sup>18</sup>  
→ Arrays  
BigIntegers

factorial of  
c n'

$$n = \underline{\underline{10^9}}$$

```
int idx = 0;
while(res[idx] == 0){
    idx++;
}

for(int i=idx; i<res.length; i++){
    System.out.println(res[i]);
}
```

$$i) = n3 - 1$$

$$\text{max}(n1, n2) + 1$$

$$\text{size} =$$

## Pseudo Code

```
synt []res = new int { man(n1, n2) + r ]
```

carry = 0      i1 = n1 - 1, i2 = n2 - 1, i3 = n3 - 1

while (      i3 ≥ 0      ) {

[  
d1 = arr1[i1]  
d2 = arr2[i2]

temp = c + d1 + d2

quot = temp / 10

rem = temp % 10

carry = quot

res[i3] = rem;

i1--; i2--; i3--

}

if (i1 < 0)      d1 = 0

else

d1 = arr1[i1]

if (i2 < 0)      d2 = 0

else      d2 = arr2[i2]