

QUERY

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Query

$\pi_{e.name} (\sigma_{e.salary > m.salary} (\rho_e(employee) \bowtie_{e.manager = m.name} \rho_m(employee)))$

(a) Relational Algebra version

RANGE employee e;

RANGE employee m;

GET w (e.name): $\exists m((e.manager = m.name) \wedge (e.salary > m.salary))$

(b) Relational Calculus version

select e.name

from employee e, employee m

where e.manager = m.name and e.salary > m.salary

(c) Sequel (SQL) version



Figure 2. Three versions of the query, “Find names of employees who earn more than their managers.”

SQL É UBIQUO



Query

- ▶ ORMs
- ▶ Schemas
- ▶ Controle de Concorrência
- ▶ Manipulação de dados



Query

- ▶ 5 abordagens
 - ▶ NOMP
 - ▶ MapReduce
 - ▶ Abordagens específica do SGBD
 - ▶ Padronização
 - ▶ Derivações do SQL



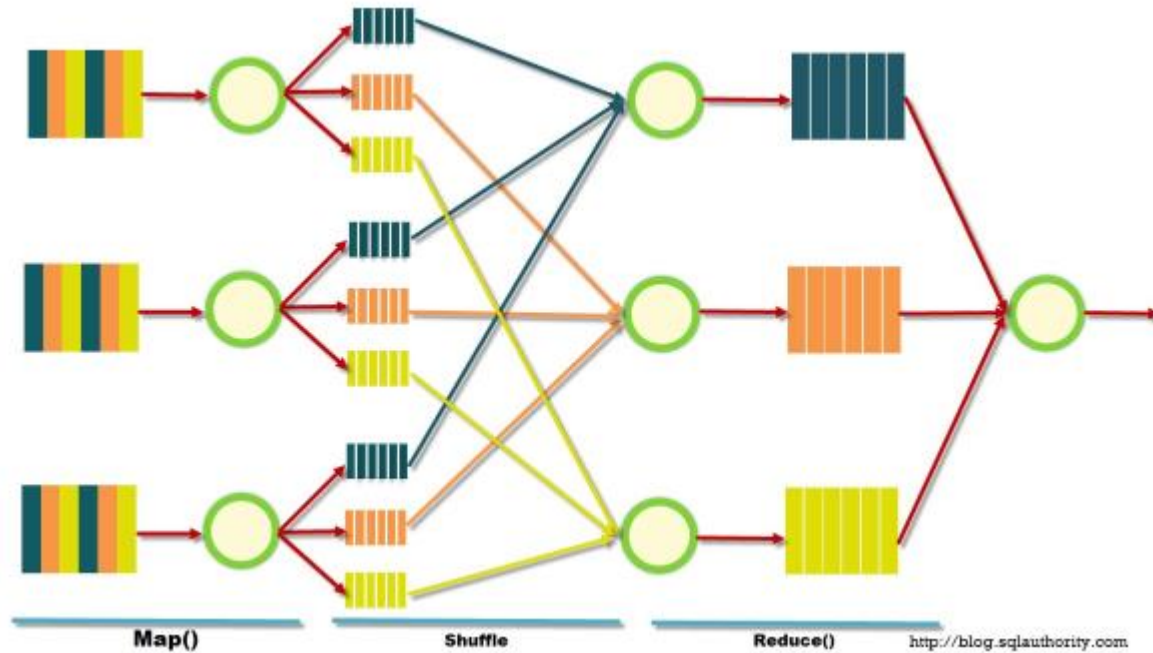
Query

- ▶ Not is My Problem
 - ▶ Key-Value
 - ▶ Schema free
 - ▶ Rest API

**HI
I DON'T CARE
THANKS**

Query

► MapReduce



Query

► SGBD

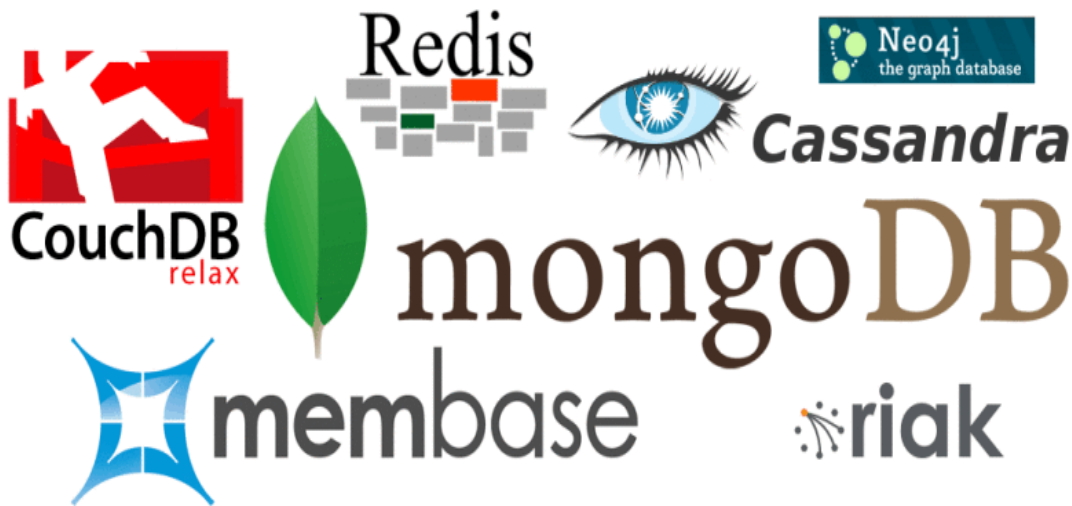
```
db.collection.command(data)
```

```
db.restaurants.find( { "cuisine": "Fish and chips" } )
```



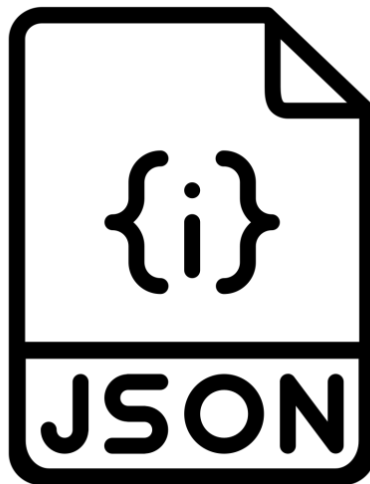
Query

► SGBD



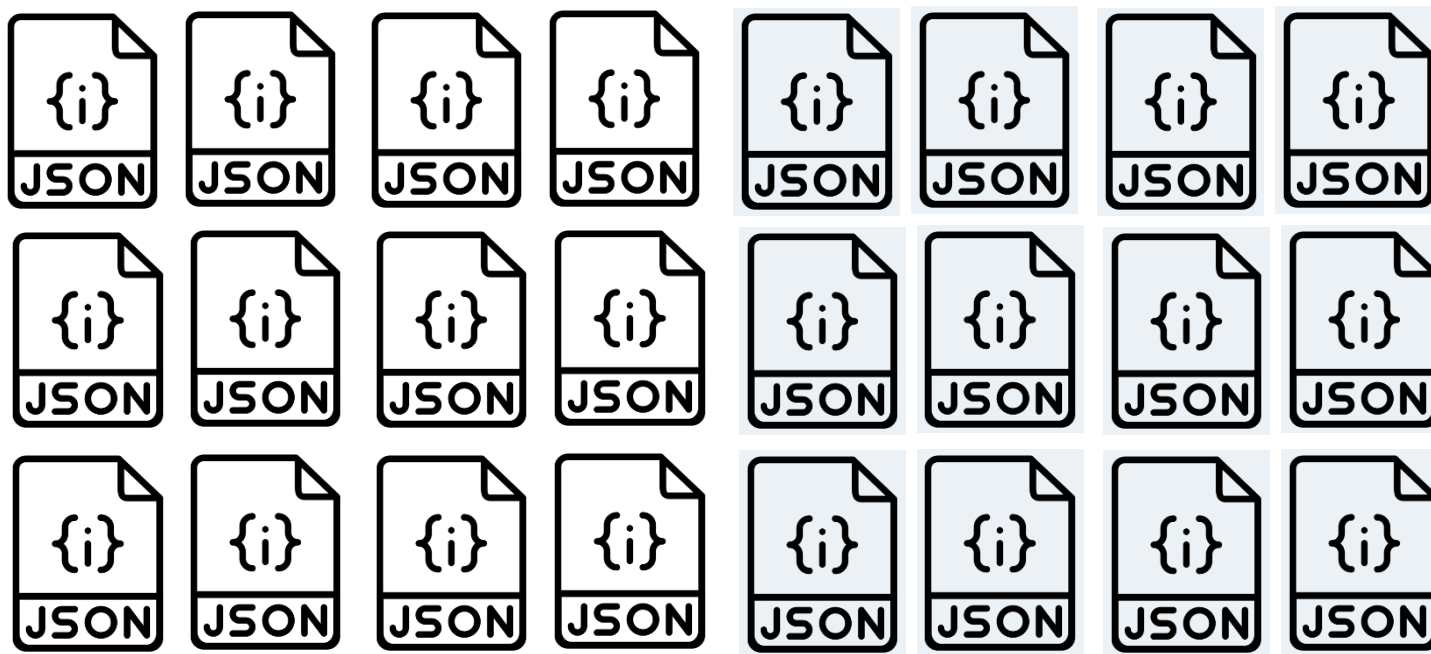
Query

► Padronização



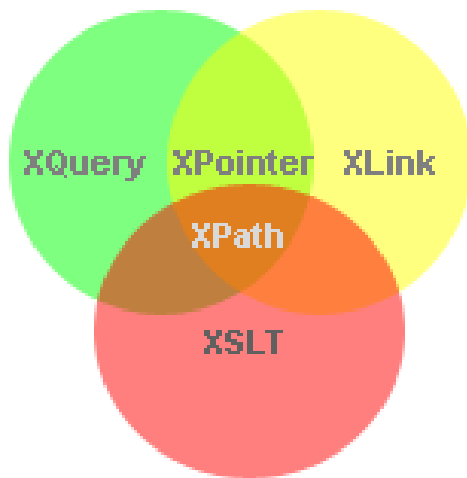
Query

► Padronização



Query

► Padronização



Query

► Padronização

```
xquery version "1.0-ml";
declare namespace rest =
"http://example.com/namespace/restaurants";
{
  for $restaurant in
  doc("restaurants.xml")/rest:restaurants/rest:restaurant
  return
  {
    $restaurant/rest:name/text()
  }
}
```

Query

► Derivações do SQL

$\pi_{e.name} (\sigma_{e.salary > m.salary} (\rho_e(employee) \bowtie_{e.manager = m.name} \rho_m(employee)))$

► CQL

```
SELECT name FROM restaurants WHERE cuisine='Fish and chips';
```

► <https://www.couchbase.com/products/n1ql>



Query



Query

- ▶ Create Database

- ▶ Couch → REST

- ▶ PUT → Create
 - ▶ POST → Insert
 - ▶ GET → Select
 - ▶ DELETE → Delete

curl -X PUT http://127.0.0.1:5984/fruit

curl -X POST -D '{"Abacaxi":190}' http://127.0.0.1:5984/fruit



Query



mongoDB

mysql

MongoDB

SELECT

```
Dim1, Dim2,
SUM(Measure1) AS MSum,
COUNT(*) AS RecordCount,
AVG(Measure2) AS MAvg,
MIN(Measure1) AS MMin
MAX(CASE
  WHEN Measure2 < 100
  THEN Measure2
END) AS MMax
FROM DenormAggTable
WHERE (Filter1 IN ('A', 'B'))
AND (Filter2 = 'C')
AND (Filter3 > 123)
GROUP BY Dim1, Dim2
HAVING (MMin > 0)
ORDER BY RecordCount DESC
LIMIT 4, 8
```

- ① Grouped dimension columns are pulled out as keys in the map function, reducing the size of the working set.
- ② Measures must be manually aggregated.
- ③ Aggregates depending on record counts must wait until finalization.
- ④ Measures can use procedural logic.
- ⑤ Filters have an ORM/ActiveRecord-looking style.
- ⑥ Aggregate filtering must be applied to the result set, not in the map/reduce.
- ⑦ Ascending: 1; Descending: -1

```
db.runCommand({
  mapreduce: "DenormAggCollection",
  query: {
    filter1: { '$in': [ 'A', 'B' ] },
    filter2: 'C',
    filter3: { '$gt': 123 }
  },
  map: function() { emit(
    { d1: this.Dim1, d2: this.Dim2 },
    { msum: this.measure1, recs: 1, mmin: this.measure1,
      mmax: this.measure2 < 100 ? this.measure2 : 0 }
  ); },
  reduce: function(key, vals) {
    var ret = { msum: 0, recs: 0, mmin: 0, mmax: 0 };
    for(var i = 0; i < vals.length; i++) {
      ret.msum += vals[i].msum;
      ret.recs += vals[i].recs;
      if(vals[i].mmin < ret.mmin) ret.mmin = vals[i].mmin;
      if((vals[i].mmax < 100) && (vals[i].mmax > ret.mmax))
        ret.mmax = vals[i].mmax;
    }
    return ret;
  },
  finalize: function(key, val) {
    val.mavg = val.msum / val.recs;
    return val;
  },
  out: 'result1',
  verbose: true
});
db.result1.
  find({ mmin: { '$gt': 0 } }).
  sort({ recs: -1 }).
  skip(4).
  limit(8);
```

OBRIGADO.