Optimizing compute graphs

April 8, 2015

1 Enumeration

Assumption: Only care about functions f that are multivariate polynomials, e.g.,

$$f = f(X_1, X_2, \cdots, X_n).$$

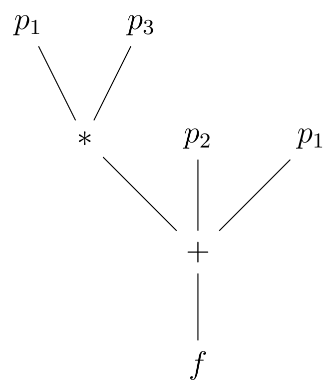
I am going to ignore division for now. I also don't really know how to handle the coefficients, I am somehow magically assuming they don't exist (you will see later).

Definition: Compute graph $\mathcal{G}$ is a "representation" of some function f. Each function f is "represented" by multiple compute graphs. Consider the following "recursive" definition of $\mathcal{G}$: A compute graph $\mathcal{G}$ is said to be k-computation of a multivariate polynomial f, if f satisifes

$$f = g(p_1, p_2, \cdots, p_k)$$

where i) g is a k-variate polynomial, and ii) p_i are polynomials in $X_1, \cdots, X_n$ that have k-computation representations by "children" graphs of $\mathcal{G}$. The motivation for k-computations from practical standpoints, is that the maximum number of mults/adder is k. You can also limit the number of adds but that is a simple generalization. Since I ignore division the polynomials p_i must have smaller degree than f.

The below computation graph is a 2-computation representation of f.



I am interested enumerating all possible k-computations of some multivariate polynomial f. If we can enumerate, we can pick the one with some smallest cost.

General idea: Maintain a pool of multivariate polynomials. Build larger degree polynomials by "combining" them via k-variate polynomials. We have to limit the pool by quickly eliminating polynomials that we know cannot build f. We can use the following observation to eliminate candidates.

Observation: If a polynomial p can be used to build f, then any monomial of p must divide some monomial of f.

Example: A trite one, $f = X_1^n$. In this case, any polynomial p that is a function of any $X_2, \cdots$ is eliminated.

Example: Consider $f = X_1X_2 + X_2X_3 + X_1X_3$. In this case, any polynomial p that has multiple variates, e.g. $X_2^2X_3$, is eliminated.