Package 'parallel'

R-core

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1 Introduction

Package **parallel** was rst included in R 2.14.0. It builds on the work done for CRAN packages **multicore** (Urbanek, 2009{2014) and **snow** (Tierney $et\ al.$, 2003{present) and provides drop-in replacements for most of the functionality of those packages, with integrated handling of random-number generation.

Parallelism can be done in computation at many dierent levels: this package is principally concerned with `coarse-grained parallelization'. At the lowest level, modern CPUs can do several basic operations simultaneously (e.g. integer and oating-point arithmetic), and several implementations of external BLAS libraries use multiple threads to do parts of basic vector/matrix operations in parallel. Several contributed R packages use multiple threads at C level via OpenMP or pthreads.

This package handles running much larger chunks of computations in parallel. A typical example is to evaluate the same R function on many di erent sets of data: often simulated data as in bootstrap computations (or with `data' being the random-number stream). The crucial point is that these chunks of computation are unrelated and do not need to communicate in any way. It is often the case that the chunks take approximately the same length of time. The basic computational model is

- (a) Start up M `worker' processes, and do any initialization needed on the workers.
- (b) Send any data required for each task to the workers.
- (c) Split the task into M roughly equally-sized chunks, and send the chunks (including the R code needed) to the workers.
- (d) Wait for all the workers to complete their tasks, and ask them for their results.
- (e) Repeat steps (b{d) for any further tasks.
- (f) Shut down the worker processes.

Amongst the initializations which may be needed are to load packages and initialize the randomnumber stream.

There are implementations of this model in the functions mcl apply and parLapply as near-drop-in replacements for lapply.

A slightly di erent model is to split the task into $M_1 > M$ chunks, send the rst M chunks to the workers, then repeatedly wait for any worker to complete and send it the next remaining task: see the section on `load balancing'.

In principle the workers could be implemented by threads¹ or lightweight processes, but in the current implementation they are full processes. They can be created in one of three ways:

1. Via system("Rscript") or similar to launch a new process on the current machine or a similar machine with an identical R installation. This then needs a way to communicate between master and worker processes, which is usually done via sockets.

This should be available on all R platforms, although it is conceivable that zealous security measures could block the inter-process communication via sockets. Users of Windows and macOS may expect pop-up dialog boxes from the rewall asking if an R process should accept incoming connections.

Following \mathbf{snow} , a pool of worker processes listening via sockets for commands from the master is called a `cluster' of nodes.

2. Via forking. Fork is a concept² from POSIX operating systems, and should be available on all R platforms except Windows. This creates a new R process by taking a complete copy of the master process, including the workspace and state of the random-number stream. However, the copy will (in any reasonable OS) share memory pages with the master until modi ed so forking is very fast.

The use of forking was pioneered by package multicore.

Note that as it does share the complete process, it also shares any GUI elements, for example an R console and on-screen devices. This can cause havoc.³

There needs to be a way to communicate between master and worker. Once again there are several possibilities since master and workers share memory. In **multicore** the initial fork sends an R expression to be evaluated to the worker, and the master process opens a pipe for reading that is used by the worker to return the results. Both that and creating a cluster of nodes communicating *via* sockets are supported in package **parallel**.

3. Using OS-level facilities to set up a means to send tasks to other members of a group of machines. There are several ways to do that, and for example package **snow** can make use of MPI ('message passing interface') using R package **Rmpi**. Communication overheads can dominate computation times in this approach, so it is most often used on tightly-coupled networks of computers with high-speed interconnects.

CRAN packages following this approach include **GridR** (using Condor or Globus) and **Rsge** (using SGE, currently called `Oracle Grid Engine').

It will not be considered further in this vignette, but those parts of **parallel** which provide **snow**-like functions will accept **snow** clusters including MPI clusters.

The landscape of parallel computing has changed with the advent of shared-memory computers with multiple (and often many) CPU cores. Until the late 2000's parallel computing was mainly done on clusters of large numbers of single- or dual-CPU computers: nowadays even laptops have two or four cores, and servers with 8, 32 or more cores are commonplace. It is such hardware that package **parallel** is designed to exploit. It can also be used with several computers running the same version of R connected by (reasonable-speed) ethernet: the computers need not be running the same OS.

Note that all these methods of communication use serial ize/unserial ize to send R objects between processes. This has limits (typically hundreds of millions of elements) which a well-designed parallelized algorithm should not approach.

¹only 'in principle' since the R interpreter is not thread-safe.

²https://en.wikipedia.org/wiki/Fork_(operating_system)

³Some precautions are taken on macOS: for example the event loops for R.app and the quartz device are inhibited in the child. This information is available at C level in the Rboolean variable R_isForkedChild.

2 Numbers of CPUs/cores

In setting up parallel computations it can be helpful to have some idea of the number of CPUs or cores available, but this is a rather slippery concept. Nowadays almost all physical CPUs contain two or more cores that run more-or-less independently (they may share parts of the cache memory, and they do share access to RAM). However, on some processors these cores may themselves be able to run multiple tasks simultaneously, and some OSes (e.g. Windows) have the concept of *logical* CPUs which may exceed the number of cores.

Note that all a program can possibly determine is the total number of CPUs and/or cores available. This is not necessarily the same as the number of CPUs available *to the current user* which may well be restricted by system policies on multi-user systems. Nor does it give much idea of a reasonable number of CPUs to use for the current task: the user may be running many R processes simultaneously, and those processes may themselves be using multiple threads through a multi-threaded BLAS, compiled code using OpenMP or other low-level forms of parallelism. We have even seen instances of **multicore**'s mcl appl y being called recursively, 4 generating $2n+n^2$ processes on a machine estimated to have n=16 cores.

But in so far as it is a useful guideline, function detectCores() tries to determine the number of CPU cores in the machine on which R is running: it has ways to do so on all known current R platforms. What exactly it measures is OS-speci c: we try where possible to report the number of logical cores available.

On Windows the default is to report the number of logical CPUs. On modern hardware (e.g. Intel $Core\ i7$) the latter may not be unreasonable as hyper-threading does give a signi cant extra throughput. What detectCores(logical = FALSE) reports is OS-version-dependent: on recent versions of Windows it reports the number of physical cores but on older versions it might report the number of physical CPU packages.

3 Analogues of apply functions

By far the most common direct applications of packages **multicore** and **snow** have been to provide parallelized replacements of lapply, sapply, apply and related functions.

As analogues of Lapply there are

```
parLapply(cl, x, FUN, ...)
mclapply(X, FUN, ..., mc.cores)
```

where mcl apply is not available on Windows and has further arguments discussed on its help page. They di er slightly in philosophy: mcl apply sets up a pool of mc. cores workers just for this computation, whereas parLapply uses a less ephemeral pool speci ed by the object cl created by a call to makeCluster (which $inter\ alia$ speci es the size of the pool). So the work ow is

```
cl <- makeCluster(<size of pool>)
# one or more parLapply calls
stopCluster(cl)
```

For matrices there are the rarely used parAppl y and parCappl y functions, and the more commonly used parRappl y, a parallel row appl y for a matrix.

⁴parallel::mclapply detects this and runs nested calls serially.

⁵except as a stub which simply calls lapply.

4 SNOW Clusters

The package contains a slightly revised copy of much of **snow**, and the functions it contains can also be used with clusters created by **snow** (provided the package is on the search path).

Two functions are provided to create SNOW clusters, makePS0CKcl uster (a streamlined version of snow::makeS0CKcl uster) and (except on Windows) makeForkCl uster. They di er only in the way they spawn worker processes: makePS0CKcl uster uses Rscri pt to launch further copies of R (on the same host or optionally elsewhere) whereas makeForkCl uster forks the workers on the current host (which thus inherit the environment of the current session).

These functions would normally be called *via* makeCluster.

Both stdout() and stderr() of the workers are redirected, by default being discarded but they can be logged using the outfile option. Note that the previous sentence refers to the *connections* of those names, not the C-level le handles. Thus properly written R packages using Rprintf will have their output redirected, but not direct C-level output.

A default cluster can be registered by a call to setDefaul tCluster(): this is then used whenever one of the higher-level functions such as parApply is called without an explicit cluster. A little care is needed when repeatedly re-using a pool of workers, as their workspaces will accumulate objects from past usage, and packages may get added to the search path.

If clusters are to be created on a host other than the current machine ('I ocal host'), makeCl uster may need to be given more information in the shape of extra arguments.

- If the worker machines are not set up in exactly the same way as the master (for example if they are of a di erent architecture), use homogeneous = FALSE and perhaps set rscript to the full path to Rscript on the workers.
- The worker machines need to know how to communicate with the master: normally this can be done using the hostname found by Sys. info(), but on private networks this need not be the case and master may need to be supplied as a name or IP address, e.g. master = "192.168.1.111".
- By default ssh is used to launch R on the workers. If it is known by some other name, use e.g. rshcmd = "pl i nk. exe" for a Windows box using PUTTY. SSH should be set up to use silent authentication: setups which require a password to be supplied may or may not work.
- Socket communication is done over port a randomly chosen port in the range 11000: 11999: site policies might require some other port to be used, in which case set argument port or environment variable R PARALLEL PORT.

5 Forking

Except on Windows, the package contains a copy of **multicore**: there a few names with the added pre x mc, e.g. mccollect and mcparallel. (Package **multicore** used these names, but also the versions without the pre x which are too easily masked: e.g. package **lattice** used to have a function parallel.)

The low-level functions from multicore are provided but not exported from the namespace.

There are high-level functions mcl appl y and pvec: unlike the versions in **multicore** these default to 2 cores, but this can be controlled by setting opti ons ("mc. cores"), and that takes its default from environment variable MC_CORES when the package is loaded. (Setting this to 1 inhibits parallel operation: there are stub versions of these functions on Windows which force mc. cores = 1.)

Functions mcmapply and mcMap provide analogues of mapply and Map.

Note the earlier comments about using forking in a GUI environment.

The parent and forked R processes share the per-session temporary directory tempdir(), which can be a problem as a lot of code has assumed it is private to the R process. Further, prior to R 2.14.1 it was possibly for tempfile in two processes to select the same lename in that temporary directory, and do it su ciently simultaneously that neither saw it as being in use.

The forked workers share le handles with the master: this means that any output from the worker should go to the same place as 'stdout' and 'stderr' of the master process. (This does not work reliably on all OSes: problems have also been noted when forking a session that is processing batch input from 'stdin'.) Setting argument mc. silent = TRUE silences 'stdout' for the child: 'stderr' is not a ected.

Sharing le handles also impacts graphics devices as forked workers inherit all open graphics devices of the parent: they should not attempt to make use of them.

6 Random-number generation

Some care is needed with parallel computation using (pseudo-)random numbers: the processes/threads which run separate parts of the computation need to run independent (and preferably reproducible) random-number streams. One way to avoid any di-culties is (where possible) to do all the randomization in the master process: this is done where possible in package **boot** (version 1.3-1 and later).

When an R process is started up it takes the random-number seed from the object . Random. seed in a saved workspace or constructs one from the clock time and process ID when random-number generation is rst used (see the help on RNG). Thus worker processes might get the same seed because a workspace containing . Random. seed was restored or the random number generator has been used before forking: otherwise these get a non-reproducible seed (but with very high probability a di erent seed for each worker).

The alternative is to set separate seeds for each worker process in some reproducible way from the seed in the master process. This is generally plenty safe enough, but there have been worries that the random-number streams in the workers might somehow get into step. One approach is to take the seeds a long way apart in the random-number stream: note that random numbers taken a long (xed) distance apart in a single stream are not necessarily (and often are not) as independent as those taken a short distance apart. Yet another idea (as used by e.g. **JAGS**) is to use di erent random-number generators for each separate run/process.

Package **parallel** contains an implementation of the ideas of L'Ecuyer $et\ al.\ (2002)$: this uses a single RNG and make streams with seeds 2^{127} steps apart in the random number stream (which has period approximately 2^{191}). This is based on the generator of L'Ecuyer (1999); the reason for choosing that generator⁶ is that it has a fairly long period with a small seed (6 integers), and unlike R's default "Mersenne-Twi ster" RNG, it is simple to advance the seed by a xed number of steps. The generator is the combination of two:

```
x_n = 1403580 \times x_{n-2} - 810728 \times x_{n-3} \mod (2^{32} - 209)

y_n = 527612 \times y_{n-1} - 1370589 \times y_{n-3} \mod (2^{32} - 22853)

z_n = (x_n - y_n) \mod 4294967087

u_n = z_n = 4294967088 \text{ unless } z_n = 0
```

⁶apart from the commonality of authors!

The `seed' then consists of $(x_{n-3}; x_{n-2}; x_{n-1}; y_{n-3}; y_{n-2}; y_{n-1})$, and the recursion for each of x_n and y_n can have pre-computed coe—cients for k steps ahead. For $k = 2^{127}$, the seed is advanced by k steps by R call . Random. seed <-nextRNGStream(. Random. seed).

The L'Ecuyer (1999) generator is available via RNGki nd("L'Ecuyer-CMRG"). Thus using the ideas of L'Ecuyer $et\ al.\ (2002)$ is as simple as

and this is is implemented for SNOW clusters in function clusterSetRNGStream, and as part of mcparallel and mclapply (by default).

Apart from streams (2¹²⁷ steps apart), there is the concept of sub-streams starting from seeds 2⁷⁶ steps apart. Function nextRNGSubStream advances to the next substream.

A direct R interface to the (clunkier) original C implementation is available in CRAN package **rlecuyer** (Sevcikova and Rossini, 2004{present). That works with named streams, each of which have three 6-element seeds associated with them. This can easily be emulated in R by storing . Random. seed at suitable times. There is another interface using S4 classes in package **rstream** (Leydold, 2005{present).

7 Load balancing

The introduction mentioned a di erent strategy which dynamically allocates tasks to workers: this is sometimes known as `load balancing' and is implemented in mcl appl y(mc. preschedul e = FALSE), cl usterAppl yLB and wrappers.

Load balancing is potentially advantageous when the tasks take quite dissimilar amounts of computation time, or where the nodes are of disparate capabilities. But some *caveats* are in order:

- (a) Random number streams are allocated to nodes, so if the tasks involve random numbers they are likely to be non-repeatable (as the allocation of tasks to nodes depends on the workloads of the nodes). It would however take only slightly more work to allocate a stream to each task.
- (b) More care is needed is allocating the tasks. If 1000 tasks need to be allocated to 10 nodes, the standard approach send chunks of 100 tasks to each of the nodes. The load-balancing approach sends tasks one at a time to a node, and the communication overhead may be high. So it makes sense to have substantially more tasks than nodes, but not by a factor of 100 (and maybe not by 10).

8 Setting the CPU Affinity with mclapply

The parameter affinity. I ist of the mclapply function can be used to run elements of the input vector X of mclapply on speci c CPUs. affinity. I ist is a vector (atomic or list) containing the CPU a nity mask for each element of X, it describes on which CPU (core or hyperthread unit) a given item is allowed to run, see? mcaffinity. This can be helpful, if

the elements of X (parallel jobs) have a high variance of completion time or if the hardware architecture is heterogeneous. It also enables the development of scheduling strategies for optimizing the overall runtime of independent parallel jobs. If affi ni ty. I i st is set, the mc. cores parameter is replaced with the number of CPU ids used in the a nity masks. To use this parameter prescheduling has to be deactivated (mc. preschedul e = FALSE). For each value of the input vector X a separate job is forked. The master process only forks one child process for each selected CPU at once, to ensure that the number of forked jobs does not exceed the number of selected CPUs. As soon as a child process has nished, the next available job for the CPU is forked.

The following code example demonstrates how the execution time can be reduced by assigning the elements of X to speci c CPUs, not on Windows, where mc. cores must remain at 1.

```
> ## Exemplary variance filter executed on three different matrices in parallel.
> ## Can be used in gene expression analysis as a prefilter
> ## for the number of covariates.
> library(parallel)
> n <- 300 # observations
> p <- 20000 # covariates
> ## Different sized matrices as filter inputs
> ## Matrix A and B form smaller work loads
> ## while matrix C forms a bigger workload (2*p)
> library(stats)
> A <- matrix(replicate(p, rnorm(n, sd = runif(1, 0.1, 10))), n, p)
> B < -matrix(replicate(p, rnorm(n, sd = runif(1, 0.1, 10))), n, p)
> C < -matrix(replicate(2*p, rnorm(n, sd = runif(1, 0.1, 10))), n, 2*p)
> varFilter <- function (X, nSim = 20) {</pre>
    for (i in 1: nSim) {
      train <- sample(nrow(X), 2 / 3 * nrow(X))
      colVars <- apply(X[train, ], 2, var)
      keep <- names(head(sort(col Vars, decreasing = TRUE), 100))</pre>
      # myAlgorithm(X[, keep])
+ }
> ## Runtime comparison -----
> ## mclapply with affinity.list
> ## CPU mapping: A and B run on CPU 1 while C runs on CPU 2:
> affinity <- c(1, 1, 2)
> system. time(
    mclapply(X = list(A, B, C), FUN = varFilter,
             mc. preschedule = FALSE, affinity.list = affinity))
            system elapsed
      user
> ## 34, 909
            0. 873 36. 720
> ## mclapply without affinity.list
> system. time(
    mclapply(X = list(A, B, C), FUN = varFilter, mc. cores = 2,
             mc. preschedul e = FALSE) )
     user
            system elapsed
> ## 72.893
             1.588 55.982
```

```
> 
> 
## mclapply with prescheduling
> system.time(
+ mclapply(X = list(A, B, C), FUN = varFilter, mc.cores = 2,
+ mc.preschedule = TRUE) )
> ## user system elapsed
> ## 53.455   1.326   53.399
```

Instead of using affi ni ty. I i st for runtime optimization, it can also be used to simply restrict the computation to speci c CPUs as in the following example.

```
> ## Restricts all elements of X to run on CPU 1 and 2.
> X <- list(1, 2, 3)
> affinity.list <- list(c(1,2), c(1,2), c(1,2))
> mclapply(X = X, FUN = function (i) i*i,
+ mc.preschedule = FALSE, affinity.list = affinity.list)
```

9 Portability considerations

People wanting to provide parallel facilities in their code need to decide how hard they want to try to be portable and e cient: no approach works optimally on all platforms.

Using mcl apply is usually the simplest approach, but will run serial versions of the code on Windows. This may su ce where parallel computation is only required for use on a single multi-core Unix-alike server | for mcl apply can only run on a single shared-memory system. There is fallback to serial use when needed, by setting mc. cores = 1.

Using parLappl y will work everywhere that socket communication works, and can be used, for example, to harness all the CPU cores in a lab of machines that are not otherwise in use. But socket communication may be blocked even when using a single machine and is quite likely to be blocked between machines in a lab. There is not currently any fallback to serial use, nor could there easily be (as the workers start with a di erent R environment from the one current on the master).

An example of providing access to both approaches as well as serial code is package **boot**, version 1.3-3 or later.

10 Extended examples

```
> library(parallel)
```

Probably the most common use of coarse-grained parallelization in statistics is to do multiple simulation runs, for example to do large numbers of bootstrap replicates or several runs of an MCMC simulation. We show an example of each.

Note that some of the examples will only work serially on Windows and some actually are computer-intensive.

10.1 Bootstrapping

Package **boot** (Canty and Ripley, 1999{present) is support software for the monograph by Davison and Hinkley (1997). Bootstrapping is often used as an example of easy parallelization, and some methods of producing condence intervals require many thousands of bootstrap samples.

As from version 1. 3-1 the package itself has parallel support within its main functions, but we illustrate how to use the original (serial) functions in parallel computations.

We consider two examples using the cd4 dataset from package **boot** where the interest is in the correlation between before and after measurements. The rst is a straight simulation, often called a *parametric bootstrap*. The non-parallel form is

To do this with mcl apply we need to break this into separate runs, and we will illustrate two runs of 500 simulations each:

There are many ways to program things like this: often the neatest is to encapsulate the computation in a function, so this is the parallel form of

```
> do. call(c, lapply(seq_len(mc), run1))
```

To run this with parLappl y we could take a similar approach by

Note that whereas with mcl appl y all the packages and objects we use are automatically available on the workers, this is not in general⁷ the case with the parLappl y approach. There is often a delicate choice of where to do the computations: for example we could compute cd4. ml e on

⁷it is with clusters set up with makeForkCluster.

the workers (as above) or on the master and send the value to the workers. We illustrate the latter by the following code

```
> cl <- makeCluster(mc)</pre>
> cd4.rg <- function(data, mle) MASS::mvrnorm(nrow(data), mle$m, mle$v)
> cd4.mle <- list(m = colMeans(cd4), v = var(cd4))
> clusterExport(cl, c("cd4.rg", "cd4.mle"))
> junk <- clusterEvalQ(cl, library(boot)) # discard result
> clusterSetRNGStream(cl, 123)
> res <- clusterEvalQ(cl, boot(cd4, corr, R = 500,
                        sim = "parametric", ran.gen = cd4.rg, mle = cd4.mle))
> library(boot) # needed for c() method on master
> cd4. boot <- do. call(c, res)
> boot.ci(cd4.boot, type = c("norm", "basic", "perc"),
          conf = 0.9, h = atanh, hinv = tanh)
> stopCluster(cl)
Running the double bootstrap on the same problem is far more computer-intensive. The stan-
dard version is
> R <- 999; M <- 999 ## we would like at least 999 each
> cd4.nest <- boot(cd4, nested.corr, R=R, stype="w", t0=corr(cd4), M=M)</pre>
> ## nested.corr is a function in package boot
> op <- par(pty = "s", xaxs = "i", yaxs = "i")
> qqplot((1:R)/(R+1), cd4.nest$t[, 2], pch = ".", asp = 1,
           xlab = "nominal", ylab = "estimated")
> abline(a = 0, b = 1, col = "grey")
> abline(h = 0.05, col = "grey")
> abline(h = 0.95, col = "grey")
> par(op)
> nominal <- (1:R)/(R+1)
> actual <- cd4.nest$t[, 2]</pre>
> 100*nominal[c(sum(actual <= 0.05), sum(actual < 0.95))]
which took about 55 secs on one core of an 8-core Linux server.
Using mcl apply we could use
> mc <- 9
> R <- 999; M <- 999; RR <- floor(R/mc)
> run2 <- function(...)</pre>
      cd4.nest <- boot(cd4, nested.corr, R=RR, stype="w", t0=corr(cd4), M=M)
> cd4. nest <- do. call(c, mclapply(seq_len(mc), run2, mc. cores = mc) )</pre>
```

which ran in 11 secs (elapsed) using all of that server.

> 100*nominal[c(sum(actual <= 0.05), sum(actual < 0.95))]

10.2 MCMC runs

> nomi nal <- (1:R)/(R+1)
> actual <- cd4.nest\$t[, 2]</pre>

Ripley (1988) discusses the maximum-likelihood estimation of the Strauss process, which is done by solving a moment equation

$$E_c T = t$$

where T is the number of R-close pairs and t is the observed value, 30 in the following example. A serial approach to the initial exploration might be

```
> library(spatial)
> towns <- ppinit("towns.dat")</pre>
> tget <- function(x, r=3.5) sum(dist(cbind(x$x, x$y)) < r)
> t0 <- tget(towns)</pre>
> R <- 1000
> c < - seq(0, 1, 0.1)
> ## res[1] = 0
> res <- c(0, sapply(c[-1], function(c))
      mean(replicate(R, tget(Strauss(69, c=c, r=3.5))))))
> plot(c, res, type="l", ylab="E t")
> abline(h=t0, col="grey")
which takes about 20 seconds today, but many hours when rst done in 1985. A parallel version
might be
> run3 <- function(c) {</pre>
      library(spatial)
      towns <- ppinit("towns.dat") # has side effects
      mean(replicate(R, tget(Strauss(69, c=c, r=3.5))))
+ }
> cl <- makeCluster(10, methods = FALSE)</pre>
> clusterExport(cl, c("R", "towns", "tget"))
> res < c(0, parSapply(cl, c[-1], run3)) # 10 tasks
> stopCluster(cl)
which took about 4.5 secs, plus 2 secs to set up the cluster. Using a fork cluster (not on
Windows) makes the startup much faster and setup easier:
> cl <- makeForkCluster(10) # fork after the variables have been set up
> run4 <- function(c) mean(replicate(R, tget(Strauss(69, c=c, r=3.5))))</pre>
> res <- c(0, parSapply(cl, c[-1], run4))
> stopCluster(cl)
As one might expect, the mcl apply version is slightly simpler:
> run4 <- function(c) mean(replicate(R, tget(Strauss(69, c=c, r=3.5))))
> res < c(0, unlist(mclapply(c[-1], run4, mc.cores = 10)))
```

If you do not have as many as 10 cores, you might want to consider load-balancing in a task like this as the time taken per simulation does vary with c. This can be done using mcl apply (mc. preschedul e = FALSE) or parSapplyLB. The disadvantage is that the results would not be reproducible (which does not matter here).

10.3 Package installation

With over 4000 R packages available, it is often helpful to do a comprehensive install in parallel. We provide facilities to do so in install. packages via a parallel make, but that approach may not be suitable. Some of the tasks take many times longer than others, and we do not know in advance which these are.

We illustrate an approach using package **parallel** which is used on part of the CRAN check farm. Suppose that there is a function do_one(pkg) which installs a single package and then returns. Then the task is to run do_one on as many of the M workers as possible whilst ensuring that all of the direct and indirect dependencies of pkg are installed before pkg itself. As the installation

⁸A parallel make might not be available, and we have seen a couple of instances of package installation not working correctly when run from make.

it refers to objects within the evaluation frame of boot and hence the environment of fn which will therefore be sent to the workers along with fn.

11 Differences from earlier versions

The support of parallel RNG di ers from snow: multicore had no support.

11.1 Differences from multicore

multicore had quite elaborate code to inhibit the Aqua event loop for R. app and the event loop for the quartz graphics device. This has been replaced by recording a ag in the R executable for a child process.

Functions fork and kill have mc pre xes and are not exported. This avoids confusion with other packages (such as package fork), and note that mckill is not as general as tool s: pskill.

Aliased functions collect and parallel are no longer provided.

11.2 Differences from snow

snow set a timeout that exceeded the maximum value required by POSIX, and did not provide a means to set the timeout on the workers. This resulted in process interlocks on Solaris.

makeCluster creates MPI or NWS clusters by calling snow.

makePSOCKcl uster has been streamlined, as package **parallel** is in a known location on all systems, and Rscript is these days always available. Logging of workers is set up to append to the le, so multiple processes can be logged.

parSapply has been brought into line with sapply.

clusterMap() gains SIMPLIFY and USE. NAMES arguments to make it a parallel version of mapply and Map.

The timing interface has not been copied.

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