

Project 5: Sheldon Air

Repository at:

https://github.com/fabriziorocco/Project-Group-Luiss.git

The Problem

SheldonAir is an emerging airline in the airline business.

Dr. Cooper, the CEO of SheldonAir, would like to know the minimum number of legs (flight segments) required to go from one city to another city.

This problem can be represented by an undirected graph where:

- vertices correspond to cities where SheldonAir is operating
- an <u>edge</u> between two vertices corresponds to a direct SheldonAir flight between the two corresponding cities.

In other terms, given two vertices u and v, you are required to compute the length of a shortest path between u and v (i.e., the minimum number of edges in a path from u to v).

Idea

Goal: Find the shortest path between vertices of a graph.

Assumption

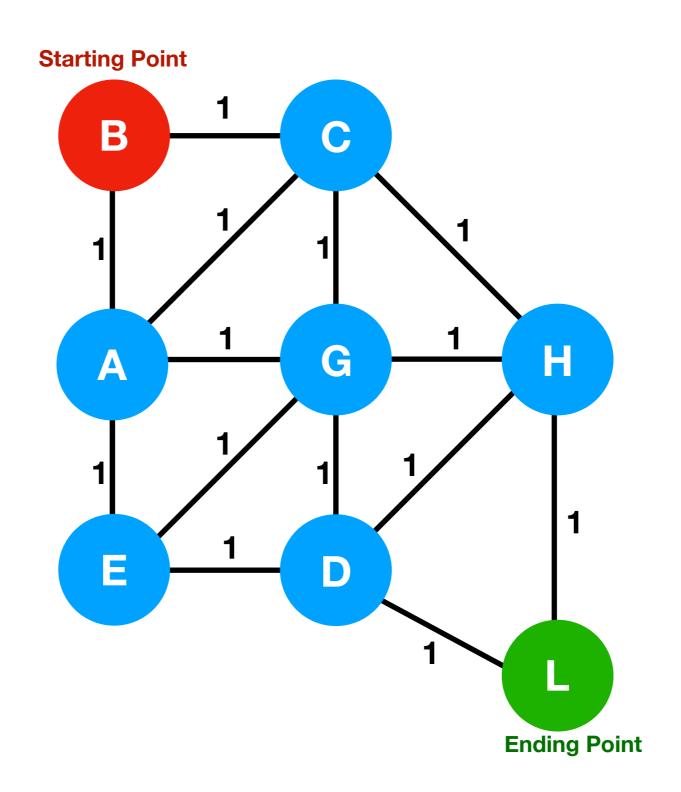
We have to initialize all the nodes with a placeholder value, let's suppose ∞, and we have to attribute to each edge a value.

There will be a set S containing the nodes already visited, $S = \{ \}$ and a set T, containing the nodes to analyze yet. $T = \{ \}$.

At the beginning of the execution the two sets are so defined:

- S = {a} # Only the starting vertex a
- T = { All the other nodes }

We give to each edge the weight 1.



Intuition: Dijkstra Algorithm

This algorithm is also known as "Dijkstra Algorithm"
 "Dijkstra's algorithm (or Dijkstra's Shortest Path First algorithm, SPF algorithm) is an algorithm for finding the shortest paths between nodes in a graph. It was conceived by computer scientist Edsger W. Dijkstra in 1956 and published three years later."

Source: Wikipedia

Key Points for "constant weight" Dijkstra Algorithm

- Every edge has weight 1 and we're just interested in the shortest path of segments.
- Once we've moved to the node we're going to visit, we will check each of its neighboring nodes.
- For each neighboring node, we'll calculate the distance/cost for the neighboring nodes by summing the cost of the edges that lead to the node we're checking from the starting vertex.
- Finally, if the distance/cost to a node is less than a known distance, we'll update the shortest distance that we have on file for that vertex.

Running Time of Dijkstra Algorithm

Complexity can expressed by |V| and |E|, which are respectively the number of vertices and the number of edges.

	m Decrease Key	n DeleteMin	n Insert	DIJKSTRA
ARRAYS	Θ(1)	O(V)	Θ(V)	$O(V ^2 + E)$
BINARY HEAP	$O(log_2 V)$	$O(log_2 V)$	Θ(V)	$O((V + E)\log_2 V)$
FIBONACCI HEAP	Θ(1)	O(log ₂ V)	Θ(V)	O(V log ₂ V + E)
QUEUE	Θ(1)	Θ(1)	Θ(1)	O(V + E)

Since each edge has weight 1(constant), our running time is O(n).(linear)

To do this we use a queue.

Queue works correctly only if the edges have constant weight.

Implementation Choice

- We decided to implement this algorithm using C++ in order to perform a benchmark as accurate as possible.
- For the test environment we used a Ryzen 2700X.
- Task pinned to a physical core and hyper threading disabled.
- Our implementation is single threaded.

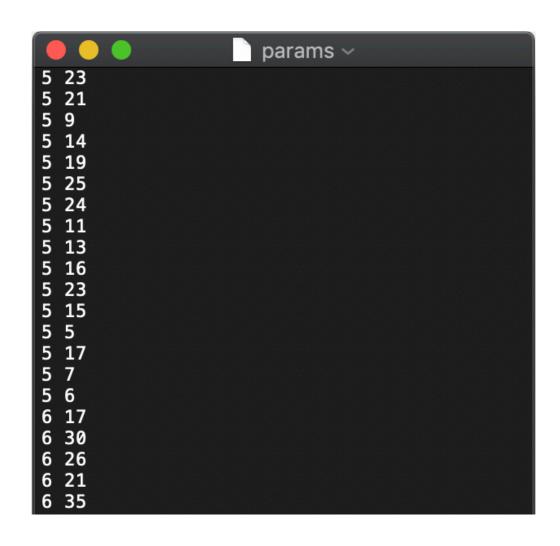
```
#pragma GCC optimize("Ofast")
   #include <bits/stdc++.h>
   using namespace std;
  vector<long long int> enGraph[123456];
   long long int curNodo;
   int main() {
       FILE* fin = fopen("input.txt", "r");
       FILE* fout = fopen("output.txt", "w");
10
       long long int N, V;
       fscanf(fin, "%lli %lli", &N, &V);
       long long int startNodo = 0;
       long long int endNodo = N-1;
       long long int a, b;
       for (int i = 0; i < V; i++) {
16
           fscanf(fin, "%lli %lli", &a, &b);
           enGraph[a].push_back(b);
           enGraph[b].push_back(a);
       vector<bool> visited;
21
       visited.resize(N);
       vector<long long int> cost;
       cost.resize(N);
24
       memset(&cost[0], -1, N*sizeof(long long int));
25
       cost[startNodo] = 0;
       queue<long long int> q;
       q.push(startNodo);
       while ((!q.empty()) && (cost[endNodo] == -1)) {
           curNodo = q.back();
31
           q.pop();
           for (auto e : enGraph[curNodo]) {
               if (visited[e])
34
                   continue;
               cost[e] = cost[curNodo] + 1;
36
               q.push(e);
37
               visited[e] = 1;
           }
40
       fprintf(fout, "%11i\n", cost[endNodo]);
42 }
```

Algorithm

- The full project is uploaded in the repository on Git-Hub.
- The code is written with a CamelCase style
- Extra comment are on the full project

User Guide

1. To run the code, all we need is just to run the file run.sh.



```
1 #!/bin/bash
2
3 echo "Compiling..."
4 g++ ./gen.cpp -Ofast -o ./gen.o
5 g++ ./sol.cpp -Ofast -o ./sol.o
6
7 echo "Generating parameters..."
8 python3 genParams.py
9
10 echo "Testing..."
11 python3 test.py < params > res
12 echo "Done!"
13 echo 'Run "cat res" to view test results'
14
```

bash /path/to/script.sh

2. This script is able to compile the generator of parameters, it tests it and at the end it returns the results on the file res.

Will be created also some **plots** based on the test results.

CMS Platform

- We developed a platform hosted on <u>cms.lizzit.it</u>, which aims is to allow everyone to try to solve the problem.
- Each solution is tested against a group of testcase.
- The score is assigned according to how many testcase are correctly solved.
- User and password to login are guest and guest.

Test Results

Nodes	Edges	Time (sec)
5	9	0.67
5	14	0.66
5	19	0.66
5	25	0.67
5	24	0.72
5	11	0.66
5	13	0.67
5	16	0.66
5	23	0.66
5	15	0.67
5	5	0.66
5	17	0.66
5	7	0.66
5	6	0.67
6	17	0.67
6	30	0.67
6	26	0.66
6	21	0.67
6	35	0.7
6	22	0.73
6	22	0.72
7	41	0.7
7	14	0.73
7	21	0.69
7	37	0.67

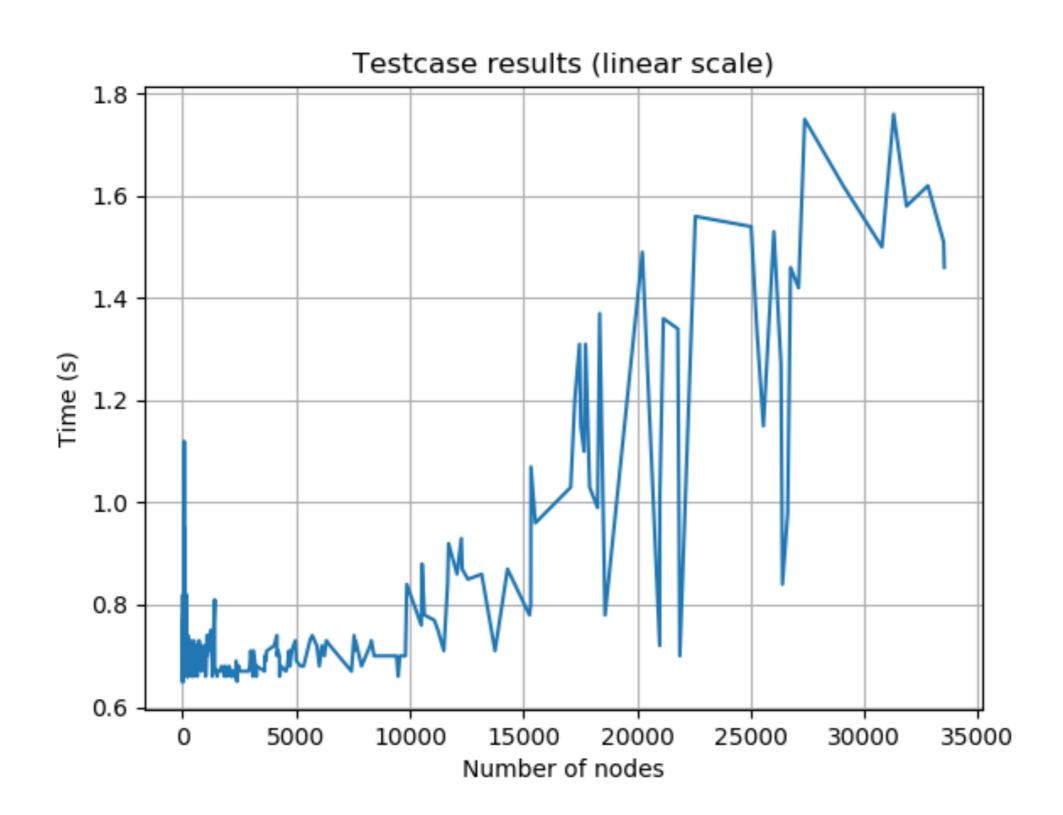
Nodes	Edges	Time (sec)
17	92	0.66
17	21	0.67
18	57	0.67
19	319	0.67
19	278	0.67
20	370	0.73
20	21	0.72
20	229	0.75
21	325	0.73
21	238	0.68
21	361	0.69
21	234	0.7
22	450	0.7
22	114	0.7
22	101	0.69
23	441	0.73
23	261	0.7
23	292	0.69
23	354	0.71
23	235	0.68
24	318	0.76
24	186	0.77
24	277	0.75
24	387	0.72
25	423	0.68
25	36	0.72
25	427	0.69

Nodes	Edges	Time (sec)
10540	94832526	0.88
10643	47163130	0.78
11078	42671090	0.77
11242	42221146	0.75
11499	7842265	0.71
11664	70389426	0.84
11708	10315874	0 0.92
12079	39271909	0.86
12270	11754570	6 0.93
12305	15750155	0.87
12559	688788 0	.85
13165	70279860	0.86
13751	457513 0	.71
14302	79736328	0.87
15267	39158667	0.78
15323	34486807	0.8
15334	13461121	4 1.07
15526	93393275	0.96
17083	11710759	8 1.03
17263	19493223	9 1.2
17471	28167688	7 1.31
17527	18526678	6 1.15
17671	16735519	9 1.1

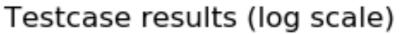
Nodes	Edges	Time (sec)
21001	20246546	0.72
21007	111205845	1.01
21163	421451894	1.36
21798	295346015	1.34
21886	11983290	0.7
22569	356391531	1.56
25027	210658405	1.54
25275	254782217	1.34
25567	170688843	1.15
26027	640914604	1.53
26340	480156665	1.27
26403	55301082	0.84
26635	107826250	0.98
26761	416957058	1.46
27114	386686413	1.42
27383	543328247	1.75
29077	205012326	1.62
30784	753721269	1.5
31297	661641082	1.76
31865	522173315	1.58
32805	853377074	1.62
33499	511916368	1.51
33527	800836659	1.46

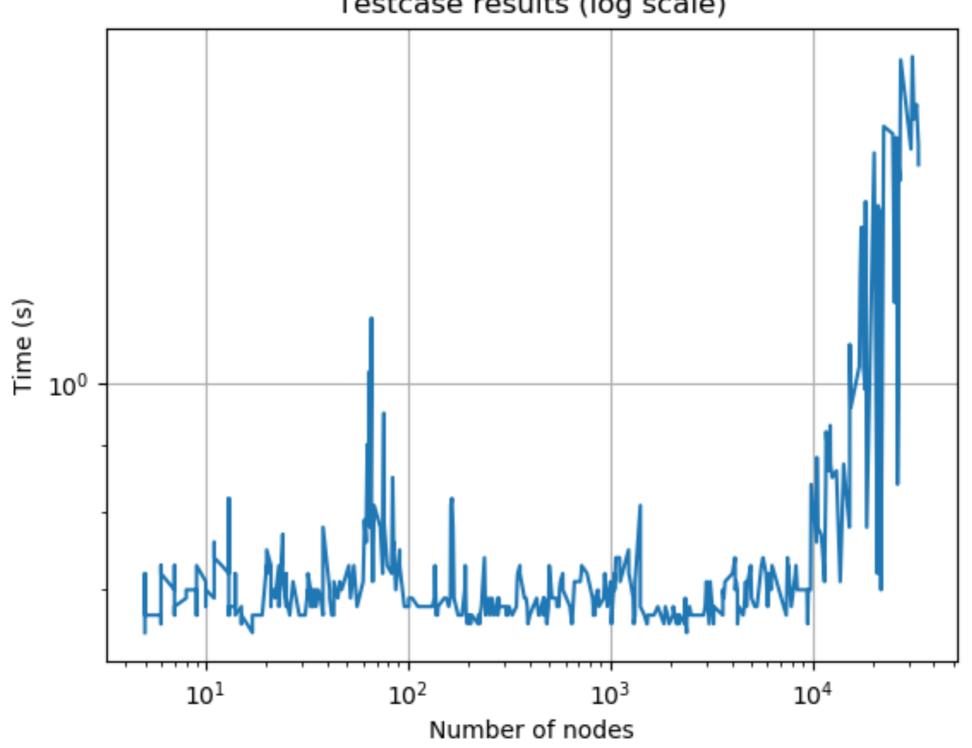
The entire set of results is contained in the res file

Plots



Plots





Critical Analysis

- From the plots we can see that our experimental data is compatible with our theoretical complexity analysis.
- We can understand that the time complexity of Dijkstra constant weight algorithm, in our case, is linear and implemented with a queue is a Breadth First Search.
- In general, some implementations of Dijkstra Algorithm are O(n) when the graph has constant cost for each single edge.