High Availability For Key-Value Stores Using Checkpoint/Restore

Fadhil Abubaker, Hussain Sadiq Abuwala

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- Rely on an external layer to provide replication?

Push replication outside of the database system, delegating it to an external layer.

Shared-Disk

¹J. Kim, K. Salem, K. Daudjee, A. Aboulnaga, and X. Pan. Database High Availability Using SHADOW Systems.

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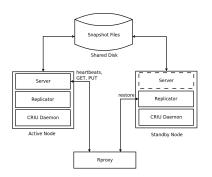
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- ...Checkpoint/Restore?

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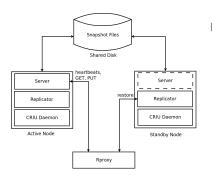
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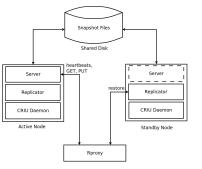
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- Integrated with container runtimes such as OpenVZ, LXC/LXD, Docker, Podman, etc.
- Commonly used for container live-migration.



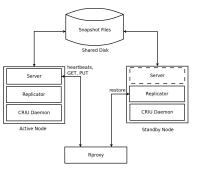
Normal Operation



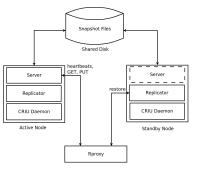
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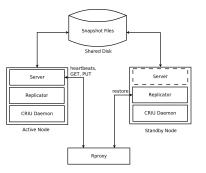


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- server calls checkpoint endpoint on the replicator after n updates.



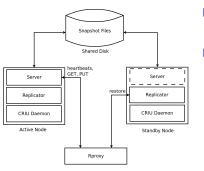
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- replicator interfaces with the C/R daemon and saves a snapshot to the shared-disk.

Failover Process



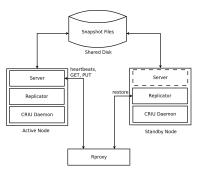
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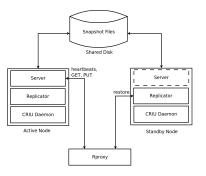
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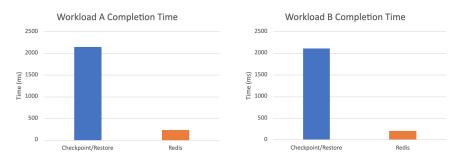


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- replicator restores the server from the latest snapshot on the shared-disk.
- Normal operations resume, with requests forwarded to the standby.

Methodology

- ► Yahoo! Cloud Serving Benchmark.
- Implemented a custom YCSB interface to support our key-value store.
- Each workload consists of 1000 operations.
- Workload A: 50/50 read/write mix.
- ► Workload B: 95/5 read/write mix.
- Snapshots were captured after every 250 updates.

Storage Performance vs Redis (No Replication)



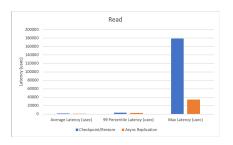
Storage performance likely suffers due to HTTP and JSON overhead.

Replication Performance vs Asynchronous Replication



Requests are blocked while snapshots happen, leading to slightly higher completion times.

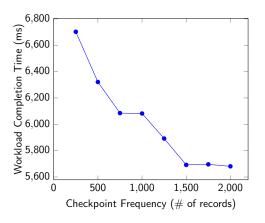
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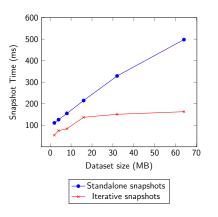
Discrepancy exists for max latency on reads due to snapshots blocking unlucky read requests.

Impact of Checkpoint Frequency on Workload Completion



Custom workload with 5000 insert operations with varying checkpoint intervals.

Standalone vs Iterative Snapshots



Iterative snapshots track changes in memory pages across consecutive checkpoints, making it more efficient.

Demo

Questions?