

3.1 From basis pursuit program

Suppose that $(y, \mathbf{X}) \in \mathbb{R}^n \times \mathbb{R}^{n \times d}$. we now consider the model $y = \mathbf{X}\theta^* + w$ with the noise vector $w \in \mathbb{R}^n$.

Extension of the basis pursuit program. The extension relaxes the constraints of the basis pursuit program, i.e., y does not have to $\mathbf{X}\theta$ for some θ . The extended program can be written as

$$\min_{\theta \in \mathbb{R}^d} \|\theta\|_1 \quad \text{such that} \quad \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 \leq b^2 \quad (3.1)$$

for some noise tolerance $b > 0$.

The program above can be shown as equivalent as such a program:

$$\min_{\theta \in \mathbb{R}^d} \left\{ \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 \right\} \quad \text{such that} \quad \|\theta\|_1 \leq R \quad (3.2)$$

for some radius $R > 0$.

Lasso program. To eliminate the constraint, one can consider the lasso program as well:

$$\min_{\theta \in \mathbb{R}^d} \left\{ \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 + \lambda_n \|\theta\|_1 \right\}. \quad (3.3)$$

Here $\lambda_n > 0$ is a regularization parameter to be chosen by the user.

Note: By Lagrangian duality theory, all three families of convex programs are equivalent, which is given in the next section.

3.2 Equivalent programs

Consider such a convex optimization:

$$p^* := \min_x f_0(x) : \begin{aligned} f_i(x) &\leq 0, & i = 1, \dots, m, \\ h_i(x) &= 0, & i = 1, \dots, p, \end{aligned}$$

where the functions f_0, f_1, \dots, f_m are convex, and h_1, \dots, h_p are affine.

Let \mathcal{D} be the domain of the problem and $\mathcal{X} \subseteq \mathcal{D}$ be its feasible set.

Define the Lagrangian $\mathcal{L} : \mathbf{R}^n \times \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$, with values

$$\mathcal{L}(x, \lambda, \nu) := f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x).$$

The dual function is $g : \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$, with values

$$g(\lambda, \nu) := \min_x \mathcal{L}(x, \lambda, \nu).$$

The dual problem is

$$d^* = \max_{\lambda \geq 0, \nu} g(\lambda, \nu).$$

Strong duality. The convex program is said to satisfy strong duality if $p^* = d^*$.

Theorem 3.1 (Strong duality via Slater condition) *We say that the problem satisfies Slater's condition if it is strictly feasible, that is:*

$$\exists x_0 \in \mathcal{D} : f_i(x_0) < 0, \quad i = 1, \dots, m, \quad h_i(x_0) = 0, \quad i = 1, \dots, p.$$

If the primal problem (8.1) is convex, and satisfies the weak Slater's condition, then strong duality holds, that is, $p^ = d^*$.*

We also need minimax theorem to help change the order of minimization and maximization:

Theorem 3.2 (Minimax theorem) *Let $\phi(x, y) : \mathcal{X} \times \mathcal{Y} \rightarrow \mathbf{R}$. Define*

$$p^* = \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} \phi(x, y)$$

and

$$d^* = \min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \phi(x, y)$$

It holds that the gap $p^ - d^*$ is zero if:*

- \mathcal{X}, \mathcal{Y} are both convex, and one of them is compact.
- The function ϕ is convex-concave: $\phi(\cdot, y)$ is convex for every $y \in \mathcal{Y}$, and $\phi(x, \cdot)$ is concave for every $x \in \mathcal{X}$.
- The function ϕ is continuous.

In our case, the program is reduced to the situation that $m = 1$ and $p = 0$ for 3.1 and 3.2. By theorem 3.1 and theorem 3.2 we would have following statements.

Proposition 3.3 *Suppose that (3.1), (3.2), and (3.3) are convex programs. Then it holds that*

- (i) *For any $b > 0$, there exists $\lambda \geq 0$ such that program (3.1) and program (3.3) are equivalent;*
- (ii) *For any $R > 0$, there exists $\lambda \geq 0$ such that program (3.2) and program (3.3) are equivalent.*

Proof:

(i) The primal program is

$$\min_{\theta \in \mathbf{R}^d} \|\theta\|_1 \quad \text{such that} \quad \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 - b^2 \leq 0. \quad (3.4)$$

The corresponding dual program is

$$\max_{\lambda > 0} \min_{\theta \in \mathbf{R}^d} \|\theta\|_1 + \lambda \left\{ \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 - b^2 \right\}. \quad (3.5)$$

By theorem 3.2 we can exchange the min and max:

$$\max_{\lambda > 0} \min_{\theta \in \mathbb{R}^d} \|\theta\|_1 + \lambda \left\{ \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 - b^2 \right\} \quad (3.6)$$

$$= \min_{\theta \in \mathbb{R}^d} \max_{\lambda > 0} \|\theta\|_1 + \lambda \left\{ \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 - b^2 \right\} \quad (3.7)$$

$$= \min_{\theta \in \mathbb{R}^d} \|\theta\|_1 + \lambda^* \left\{ \frac{1}{2n} \|y - \mathbf{X}\theta\|_2^2 - b^2 \right\} \text{ for some } \lambda^* > 0. \quad (3.8)$$

Then by theorem 3.1, we know that program (3.1) and program (3.3) are equivalent.

(ii) The proof is identical to (i). ■

3.3 Proof of the minimax theorem

Lemma 3.4 *It holds that*

$$\max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} \phi(x, y) \leq \min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \phi(x, y).$$

Proof: For any $x \in \mathcal{X}$, $y \in \mathcal{Y}$

$$\min_{x' \in \mathcal{X}} \phi(x', y) \leq \phi(x, y) \leq \max_{y' \in \mathcal{Y}} \phi(x, y').$$

Since the inequality holds for any $x \in \mathcal{X}$ and any $y \in \mathcal{Y}$, we can take max on the left and take min on the right:

$$\max_{y \in \mathcal{Y}} \min_{x' \in \mathcal{X}} \phi(x', y) \leq \min_{x \in \mathcal{X}} \max_{y' \in \mathcal{Y}} \phi(x, y').$$
■

Lemma 3.5 *The following statements are equivalent:*

(1) *There exists $(x^*, y^*) \in \mathcal{X} \times \mathcal{Y}$ such that for any $x \in \mathcal{X}$, $y \in \mathcal{Y}$,*

$$\max_{y \in \mathcal{Y}} \phi(x^*, y) \leq \phi(x^*, y^*) \leq \min_{x \in \mathcal{X}} \phi(x, y^*).$$

(2) *The minimax equation holds:*

$$\max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} \phi(x, y) = \min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \phi(x, y)$$

and

$$x^* = \arg \min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \phi(x, y),$$

$$y^* = \arg \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} \phi(x, y).$$

Proof: (1) \implies (2): Take min and max on the left and right respectively:

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \phi(x^*, y) \leq \phi(x^*, y^*) \leq \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} \phi(x, y^*).$$

Together with lemma 3.4, we have

$$\phi(x^*, y^*) = \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} \phi(x, y) = \min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \phi(x, y).$$

(2) \implies (1): by the definition of x^* and y^* , we have

$$\max_{y \in \mathcal{Y}} \phi(x^*, y) = \min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \phi(x, y) = \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} \phi(x, y) = \min_{x \in \mathcal{X}} \phi(x, y^*).$$

Thus,

$$\phi(x^*, y^*) \leq \max_{y \in \mathcal{Y}} \phi(x^*, y) = \min_{x \in \mathcal{X}} \phi(x, y^*) \leq \phi(x^*, y^*),$$

which implies that $\phi(x^*, y^*) = \max_{y \in \mathcal{Y}} \phi(x^*, y) = \min_{x \in \mathcal{X}} \phi(x, y^*)$. ■

Note: Then the proof of theorem follows by proving that Lemma 3.2 (1) holds.