

# Lecture 3/4 – Control Flow Hijacking

University of Illinois  
ECE 422/CS 461

# Control Flow Hijacking

- Altering control flow of a target program to cause it to do what attacker wants
  - Identify a **value** that will be loaded into **PC** (%eip)
  - Overwrite it (e.g., to point to shellcode)
- The simple stack buffer overflow attack from last lecture overwrote return address on stack

# Goals

- By the end of this lecture you should:
  - Understand common vulnerabilities that lead to control flow hijacking
  - Understand common countermeasures to control flow hijacking and their limitations
  - Understand which countermeasure an (advanced) attack bypasses

# Defenses and Counter-Attacks

- Stack canaries
- Other forms of control flow hijacking
- Data Execution Prevention (DEP, W<sup>X</sup>)
- Return-to-libc and Return-Oriented Programming (ROP)
- Address Space Layout Randomization (ASLR)
- Heap Spray

```
void foo(int a, int b) {
    char buf1[16];
    gets(buf1);
}
```

buf1[] →

**foo:**

```
push    %ebp
mov     %esp, %ebp
sub     $16, %esp
.....
call    gets
leave
ret
```



```
void foo(int a, int b) {  
    char buf1[16];  
    gets(buf1);  
}
```

buf1[] →

**foo:**

```
push    %ebp  
mov     %esp, %ebp  
sub     $16, %esp  
.....  
call    gets  
leave  
ret
```



# Stack Canary

- **Idea:** detect return address overwrite
- Place special value (**canary**) before return address on the stack
- Check **canary** before executing ret
  - If return address is overwritten, so is canary



foo:

push %ebp

mov %esp, %ebp

push CANARY

sub \$16, %esp

.....

call gets

mov -4(%ebp), %eax

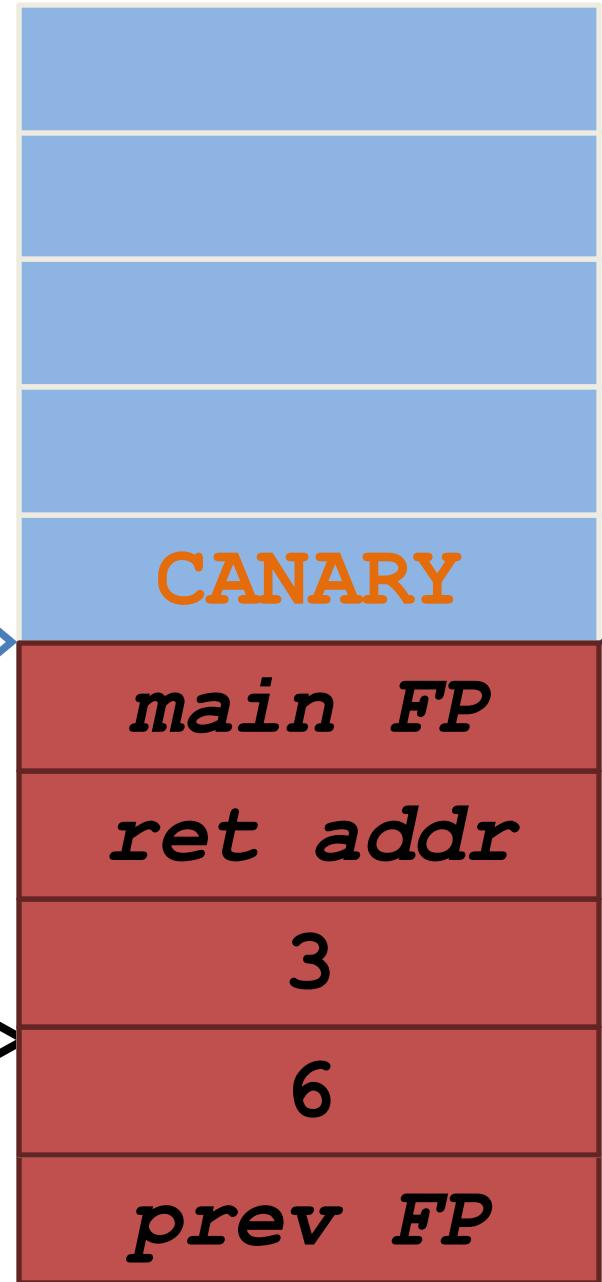
cmp CANARY, %eax

jne <stack\_chk\_fail>

leave

ret

buf1[]





foo:

push %ebp

mov %esp, %ebp

push CANARY

sub \$16, %esp

.....

call gets

mov -4(%ebp), %eax

cmp CANARY, %eax

jne <stack\_chk\_fail>

leave

ret

buf1[]

good

mor

ning

you

are

*hac*

*ked\0*

3

6

*prev FP*

User input: good morning you are hacked

# Stack Canary Value

- Exploit must contain the **canary** value to pass canary check
- **CANARY = 0**: can't strcpy past canary
- **CANARY = \n**: can't gets past canary
- **Random CANARY**: can't write past canary
  - Must not be discovered by attacker

# Stack Canary

- Low cost and modest performance penalty
- Enabled by default in GCC and Clang
  - To disable: `-fno-stack-protector`
- Requires re-compile (need source code)
- Only protects **return address** against **stack buffer overwrites**. *Does not protect against non-stack writes!*

# Control Flow Hijacking

- Altering control flow of a target program to cause it to do what attacker wants
  - Identify a **value** that will be loaded into **PC** (%eip)
  - Overwrite it (e.g., to point to shellcode)
- The simple stack buffer overflow attack from last lecture overwrote return address on stack
- Next: other control flow hijacking vulnerabilities

# Function Pointers

```
char text[128];  
void (*my_func)(int, int);  
  
my_func = &foo;  
*my_func(3, 6);    // equivalent to foo(3,6)
```

Q: Why does it defeat stack canary?

# C++ Virtual Function

```
class Shape {  
    virtual float area(void);  
};
```

```
class Circle : Shape {  
    float r;  
    Circle(float r) {this->r = r;}  
    float area() {return PI * r * r;}  
};
```

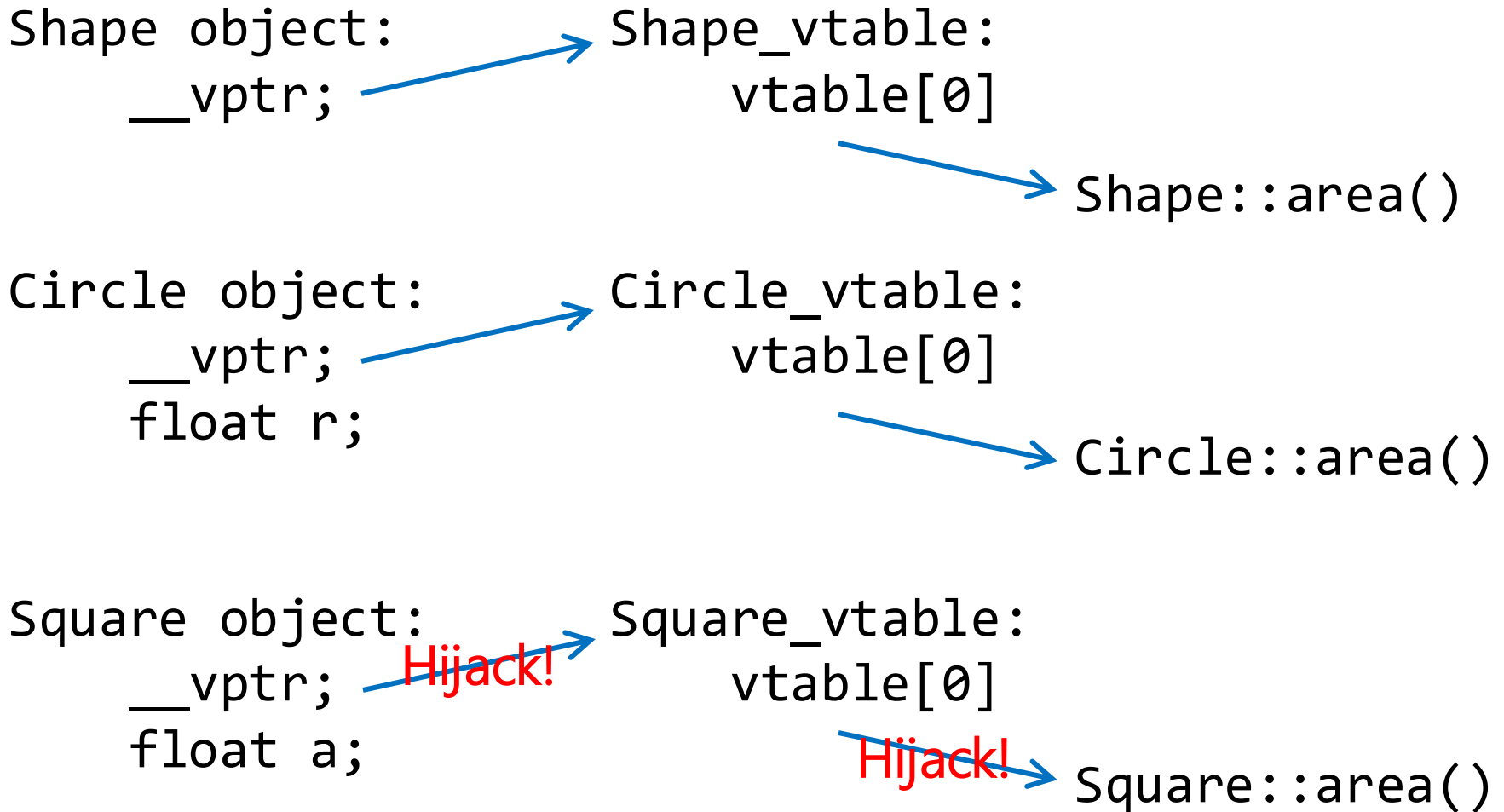
```
class Square : Shape {  
    float a;  
    Square(float a) {this->a = a;}  
    float area() {return a * a;}  
};
```

# C++ Class Polymorphism

```
Shape *s1, *s2;  
s1 = new Square(3);  
s2 = new Circle(5);  
...  
s1->area();           // calls Square::area()  
s2->area();           // calls Circle::area()
```

- How does the program know which `area()` method to call?
- Virtual classes have an invisible member variable: virtual function table pointer

# VTable: Array of Function Pointers





# Use after Free

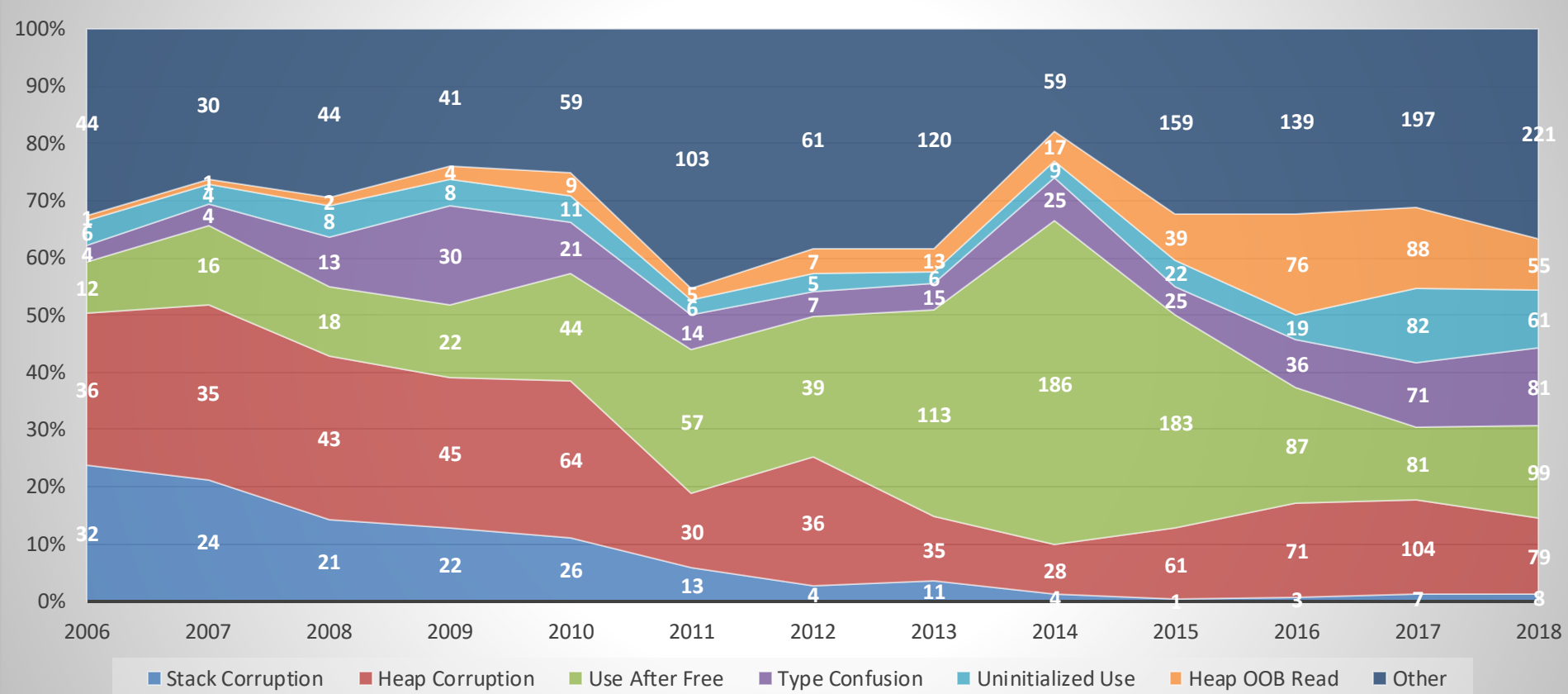
```
struct msg {                                struct student {  
    void (*my_func)(char*);                int uid;  
    char text[128];                        char name[128];  
};                                          };
```

```
student *s1 = malloc(sizeof(student));  
free(s1);  
msg *m1 = malloc(sizeof(msg));  
    // may occupy the same space as s1 (freed)  
s1->uid = updateUid();    // overwrite my_func
```

# Control Flow Hijacking Vulnerabilities

- Stack buffer overflows now less common
  - Easy-to-find bugs getting fixed
  - Use of unsafe functions deprecated
  - Stack-specific countermeasures (canary) helped
- Other variants have grown in popularity

Root cause of CVEs by patch year



Top root causes since 2016:

#1: use after free

#2: heap corruption

#3: type confusion

#4: uninitialized use

# Announcements and Review

- MP1 CP1 due today at 6 pm
  - No 24-hour auto extension for CP1
  - Autograder (v beta) is available on PrairieLearn
- Extra TA office hour every Wed 6-7 pm & Wed 1-4 pm when there is no discussion
- Last time:
  - Stack canary: detect overwrites to return address
  - Limitation: only protects return address (on stack).
  - Bypassed by heap buffer overflow, use after free, ...

# Control Flow Hijacking

- Altering control flow of a target program to cause it to do what attacker wants
- 1. Identify a **value** that will be loaded into **PC**
- 2. Overwrite it (to point to shellcode)
  - Deprecate unsafe functions and stack canaries
  - Other forms of control flow hijacking

# Control Flow Hijacking

- Altering control flow of a target program to cause it to do what attacker wants
- 1. Identify a **value** that will be loaded into **PC**
- 2. Overwrite it (to point to shellcode)
  - Deprecate unsafe functions and stack canaries
  - Other forms of control flow hijacking
- 3. Implicit step: shellcode runs
  - Can we prevent execution of injected code?

# Observation on Code Injection

- Root cause: confusion between code and data
  - Will be a recurring theme in this course
- Defense: distinguish code and data
  - Data should not be executable
  - Code need not be writable
    - \* Self-modifying code is a thing but uncommon

# Data Execution Prevention

- Make each memory region either writable or executable, never both
  - Make use of W (writable) and NX (no execute) bits in hardware page tables
  - Also called W<sup>X</sup> (write xor execute)
  - Supported by all major processors and OS
  - Malicious code can only be injected to writable regions. Attempting to run it will cause an error.





# Data Execution Prevention

- A memory region is writable or executable, never both
- No way to run shellcode ... right?
- Can still execute the program's code that is already there, but is that a problem?
- Isn't the program's code *good* code?
- Can *good* code do *bad* things?

# Return-to-libc Attacks

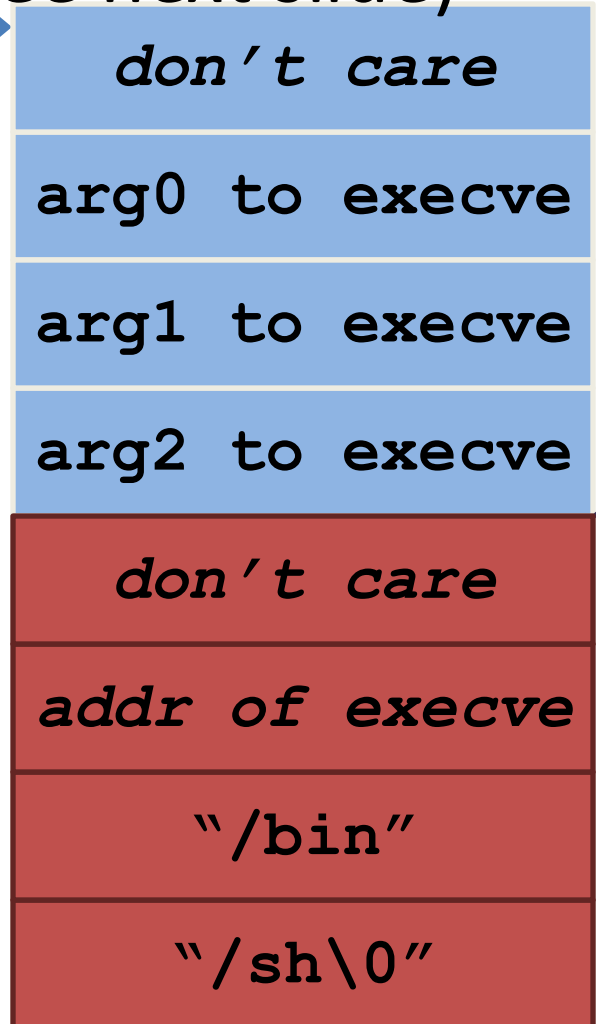
- Many programs rely on C Standard library to perform common tasks, e.g.,
  - Access files, kill processes, create users, execve, ...
- The attacker needs to (and may be able to) set up arguments on stack

# Return-to-libc Attacks

(buggy slide shown in class, see next slide)

**execve:**

```
push    %ebx
mov     0x10(%esp), %edx
mov     0xc(%esp), %ecx
mov     0x8(%esp), %ebx
mov     0xb, %eax
int     $0x80
```



# Return-to-libc Attacks

**execve:**

```
push    %ebx
mov     0x10(%esp), %edx
mov     0xc(%esp), %ecx
mov     0x8(%esp), %ebx
mov     0xb, %eax
int     $0x80
```



<i>don't care</i>
<del>ret</del> execve
<i>don't care</i>
arg0 to execve
arg1 to execve
arg2 to execve
"/bin"
"/sh\0"



# Return-to-libc Attacks

- Many programs rely on C Standard library to perform common tasks, e.g.,
  - Access files, kill processes, create users, execve, ...
- The attacker needs to (and may be able to) set up arguments on stack
  - Require precise knowledge about state of stack
- Does not work if libc is not loaded

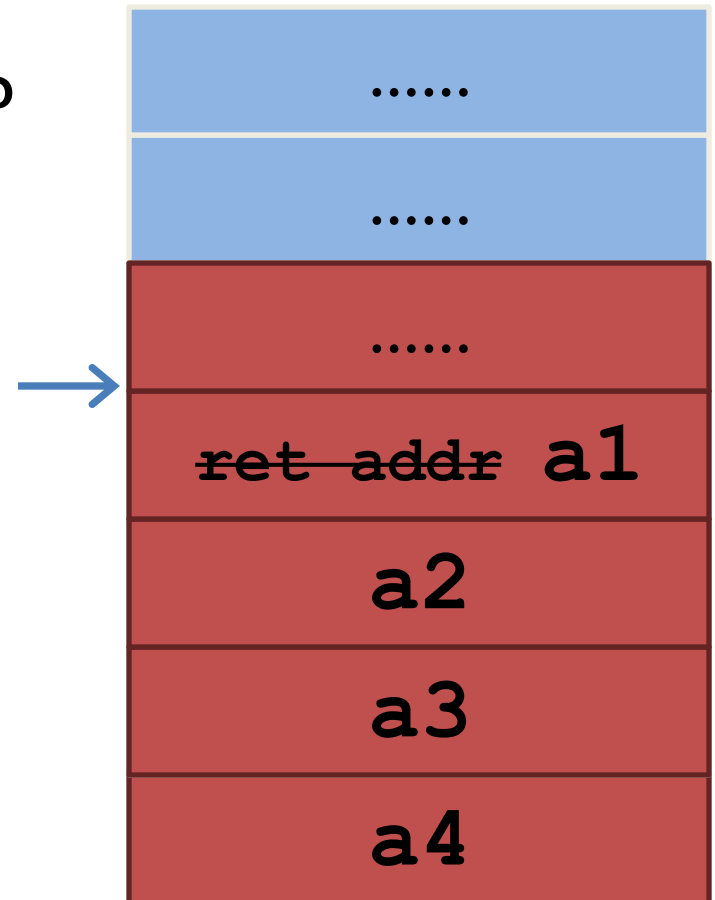
# Find Shellcode in Target Program?

6a	0b			push	\$0xb	
58				pop	%eax	
31	c9			xor	%ecx, %ecx	
31	d2			xor	%edx, %edx	
52				push	%edx	
68	2f	2f	73	68	push	\$0x68732f2f
68	2f	62	69	6e	push	\$0x6e69622f
89	e3			mov	%esp, %ebx	
cd	80			int	\$0x80	

# Another Look at Stack Buffer Overflow

- What happens after we jump to a location via ret?

```
foo:  
    ...  
    ret●
```





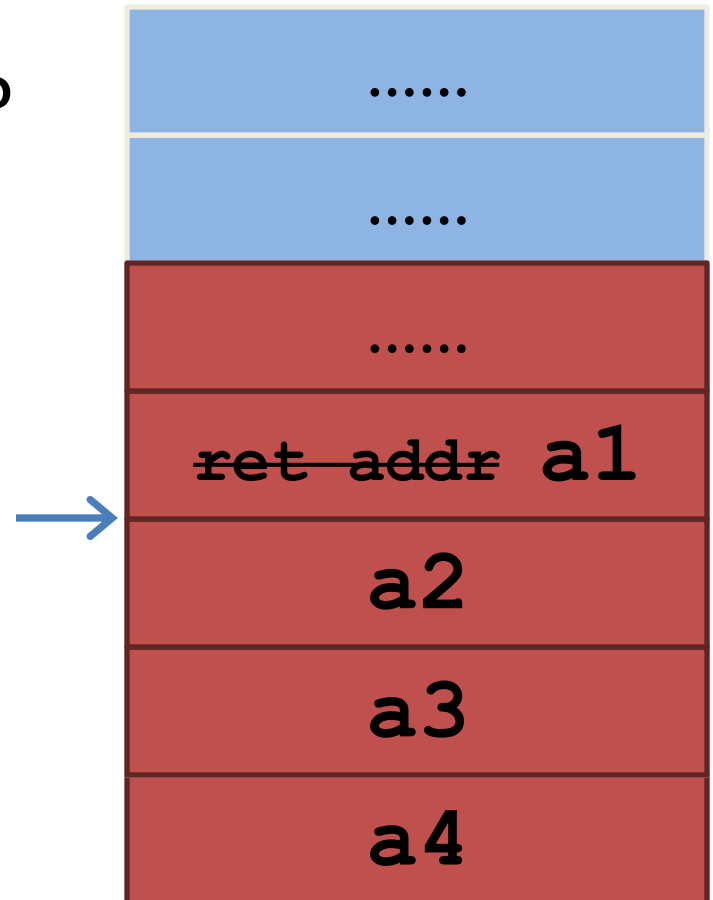
# Another Look at Stack Buffer Overflow

- What happens after we jump to a location via `ret`?
  - Jump to `a1`, execute there
  - `%ESP`  $\rightarrow$  next value on stack

`foo:`

`...  
ret`

`a1: some instr•`



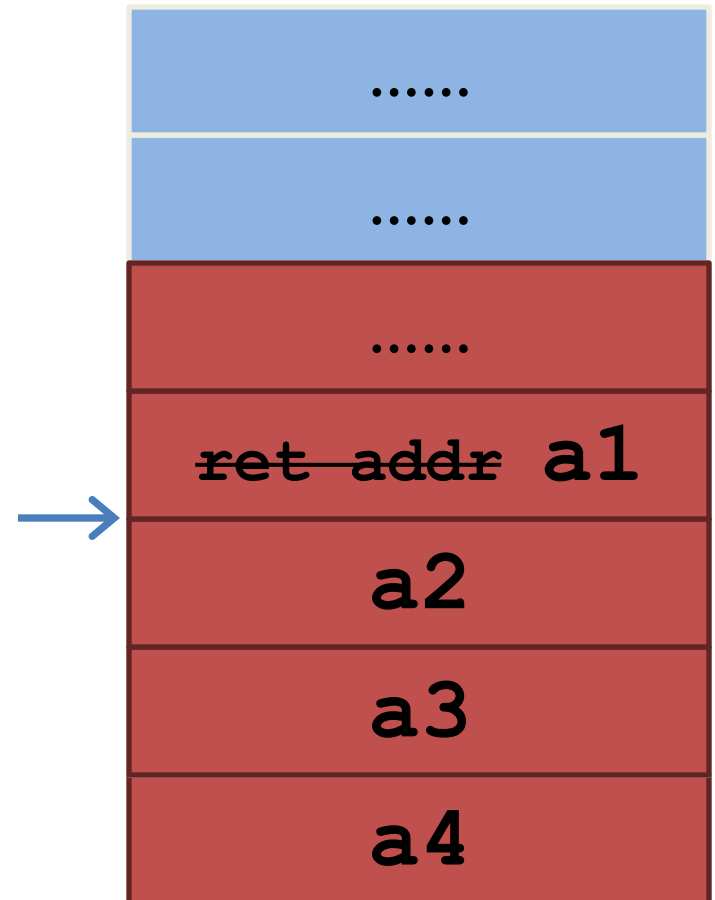
# Another Look at Stack Buffer Overflow

- What happens if the next instruction is also a ret?

```
foo:
```

```
    ...  
    ret
```

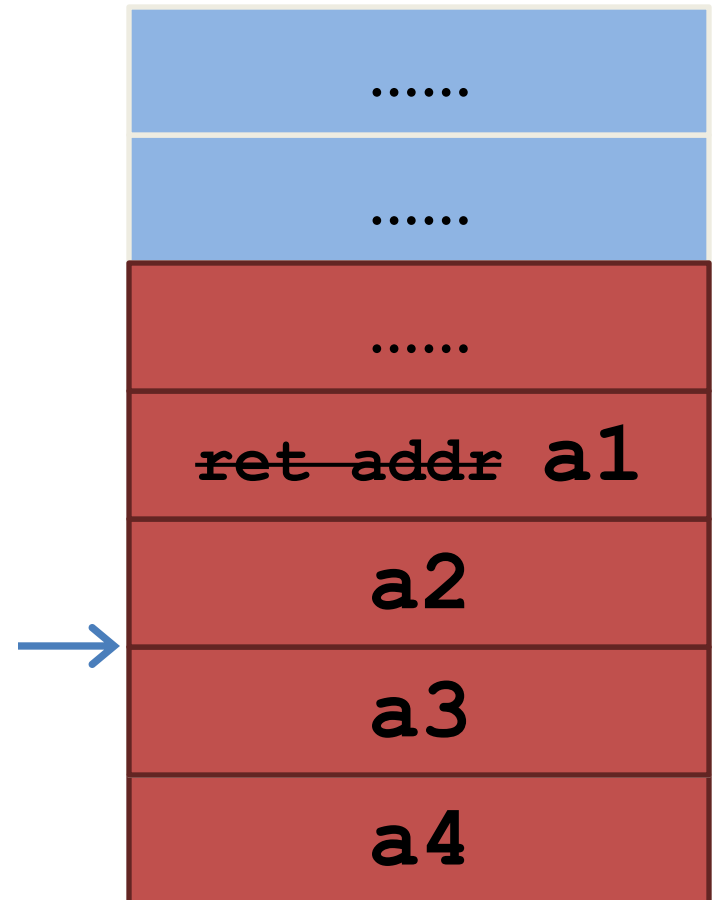
```
a1: some instr  
    ret●
```



# Another Look at Stack Buffer Overflow

- What happens if the next instruction is also a ret?
  - Jump to a2, execute there
  - %ESP → a3

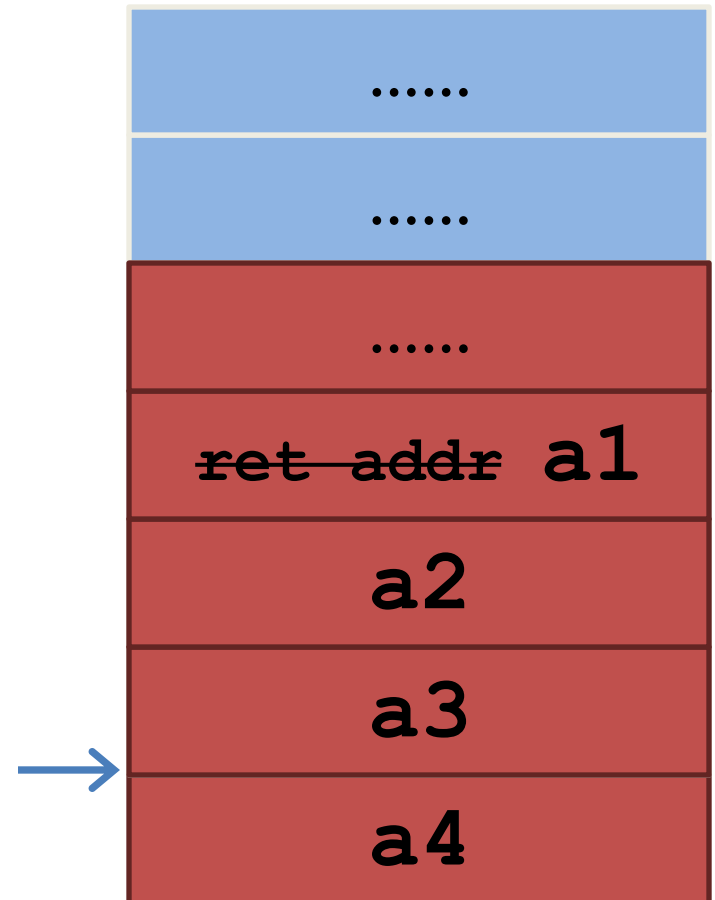
```
foo:                a2: instr●  
    ...  
    ret  
  
a1: some instr  
    ret
```



# Another Look at Stack Buffer Overflow

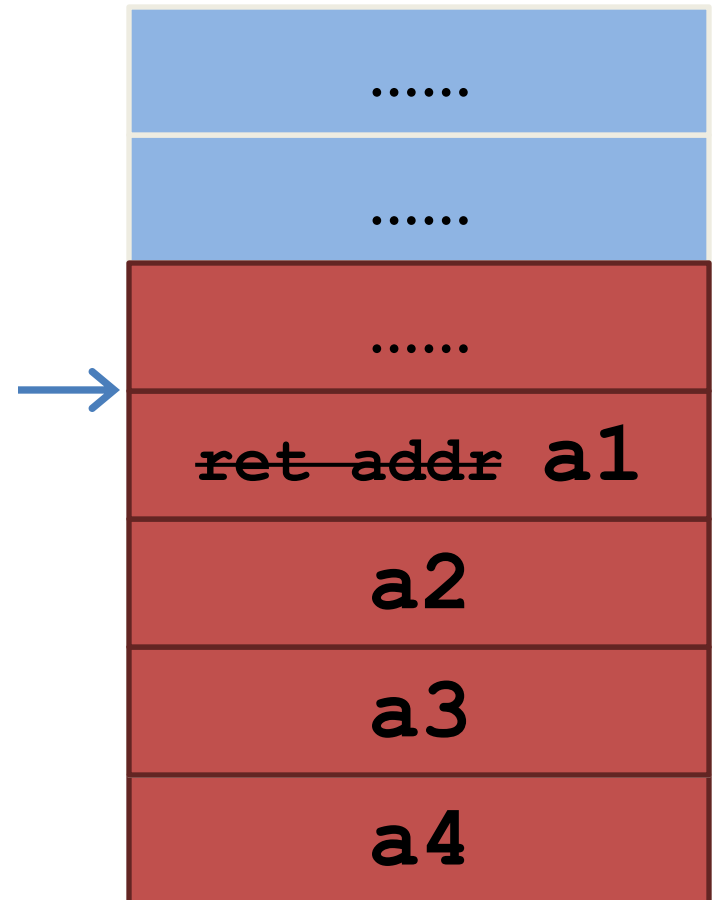
- What happens if the next instruction is also a ret?
  - Jump to a3, execute there
  - %ESP → a4

```
foo:                a2: instr
    ...              ret
    ret              a3: instr•
a1: some instr      ret
    ret
```



# Another Look at Stack Buffer Overflow

- Overwrite stack with a sequence of addresses
  - Jump (ret) to an address
  - Do some useful work
  - Repeat



# Return-Oriented Programming (ROP)

- Workflow
  - Dump executable portions of target program
  - Identify byte sequences ending in 0xC3 (ret)
    - Such a code fragment is called a **gadget**
  - Figure out what each gadget does — use a disassembler, e.g., <https://onlinedisassembler.com/>
  - Chain together useful gadgets

# Finding Gadgets

```
compy$ objdump -s /bin/ls
```

```
...
```

```
Contents of section .text:
```

```
...
```

```
804a530: 2404a120 430608c7 44241800 000000c7 $. . C...D$. . . . .
```

```
804a540: 442414c3 b90508c7 442410d3 b9050889 D$. . . . . D$. . . . .
```

```
...
```

```
804ad80: 5c240489 0424e855 8b000085 c00f8462 \$. . . $. U. . . . . b
```

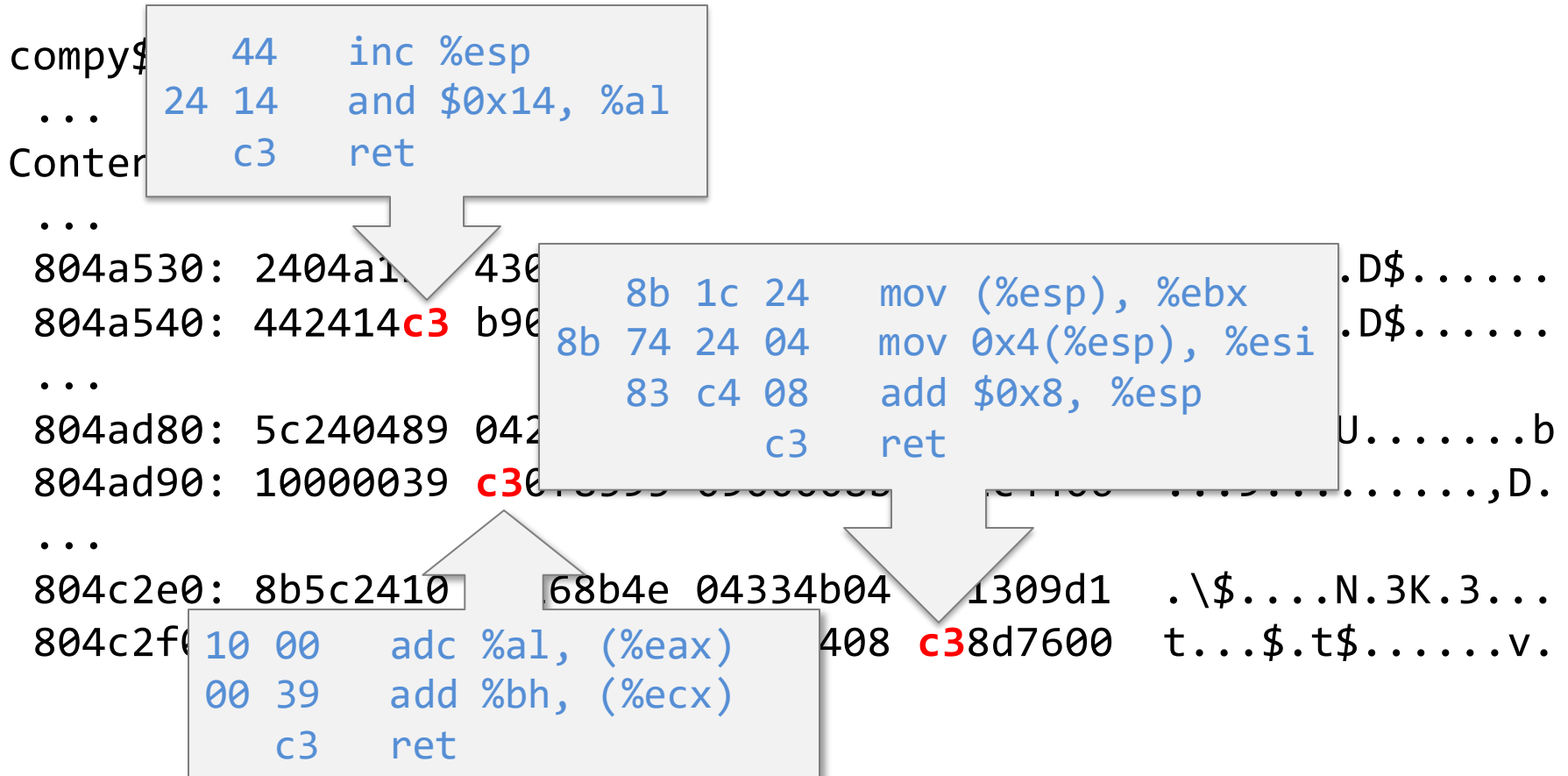
```
804ad90: 10000039 c30f8595 0900008b 0d2c4406 ...9. . . . . ,D.
```

```
...
```

```
804c2e0: 8b5c2410 8b168b4e 04334b04 331309d1 . \$. . . . N.3K.3. . .
```

```
804c2f0: 740e8b1c 248b7424 0483c408 c38d7600 t...$.t$. . . . . v.
```

# Finding Gadgets





# Tips on Finding Gadgets

- Suffix of a gadget is also a gadget
- Gadgets may not be “intended” code

```
89 5c 24 04      mov %ebx, 4(%esp)
```

```
89 c3           mov $eax, %ebx
```

```
04 89           add $89, %al
```

```
c3             ret
```

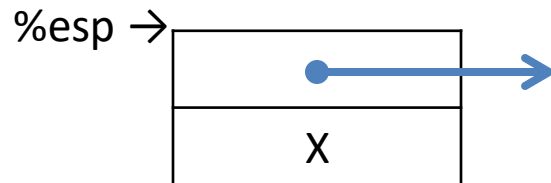
# Exercise with gadgets from /bin/l

- What can I do with these?

8b 1c 24	mov (%esp), %ebx	89 ea	mov %ebp, %edx
8b 74 24 04	mov 0x4(%esp), %esi	31 c0	xor %eax, %eax
83 c4 08	add \$0x8, %esp	8b 7c 24 04	mov 0x4(%esp), %edi
c3	ret	89 1f	mov %ebx, (%edi)
		89 0a	mov %ecx, (%edx)
b8 01 00 00 00	mov \$0x1, %eax	8b 54 24 20	mov 0x20(%esp), %edx
8b 34 24	mov (%esp), %esi	89 32	mov %esi, (%edx)
8b 7c 24 04	mov 0x4(%esp), %edi	83 c4 0c	add \$0xc, %esp
83 c4 08	add \$0x8, %esp	5b	pop %ebx
c3	ret	5e	pop %esi
		5f	pop %edi
89 eb	mov %ebp, %ebx	5d	pop %ebp
5d	pop %eax	c3	ret
c3	ret		

# Exercise with gadgets from /bin/ls

- Set EBP to X



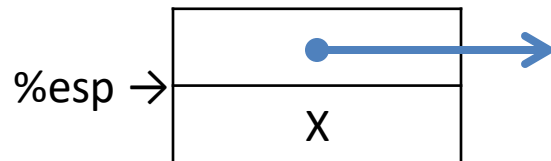
•c3

ret

```
89 ea      mov %ebp, %edx
31 c0      xor %eax, %eax
8b 7c 24 04 mov 0x4(%esp), %edi
89 1f      mov %ebx, (%edi)
89 0a      mov %ecx, (%edx)
8b 54 24 20 mov 0x20(%esp), %edx
89 32      mov %esi, (%edx)
83 c4 0c   add $0xc, %esp
5b        pop %ebx
5e        pop %esi
5f        pop %edi
5d        pop %ebp
c3        ret
```

# Exercise with gadgets from /bin/l

- Set EBP to X

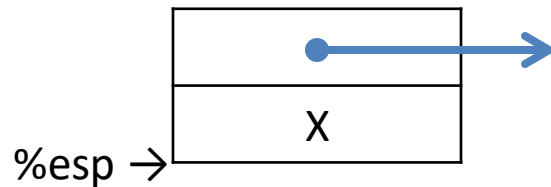


c3	ret
89 ea	mov %ebp, %edx
31 c0	xor %eax, %eax
8b 7c 24 04	mov 0x4(%esp), %edi
89 1f	mov %ebx, (%edi)
89 0a	mov %ecx, (%edx)
8b 54 24 20	mov 0x20(%esp), %edx
89 32	mov %esi, (%edx)
83 c4 0c	add \$0xc, %esp
5b	pop %ebx
5e	pop %esi
5f	pop %edi
• 5d	pop %ebp
c3	ret

# Exercise with gadgets from /bin/l

- Set EBP to X

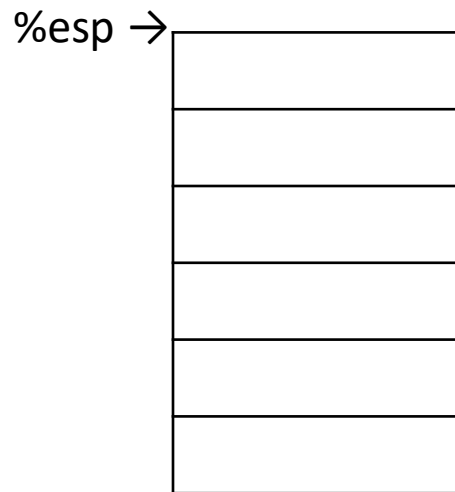
`%ebp = X`



c3	ret
89 ea	mov %ebp, %edx
31 c0	xor %eax, %eax
8b 7c 24 04	mov 0x4(%esp), %edi
89 1f	mov %ebx, (%edi)
89 0a	mov %ecx, (%edx)
8b 54 24 20	mov 0x20(%esp), %edx
89 32	mov %esi, (%edx)
83 c4 0c	add \$0xc, %esp
5b	pop %ebx
5e	pop %esi
5f	pop %edi
5d	pop %ebp
• c3	ret

# Exercise with gadgets from /bin/ls

- Set EBP to X and EBX to Y



• c3

ret

...

5b

5e

5f

5d

c3

89 eb

5d

c3

...

pop %ebx

pop %esi

pop %edi

pop %ebp

ret

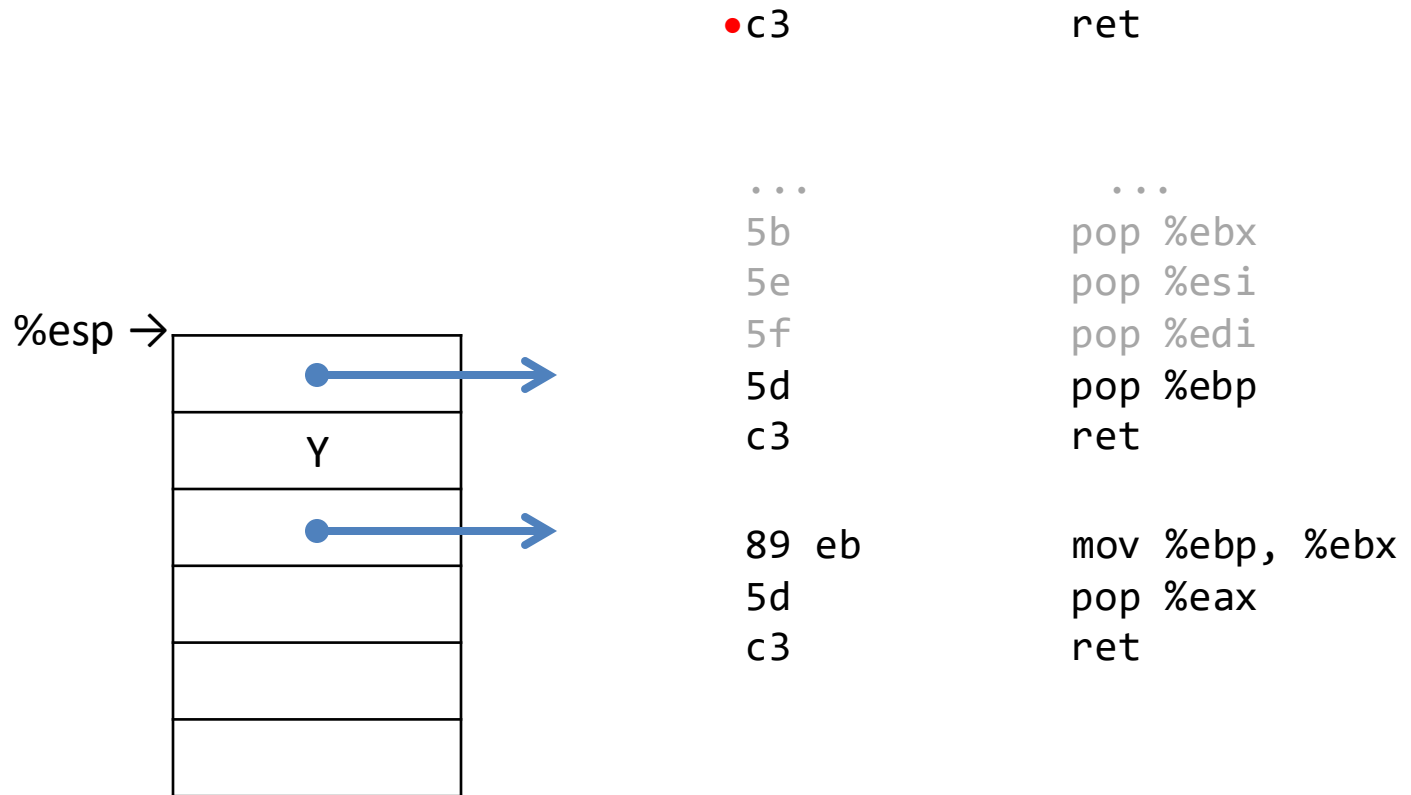
mov %ebp, %ebx

pop %eax

ret

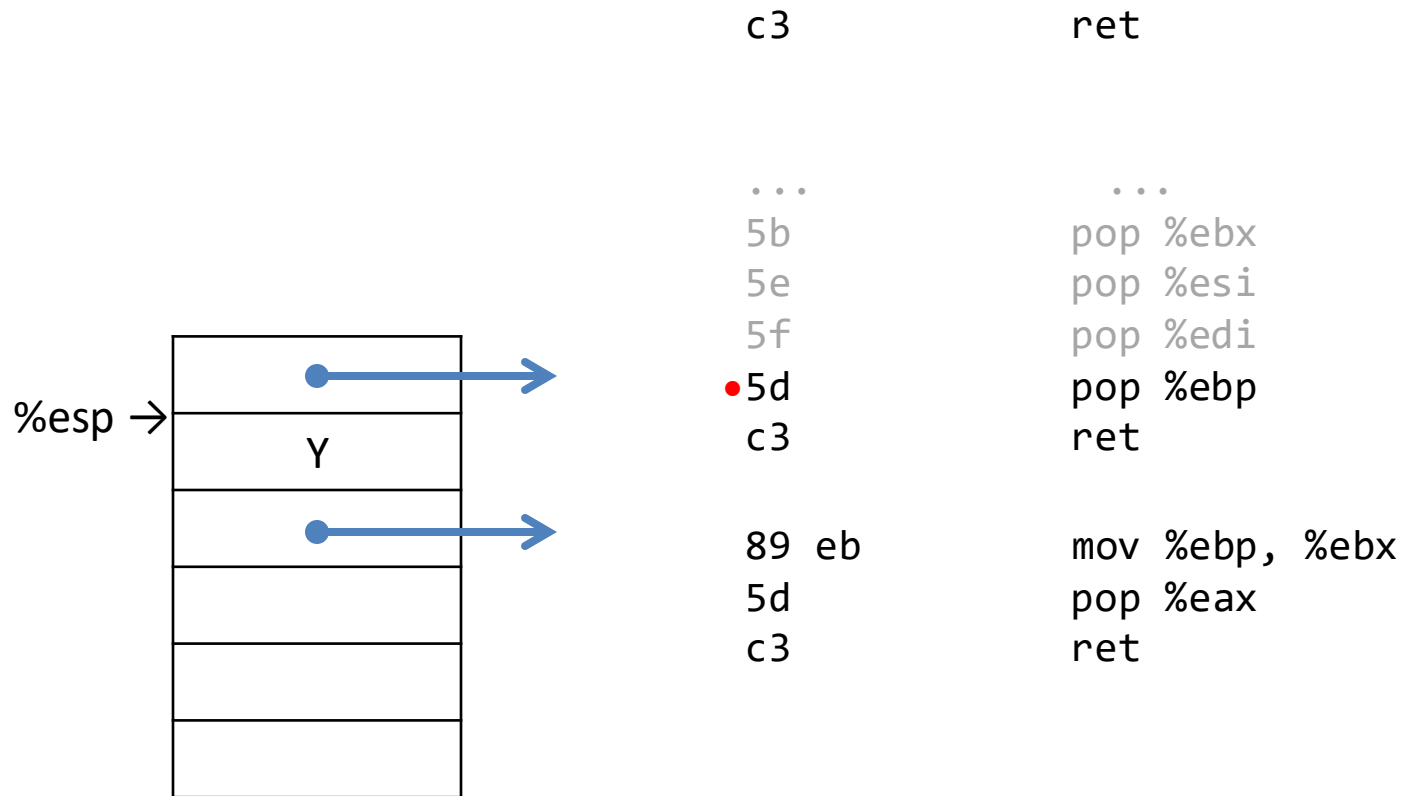
# Exercise with gadgets from /bin/l

- Set EBP to X and EBX to Y



# Exercise with gadgets from /bin/l

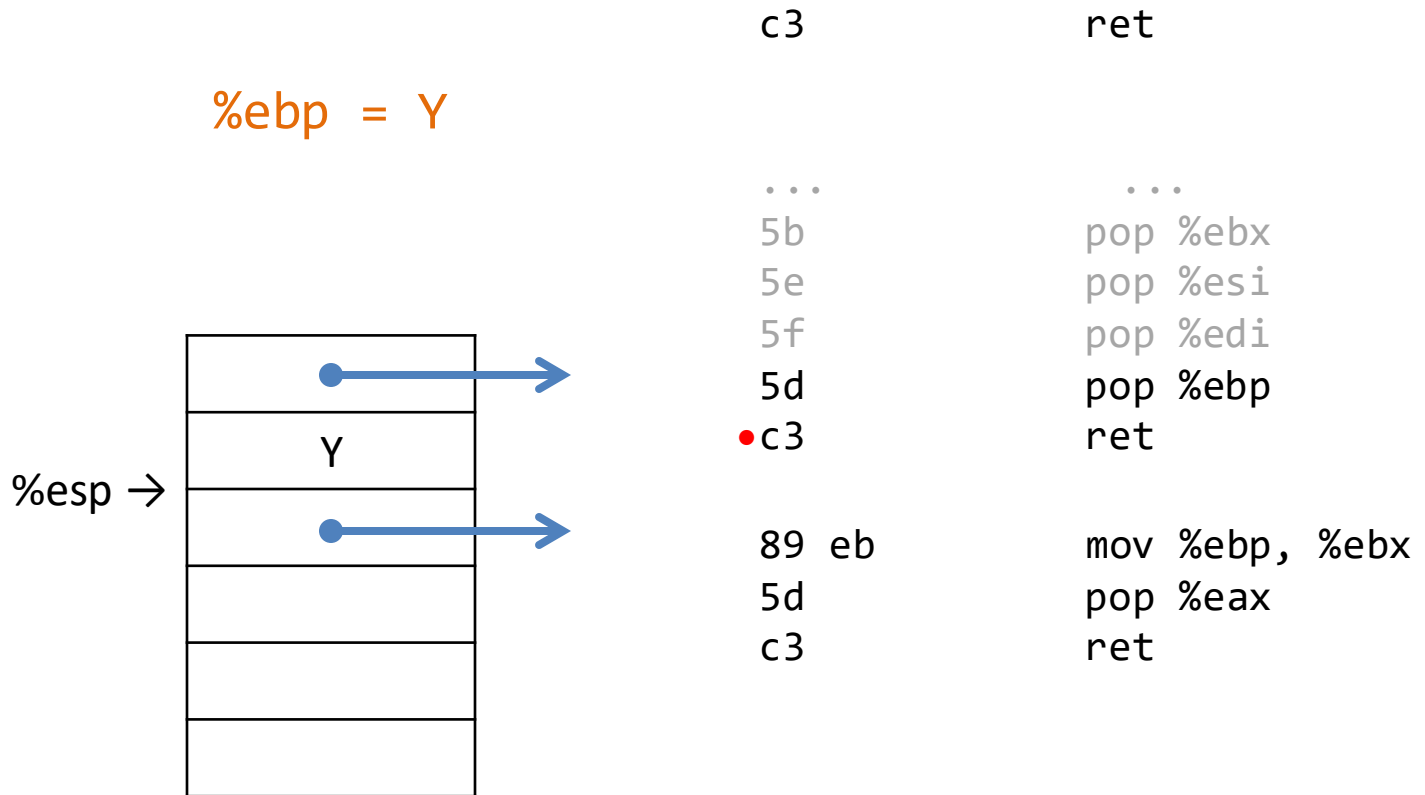
- Set EBP to X and EBX to Y





# Exercise with gadgets from /bin/l

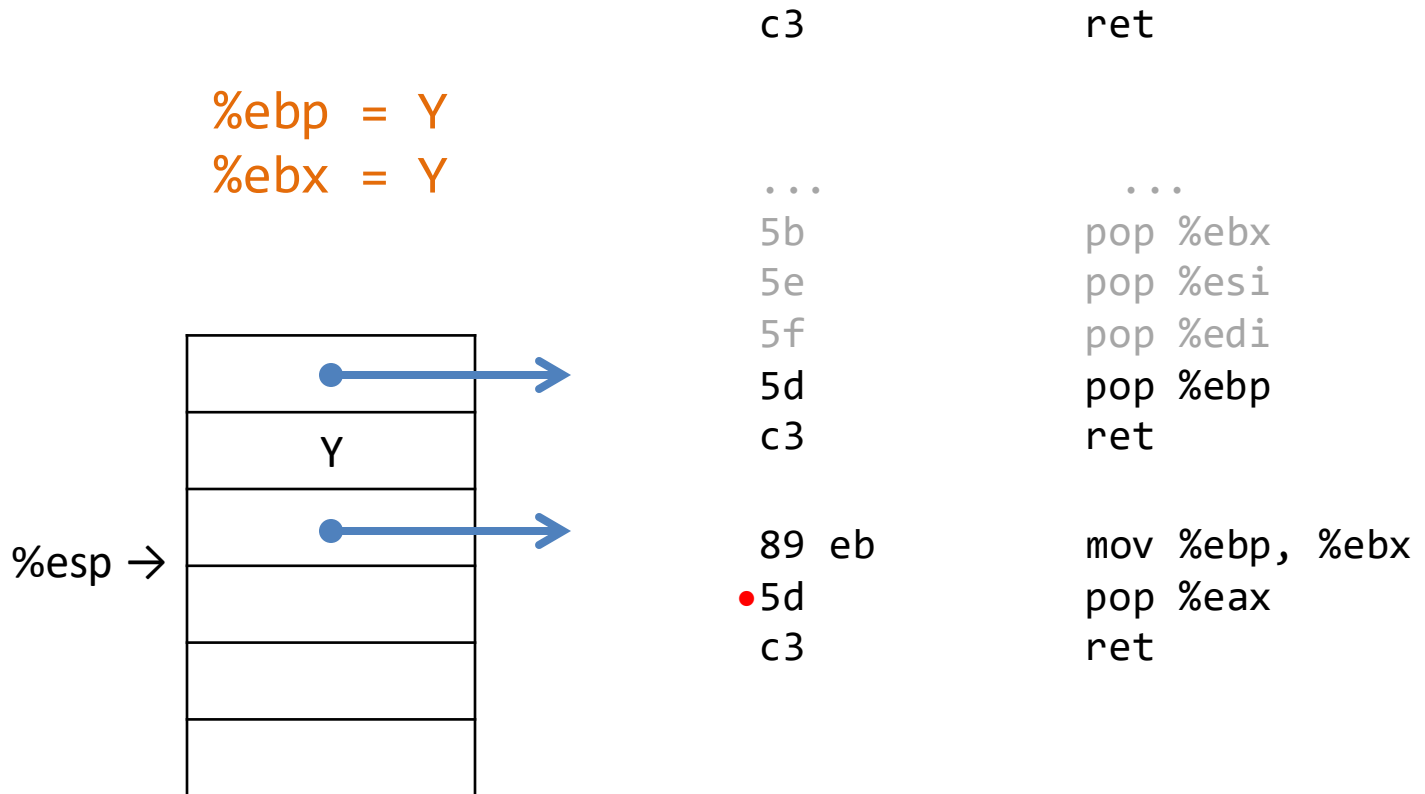
- Set EBP to X and EBX to Y





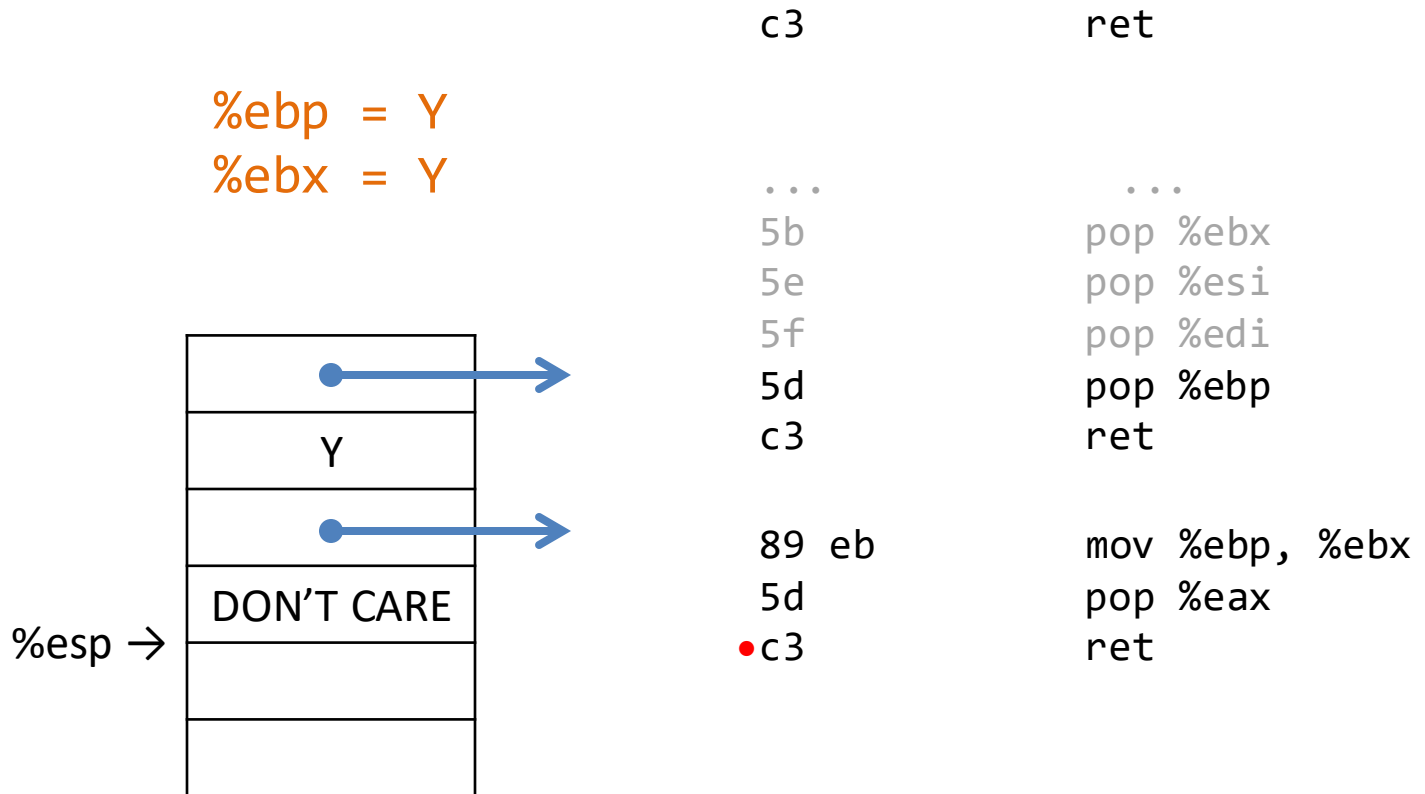
# Exercise with gadgets from /bin/ls

- Set EBP to X and EBX to Y



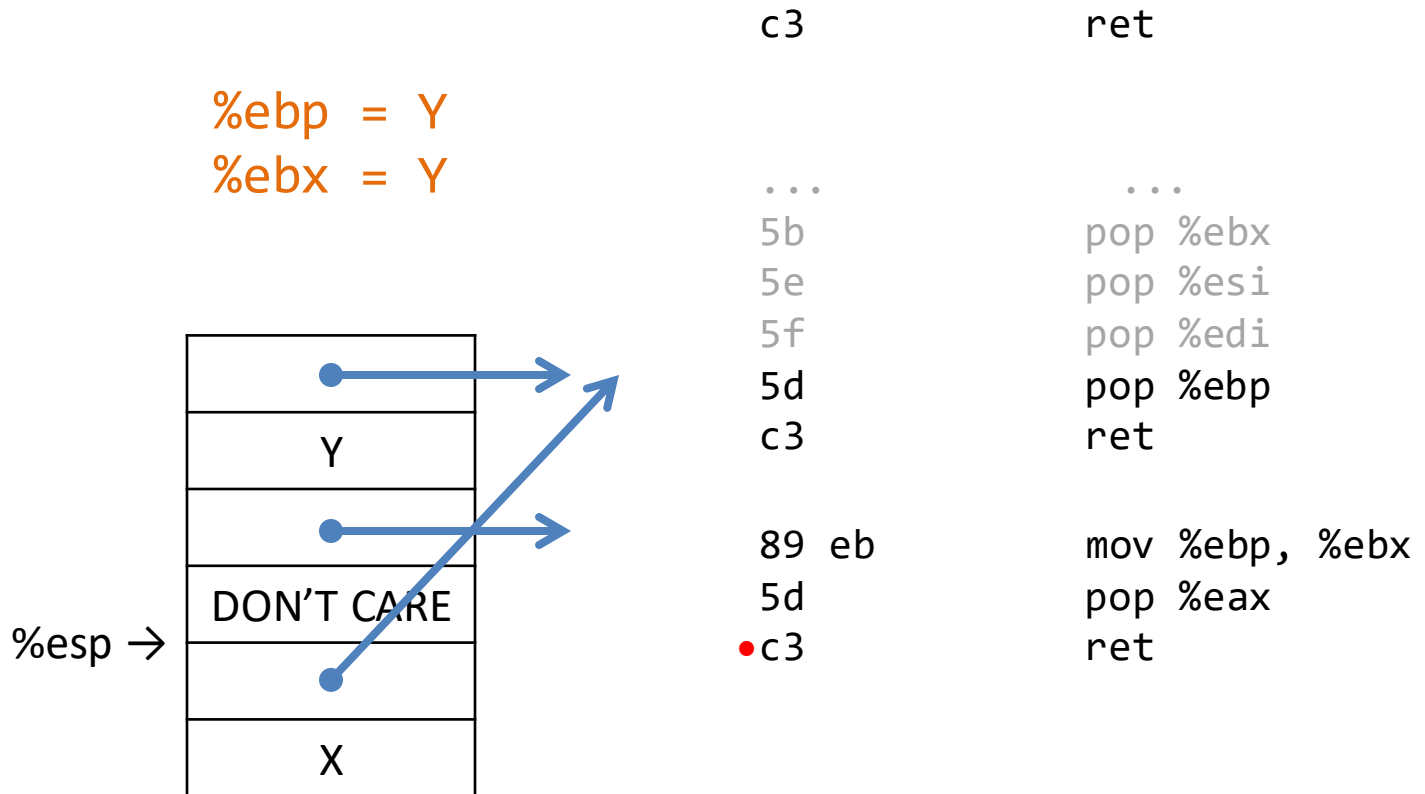
# Exercise with gadgets from /bin/l

- Set EBP to X and EBX to Y



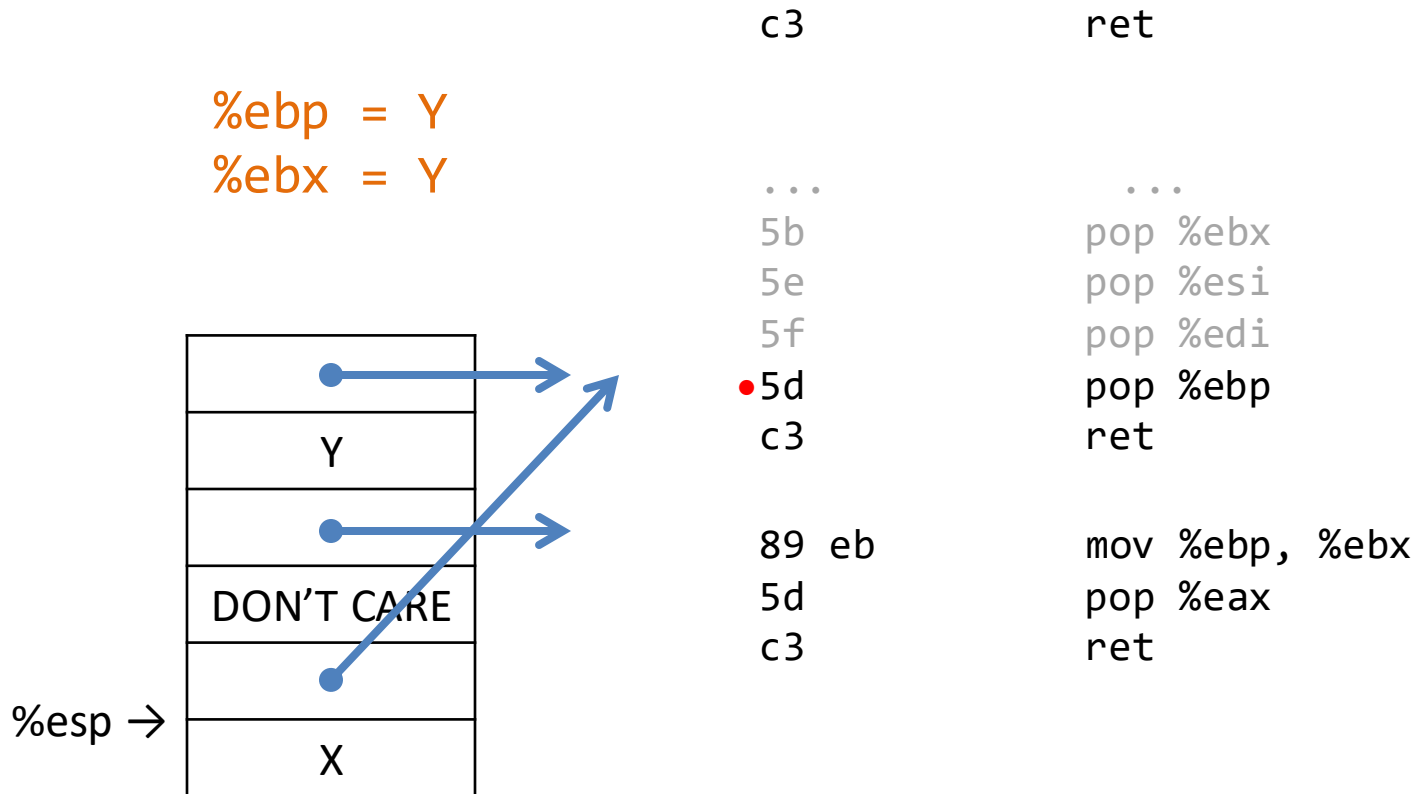
# Exercise with gadgets from /bin/l

- Set EBP to X and EBX to Y



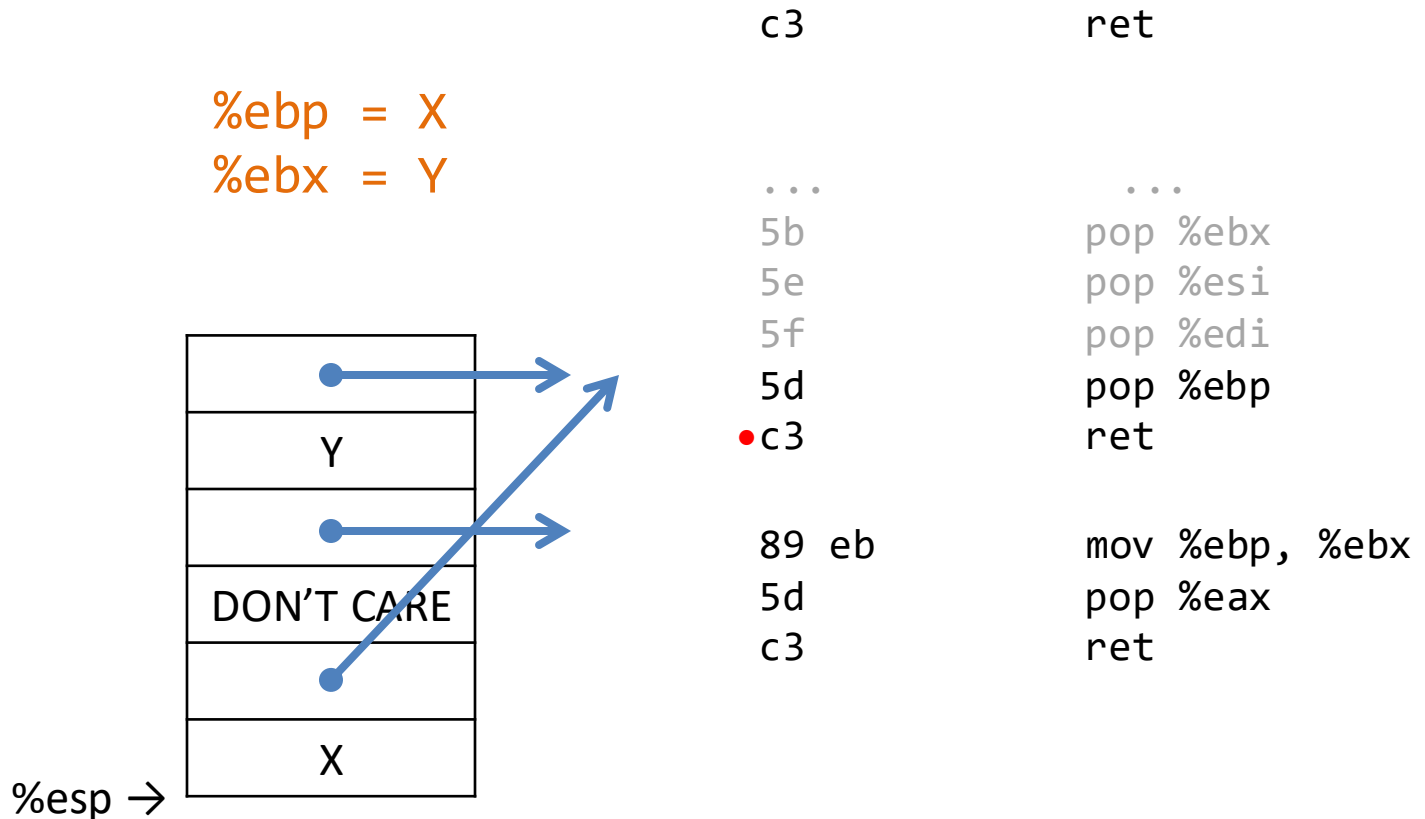
# Exercise with gadgets from /bin/l

- Set EBP to X and EBX to Y



# Exercise with gadgets from /bin/l

- Set EBP to X and EBX to Y



# Exercise with gadgets from /bin/l

- Write value X to location Y?

8b 1c 24	mov (%esp), %ebx	89 ea	mov %ebp, %edx
8b 74 24 04	mov 0x4(%esp), %esi	31 c0	xor %eax, %eax
83 c4 08	add \$0x8, %esp	8b 7c 24 04	mov 0x4(%esp), %edi
c3	ret	89 1f	mov %ebx, (%edi)
		89 0a	mov %ecx, (%edx)
b8 01 00 00 00	mov \$0x1, %eax	8b 54 24 20	mov 0x20(%esp), %edx
8b 34 24	mov (%esp), %esi	89 32	mov %esi, (%edx)
8b 7c 24 04	mov 0x4(%esp), %edi	83 c4 0c	add \$0xc, %esp
83 c4 08	add \$0x8, %esp	5b	pop %ebx
c3	ret	5e	pop %esi
		5f	pop %edi
89 eb	mov %ebp, %ebx	5d	pop %ebp
5d	pop %eax	c3	ret
c3	ret		



# Exercise with gadgets from /bin/l

- Write value X to location Y?

8b 1c 24	mov (%esp), %ebx	89 ea	mov %ebp, %edx
8b 74 24 04	mov 0x4(%esp), %esi	31 c0	xor %eax, %eax
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c3	ret	89 1f	mov %ebx, (%edi)
		89 0a	mov %ecx, (%edx)
b8 01 00 00 00	mov \$0x1, %eax	8b 54 24 20	mov 0x20(%esp), %edx
8b 34 24	mov (%esp), %esi	89 32	mov %esi, (%edx)
8b 7c 24 04	mov 0x4(%esp), %edi	83 c4 0c	add \$0xc, %esp
83 c4 08	add \$0x8, %esp	5b	pop %ebx
c3	ret	5e	pop %esi
		5f	pop %edi
89 eb	mov %ebp, %ebx	5d	pop %ebp
5d	pop %eax	c3	ret
c3	ret		

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then **G110**

**G116**

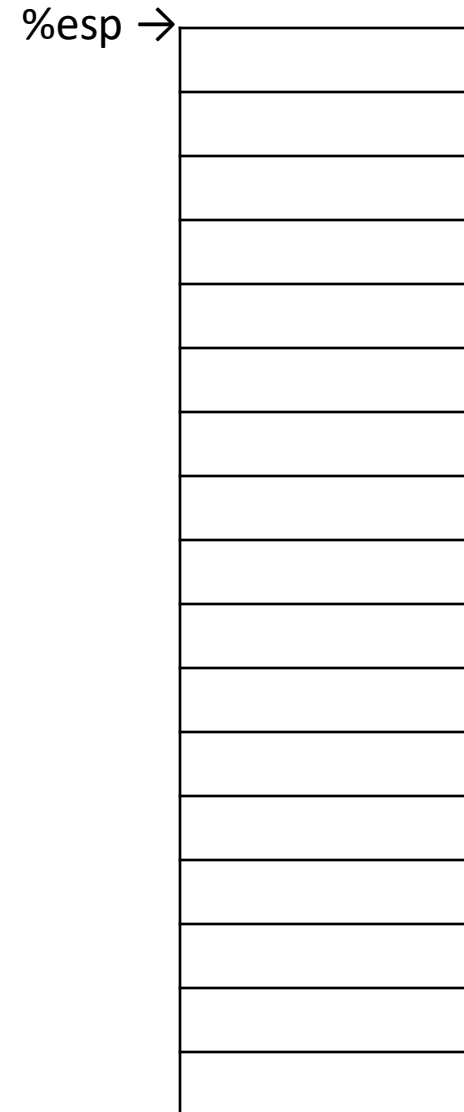
**G10c**

<b>G100:</b>	89 ea	mov %ebp, %edx
<b>G102:</b>	31 c0	xor %eax, %eax
<b>G104:</b>	8b 7c 24 04	mov 0x4(%esp), %edi
<b>G108:</b>	89 1f	mov %ebx, (%edi)
<b>G10a:</b>	89 0a	mov %ecx, (%edx)
<b>G10c:</b>	8b 54 24 20	mov 0x20(%esp), %edx
<b>G110:</b>	89 32	<b>mov %esi, (%edx)</b>
<b>G112:</b>	83 c4 0c	add \$0xc, %esp
<b>G115:</b>	5b	pop %ebx
<b>G116:</b>	5e	pop %esi
<b>G117:</b>	5f	pop %edi
<b>G118:</b>	5d	pop %ebp
<b>G119:</b>	c3	ret

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then **G110**

	G116	G10c	
G100:	89 ea		mov %ebp, %edx
G102:	31 c0		xor %eax, %eax
G104:	8b 7c 24 04		mov 0x4(%esp), %edi
G108:	89 1f		mov %ebx, (%edi)
G10a:	89 0a		mov %ecx, (%edx)
G10c:	8b 54 24 20		mov 0x20(%esp), %edx
G110:	89 32		mov %esi, (%edx)
G112:	83 c4 0c		add \$0xc, %esp
G115:	5b		pop %ebx
G116:	5e		pop %esi
G117:	5f		pop %edi
G118:	5d		pop %ebp
G119:	c3		ret



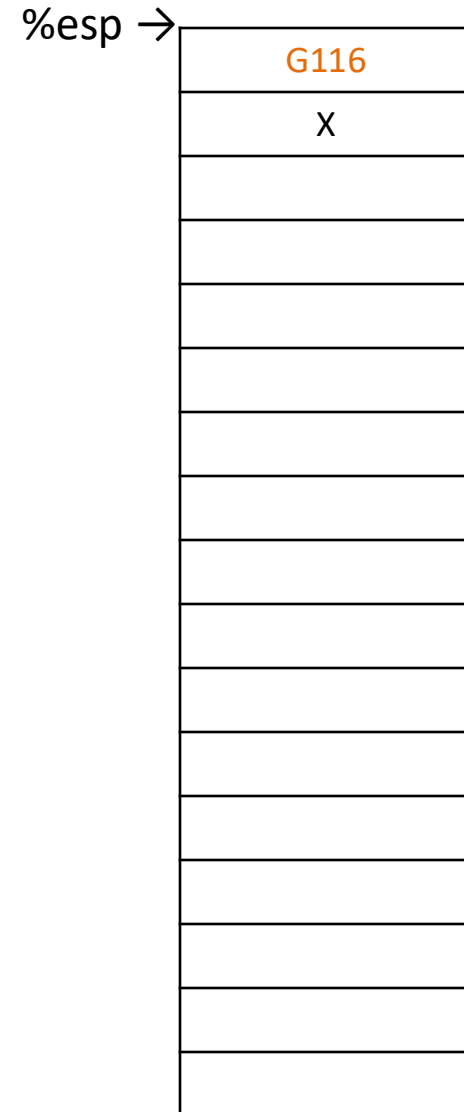
# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then **G110**

# G116

# G10c

G100:	89	ea			mov %ebp, %edx
G102:	31	c0			xor %eax, %eax
G104:	8b	7c	24	04	mov 0x4(%esp), %edi
G108:	89	1f			mov %ebx, (%edi)
G10a:	89	0a			mov %ecx, (%edx)
G10c:	8b	54	24	20	mov 0x20(%esp), %edx
G110:	89	32			mov %esi, (%edx)
G112:	83	c4	0c		add \$0xc, %esp
G115:	5b				pop %ebx
G116:	5e				pop %esi
G117:	5f				pop %edi
G118:	5d				pop %ebp
G119:	c3				ret



# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then **G110**

# G116

# G10c

G100:	89	ea			mov %ebp, %edx
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G116:	5e				pop %esi
G117:	5f				pop %edi
G118:	5d				pop %ebp
G119:	c3				ret

%esp →

[illegible]

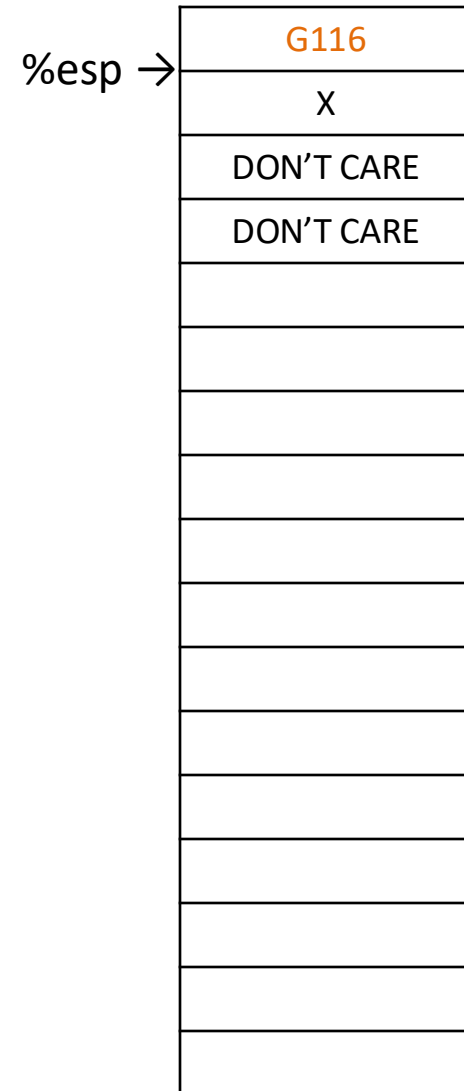
# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then **G110**

# G116

# G10c

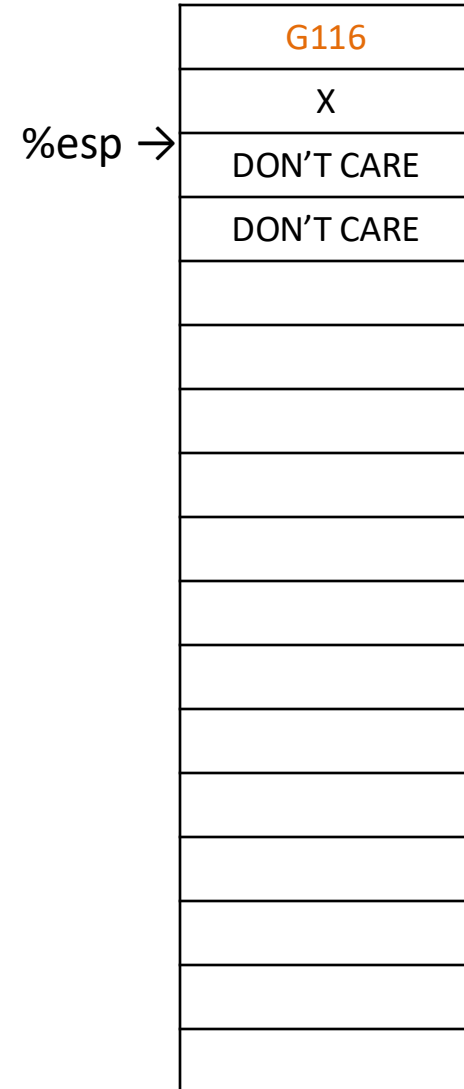
G100:	89	ea			mov %ebp, %edx
G102:	31	c0			xor %eax, %eax
G104:	8b	7c	24	04	mov 0x4(%esp), %edi
G108:	89	1f			mov %ebx, (%edi)
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G110:	89	32			mov %esi, (%edx)
G112:	83	c4	0c		add \$0xc, %esp
G115:	5b				pop %ebx
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G117:	5f				pop %edi
G118:	5d				pop %ebp
G119:	c3				ret



# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then `G110`  
`G10c`

G100:	89	ea			mov %ebp, %edx
G102:	31	c0			xor %eax, %eax
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# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110

# G10c

%esp →

G100:	89	ea			mov %ebp, %edx
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[illegible]



# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110

- Plan: **X** in **%esi**, **Y** in **%edx**, then **G110**

# G10c

%esp →

G100:	89	ea			mov %ebp, %edx
G102:	31	c0			xor %eax, %eax
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G117:	5f				pop %edi
G118:	5d				pop %ebp
G119:	c3				ret

[illegible]

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110

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G115:	5b				pop %ebx
G116:	5e				pop %esi
G117:	5f				pop %edi
G118:	5d				pop %ebp
G119:	c3				ret

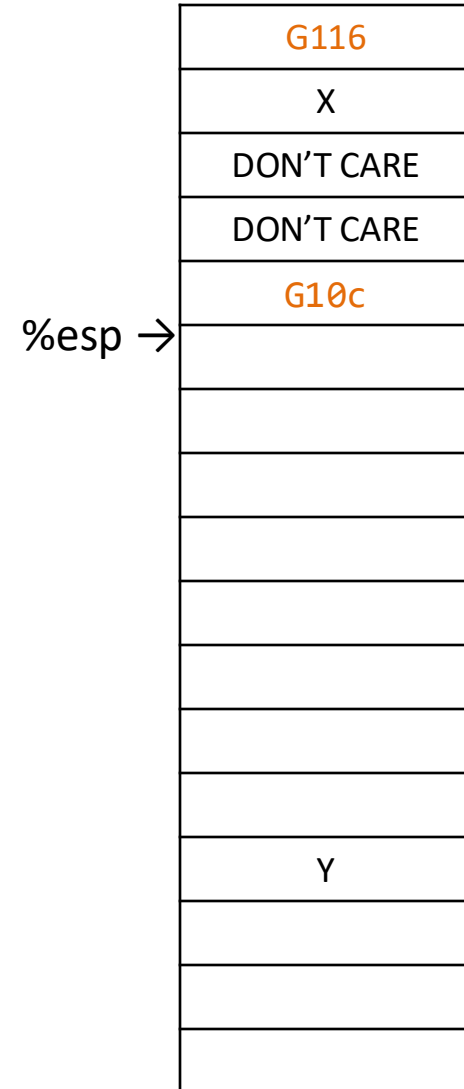
%esp →

[illegible]

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then `G110`  
`G10c`

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G117:	5f				pop %edi
G118:	5d				pop %ebp
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  - Plan: X in %esi, Y in %edx, then G110

```
G100: 89 ea    mov %ebp, %edx
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G10c: 8b 54 24 20 mov 0x20(%esp), %edx
G110: 89 32    mov %esi, (%edx)
G112: 83 c4 0c  add $0xc, %esp
G115: 5b      pop %ebx
G116: 5e      pop %esi
G117: 5f      pop %edi
G118: 5d      pop %ebp
G119: c3      ret
```

%esp →

G116
X
DON'T CARE
DON'T CARE
G10c
Y

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110

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G100: 89 ea    mov %ebp, %edx
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G10c: 8b 54 24 20  mov 0x20(%esp), %edx
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G112: 83 c4 0c    add $0xc, %esp
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G116: 5e      pop %esi
G117: 5f      pop %edi
G118: 5d      pop %ebp
G119: c3      ret
```

%esp →

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DON'T CARE
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G118: 5d      pop %ebp
G119: c3      ret
```

%esp →

G116
X
DON'T CARE
DON'T CARE
G10c
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
Y

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G100: 89 ea      mov %ebp, %edx
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G10a: 89 0a      mov %ecx, (%edx)
G10c: 8b 54 24 20  mov 0x20(%esp), %edx
G110: 89 32      mov %esi, (%edx)
G112: 83 c4 0c    add $0xc, %esp
G115: ●5b      pop %ebx
G116: 5e        pop %esi
G117: 5f        pop %edi
G118: 5d        pop %ebp
G119: c3        ret
```

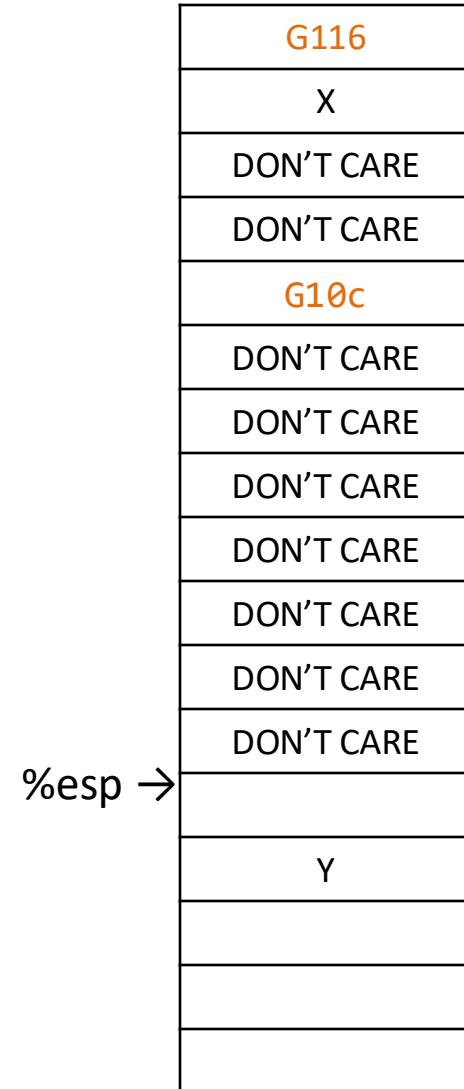
%esp →

G116
X
DON'T CARE
DON'T CARE
G10c
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
Y

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110

```
G100: 89 ea      mov %ebp, %edx
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G112: 83 c4 0c    add $0xc, %esp
G115: 5b        pop %ebx
G116: 5e        pop %esi
G117: 5f        pop %edi
G118: 5d        pop %ebp
G119: c3        ret
```





# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110

```
G100: 89 ea      mov %ebp, %edx
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G117: 5f        pop %edi
G118: 5d        pop %ebp
G119: c3        ret
```

%esp →

G116
X
DON'T CARE
DON'T CARE
G10c
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
Next gadget?
Y

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110
  - Skip Y, then next gadget

```
G100: 89 ea      mov %ebp, %edx
G102: 31 c0      xor %eax, %eax
G104: 8b 7c 24 04  mov 0x4(%esp), %edi
G108: 89 1f      mov %ebx, (%edi)
G10a: 89 0a      mov %ecx, (%edx)
G10c: 8b 54 24 20  mov 0x20(%esp), %edx
G110: 89 32      mov %esi, (%edx)
G112: 83 c4 0c    add $0xc, %esp
G115: 5b        pop %ebx
G116: 5e        pop %esi
G117: 5f        pop %edi
G118: 5d        pop %ebp
G119: c3        ret
```

%esp →

G116
X
DON'T CARE
DON'T CARE
G10c
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
%esp →
Y

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110
  - Skip Y, then next gadget

```
G100: 89 ea    mov %ebp, %edx
G102: 31 c0    xor %eax, %eax
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G10c: 8b 54 24 20  mov 0x20(%esp), %edx
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G112: 83 c4 0c    add $0xc, %esp
G115: 5b      pop %ebx
G116: 5e      pop %esi
G117: 5f      pop %edi
G118: 5d      pop %ebp
G119: c3      ret
```

%esp →

G116
X
DON'T CARE
DON'T CARE
G10c
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
G118
Y
Next gadget

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110
  - Skip Y, then next gadget

```
G100: 89 ea    mov %ebp, %edx
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G112: 83 c4 0c    add $0xc, %esp
G115: 5b      pop %ebx
G116: 5e      pop %esi
G117: 5f      pop %edi
G118: •5d    pop %ebp
G119: c3      ret
```

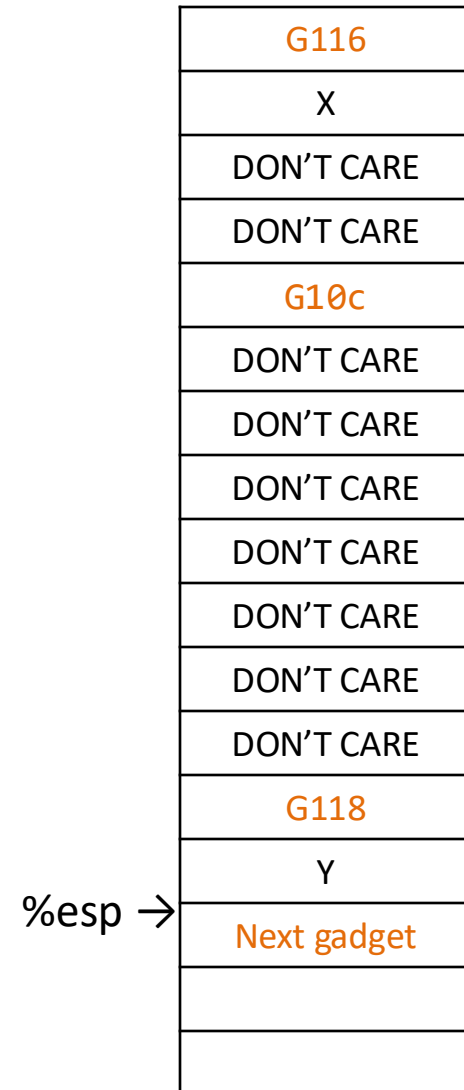
%esp →

G116
X
DON'T CARE
DON'T CARE
G10c
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
DON'T CARE
G118
Y
Next gadget

# Exercise with gadgets from /bin/l

- Write value X to location Y?
  - Plan: X in %esi, Y in %edx, then G110
  - Skip Y, then next gadget

```
G100: 89 ea    mov %ebp, %edx
G102: 31 c0    xor %eax, %eax
G104: 8b 7c 24 04  mov 0x4(%esp), %edi
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G112: 83 c4 0c    add $0xc, %esp
G115: 5b      pop %ebx
G116: 5e      pop %esi
G117: 5f      pop %edi
G118: 5d      pop %ebp
G119: c3      ret
```



# Return-Oriented Programming (ROP)

- **Gadgets** serve the role of instructions
- ROP programs assembled from gadgets
- There are ROP compilers to automate this

	Normal programming	Return-oriented programming
PC	%eip	%esp
No-op	nop	ret
Jump	jmp 4	pop %eax/%ebx/...

# Control Flow Hijacking

- Altering control flow of a target program to cause it to do what attacker wants
- 1. Identify a **value** that will be loaded into **PC**
- 2. Overwrite it (to point to shellcode)
  - Deprecate unsafe functions and stack canaries
  - Other forms of control flow hijacking
- 3. Shellcode runs
  - Data Execution Prevention (DEP/W^X)
  - Return-to-libc and Return-Oriented Programming

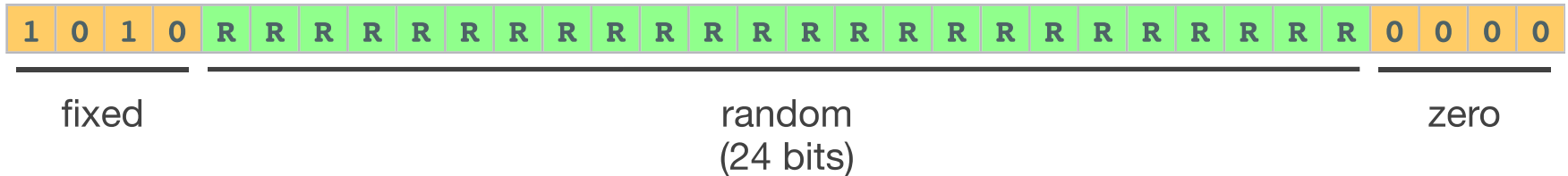
# Address Space Layout Randomization

- Randomize location of stack, heap, and code
  - So that the attacker does not know where the injected shellcode is located
  - Best to randomize on every launch
- Implemented (in some form) on most OSes
  - GCC and Clang: -fPIE
  - Code must be position independent
  - Binaries must be compiled to support ASLR

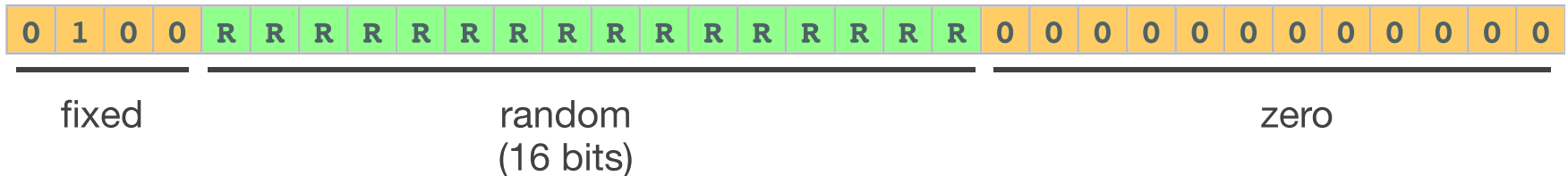


# 32-bit PaX ASLR (x86) Base Addresses

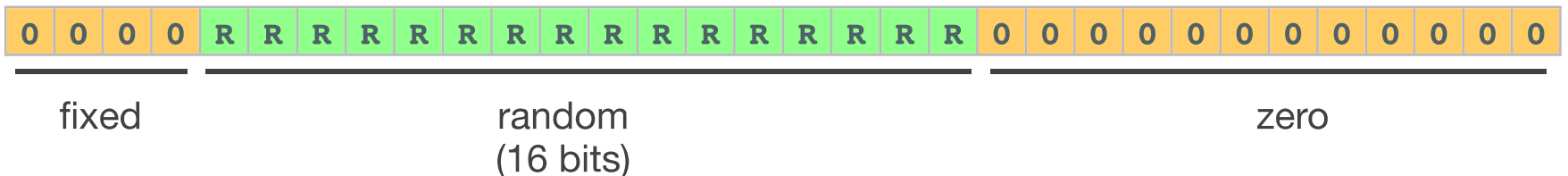
*Stack:*



*Mapped area:*



*Executable code, static variables, and heap:*

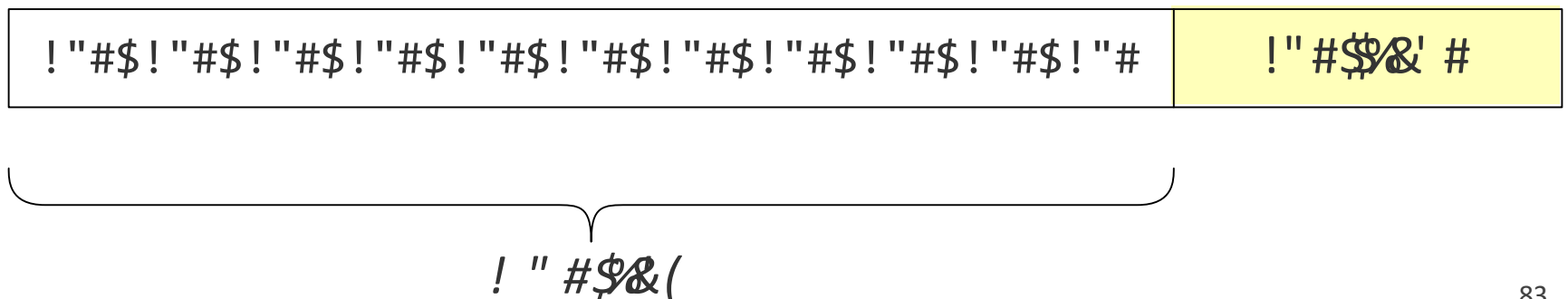


# Attacking ASLR

- Brute-force
  - Need to make sure an unsuccessful guess does not crash the target program
  - May be feasible on 32-bit systems: e.g.  $2^{16} = 65,536$  possible PaX code offsets
- Exploit other vulnerabilities to learn the random memory offset
- Heap spray

# Attacking ASLR: Heap Spray

- Suppose attacker can allocate objects on the victim's machine
- Fill the entire memory/heap with many instances of NOP sled + shellcode
- Overwrite return address / function pointer with arbitrary value → most likely land in a NOP sled



# Heap Spray

- Requires attacker to be able allocate objects on the victim's machine — how?
- Browsers are popular targets
  - Victim user visits a malicious website
  - Malicious site serves JavaScript code to browser
  - Browser is supposed to be a sandbox
  - Malicious JavaScript code performs heap spray and exploits a control flow vulnerability in browser

# Summary: Countermeasures

- Altering control flow of a target program to cause it to do what attacker wants
- 1. Identify a **value** that will be loaded into **PC**
- 2. Overwrite it (to point to shellcode)
  - Deprecate unsafe functions and stack canaries
  - Other forms of control flow hijacking
- 3. Shellcode runs
  - Data Execution Prevention (DEP/W<sup>X</sup>)
  - Return-to-libc and Return-Oriented Programming
  - Address Space Layout Randomization (ASLR)
  - Heap spray

# Summary: Countermeasures

- Combination of defenses more effective
  - DEP, to some extent, forces attacker to use ROP
  - ASLR randomizes where ROP gadgets are located
  - ROP happens on stack; stack canaries prevent some forms of stack overwrites
- To Learn More
  - Laszlo Szekeres, Mathias Payer, Tao Wei, and Dawn Song. *SoK: Eternal War in Memory*. 2013.

# References/Acknowledgements

- <https://users.ece.cmu.edu/~dbrumley/courses/18487-f14/www/powerpoint/>
- Return-to-libc demo <http://www.securitytube.net/video/258>
- <http://seclists.org/bugtraq/1997/Aug/63>
- <https://hovav.net/ucsd/dist/geometry.pdf>
- [https://www.usenix.org/legacy/publications/library/proceedings/sec98/full\\_papers/cowan/cowan.pdf](https://www.usenix.org/legacy/publications/library/proceedings/sec98/full_papers/cowan/cowan.pdf)
- <http://security.stackexchange.com/questions/20497/stack-overflows-defeating-canaries-aslr-dep-nx>
- <http://phrack.org/issues/58/4.html>