Query Auto-Completion (QAC)

Alfan Farizki Wicaksono Fakultas Ilmu Komputer, Universitas Indonesia

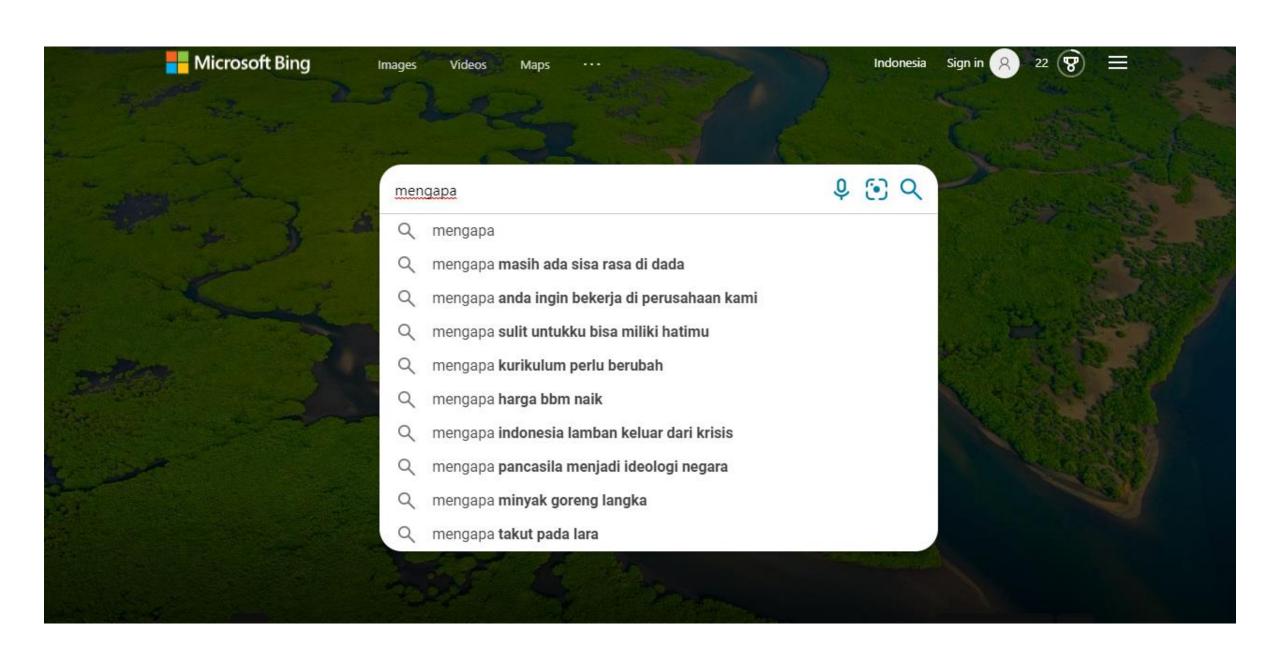


Q anak ui X

- Q anak ui
- anak ui tertabrak kereta
- Q anak uin
- Q anak ui meninggal
- anak ui sombong
- anak ui pintar
- anak ui hedon
- anak uir gantung diri
- anak ui gaji 8 juta
- anak uin viral

Penelusuran Google

Saya Lagi Beruntung



Mengapa Query Auto-Completion?

- Membantu user dalam formulasi query
- · Mengurangi keystrokes yang dibutuhkan untuk input query
- Membantu koreksi spelling error
- · Cache results! Reduce server load

High Level Algorithm

Given a query pattern P

- Retrieve set of candidates "matching" P from set 5 of possible target queries.
- Rank candidates by frequency.
- Possibly re-rank highest ranked candidates with more complex ranking
- Return the top-K highest ranking candidates as suggestions

QAC Modes

1. Prefix match

- 2. Substring match
- 3. Multi-term prefix match
- 4. Relaxed match

Contoh: Target -> "FIFA world cup 2022"

	Mode 1	Mode 2	Mode 3	Mode 4
FIFA wo	×	×	×	×
orl		×		
FI wor			×	×
FIFO walrd cu				×

Prefix Completion

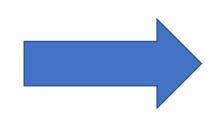
Prefix Completion

- Task: given a query prefix P, retrieve the top-K most frequent completions.
- Data: search query logs yang sudah dikumpulkan lama
- Requirements:
 - Efficient retrieval time required
 - Space efficient index

Tahap 1

Preprocess data by sorting query log in lexicographical order and based on frequency of unique queries.

bunnings bachelor in paradise bunnings bbc news bunnings bbc news bbc news big w big w bunnings



<bunnings, 4>

<bbc news, 3>

<bachelor in paradise, 1>

Tahap 2 - Trie + Frequency Table

Insert all unique queries and their frequencies into a trie data structure

Trie Data Structure:

- A tree data structure that contains a set of strings
- Edges of the tree are labeled
- Children of nodes are ordered
- A path from root to a node represents a prefix of all strings in the subtree starting at that node

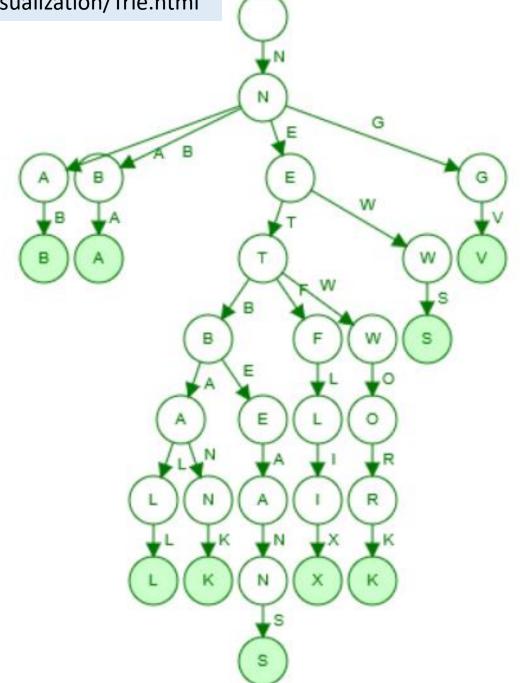
https://www.cs.usfca.edu/~galles/visualization/Trie.html

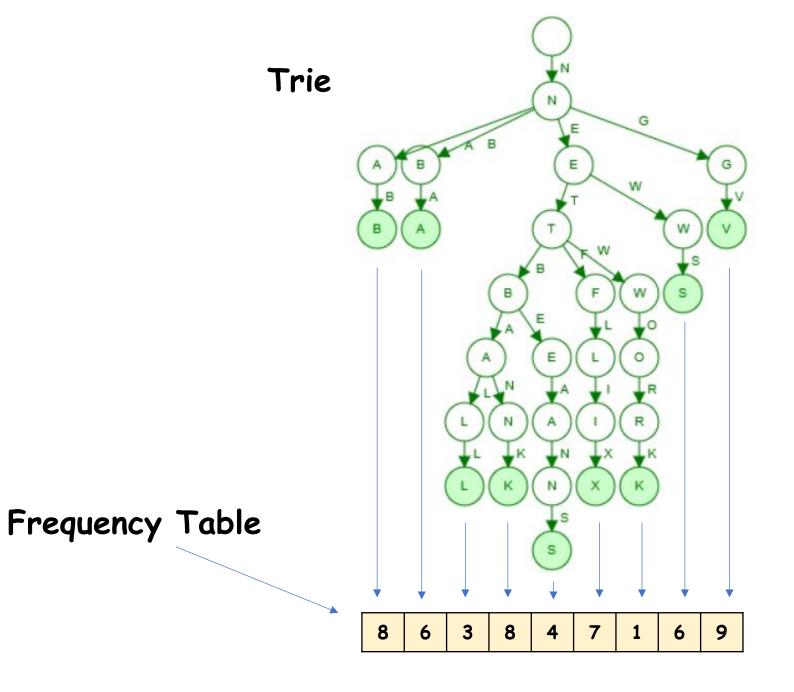
Trie Visualization

Set of strings

```
<nba, 6>
<news, 6>
<news, 6>
<nab, 8>
<ngv, 9>
<netflix, 7>
<netbank, 8>
<network, 1>
<netball, 3>
<netbeans, 4>
```

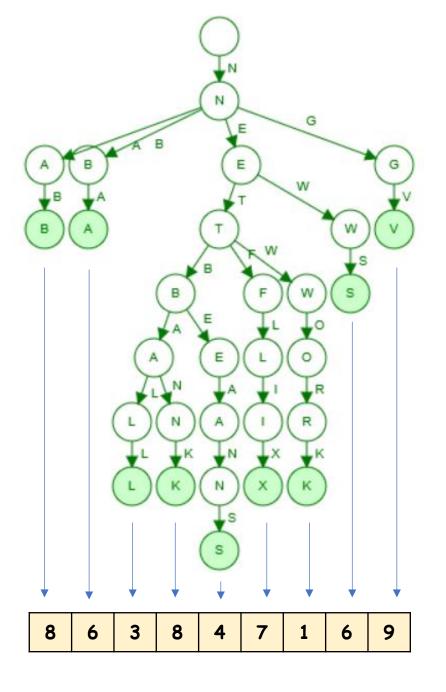






Implementasi Trie dengan Python

```
class Trie(object):
    """Abstraksi struktur data Trie"""
    def init (self):
        self.root = TrieNode("")
    def insert(self, word, freq):
        """tambahkan sebuah term pada Trie"""
        node = self.root
        # loop ke setiap karakter di word
        for char in word:
            # jika ada di salah satu anak
            # langsung node pindah ke anak tersebut
            if char in node.children:
                node = node.children[char]
            else:
                # jika tidak ditemukan, buat node baru
                new node = TrieNode(char)
                node.children[char] = new node
                node = new node
        node.freq += freq
```



net

search

Rekomendasikan Top-K queries, K = 4

net

search

netbank netflix netbeans netball

Untuk sebuah pattern P, pencarian node pada Trie yang merepresentasikan subtree dengan prefix P dapat dilakukan dengan kompleksitas O(|P|)

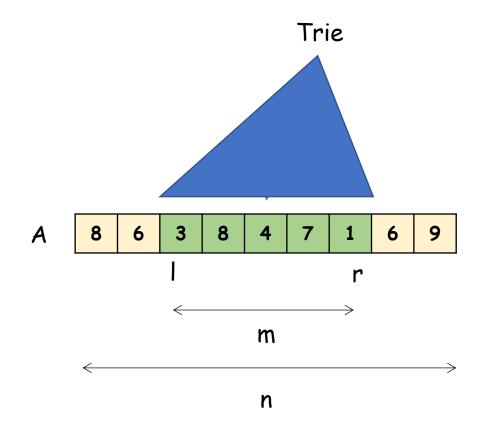
Range Maximum Queries (RMQ)

net

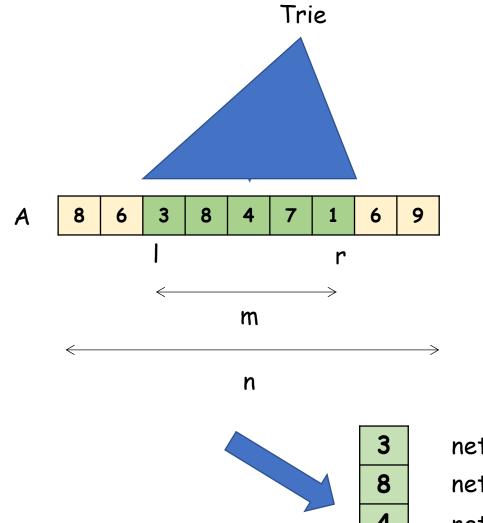
search

netbank netflix netbeans netball

Task: Given an array A of n numbers, and a range [I, r] of size m, find the positions of the K largest numbers in A[I, r]



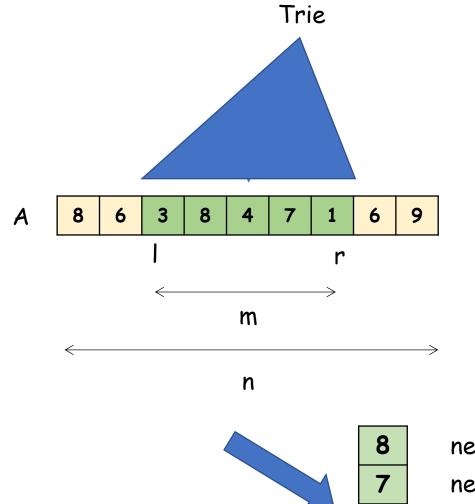
Algoritma Sederhana:



Algoritma Sederhana:

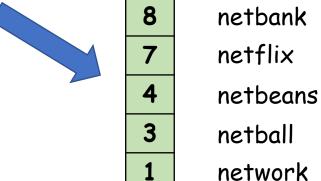
1. Copy A[I, r] ke array baru B dalam waktu O(m)

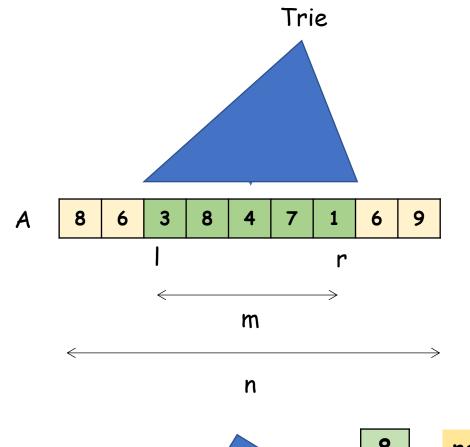
netball
netbank
netbeans
netflix
network



Algoritma Sederhana:

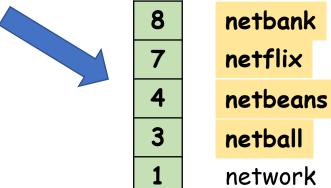
- 1. Copy A[I, r] ke array baru B dalam waktu O(m)
- 2. Sort B dalam waktu O(m log m)





Algoritma Sederhana:

- 1. Copy A[I, r] ke array baru B dalam waktu O(m)
- 2. Sort B dalam waktu O(m log m)
- 3. Kembalikan posisi K angka terbesar di A[I, r]

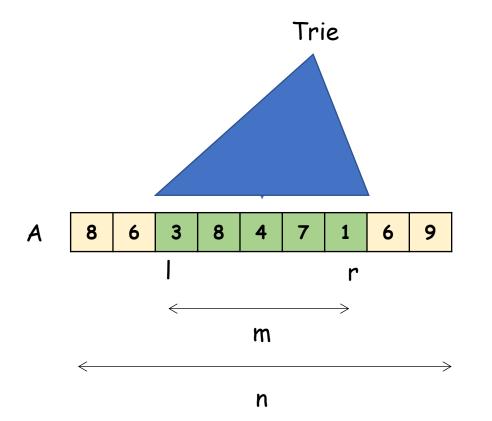


Masalah

- m bisa besar!
- Perlu O(m) space

Range Maximum Queries - Index

Precompute matrix berukuran $\mathbf{n} \times \mathbf{n}$, dimana $\mathbf{M}[\mathbf{i}, \mathbf{j}]$ berisi posisi dimana elemen terbesar pada lokasi array $\mathbf{A}[\mathbf{i}, \mathbf{j}]$ berada.



	0	1	2	3	4	5	6	7	8
0	0	0	0	0	0	0	0	0	8
1		1	1	3	3	3	3	3	8
2			2	3	3	3	3	3	8
3				3	3	3	3	3	8
4					4	5	5	5	8
5						5	5	5	8
6							6	7	8
7								7	8
8									8

Range Maximum Queries - Index

Dengan index, pencarian nilai maksimum pada range A[i,j] menjadi O(1).

Jadi, sekarang bagaimana mencari K nilai terbesar?

- Cari posisi nilai terbesar p pada A[i,j]
- Proses ke A[i,p-1] dan A[p+1,j]
- Lakukan terus hingga mendapatkan K elemen terbesar

O(K log m)