### Scheduling

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#### Scheduling Criteria

Have various, often conflicting criteria to measure success of scheduling:

- CPU utilisation
- Throughput: Number of processes completed within a given
- Turnaround time: Time it takes for each process to be executed
- Waiting time: Amount of time spent in the ready-queue
- Response time: time between submission of request and production of first response

#### Scheduling

Problem: Which process ready to execute commands gets the CPU?

key function of the operating system

Prerequisites for successful scheduling:

1.) CPU-I/O-Burst Cycle

Experience shows: I/O occurs after fixed amount of time in  $\geq 90\%$  $\Rightarrow$  appropriate time for re-scheduling

2.) Preemptive Scheduling: Processes can be forced to relinquish processor

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#### Scheduling algorithms

1.) First-Come, First-Served (FCFS)

Jobs are put in a queue, and served according to arrival time

Easy to implement but CPU-intensive processes can cause long waiting time.

FCFS with preemption is called Round-Robin standard method in time sharing systems

Problem: get the time quantum (time before preemption) right.

- too short: too many context switches
- too long: Process can monopolise CPU

#### Shortest Job First

Next job is one with shortest CPU-burst time (shortest CPU-time before next I/O-operation)

Not implementable, but this is algorithm with the smallest average waiting time

⇒ Strategy against which to measure other ones Approximation: Can we predict the burst-time? Only hope is extrapolation from previous behaviour done by weighting recent times more than older ones.

$$\tau_{n+1} = \alpha t_n + (1 - \alpha) \tau_n$$

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#### Multilevel Queue Scheduling

Applicable when processes can be partitioned into groups (eg interactive and batch processes):

Split ready-queue into several separate queues, with separate scheduling algorithm

Scheduling between queues usually implemented as pre-emptive priority scheduling

#### Possible setup of queues:

- System processes
- Interactive processes
- Interactive editing processes
- Batch processes

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#### **Priority Scheduling**

Assumption: A priority is associated with each process CPU is allocated to process with highest priority Equal-priority processes scheduled according to FCFS

#### Two variations:

- With preemption: newly-arrived process with higher priority may gain processor immediately if process with lower priority is running
- Without preemption: newly arrived process always waits Preemption good for ensuring quick response time for high-priority processes

Disadvantage: Starvation of low-priority processes possible Solution: Increase priority of processes after a while (Ageing)

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Other w	ay of organis	sing queues	according	to length o	of CPU-burst
Burst time	e				
					1
Burst time	e				
Burst time  4ms	2				
		1			1
FCFS					

#### Scheduling for Multiprocessor Systems

CPU scheduling more complex when multiple CPU's are available

Most common case: Symmetric multiprocessing (SMP):

- all processors are identical, can be scheduled independently
- have separate ready-queue for each processor (Linux), or shared ready-queue

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## Load Balancing

Idea: use all CPU's equally (goes against processor affinity)

- Push migration: periodically check load, and push processes to less loaded CPU's
- Pull migration: idle CPU's pull processes from busy CPU's

#### **Processor Affinity**

Process affinity for CPU on which it is currently running

- Soft Affinity current CPU only preferred when re-scheduled
- Hard Affinity Process may be bound to specific CPU Advantage: caches remain valid, avoiding time-consuming cache invalidation and recovery

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#### Linux Implementation

Several schedulers may co-exist, assign fixed percentage of CPU-time to each scheduler

Important schedulers:

- Round-robin scheduler with priorities (the default scheduler)
- Real-time scheduler (process needs to be assigned explicitly to this one) (typically FIFO)

# Round-Robin Scheduler with priorities

implemented in an interesting way: maintain tree of processed ordered by runtime allocated so far pick next process as one with least runtime allocated so far insert new process in ready queue at appropriate place in tree Priorities handled by giving weights to run-times.

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