Security

At least three different categories:

- *Data integrity*: Backup etc., taken care by normal OS functions
- *Protection against user errors*: done by separating users and processes
- *Protection against malicious users*: more complicated, involves trade-offs between ease of use and level of protection

Major problem: *Identification of users*
Most common scheme: Passwords
+ : Easy to use and understand
- : All too often too easily guessable
- : Exposure problem
- : Where to store passwords

Other problems arise from vicious programs.
Most famous example so far : *Internet Worm* (1988)
Very clever use of weaknesses in OS design/program bugs

Protection mechanisms

Access to shared resources must be controlled

Two aims:
- Protection against users' mistakes
- Increase in reliability

Principles:

- separate policy (*what*) from mechanism (*how*)
  Important for flexibility
- Users should have as much privilege as necessary to get job done

Standard way: each user separate domain of protection
Need *trusted way* of making system privileges available
Disadvantage: primary target for breakins (`setuid`-programs in UNIX)

Access matrix

abstract view of protections
have row for each domain and column for object
Entry indicates access right
Example:

	F_1	F_2	F_3	Printer
D_1	read		write	
D_2				print

Have capabilities like read, write, copy, owner, control

Unsuitable for implementation because matrix far too large

Implementation Issues

- Access lists for objects (store columns)
- Capability lists for domains (store rows)

Capabilities allow great flexibility
Example: Hydra

- Auxiliary rights: each process can pass access to procedure to other processes
  $\Rightarrow$ dynamic change of access rights
- Rights amplifications: procedure can act on specified type on behalf of any process which is allowed to execute it
  Rights *cannot* be passed on
  $\Rightarrow$ flexible and more secure mechanism for granting higher privileges temporarily

coarser than access lists

for each file, have three categories of possible users

- owner
- group (pre-defined set of users)
- all others

owner can grant permission to

- read
- write
- execute program/find files in directory

Aim: authentication in insecure networks

Two possible solutions:
- Central authentication server (Kerberos):

Have following protocol:

- Client → Auth Server: Credentials, please?
- Auth Server → Client: ((Client Id, Session Key)$_S$, Session Key)$_C$
- Client decrypts message and keeps Session Key
- Client → Server: (Client Id, Session Key)$_S$
- Server decrypts message and obtains Client Id and Session Key

Authentication Server must be trusted
Protocol vulnerable to replay attacks

- Newer Alternative: PGP (Pretty Good Privacy)

Relies on so-called Public Key Cryptography:

Have two different keys, *public key* K_p and *secret key* K_s
clear text encrypted via K_p can be decrypted only via K_s and vice versa

Assume every host in network has public and secret key

Now communication can take place by sending message from A to B as $(m_{A_s})_{B_p}$. Only B can decode the message, and it can verify that it came from A.

Certification of keys still necessary.