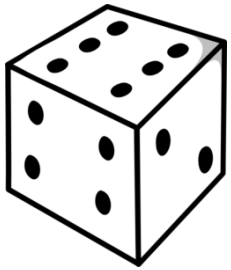


DICE - Deception of InterCommunication to Enemies

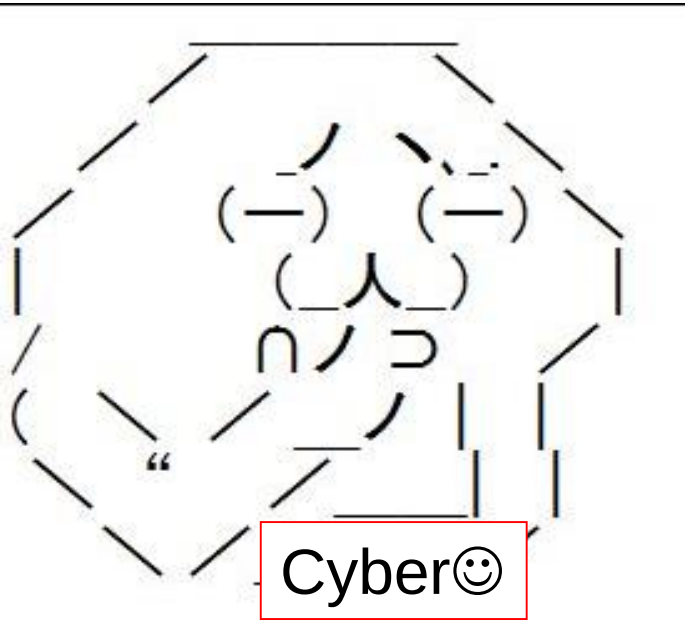
wakatono

wakatono@gmail.com

Twitter: @wakatono

Facebook: www.facebook.com/wakatono

Other side:



It’s a Joke ☺

Today's Presentation

(Almost) All you need is…

- TCP/IP basic(knowledge of IP, UDP, TCP headers)
- DNS basic(knowledge of query – response)
- Speed of Light (300,000km/sec)

And some knowledge to develop, deploy

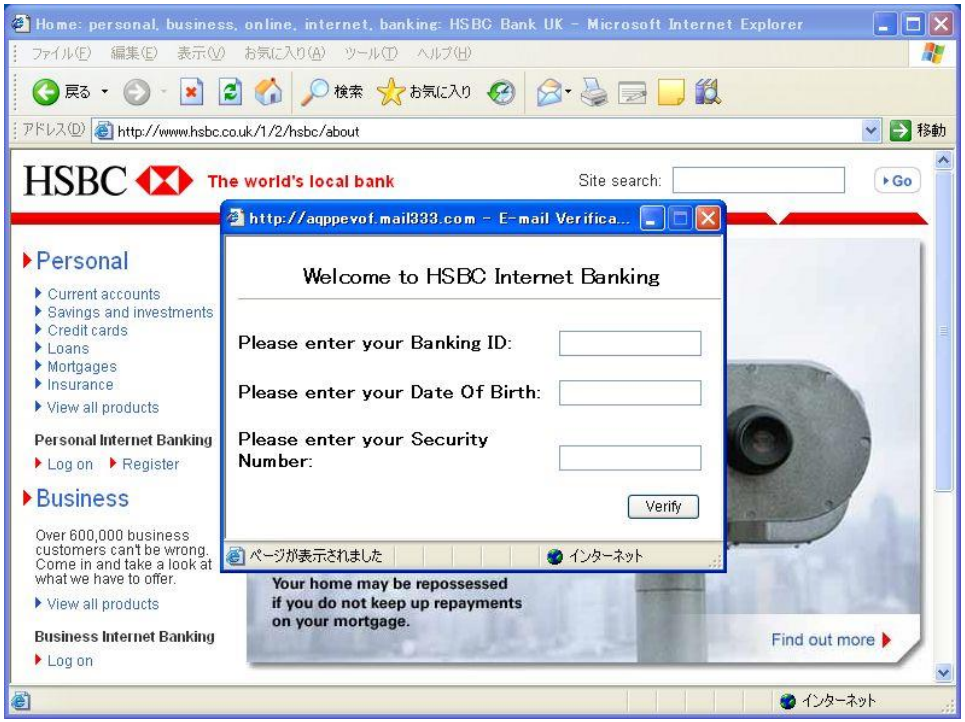
Table of Content

- Various of attacks on internet
- Popular solutions and issues
- Focus to DNS query and response
- Decept the Enemies! – DICE architecture, keys of implementation, and deployment
- Protocol Deception – Joyful and Useful Decepting Technology
- DICE Applications, Issues, and future

Various of attacks on internet

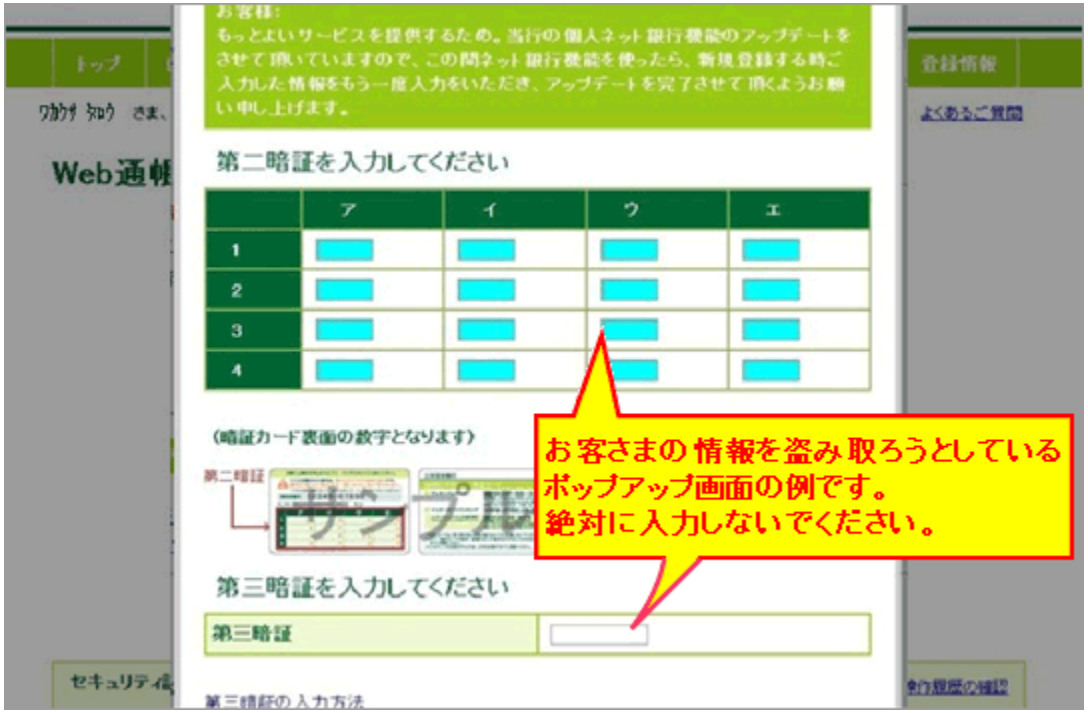
Assumption: Attackers use Malicious FQDNs, Domains, and FastFlus combination

Phishing makes malicious sites like a real service sites.



Reference:
http://www.atmarkit.co.jp/fsecurity/special/65phishing/phishing01.html

MITB viruses inject malicious forms to real contents

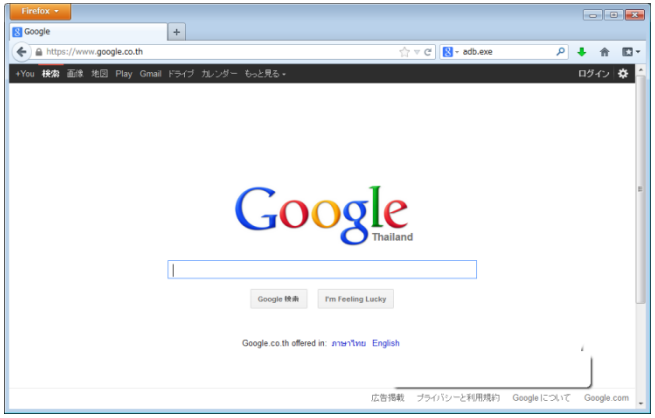


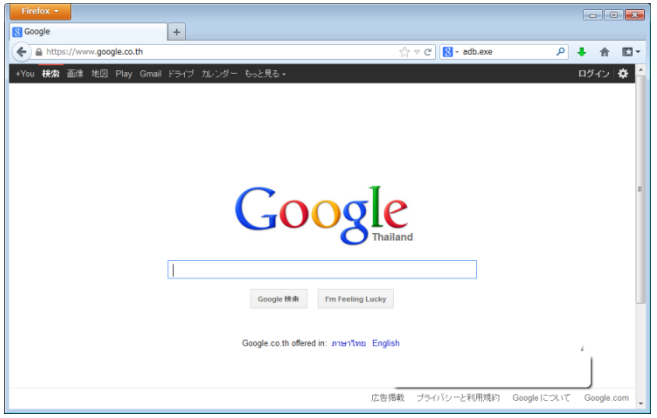
Reference: http://www.smbc.co.jp/security/popup.html

Common Spec:malicious hosts exist

- Attacker prepares the host to receive data of victims' like banking information.
- Most of malicious host has own FQDN
- IP addresses is changed due to their lifecycle
  - Stopping access to malicious hosts that have fixed IP addresses is easy due to many technology to take down.

Modern Attacks triggered by Web Access

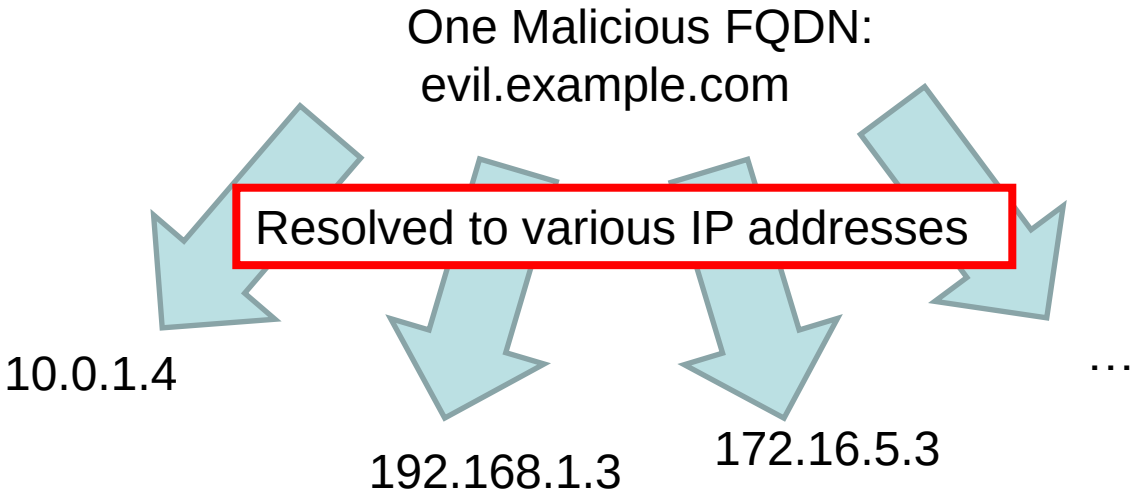




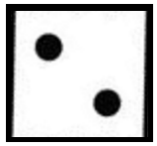
Which is better, left “Google” or right “Google” ?

Both sites are better(correct) web site ☺

FastFlux – common tech for attackers

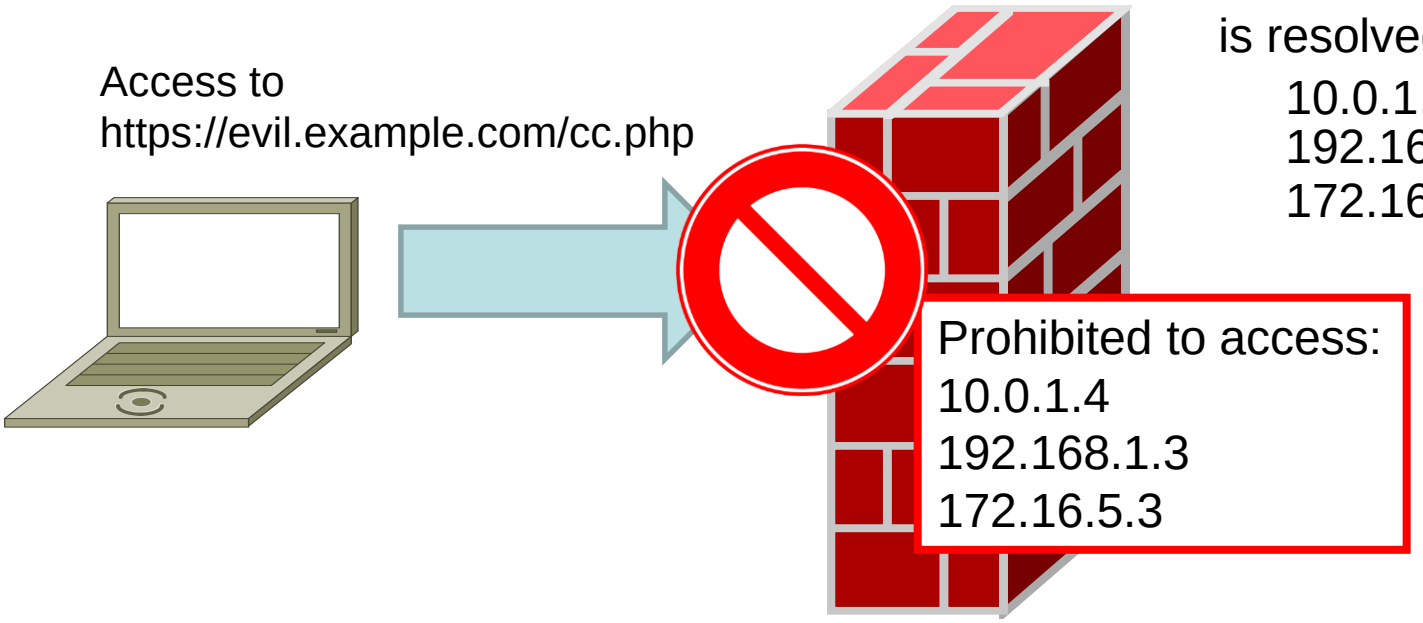


- One malicious FQDN has multiple IP addresses ☹

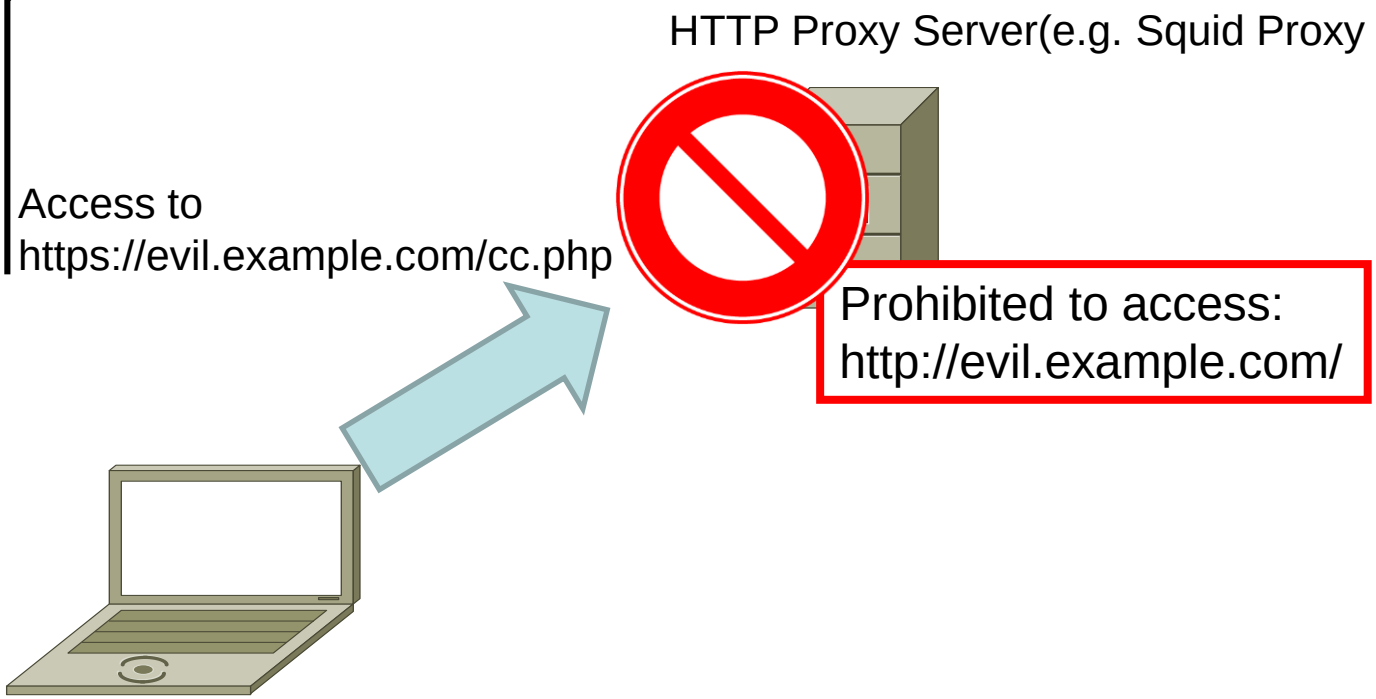


Popular solutions and issues

Firewall


Access to
https://evil.example.com/cc.php
evil.example.com
is resolved to:
10.0.1.4
192.168.1.3
172.16.5.3
Prohibited to access:
10.0.1.4
192.168.1.3
172.16.5.3

URL Filtering by HTTP(S) Proxy



IDS / IPS

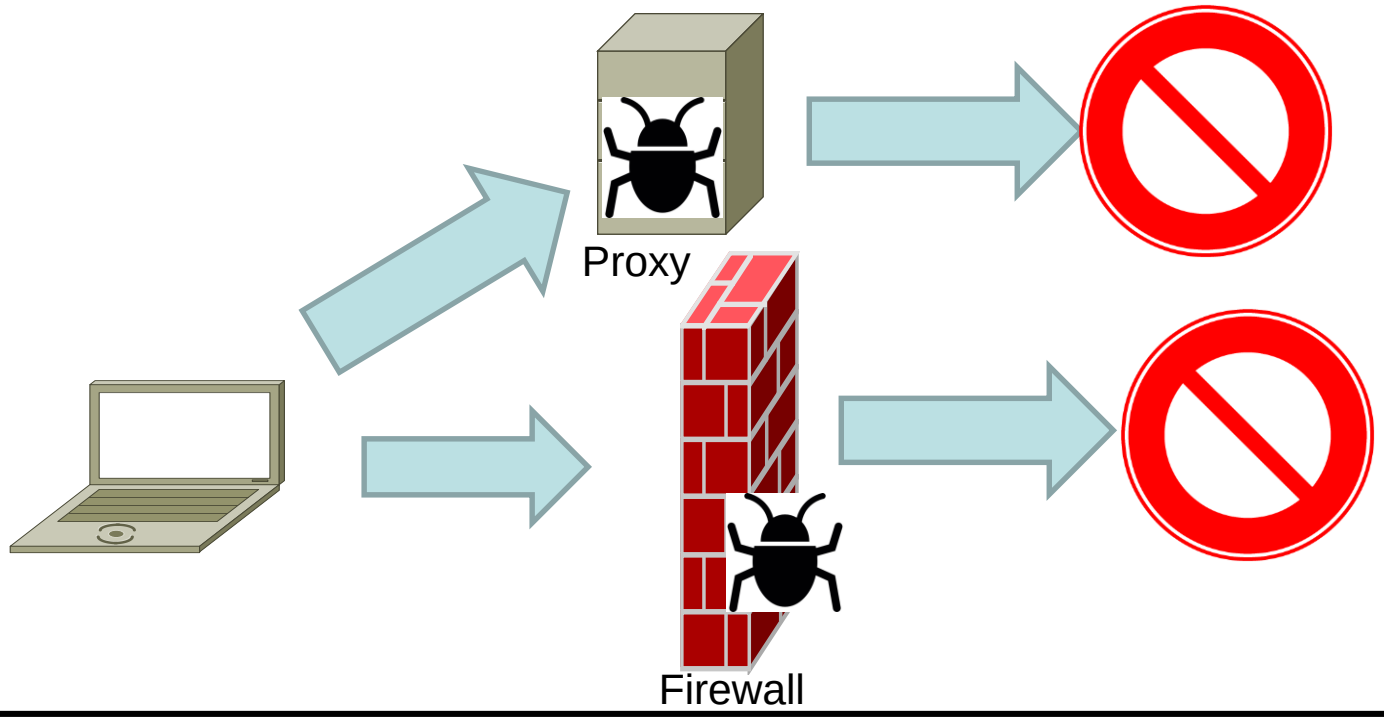
- (snip)

RPZ(Response Policy Zone)

- RPZ is the DNS blocking technology
- Implemented in newer BIND releases

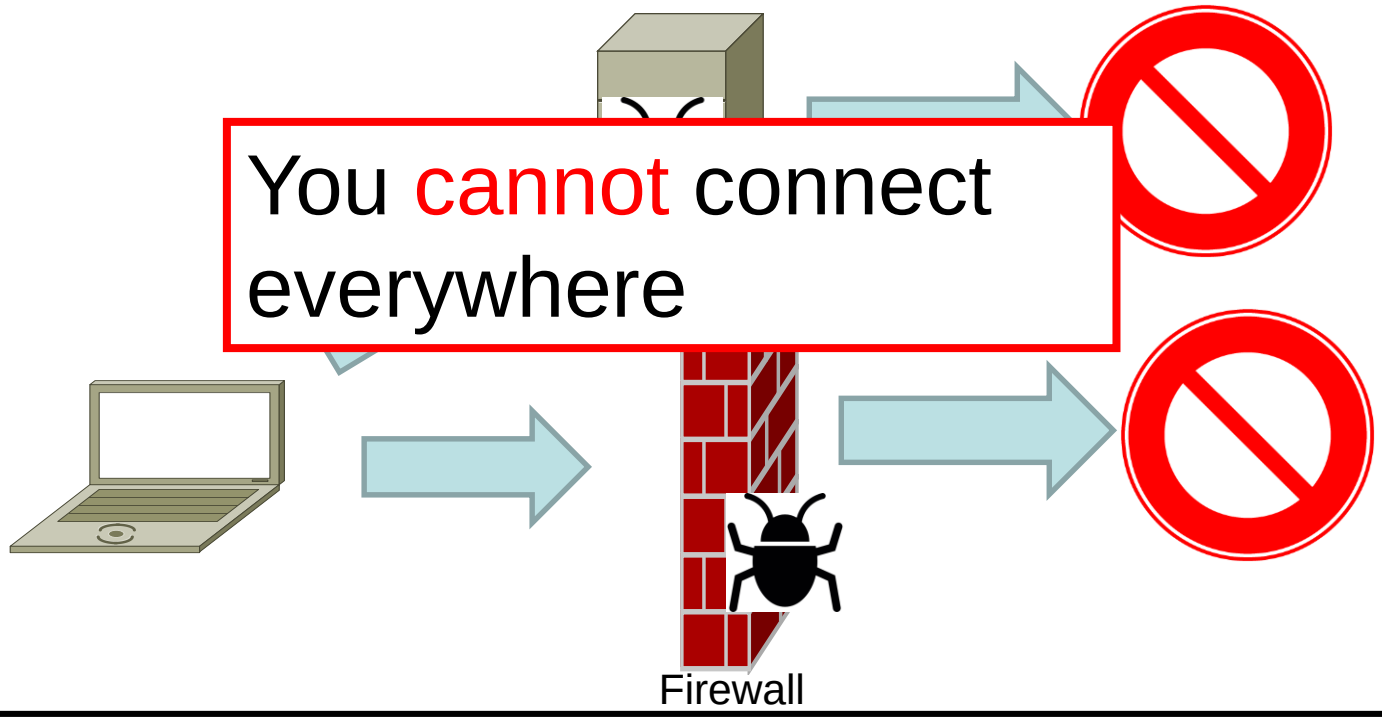
Generic Issues of Solutions

- Performance Bottleneck
- Reliability / Availability



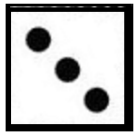
Generic Issues of Solutions

- Performance Bottleneck
- Reliability / Availability



How to avoid accessing to malicious host?

- HTTP/HTTPS Proxy Server access block by using Blacklist and/or Firewall
  - Load of Proxy (Servers & Operators) and/or Firewall(Servers & Operators) may be High!
- Takedown by ISP and Various Service Provider
  - Sometimes Long Term discussion is needed
- Temporarily:
  - Stop by using DNS response deception(example)
    - I assume this to use edge network(response from nearest DNS Cache Server is decepted)



Focus to DNS query and response

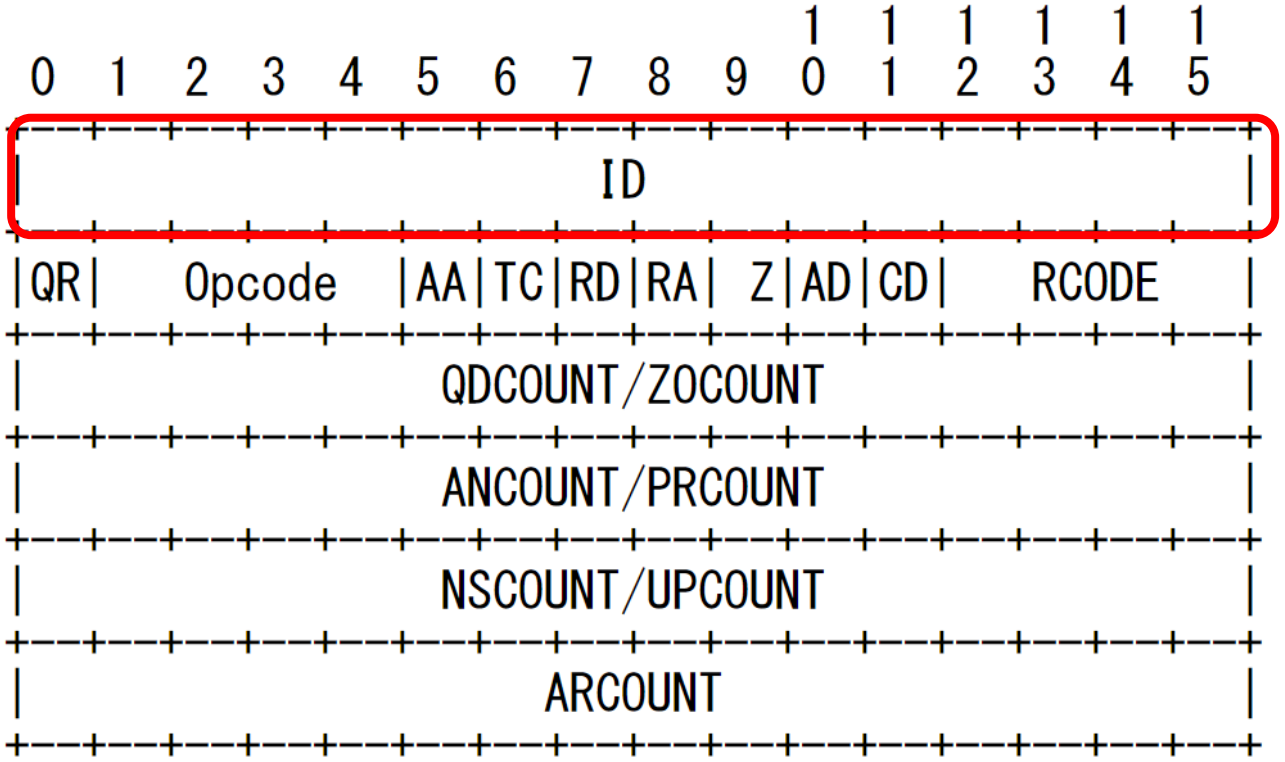
Assumption

- We have technology to identify malicious FQDN
  - By using SIEM, Domain Blacklist, URL Blacklist, etc…
- Network Operator and Security Operator are independent

DNS Response
- easily Deceptable Protocol Response

- Normal DNS Response can be decepted easily
  - Signed DNS Response (e.g. DNSSEC specification) is hard to be decepted.
  - TCP DNS Interaction is little a bit hard to be decepted.
- Applicable to various of deception
  - Various RR Type of Response(e.g. NS, MX) can be used for applying.

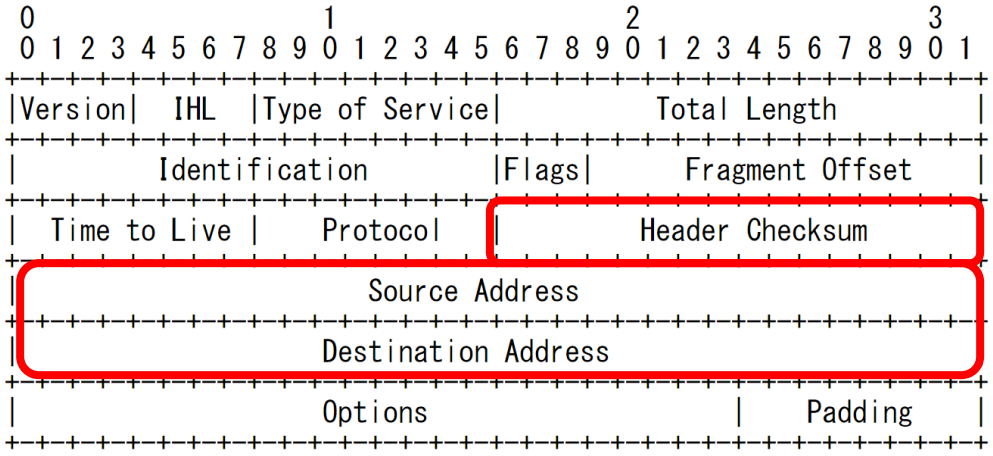
DNS Query/Response Header Format



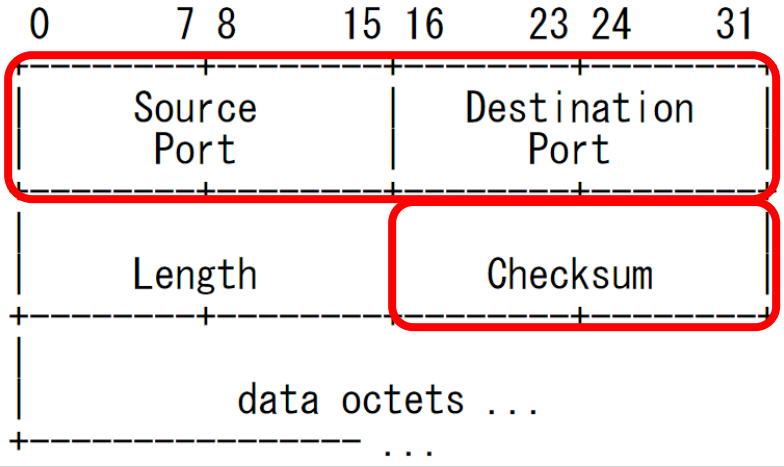
The ID field identifies the query and is echoed in the response so they can be matched.

Reference: RFC6895 Domain Name System (DNS) IANA Considerations

UDP and IP Header Format

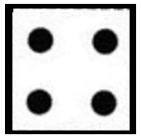


References:
RFC768 User Datagram Protocol
RFC791 INTERNET PROTOCOL



DNS Response Deception Summary

- Response Packet is easy to decept
  - Know pair of IP addresses(IP), Port numbers(UDP), (Transaction) ID(DNS)
  - Send the decepted response packet to client(s) faster than “true” DNS response is sent
  - Client(s) will access host(s) according to the decepted response



Decept the Enemies! – DICE architecture, keys of implementation, and deployment

DICE architecture is simple, but implementation and deployment is difficult little a bit

What is DICE?

- Abbrev of:
  Deception of InterCommunication to Enemies

Gartner Security & Risk Management Summit 2016

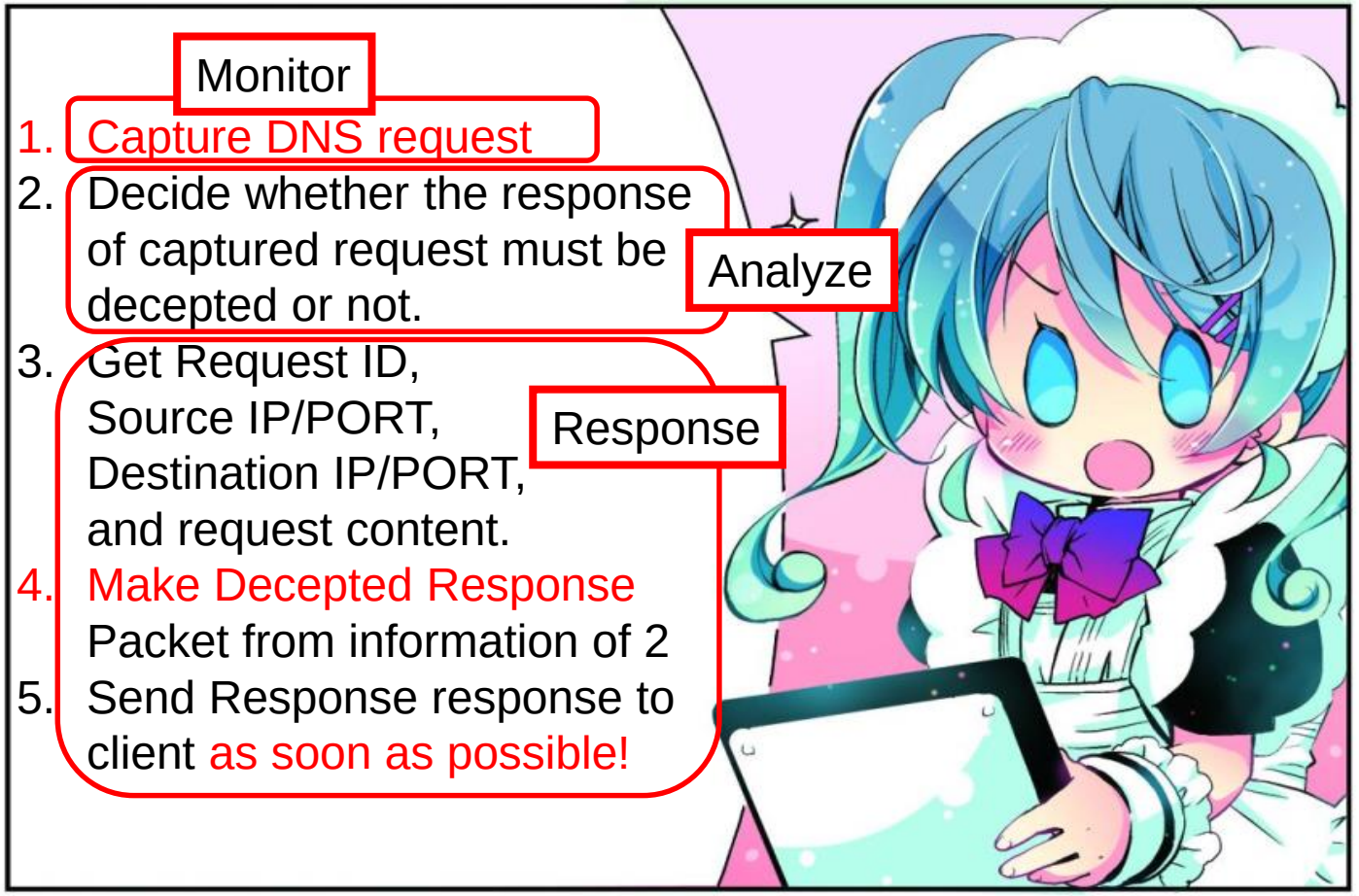
To Decept Protocol(s)

- Understand Protocol specification(s)
  - Connection sequence and requirement(TCP)
  - Communication sequence and fields of payloads
  - Transaction Relation Information(e,g,ID on DNS)
- Understand Protocol implementation(s)
  - TCP stack implementation UDP stack implementation, BIND on DNS, etc…
- Understand Target Application(s)
  - (snip)

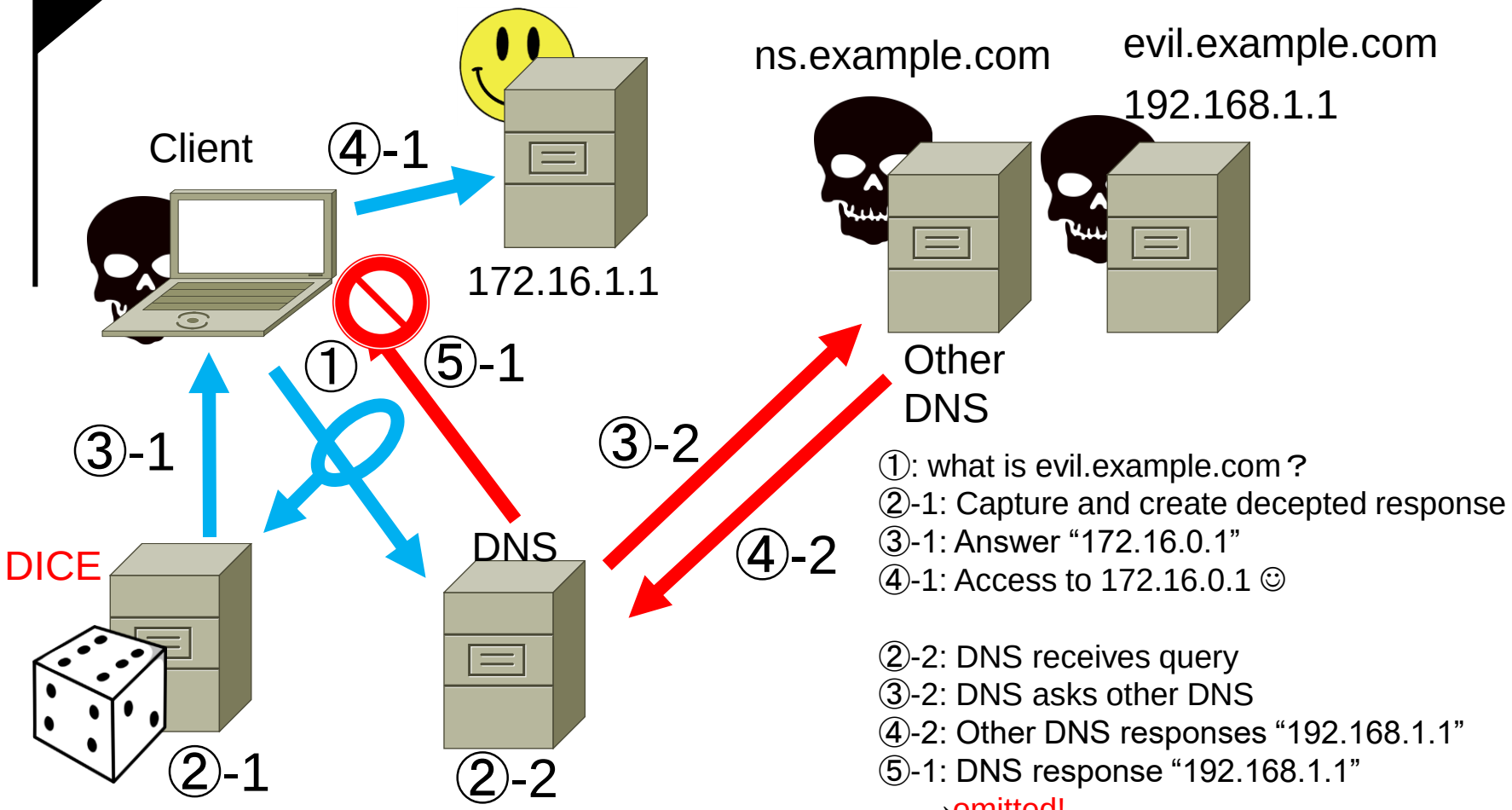
DICE architecture(simple!)

- Monitor subsystem
  - High-speed and lossless monitoring
  - Like tcpdump, wireshark, etc…(but simple)
- Analyze subsystem
  - Decede to Response
- Response Subsystem
  - High-speed and synchronous processing
  - Works like hping(but simple)

Applied DICE architecture to DNS client deception



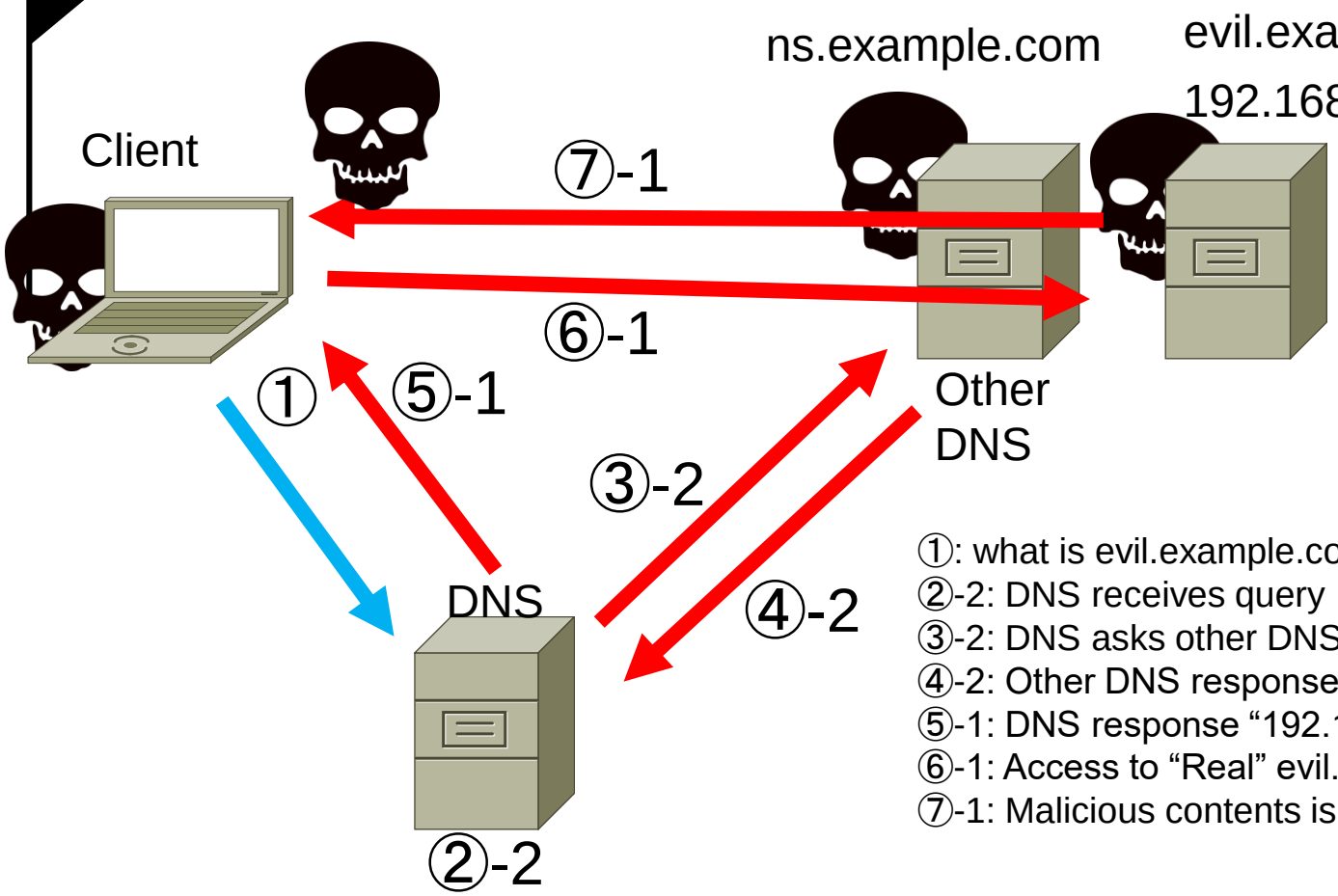
Concept Diagram



①: what is evil.example.com？
②-1: Capture and create decepted response
③-1: Answer “172.16.0.1”
④-1: Access to 172.16.0.1 ☺

②-2: DNS receives query
③-2: DNS asks other DNS
④-2: Other DNS responses “192.168.1.1”
⑤-1: DNS response “192.168.1.1”
→omitted!

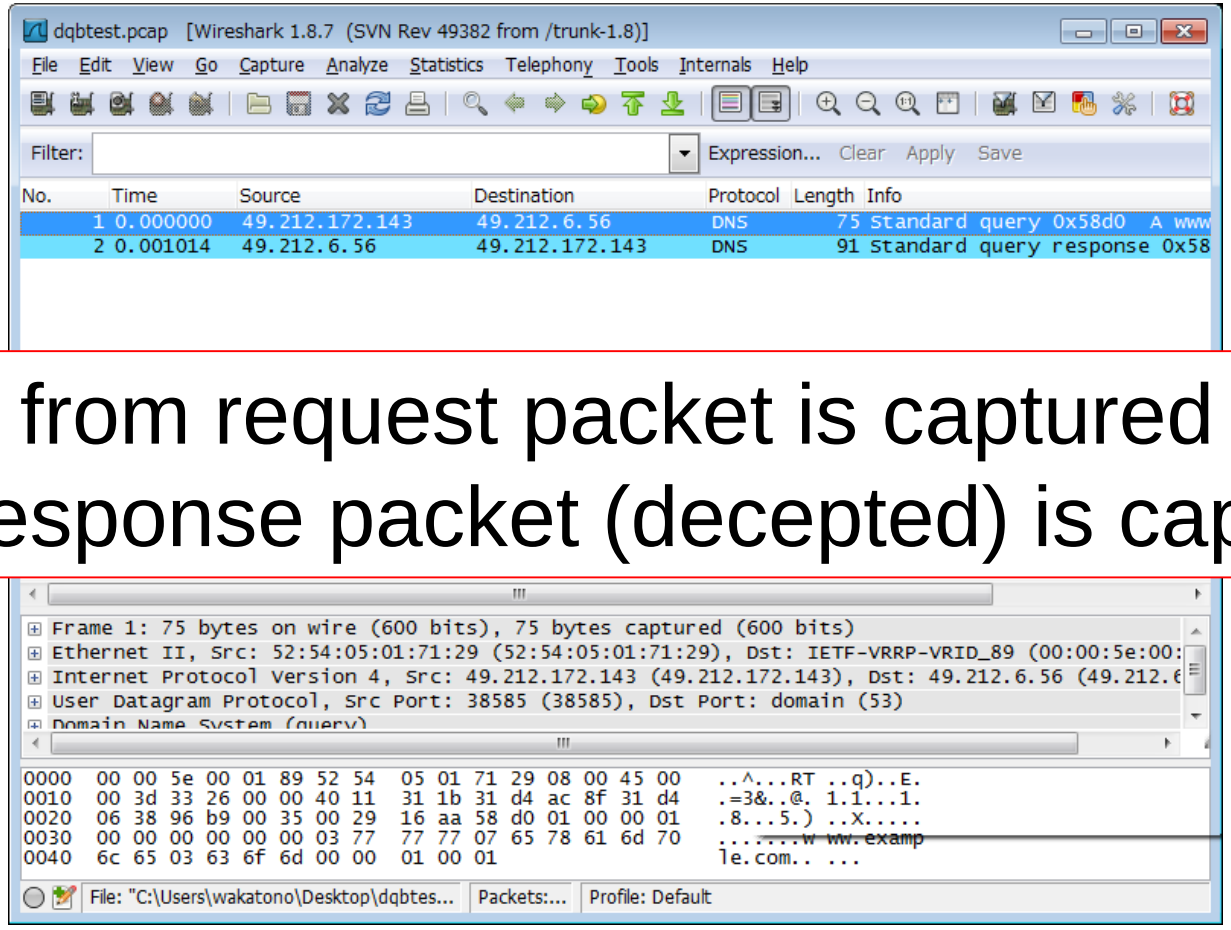
If there is not DICE…



①: what is evil.example.com？
②-2: DNS receives query
③-2: DNS asks other DNS
④-2: Other DNS responses “192.168.1.1”
⑤-1: DNS response “192.168.1.1”
⑥-1: Access to “Real” evil.example.com
⑦-1: Malicious contents is downloaded by client☹

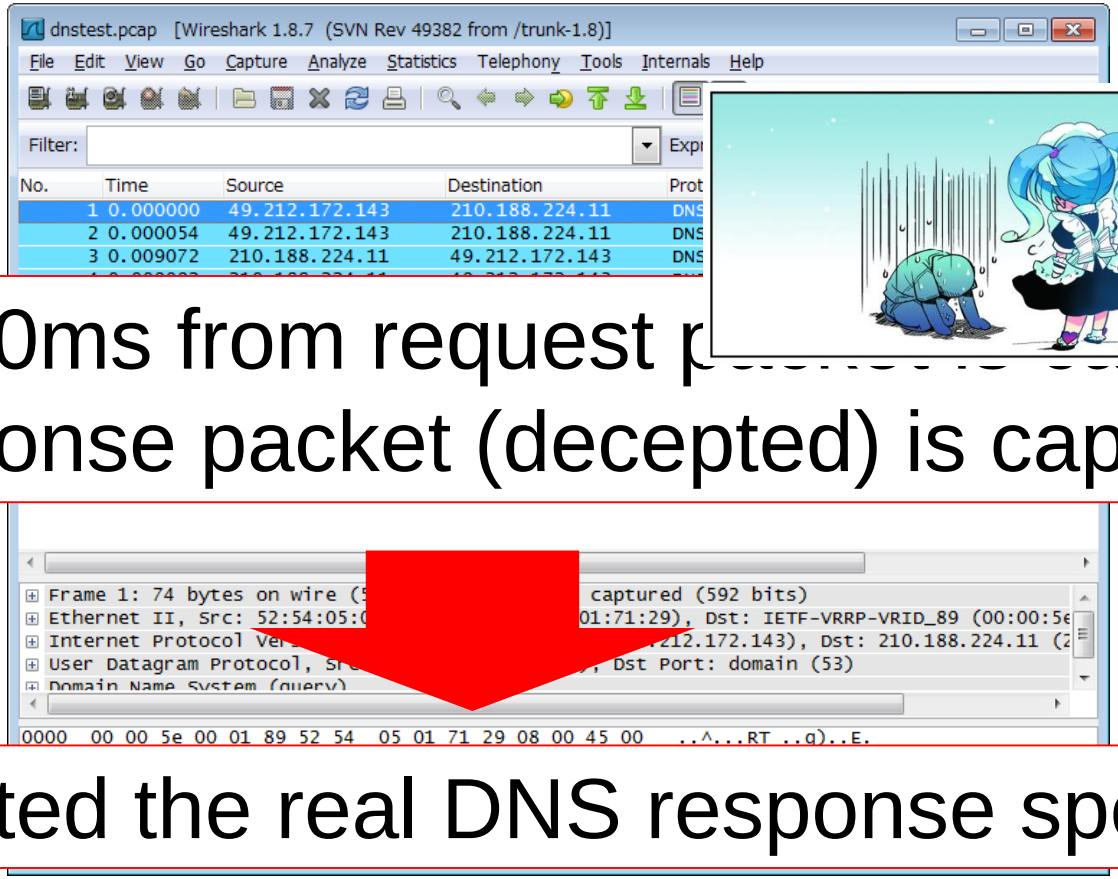
Proof of Concept:

About 1ms



1ms from request packet is captured
to response packet (decepted) is captured

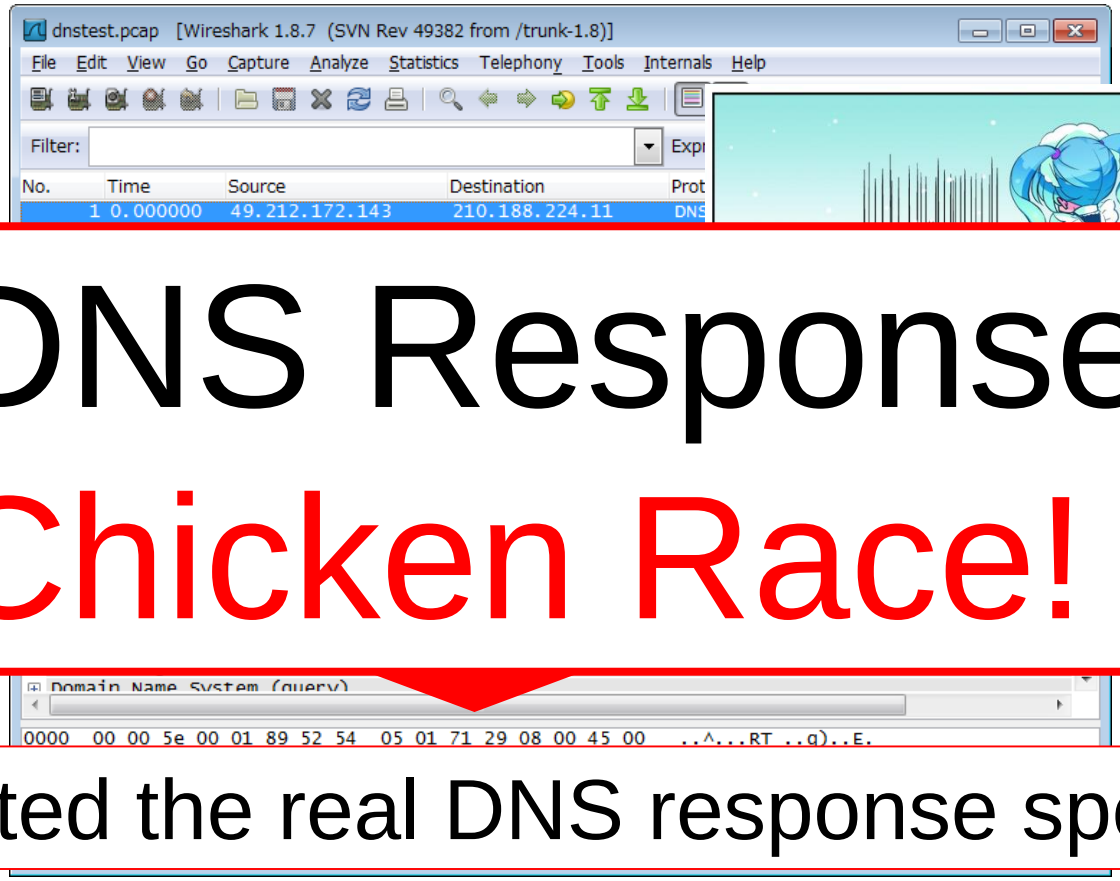
Normal Request/Response



about 10ms from request p to response packet (decepted) is captured

I defeated the real DNS response speed ☺

Normal Request/Response



about

to res

ured

DNS Response

Chicken Race!

I defeated the real DNS response speed ☺

Name Resolution Step Summary Decepted by DICE

- 1. DNS Request is sent by client
- 2. Decepted DNS Response is sent to client by DICE
- 3. Real DNS Response is sent by DNS Cache

Name Resolution Step Summary Decepted by DICE

- 1. DNS Request is sent by client
- 2. Decepted DNS Response is sent to

One Request for Two Response!

- 3. Real DNS Re is sent by DNS Cache

Omitted

Real DNS response too slow?

- Slow (most case of oversea)
  - e.g. Tokyo – San Francisco
    - Round trip: 18,000km
  - Speed of Light: 300,000km/s
  - At least, about 60msec is required as a time between IP packet round trip - Traffic initiated from Tokyo to San Francisco and Response sent from San Francisco to Tokyo reaches to traffic initiator(in Tokyo)

To develop like DICE

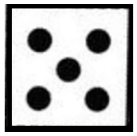
- Use Linux Socket API（～GbE）
  - In fact, GbE wirespeed capture is too hard for (old) Socket API(!)
- Use Intel DPDK（～10GbE）
  - DPDK is used for various development of NW appliance
    - e.g. Lagopas (SDN Switch implementation by NTT and IIJ)

DICE Deployment principle

- 1st: Understand your network topology
- 2nd: Identify DNS server(or Server to Decept) and Gateway to internet
- 3rd: Place the appropriate DICE unit

DICE Practical Deployment in DNS

- 1st: Understand your network topology
- 2nd: Identify Top DNS “Cache” server
  - And NW Switch that DNS “Cache” Server is connected
- 3rd: Place the DICE unit near the Top DNS “Cache server”
  - And change configuration of the NW switch above.

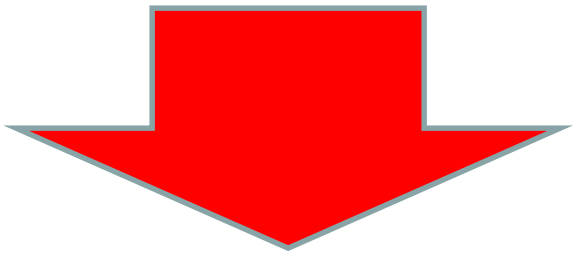


Protocol Deception – Joyful and Useful Decepting Technology

Plain Text Communications without sign can be decepted easily ☺

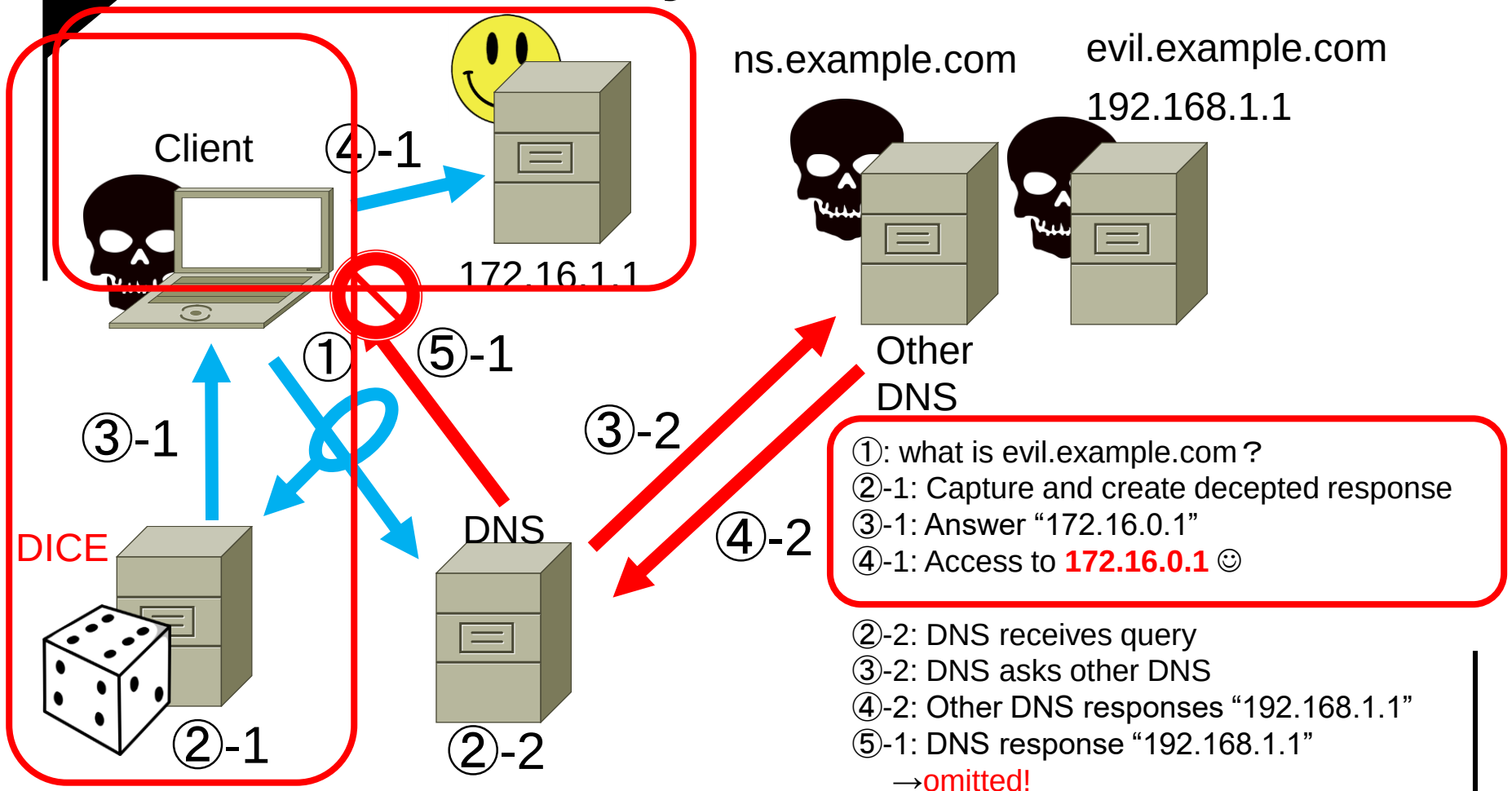
If DNS Response is decepted successfully…

- Connection can be intercepted as you like



- Malicious request parameter(s) can be obtained easily and effectively

Connection can be intercepted as you like



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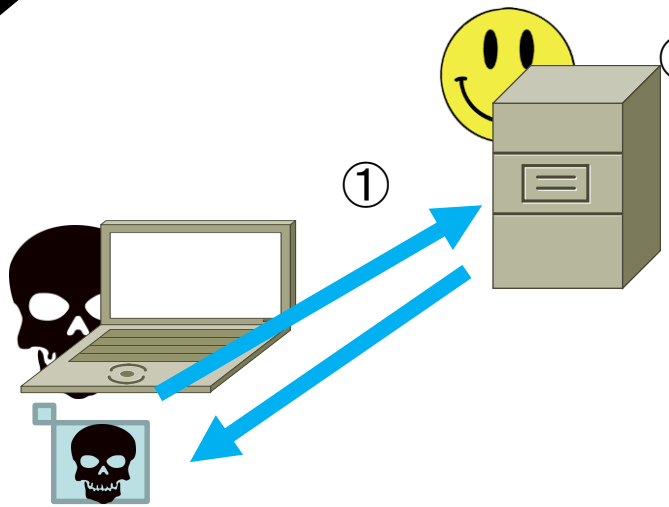
⑤-1: DNS response “192.168.1.1”
→omitted!

Malicious request parameter(s) can be obtained easily(1/2)



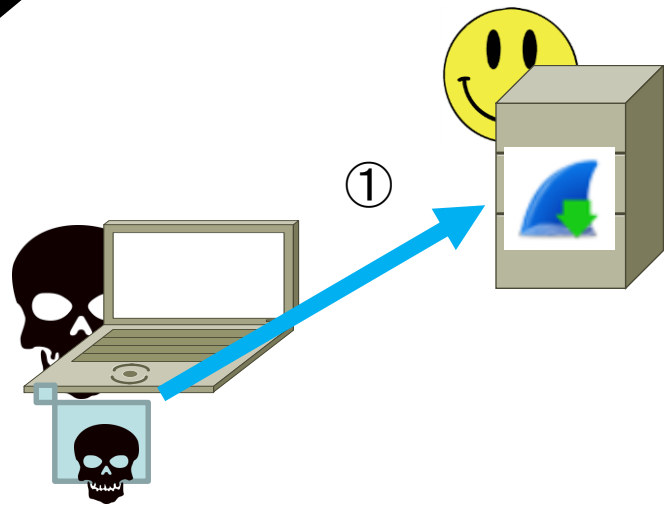
⓪ AliasMatch ^/ //var/www/html/fault.html

② logged parameter(s) is recorded in the leaded host
e.g. 172.16.0.2 - - [22/Jul/2016:13:21:24 +0900] "GET /cc.php?init=succeed&target=inhouse&..." 200 36 "-" "-"

① send request to “evil.example.com”
e.g. Access to URL
http://evil.example.com/cc.php?init=succeed&target=inhouse&...

② logged parameter(s) is recorded in the leaded host
172.16.0.2 - - [22/Jul/2016:13:21:24 +0900] "GET /cc.php?init=succeed&target=inhouse&..." 200 36 "-" "-"

Malicious request parameter(s) can be obtained easily(2/2)



⓪ AliasMatch ^/ //var/www/html/fault.html (Apache2 configuration)

② Packet data including POST data is available

① send request to “evil.example.com” by POST method
e.g. Access to URL by POST method http://evil.example.com/cc.php
POST init=succeed, target=inhouse&...

Normally, these info are not recorded in Web Server log

②detail: Packet data including POST parameter data is available
packet capture works always but capture really works only when malicious communication is identified effectively

DICE Applications, Issues, and future

In fact, DICE is the grand design to interfare various of communications between victims and attackers

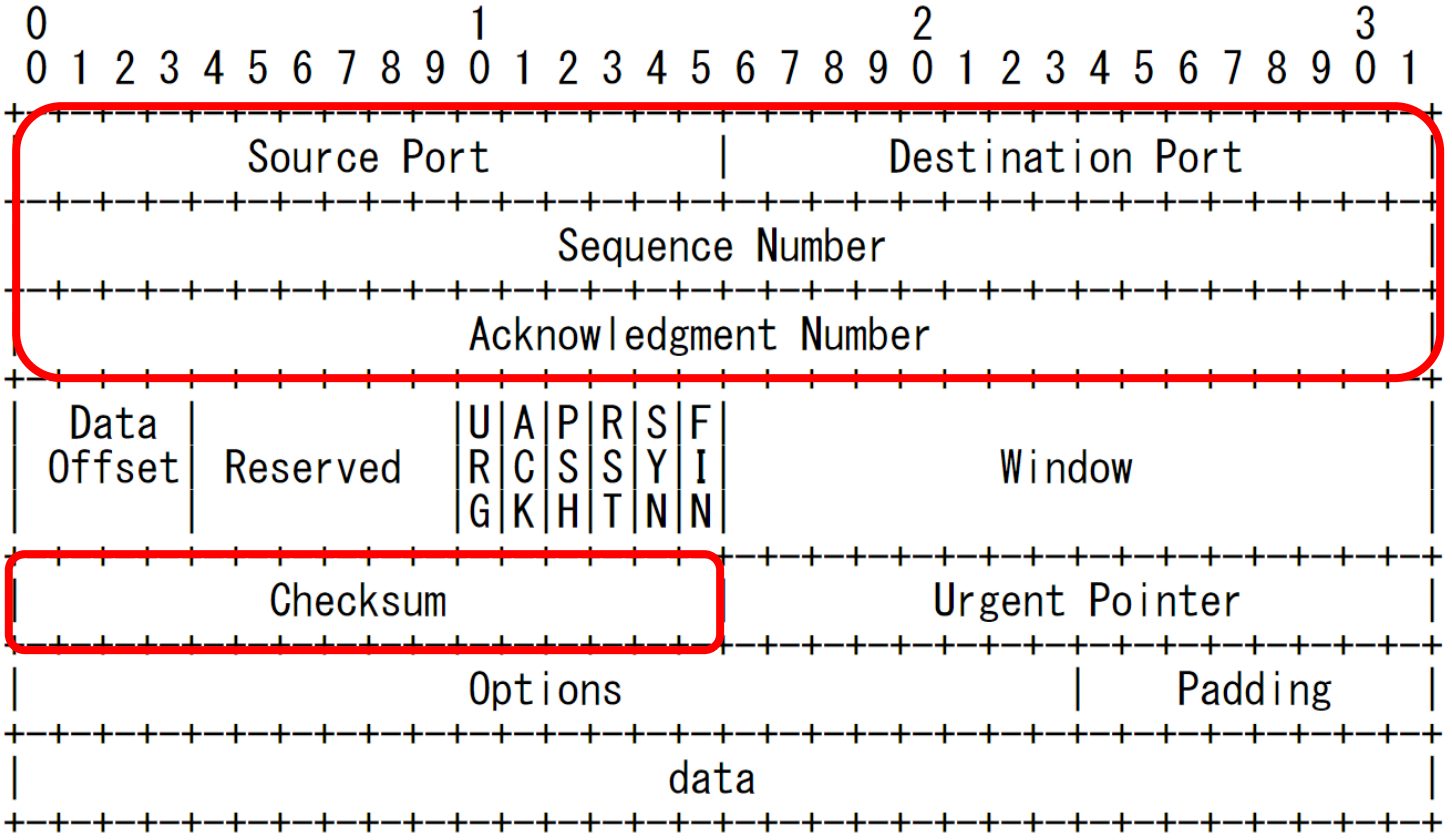
Applications:
1st step→DNS deception,
2nd step→TCP deception

TCP deception is easy to implement(except on performance)

TCP 3way handshake Response - easily Deceptable Protocol Response

- Normal TCP Response of Connection Initiation can be decepted easily
  - Signed and/or encrypted Packet (e.g. IPsec) is hard to be decepted.
- Applicable to various of deception
  - After Decepted, connection is “hijacked” ☺
  - After Connection Decepted, we can decept interaction to enemies by using Hijacked connection ☺

TCP Connection Interaction
- easily Deceptable Protocol Response



Reference: RFC793 TRANSMISSION CONTROL PROTOCOL

I want to terminate malicious connection to specific “client(s)”

- RST packet may be filtered(and connection is still alive ☹)
  - Many IPS have function of sending RST packet(and may be filtered).

- ACK, SYN+ACK packet of Connection initiation state to proper port must not be filtered ☺
  - If filtered every packet, that computer turned to useless object ☺

Protocol Condition
Easy to Decept

- Fields to be used for relation between request and response is identified from request easily
  - In case of DNS: (Transaction) ID only
- There is “No” signed field ☺
  - Other Challenge ☹
- Fulfill the General Requirement of MITM

DICE basic concept is similar to…

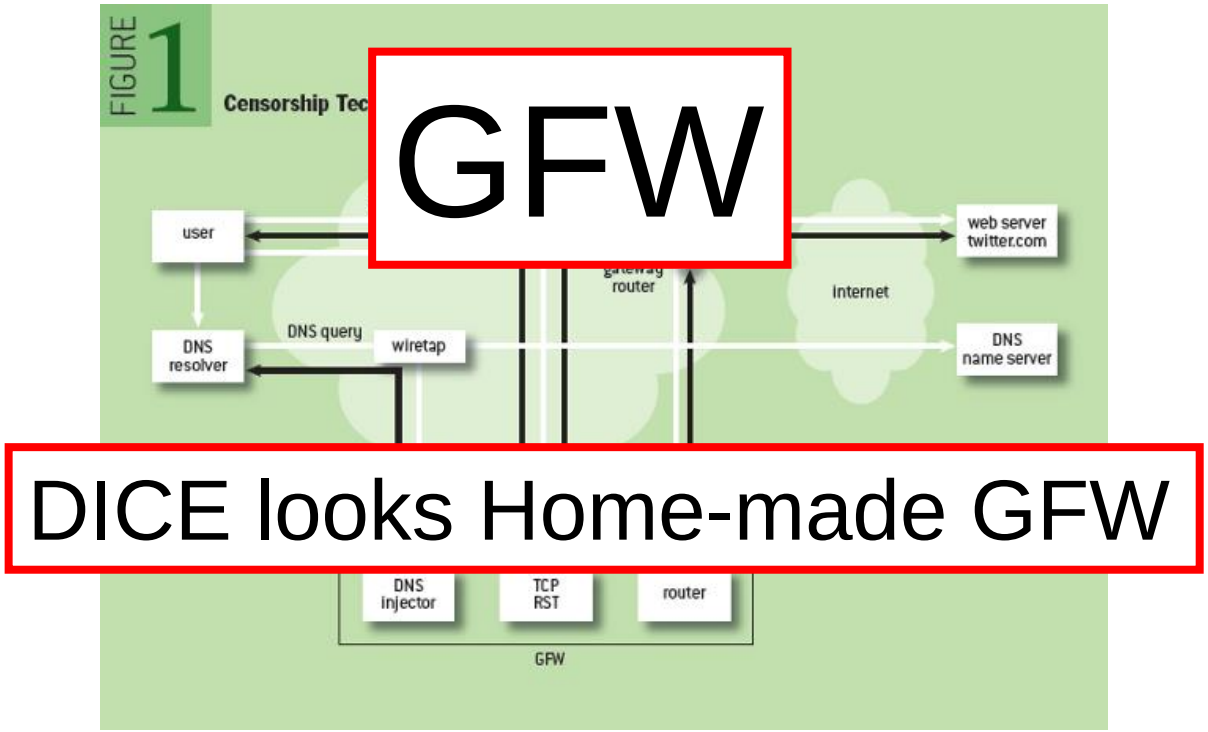


Figure From "Splinternet Behind the Great Firewall of China"
http://queue.acm.org/detail.cfm?id=2405036

Issues: Protocol Difficult to Decept

- Fields to be used for relation between request and response cannot be identified from request easily
- There is “Signed” field ☺
  - Sigh…☹
- Adaptive Payload(e.g. Result of Application Processing is included in headers and/or payloads and difficult to guess or calculate)

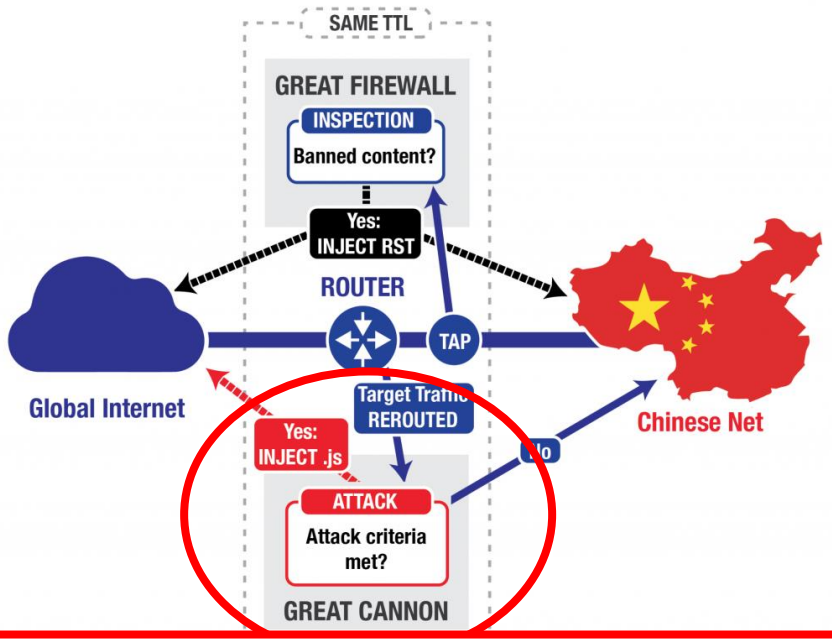
Issues: LIMITATION!

- Of course, this mechanism is not suitable for faking DNS response signed by DNSSEC mechanism.

- RFC7766 allows DNS query via TCP
  - Little a bit complex for decepting DNS response
  - RFC7766 DNS Transport over TCP - Implementation Requirements(March, 2016)
  - Fortunately, there is no implementation yet.

Issues: LIMITATION!

- Of course, this mechanism is not suitable for faking DN[illegible] DNSSEC mechanis[illegible]

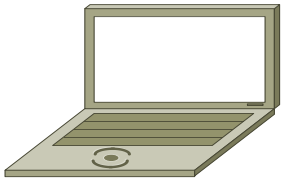


- RFC7766 [illegible] TCP
  - Little a [illegible] DNS response
  - RFC776 [illegible] -
    Impleme[illegible]h 2016)
  - Fortu[illegible] yet.

Like a GREAT CANNON…

RFC7766 Before / After

Before:

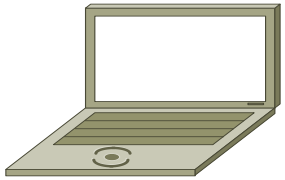


①Query by UDP

②Response by UDP with TC bit

③Query by TCP

After:



①Query by TCP

Future: Current Status of this research

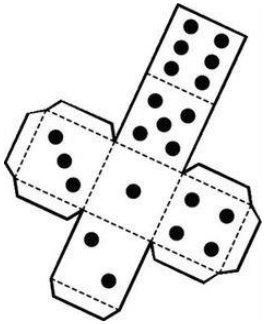
- I developed and deployed the variation of DICE
  - Malicious FQDNs find by other method(s)

- Ideas for more(for example):
  Respond by decepting response related to the request of domain name generated by DGA(Domain Generation Algorithm)

Future: DICE troubles and then stop?

- If DICE stops abnormally, no traffic can be decepted
  - These troubles don’t stop communication to outside

- DICE can be redundant by place same DICE unit to same place on network topology
  - This is effective to reduce load of each DICE unit



Conclusion

Summary of this presentation

Conclusions

- If you understand protocols, you can decept part of TCP/IP and Some Application Procotol easily
- DICE architecture is very simple(you can develop similar system(s) easily)
- Decepting Protocol makes more application to respond some kind of attacks

Thank you!

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If possible, any questions are welcome via email or Twitter.
Of course, in banquet or any networking time ☺

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