

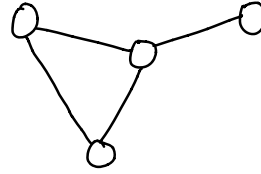
graph coloring IP formulation

Recall : input an undirected graph $G = (V, E)$

want : a color for each node s.t.

if $(u, v) \in E$, then u, v have

different colors. Want # colors as small as possible



IP formulation : decision variables :

- x_{ij} for $i \in V$, $j \in [n]$, $n = |V|$

$x_{ij} = 1$ means node i is assigned color j (max # colors = n)

- y_j for $j \in [n]$

$y_j = 1$ means color j is used

objective function :

$$\min \sum_{j=1}^n y_j \quad (\text{i.e. min \# used colors})$$

constraints :

- x_{ij}, y_j all binary

- each node is assigned one color

$$\sum_{j=1}^n x_{ij} = 1 \quad \forall i \in V$$

- if any node is assigned color j , then $y_j = 1$

$$\equiv \sum_{i \in V} x_{ij} > 0 \Rightarrow y_j = 1$$

$$\equiv n \cdot y_j \geq \sum_{i \in V} x_{ij} \quad \text{or} \quad x_{ij} \leq y_j \quad \forall i \in V, j \in [n]$$

- no two adjacent node has the same color

$$x_{uj} + x_{vj} \leq 1 \quad \forall (u, v) \in E, j \in [n]$$

$$x_{uj} + x_{vj} \leq y_j$$

$$\forall (u, v) \in E \\ j \in [n]$$

Symmetry of the LP:

if an optimal solution uses 3 colors

then there are $n \times (n-1) \times (n-2)$ "equivalent" solutions

(like color 1, 2, 5)
 \downarrow \downarrow \downarrow
 color 3, 4, 6

this slows down the solver \Rightarrow break the symmetry by adding constraint

- $y_j \geq y_{j+1}$ (always prefer colors w/ smaller indices)
- $\sum_{i \in V} x_{ij} \geq \sum_{i \in V} x_{i(j+1)}$ (always prefers to use smaller indices for colors most used)